

# Assembly code

**registers 寄存器**: They are **small memory storage areas built into the processor**.

**Instructions 指令**: operate on memory (either RAM or registers), 比如加减, mov, jump等

## Instruction

### MOV:

移动数值到寄存器中。

would use the **MOV** (move) instruction, which

```
1. MOV RAX, 0x1234
```

As you can see, the MOV instruction takes two

MOV is similar to the assignment operator in C

```
1. myVariable = 0x1234;
```

You can also move values from one register in register RAX:

```
1. MOV RAX, RBX
```

In that case, the value of the register RBX isn't c

```
1. variable1 = variable2;
```

## ADD, SUB, AND, OR, NOT, XOR

The following instruction increments RBX by 4:

```
1. ADD RBX, 4
```

*RBX + 4*

The following instruction decrements RBX by 4:

```
1. SUB RBX, 4
```

*RBX - 4*

The instructions **AND**, **OR**, **NOT**, and **XOR** are also easy to understand. For example, the following assembly performs a bitwise AND on the RAX register with the 64-bit value `0x0000000000000000`. The result is stored in RAX.

```
1. AND RAX, 0x0000000000000000
```

## MUL, DIV 乘除

通过移位bit

multiplication (**MUL**), division (**DIV**), shifting bits to the left (SHL) or to the right (SHR) in a register.

## MOVZX move with the zero extended.

但因为取1 byte的内容到 4 bytes长的EAX寄存器，就需要补0.

*my\_arr = range force*

C code:

```
char character = my_array[0];
```

Assuming my\_array is a char array with the value "range force", the value of the variable "character" is "r".

Assembly:

```
MOVZX EAX, byte [ESI]
```

The value of EAX is 0x00000072 after the operation.

*0x72 = 'r'*

*18C72*

byte [ESI + 1] , 像array [1]取值一样。

比如range force

array的起始位，就是一个字符r，第二个字符就是[ESI + 1]

## Memory Dereferencing Calculations

move the second character? Or the third? Fortunately, you can perform calculations inside memory dereferencing operations.

	addresses	values	comments
ESI →	0x12345678	0x72616e67	72 61 6e 67 = "rang"
	0x1234567C	0x65666f72	65 66 6f 72 = "efor"
	0x12345680	0x63650000	63 65 = "ce"

C code:

```
char character = my_array[1];
```

Assuming my\_array is a char array with the value "range force", the value of the variable "character" is "a".

Assembly:

```
MOVZX EAX, byte [ESI + 1]
```

The value of EAX is 0x00000061 after the operation.

To get the second character of "range force" into EAX, you can do:

