COMP 140: Computational Thinking Breadth-First Search

Names:

The breadth-first search algorithm is given here for your reference.

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Algorithm 1: Breadth-First Search
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Input: Undirected graph graph, start node start in graph

Output: $dist_{node}$, for each node in the graph: the number of steps between start and node in graph; $parent_{node}$, for each node node in graph: the parent of node in the exploration of graph starting from start

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1 Initialize queue to an empty queue;
2 Initialize dist to an empty mapping;
3 Initialize parent to an empty mapping;
4 foreach node in graph do
        dist_{node} \leftarrow \infty;
       parent_{node} \leftarrow null;
7 dist_{start} \leftarrow 0;
8 push start onto queue;
  while queue is not empty do
10
        pop node off of queue;
       foreach neighbor nbr of node do
11
            if dist_{nbr} = \infty then
12
                dist_{nbr} \leftarrow dist_{node} + 1;
13
                parent_{nbr} \leftarrow node;
14
                push nbr onto queue;
```

16 return dist, parent;

For each of the following four graphs, walk through BFS by hand. For graphs 1-3, you should use 0 as your start node; for graph 4 you should use Kevin Bacon.

As you go, you will perform two forms of bookkeeping. First, trace the edges as you traverse them with colored pencils, using a different color for each distance from the *start* node. Second, keep track of the states of the data structures as you progress through the algorithm. In particular, for each iteration of the while loop that you walk through, write down the state of *queue* (prior to the pop on line 10), *dist*, and *parent*. Also record the final states of these data structures upon reaching the return statement. This will be a lot of writing! So, you may take a few shortcuts. First, after the first iteration, you need not record the states of unvisited nodes (those with a distance of infinity) in either the *dist* or *parent* data structure. Second, if a data structure remains entirely unchanged since the previous iteration, you can simply write "no change".

You should write the state of the queue as a sequence of values $q_0, q_1, ..., q_{n-1}$, where q_0 is the head of the queue and q_{n-1} is the tail.

For Graph 4, you do not need to record the states of the data structures at every step. Instead, you may just trace the edges with colored pencils and record the final (upon return) states.

For example, consider the following graph:

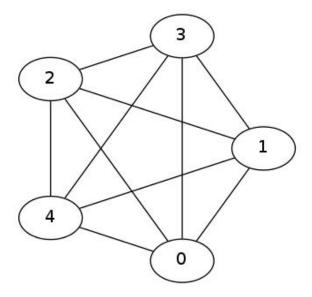


Performing BFS on this graph using node 0 as the *start* node should be recorded as follows:

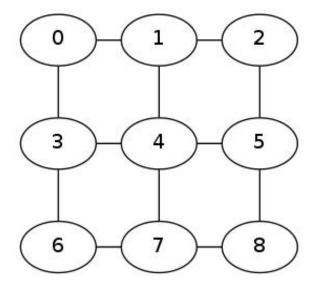
Iteration 0:	Iteration 2:
queue: 0	queue: 2
	dist:
dist:	$0 \rightarrow 0$
$0 \rightarrow 0$	$1 \rightarrow 1$
$1, 2 \to \infty$	$2 \rightarrow 2$
	parent:
parent:	$0 \rightarrow null$
$0, 1, 2 \rightarrow null$	$1 \rightarrow 0$
0, 1, 2	$2 \rightarrow 1$
	
Iteration 1:	Return
queue: 1	Return: queue: empty
queue: 1	
queue: 1 dist:	queue: empty
queue: 1 dist: $0 \rightarrow 0$	queue: empty dist:
queue: 1 dist:	queue: empty dist: $0 \rightarrow 0$
queue: 1 dist: $0 \rightarrow 0$	queue: empty dist: $0 \rightarrow 0$ $1 \rightarrow 1$ $2 \rightarrow 2$
queue: 1 dist: $0 \rightarrow 0$	queue: empty dist: $0 \rightarrow 0$ $1 \rightarrow 1$ $2 \rightarrow 2$ parent:
queue: 1 dist: $0 \rightarrow 0$ $1 \rightarrow 1$	queue: empty dist: $0 \rightarrow 0$ $1 \rightarrow 1$ $2 \rightarrow 2$

This should result in edges (0, 1) and (1, 2) each having a unique color, since (0, 1) is 1 hop from the *start* node whereas (1, 2) is 2. Note that if you had performed BFS starting from node 1, both edges would have been given the same color as they are in the same "layer". Note also that in some graphs, not all edges will be colored, as it will be possible to reach every node from the *start* node without traversing every edge in the graph.

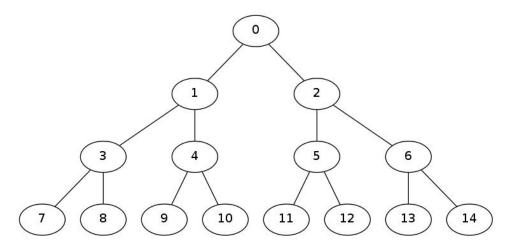
1 Graph 1



2 Graph 2



3 Graph 3



4 Graph 4

