

# Quiz 5

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## Problem 2-1

It's similar to Fibonacci sequence , for  $n \geq 2$  , every  $n$ th number will be the sum of  $n-1$ th and  $n-2$ th number .

## Problem 2-2

1. First, initialize the LFSR to length = 1 and set it to the first element of the sequence:  $LFSR = (0)$
2. Compute the discrepancy between LFSR output and the next element of the sequence , and the discrepancy is the difference between the predicted value and the actual value . In the former case , the predicted value is the only value in LFSR , making discrepancy be the same as the next element of the sequence. For example , the discrepancy between  $LFSR = (0)$  and the next element of the sequence(1) , is also 1.
3. Update the coefficients of LFSR . First , compute a new LFSR using the previous one and the discrepancy . The new LFSR'll be the sum of the previous LFSR , the product of the discrepancy , and a shifted version of previous LFSR( last element removed) .

4. Repeat steps 2 and 3 for each subsequent element in the sequence.

The following are steps for the first few elements:

- LFSR = (0)
- Discrepancy between LFSR = (0) and 1 is 1.
- Update coefficients:  $\text{LFSR} = (0) + 1 \cdot () = (0)$
- Discrepancy between LFSR = (0) and 1 is 1.
- Update coefficients:  $\text{LFSR} = (0) + 1 \cdot (0) = (0)$
- Discrepancy between LFSR = (0) and 2 is 2.
- Update coefficients:  $\text{LFSR} = (0) + 2 \cdot () = (0)$
- Discrepancy between LFSR = (0) and 3 is 3.
- Update coefficients:  $\text{LFSR} = (0) + 3 \cdot (0) = (0)$

We can see that LFSR hasn't changed after processing the first four elements of the sequence , which means that the minimum LFSR capable of generating the first four elements of the sequence has length 1 and consists of a single tap.

Continue the process for the full sequence , then we can get the following LFSR :  $[0,1,1,0,1,1,1,0,1]$  . Thus , the rule for the requested sequence is generated by LFSR with coefficients  $[0,1,1,0,1,1,1,0,1]$  .