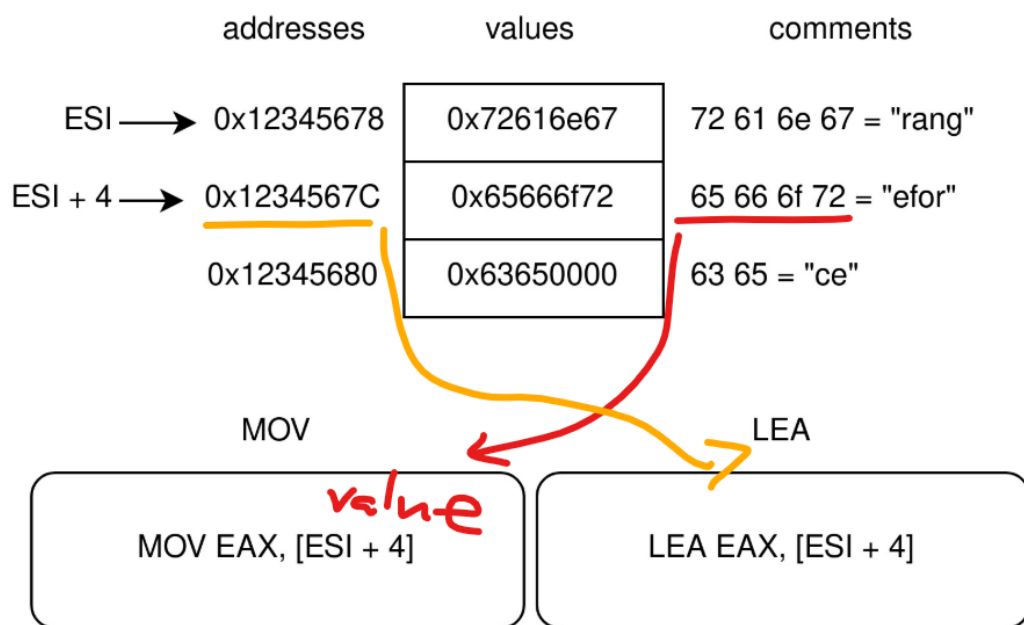
```
1. MOVZX EAX, byte [ESI + 1]
```

## LEA load effective address

获得某value的实际地址

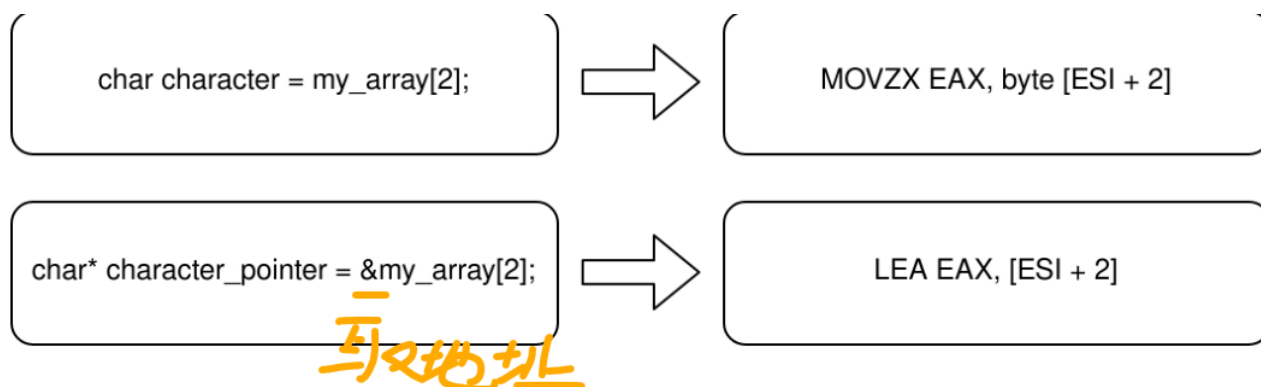
如果你像获得array中第几个字符串的地址，那就使用LEA。

Let's take the previous example with the "rangeforce" string again and see how LEA works.



The value of EAX is 0x65666f72.  
ASCII "efor"

The value of EAX is 0x1234567C.



## leave

mov ESP, EBP，就是将ebp指向的地址，给esp，这就相当于cleaning这个stack frame

## LEAVE

There's one instruction that may cause some confusion if you don't know what it does — **LEAVE**. LEAVE is a high-level instruction, which combines two other instructions into one. Executing the LEAVE instruction is exactly the same as executing the following two instructions:

In 64-bit applications:

1. MOV RSP, RBP
2. POP RBP

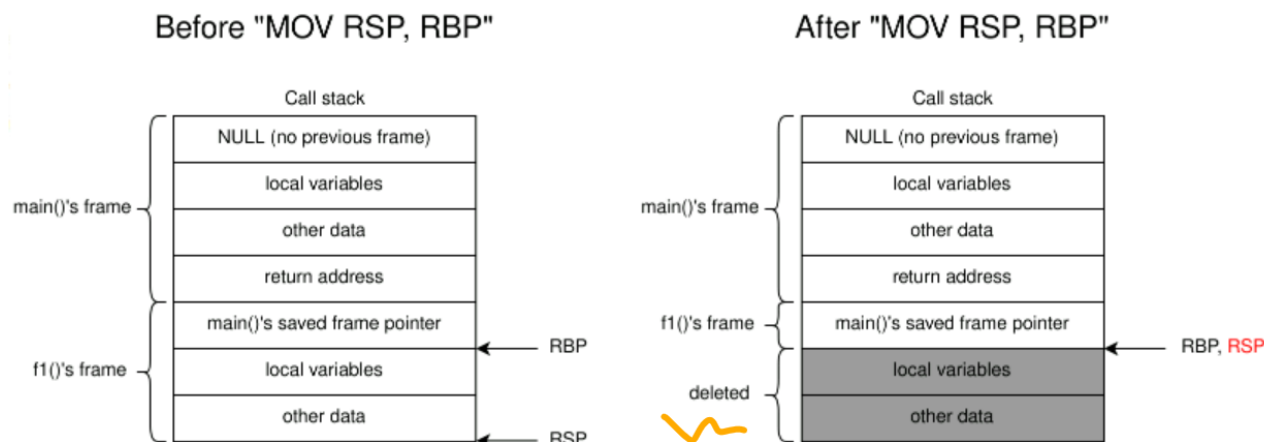
In 32-bit applications:

