Cal State Fullerton

College of Engineering and Computer Science

Project 1:

Implementing Algorithm for Sorting Disks

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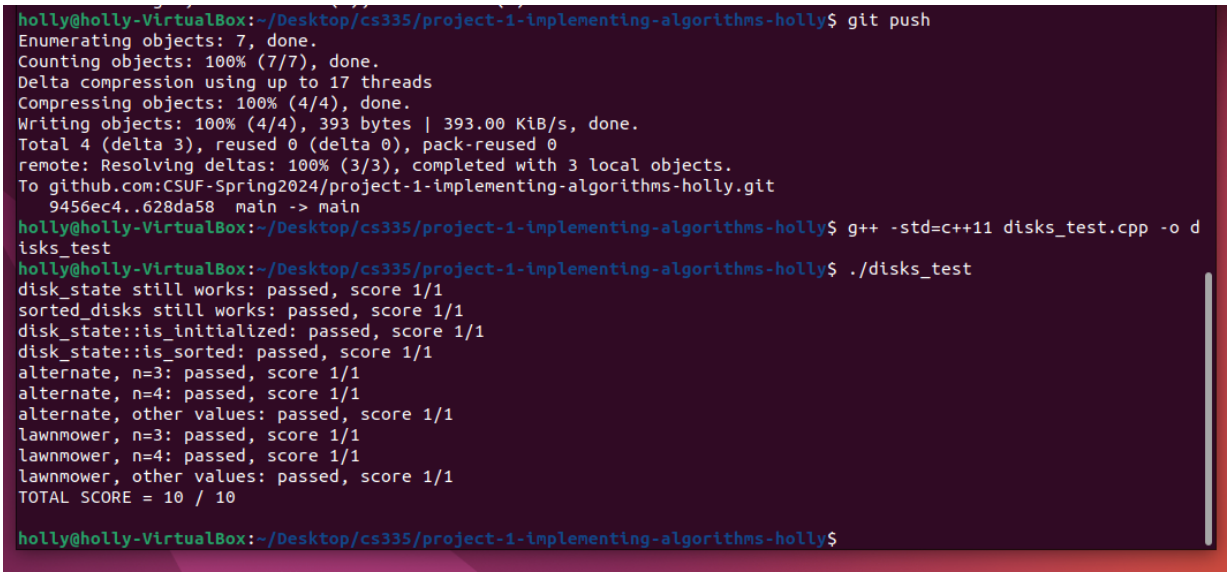
Introduction

In this first assignment for CS335 Algorithm class, we will implement two ways of sorting an alternating disks vector.

The definition of the problem is:
Input: a list of alternating light and dark disks size of 2n, with n is a positive integer.
Output: sorted vector with dark disks on the left and light disks on the right.

Screenshots



Pseudo-Code

1/
sort_lawnmower (disk_state& before)

disk_state after = before; // 1

count = 0; //1

for j=0 to (count()/2) -1 do // (count()/2) -1 +1

for i=j to (count()-1)-j-1 jump 2 steps forward do // ((count()-1)-j-1 –j) / 2 +1

if (after[i]== DISK_LIGHT && after[i+1] == DISK_DARK) // 1+1+1 = 3

{

swap; //1

count++; //1

}

else do nothing //0

endif

for i = count()-1-j to 1+j-1 jump 2 steps backward do // (1+j-1 - (count()-1-j)) /2 +1

if (after[i]== DISK_LIGHT && after[i+1] == DISK_DARK) //3

{

```
swap; //1
count++; //1
}
else do nothing
endif
endfor
endfor

 return after;

2/
sort_alternate (disk_state& before)

disk_state after = before; //1
count = 0; //1
for j=0 to (count() -2) do // (count() -2) +1
for i=j to count()-2) jump 2 steps forward do // ((count() -2) - j ) / 2 +1
if (after[i]== DISK_LIGHT && after[i+1] == DISK_DARK) //3
{
swap; //1
count++; //1
}
else do nothing
endif
endfor
endfor
return after
```

Big O Analysis

The input disks arrangement starts with a white disk and then a black disk and so on until the end of the vector. The length of this array/vector is always even because we started with an integer n number of white disks then add n number of black disks. For the algorithms to complete the entire length of the vector, we need to make sure to set the iteration numbers correctly. For example, the sort_alternate() function needs to be repeated n/2 times, which means the last pair will be at indexes (n/2)-1 and n/2. Similarly, the sort_lawnmower() swapping decision needs to be iterated until the last pair of indexes (n/2)-1 and n/2. Below is the Big O analysis of the runtime for both functions.

sort_lawnmower()

1/ if s.c = 3 + max(2,0) = 5

inner loop ≠ 1: $s.c = 5\left(\frac{c-1-j-1-j}{2}+1\right) = \left(\frac{c}{2}-j\right)\cdot 5$

second inner for loop: $s.c = 5\left(\frac{(1+j-1)-(c-1)-j)}{-2}+1\right) = \left(\frac{c}{2}-j+1\right)5$.

Total $s.c = 1+1+\sum_{j=0}^{\frac{c}{2}-1} 5\left(\frac{c}{2}-j\right) + \sum_{j=0}^{\frac{c}{2}-1} 5\left(\frac{c}{2}-j+1\right)$

$= 2 + \sum_{j=0}^{\frac{c}{2}-1} 5\left(\frac{c}{2}-j+\frac{c}{2}-j+1\right) = 2+\sum_{j=0}^{\frac{c}{2}-1} 5(c-2j+1)$

$= 2 + 5\left(\frac{c}{2}-1+1\right) - 10\,\frac{\left(\frac{c}{2}-1\right)\left(\frac{c}{2}\right)}{2} + 5c\left(\frac{c}{2}-1+1\right)$

$= 2 + 5\frac{c}{2} - 5\left(\frac{c^2}{4}-\frac{c}{2}\right) + 5c\left(\frac{c}{2}\right)$

$= 2 + \frac{5c}{2} - \frac{5}{4}c^2 + \frac{5}{2}c + \frac{5c^2}{2} = 2 + 5c + \frac{5}{4}c^2$

This belongs to big O of n^2 because largest term is n^2 (or c^2)

sort_alternate()

2| if s.c = 3 + max(2,0) = 5

inner for loop s.c = $5 \cdot \left[\left(\frac{c-2-j}{2}\right)+1\right] = 5\left(\frac{c-j}{2}\right)$

Total s.c = $1+1+\sum_{j=0}^{c-2} \frac{5}{2}c - \sum_{j=0}^{c-2} \frac{5}{2}j$

$= 2 + \frac{5}{2}c.(c-2+1) - \frac{5}{2}\frac{(c-2)(c-1)}{2}$

$= 2 + \frac{5}{2}c(c-1) - \frac{5}{4}(c^2 - c - 2c + 2)$

$= 2 + \frac{5}{2}c^2 - \frac{5}{2} - \frac{5}{4}(c^2 - 3c + 2)$

$= 2 + \frac{5}{2}c^2 - \frac{5}{2} - \frac{5}{4}c^2 + \frac{15}{4}c - \frac{10}{4}$

$= \frac{5}{4}c^2 + \frac{15}{4}c - 5 + 2 = \frac{5}{4}c^2 + \frac{15}{4}c - 3$

This belongs to big O of n^2 (or c^2)