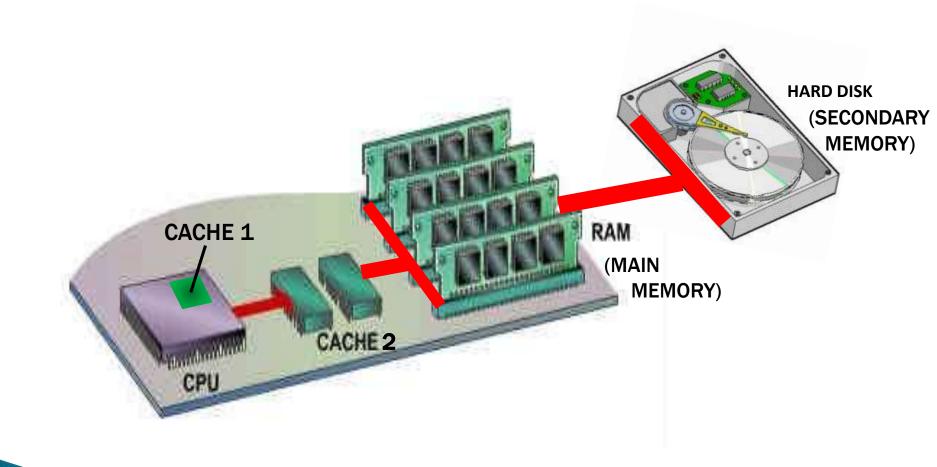
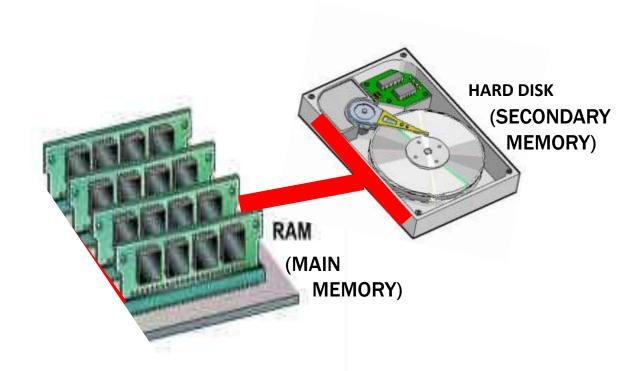
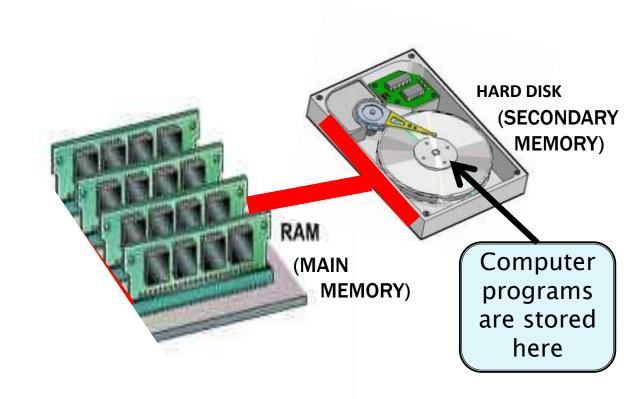
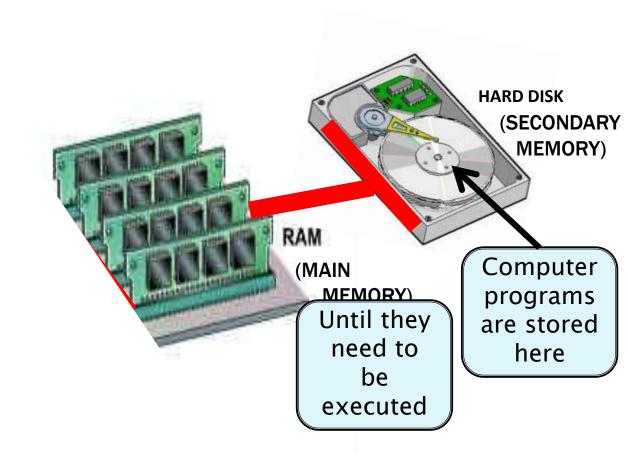
Memory Management: Virtual Memory