1. MOV ESP, EBP
2. POP EBP

*Handwritten note: "结合" (combine) with an arrow pointing to the LEAVE instruction.*

*Handwritten note: "离开一个 stack frame" (leave a stack frame).*

## 解释



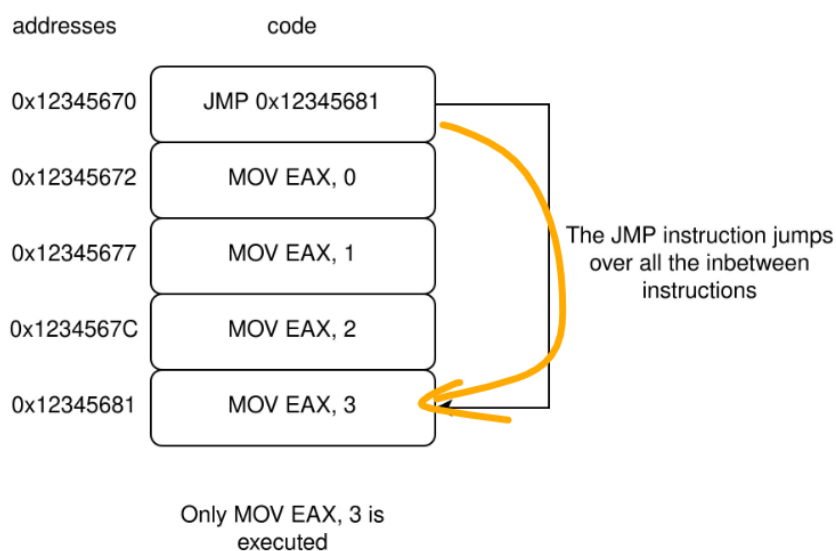
## CMP

CMP就是SUB，区别在于不会真的改变其值。

CMP works exactly like **SUB** (subtracts one operand from another), but with one exception — **it doesn't change the value you're comparing.**

## jmp, JUMP

The JMP instruction simply changes the **instruction pointer** to point to the specified address. instruction is executed next, the code will "jump" to wherever you point the instruction pointer



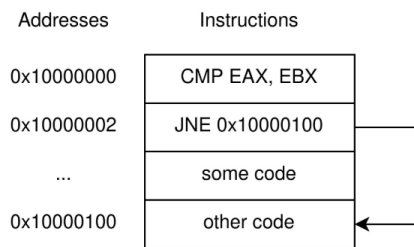
## Conditional jmp, JNE, JE, JL,JG, JLE, JGE, JZ

通常配合**CMP** compare指令

条件，比如if语句

如果不相同，则jump到指定地址。

In assembly, the conditional jump **JNE (jump if not equal)** is used along with the **CMP** (compare) instruction. Pay attention to how `some code` is not executed when the two variables are not the same:



Explanation: If EAX is **not equal** to EBX, then the JNE will **jump over** `some code`. `some code` will only get executed if EAX is equal to EBX. The assembly code example assumes that EAX contains the value of the `a` variable, and EBX contains the value of the `b` variable.

L: less than

G: greater than

E: Equal to

Z: zero, JZ, jump if zero

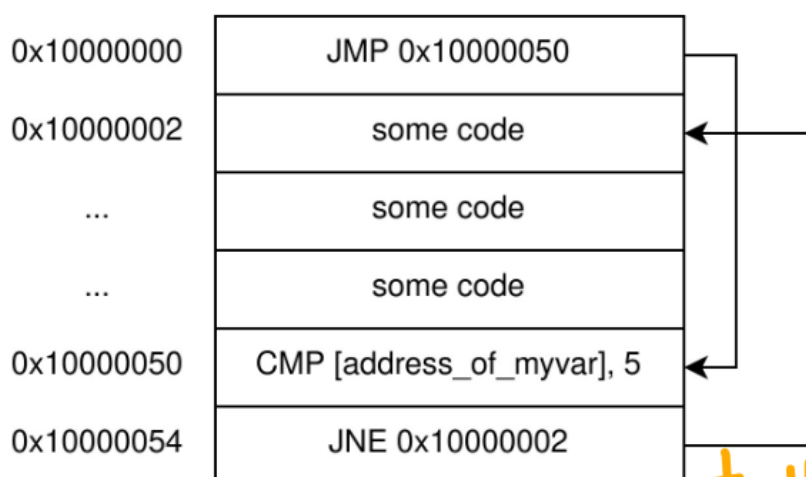
- If the conditional statement was `if (a != b)`, then instead of JNE, the **JE (jump if equals)** instruction would have been used.
- If the statement was `if (a >= b)`, then the **JL (jump if less than)** instruction would have been used.
- If the statement was `if (a <= b)`, then the **JG (jump if greater than)** instruction would have been used.
- If the statement was `if (a > b)`, then the **JLE (jump if less than or equals)** instruction would have been used.
- If the statement was `if (a < b)`, then the **JGE (jump if greater than or equals)** instruction would have been used.

