## Count prime

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- Given an integer n, return the number of prime numbers that are strictly less than n.
- Input: n = 10
- Output: 4
- Explanation: There are 4 prime numbers less than 10, they are 2, 3, 5, 7.

## 10 points out of a total 25 points that your program can run n<10<sup>4</sup>

- iterated from i = 2 to i<n</p>
  - iterated from j = 2 to j<i, check that whether i is perfectly divisible by j
    - If i is perfectly divisible by j, i is not a prime number, then terminated the loop
  - Otherwise, i is a prime number.

# 15 points out of a total 25 points that your program can run 10^5<=n<2.5\*10^6 in one second.

- iterated from i = 2 to i<n</p>
  - iterated from j = 2 to j\*j<=i, check that whether n is perfectly divisible by j</li>
    - If i is perfectly divisible by j, i is not a prime number, then terminated the loop
  - Otherwise, i is a prime number.

若N為合數  $\cdot$  且有r個因數 $(r \geq 2)$ 反證明  $P_1, P_2, \cdots, P_r > \sqrt{N}$ 

$$N = P_1 imes P_2 imes \cdots imes P_r > N^{rac{r}{2}} \ N^{rac{r}{2}} \geq N \ (\because rac{r}{2} \geq 1)$$

 $\because N > N^{rac{r}{2}}$ 和  $N^{rac{r}{2}} \geq N$  矛盾

:. 反證成立

If it is a composite number, it must have at least one factor  $<=\sqrt{n}$ 

## 25 points out of a total 25 points that your program can run 3.3\*10^6<=n<=5\*10^6 in one second.

#### Sieve of Eratosthenes

algorithm Sieve of Eratosthenes is

input: an integer n > 1.

output: all prime numbers count from 2 through n-1.

let A be an array of Boolean values, indexed by integers 2 to n-1, initially all set to true.

for i = 2, 3, 4, ..., not exceeding 
$$\sqrt{n}$$
 do if A[i] is true for j =  $i^2$ ,  $i^2$  +i,  $i^2$  +2i,  $i^2$  +3i, ..., not exceeding n do A[j] := false

return all i such that A[i] is true.

10 20 30
20 30
30
40
40
50
60
70
80
90
100
110
120

Prime numbers

### Questions?