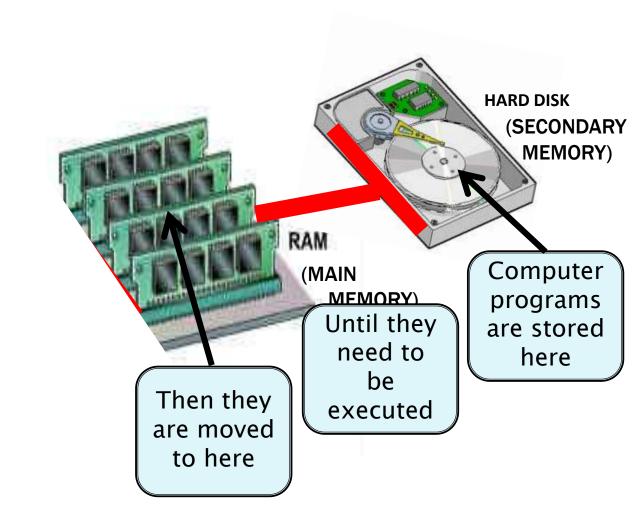
Damian Gordon









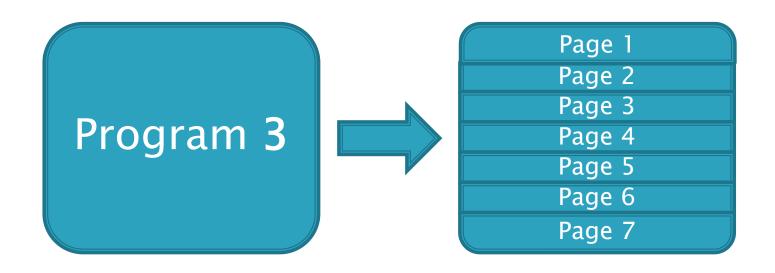


In modern operating systems, before a program is loaded into main memory, it is divided into chunks, called **PAGES**.

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• Each **PAGE** is loaded into memory locations called **PAGE FRAMES**.

▶ Each **PAGE** is loaded into memory locations

called **PAGE FRAMES**.

MAIN MEMORY

200K available

▶ Each **PAGE** is loaded into memory locations

called **PAGE FRAMES**.

Page Frame 1
Page Frame 2
Page Frame 3
Page Frame 4
Page Frame 5
Page Frame 6
Page Frame 7
Page Frame 8
Page Frame 9
Page Frame 10

200K available

If the PAGES are the exact same size as the PAGE FRAMES (and the same size as the disk sectors), this scheme works very well.

Page Frame 1
Page Frame 2
Page Frame 3
Page Frame 4
Page Frame 5
Page Frame 6
Page Frame 7
Page Frame 8
Page Frame 9
Page Frame 10

200K available

- The Memory Manager prepares a program for execution by doing the following:
 - 1. Determine the number of pages in the program
 - 2. Locate enough empty page frames in main memory
 - 3. Load all the program's pages into them

The empty page frame does not have to be contagious.

Consider a program that 350 bytes, and the page size is 100 bytes.

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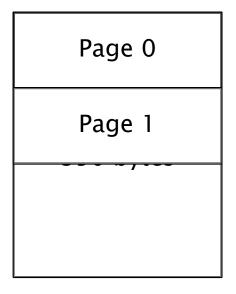
Program 4: 350 bytes

Consider a program that 350 bytes, and the page size is 100 bytes.

Page 0

Program 4: 350 bytes

Consider a program that 350 bytes, and the page size is 100 bytes.



Consider a program that 350 bytes, and the page size is 100 bytes.