## while loop 的jump使用

```
1. while (myvar != 5) {  
2.     // some code  
3. }
```

while loop

And this is what it looks like in assembly:



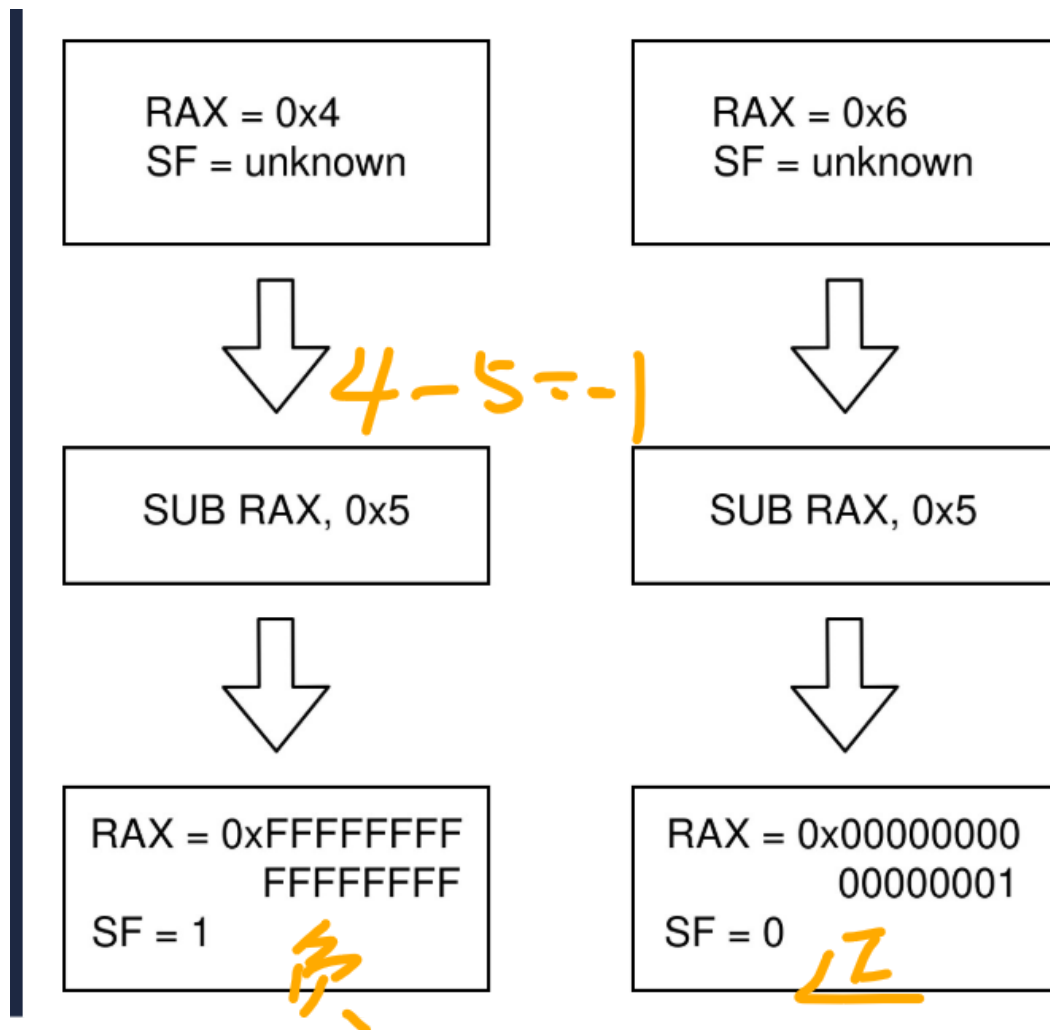
未达条件, jump back

## flags

## Sign flag

如果计算结果是positive, SF= 0, 如果是negative, SF=1

The sign flag is used for conditional jumps, such as **JGE** and **JL**. 配合JGE, JL jump



### Overflow flag

通常用来检测有符号的结果，是否overflow。

if the overflow flag is set, then that indicates an overflow condition for **signed integers**.