Page 0
Page 1
Page 2

Consider a program that 350 bytes, and the page size is 100 bytes.

Page 0
Page 1
Page 2
Page 3

Consider a program that 350 bytes, and the

page size is 100 bytes.

Page 0

Page 1

Page 2

Page 3

Main Memory

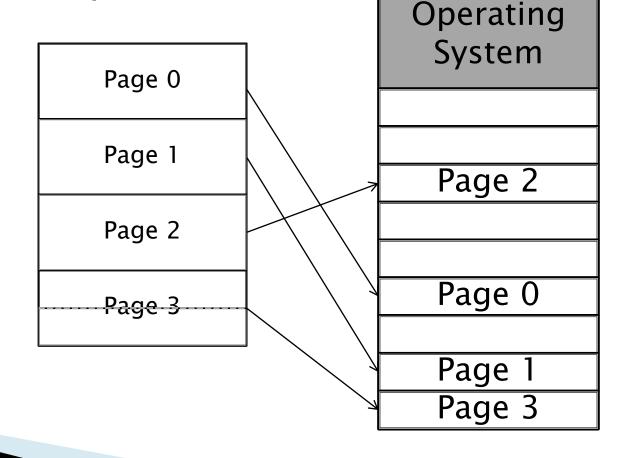
Consider a program that 350 bytes, and the page size is 100 bytes.

Page 0
Page 1
Page 2
Page 3

Operating System

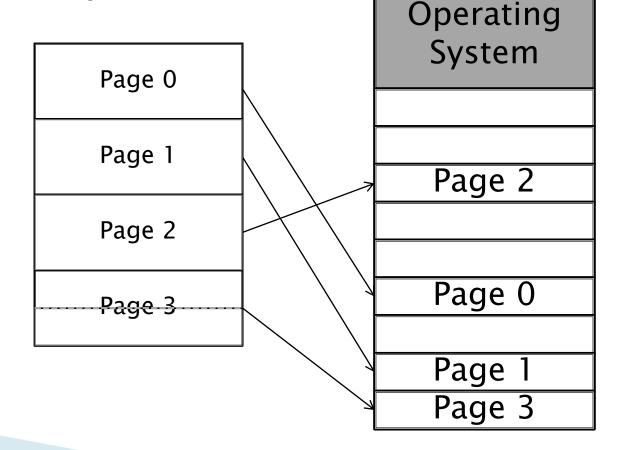
Main Memory

Consider a program that 350 bytes, and the page size is 100 bytes.



Consider a program that 350 bytes, and the page size is 100 bytes.

A little bit of internal fragmentation



Demand Paging

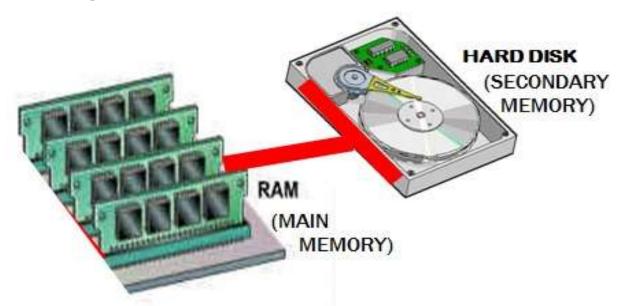
- A useful extension to the notion of PAGING is DEMAND PAGING.
- Demand Paging introduces the notion that you don't have to load the whole program into memory, just part of it.
- Because not all of the program needs to be in memory at the same time.

Demand Paging

This means that lots of programs can be run at the same time, and there is an illusion of a significantly larger amount of memory than with regular paging.

Demand Paging

To make this work, pages have to be moved very quickly from Secondary Storage to Main Memory and back again (this is called "swapping").



- This leads to the concept of VIRTUAL MEMORY
- The size of main memory appears much larger than the actual size, since many programs can appear to be fully loaded into main memory at the same time, when in actual fact, only part of many programs are loaded into main memory.