举个例子：if you subtract 1 from 0x00000000 and end up with 0xFFFFFFFF, then that is **not considered an overflow condition** for **signed integers**, but it is an overflow for unsigned integers.

When will a JGE jump be taken?



- ☒ **SF = 0, and there is no overflow (OF = 0)**
- ☐ SF = 0, but there was an overflow (OF = 1)
- ☐ SF = 1, and there is no overflow (OF = 0)
- ☒ **SF = 1, but there was an overflow (OF = 1)**

负

## Test

做的是And运算，只不过只改变对应的flags，不会改变其值。

TEST is the same as **AND**, but it only sets flags and doesn't change the value of any operands.

如果RAX是0，那么test rax, rax, 自己和自己比较只能是0，因为AND运算需要  $1 \& 1 = 1$ ，只要是0，其结果都是0。因此ZERO flag =1，因为结果为0。

```
1. TEST RAX, RAX
```

The above instruction sets the zero flag if the value of RAX is 0. That's because if you perform a logical AND on a number with that same number, then the result will be zero only if the number itself was 0 (  $0 \text{ AND } 0$  ). Otherwise, the result is a nonzero number (the result of  $2 \text{ AND } 2$  is 2).

A simpler way of checking whether RAX is 0 is by using the CMP instruction:

```
1. CMP RAX, 0
```

## machine code

汇编只是human readable代码，machine code才是机器能read的

- An **opcode**, which specifies the operation to be performed
- The operands of the instruction

For example, the machine code for `MOV EAX, 0x12345678` is `b878563412` :

- `b8` is the opcode.
- `78563412` is the operand — 0x12345678 in little endian.

1. The leftmost column shows the **address** where the instruction resides in memory.
2. The middle column shows the **machine code** that the computer understands.
3. The rightmost column shows the **human-readable assembly** that corresponds to the opcode.

- Radare2 may add comments next to some instructions. These start with `;`.

```
27: int main (int argc, char **argv, char **envp);
0x5617fdc4b149 f30f1efa endbr64
0x5617fdc4b14d 55 push rbp
0x5617fdc4b14e 4889e5 mov rbp, rsp
0x5617fdc4b151 488d3dac0e00 lea rdi, str.Hello_World_ ; 0x5617fdc4c004 ; "Hello World!"
0x5617fdc4b158 e8f3feffff call sym.imp.puts ; int puts(const char *s)
0x5617fdc4b15d b800000000 mov eax, 0
0x5617fdc4b162 5d pop rbp
0x5617fdc4b163 c3 ret
```

address      machine code      Assembly

## 寄存器

64位系统中，有16个寄存器。每一个寄存器是8 bytes(64 bits)

32位系统中，每一个寄存器是4 bytes (32 bits)



RIP 指向下一个要执行的instruction的地址

RSP和RBP来表示一个stack 范围

## Register Conventions

There are 16 general-purpose registers plus some special ones (notably the instruction pointer register). In 64-bit systems, each register is 8 bytes (64 bits) long. In 32-bit systems, each register is 4 bytes long.

64-bit general purpose registers

RAX	R8
RBX	R9
RCX	R10
RDY	R11
RSI	R12
RDI	R13
RSP	R14
RBP	R15

Instruction pointer  
RIP

Of those 16 + 1 registers, three are special:

- The **instruction pointer** register (RIP) always points to the next instruction that needs to be executed.
- Whenever an instruction is executed, the instruction pointer is automatically changed to point to the next instruction.
- The **stack pointer** register (RSP) keeps track of the top of the stack. This register is explained in the **Stack and Heap Basics** module.
- The **stack frame pointer** register (RBP) keeps track of the current stack frame. This register is explained in the **Stack Frames** module.

## 各个Pointer的主要作用。

RCX主要用于loop运算

- RAX (Accumulator register): Stores the return value of a function
- RCX (Counter register): Counter for string and loop operations
- RSI (Source Index register): Source pointer for string operations
- RDI (Destination Index register): Destination pointer for string operations

Keep in mind that these are just conventions. While you **will** often see the registers data operation purposes. The other legacy registers also have names and special conventions for now:

- RBX (Base register)
- RDX (Data register)

## RIP, EIP, IP

在16bits system中, ip就是一个16bit的寄存器

32bits, EIP, E for extended IP, 32 bit 寄存器

到64 bit, R 就for 64 bit寄存器

## 在实际情况中

在一个64bit的软件中, 可能会看到32 bit的寄存器使用。

原因是: 比如你只想取RAX内容中的其中32bit内容, 那么就可以使用EAX。其他寄存器也可以以此类推。

之所以这样做是因为原64 bit寄存器需要保存64 bit的return value .

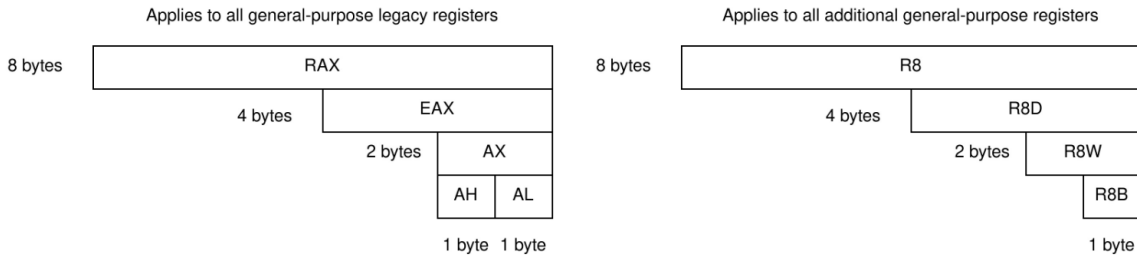
学习

What's going on is that if you only want to access the last 32 bits of the RAX register, then you can do so by using EAX instead of RAX. MOV EAX, 0 moves the 32-bit value 0x00000000 to the last 32 bits (4 bytes) of the RAX register. In this case, EAX is used instead of RAX because the value it needs to contain is 32 bits long (32-bit **int** data type). Therefore, only the last 4 bytes of RAX need to be used to contain the desired value (0x00000000).

保存原64 bit 的 value

The reason why the compiler assigns the value 0 to EAX in this "hello world" example is due to the convention that RAX should be used to store the function's return value. This topic will be covered in more depth in the **x86 Calling Conventions** modules.

Just as you can use EAX to address the last 32 bits of the RAX register, you can also use AX to address the last 16 bits. Furthermore, you can use AH to address the first 8 bits of the AX section, and you can use AL to address the last 8 bits. This naming scheme applies to all legacy registers. For example, to address the last 32 bits of the RBX register, you can use EBX.



R1015

As you can see, the naming for the additional registers is a bit different. R8 addresses the whole 8-byte register (8 bytes is a **QUADWORD**). **R8D** addresses the last 4 bytes of R8 (4 bytes is a **DWORD**). **R8W** addresses the last 2 bytes, and its length is 1 **WORD**. R8B addresses the last byte.

## 汇编语句 Intel, AT&T

# Intel vs. AT&T

Intel

1.	0000000000001149	<main>:		
2.	1149:	f3 0f 1e fa	endbr64	
3.	114d:	55	push	rbp
4.	114e:	48 89 e5	mov	rbp, rsp
5.	1151:	48 8d 3d ac 0e 00 00	lea	rdi, [rip+0xeac]
6.	1158:	e8 f3 fe ff ff	call	1050 <puts@plt>
7.	115d:	b8 00 00 00 00	mov	eax, 0x0
8.	1162:	5d	pop	rbp
9.	1163:	c3	ret	

And here is the same code in AT&T syntax:

AT&T

1.	0000000000001149	<main>:		
2.	1149:	f3 0f 1e fa	endbr64	
3.	114d:	55	push	%rbp
4.	114e:	48 89 e5	mov	%rsp, %rbp
5.	1151:	48 8d 3d ac 0e 00 00	lea	0xeac(%rip), %rdi
6.	1158:	e8 f3 fe ff ff	callq	1050 <puts@plt>
7.	115d:	b8 00 00 00 00	mov	\$0x0, %eax
8.	1162:	5d	pop	%rbp
9.	1163:	c3	retq	

There are many differences but the easy-to-notice ones are that in AT&T syntax:

- There are % signs in front of register names.
- The order of the operands is swapped.
- The instructions look a bit different.
- For example, `ret` in Intel syntax might be written as `retq` in AT&T syntax.

## Memory dereferencing

一个寄存器，你可以获取它的地址，或取得它包含的value。

[ ]表示取寄存器中的value。

Let's assume ESI contains the memory address of some value. Recapping the information you have learned so far, the instruction MOV EAX, ESI moves the address into the EAX register:

1. MOV EAX, ESI    取ESI地址

When the square brackets are used, the value at that address is accessed and moved into the EAX register:

1. MOV EAX, [ESI]    取ESI value