

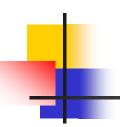
## CHAPTER 5: DISCRETE FOURIER TRANSFORM (DFT)

Lecture 9: DFT and Inverse DFT

Lecture 10: Fast Fourier Transform (FFT)

**Duration:** 4 hrs





## Lecture 9 DFT and Inverse DFT

- Duration: 2 hrs
- Outline:
  - 1. Review of DTFT of DT periodic signals
  - 2. DFT and Inverse DFT
  - 3. Frequency resolution
  - 4. DFT properties

#### Procedure to calculate DTFT of periodic signals

#### Step 1:

Start with  $x_0(n)$  – one period of x(n), with zero everywhere else

#### Step 2:

Find the DTFT  $X_0(\Omega)$  of the signal  $x_0[n]$  above

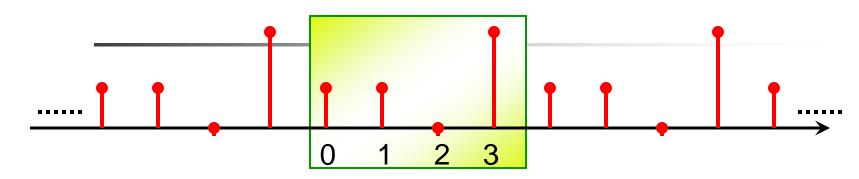
#### **Step 3:**

Find  $X_0(\Omega)$  at **N** equally spacing frequency points  $X_0(k2\pi/N)$ 

#### **Step 4:**

Obtain the DTFT of x(n): 
$$X(\Omega) = \frac{2\pi}{N} \sum_{k} X_0(k \frac{2\pi}{N}) \delta(\Omega - k \frac{2\pi}{N})$$

## Example of calculating DTFT of periodic signals



$$X_0(\Omega) = \sum_{n=0}^{3} x_0(n)e^{-j\Omega n} = 1 + e^{-j\Omega} + 2e^{-j3\Omega}$$

$$X_{o}(\frac{2\pi k}{4}) = 1 + e^{-j\frac{2\pi k}{4}} + 2e^{-3j\frac{2\pi k}{4}} \qquad k = 0, 1, 2, 3$$

$$k = 0 \rightarrow X_0(0) = 4;$$
  $k = 1 \rightarrow X_0(1) = 1+j$ 

$$k = 2 \rightarrow X_0(2) = -2;$$
  $k = 3 \rightarrow X_0(3) = 1-j$ 

#### **Example (cont)**

$$X(\Omega) = \frac{2\pi}{N} \sum_{k=-\infty}^{\infty} X_o(\frac{2\pi k}{N}) \delta(\Omega - \frac{2\pi k}{N})$$
$$= \frac{2\pi}{4} \sum_{k=-\infty}^{\infty} X_o(\frac{2\pi k}{4}) \delta(\Omega - \frac{2\pi k}{4})$$

For one period  $0 \le \Omega < 2\pi$ 

$$\frac{\pi}{2} \left\{ 4\delta(\Omega) + (1+j)\delta(\Omega - \frac{\pi}{2}) - 2\delta(\Omega - \pi) + (1-j)\delta(\Omega - \frac{3\pi}{2}) \right\}$$





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  - 4. Applications

#### **DFT** to the rescue!

Could we calculate the **frequency spectrum** of a signal using a **digital computer** with **CTFT/DTFT**?

- ➤ Both CTFT and DTFT produce continuous function of frequency → can't calculate an infinite continuum of frequencies using a computer
- $\triangleright$  Most real-world data is not in the simple form such as  $a^n u(n)$

DFT can be used as a FT approximation that can calculate a **finite set of discrete-frequency spectrum values** from a finite set of discrete-time samples of an analog signal

#### **Building the DFT formula**

Continuous time

signal **x(t)** |

sample

Discrete time signal x(n)

window

Discrete time signal x<sub>0</sub>(n)

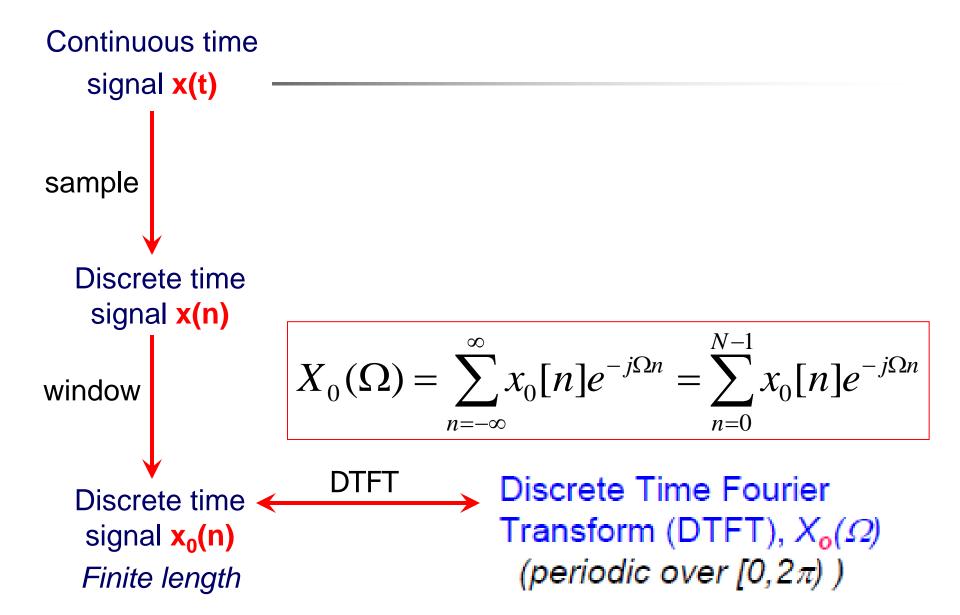
Finite length

"Window" x(n) is like multiplying the signal by the finite length rectangular window

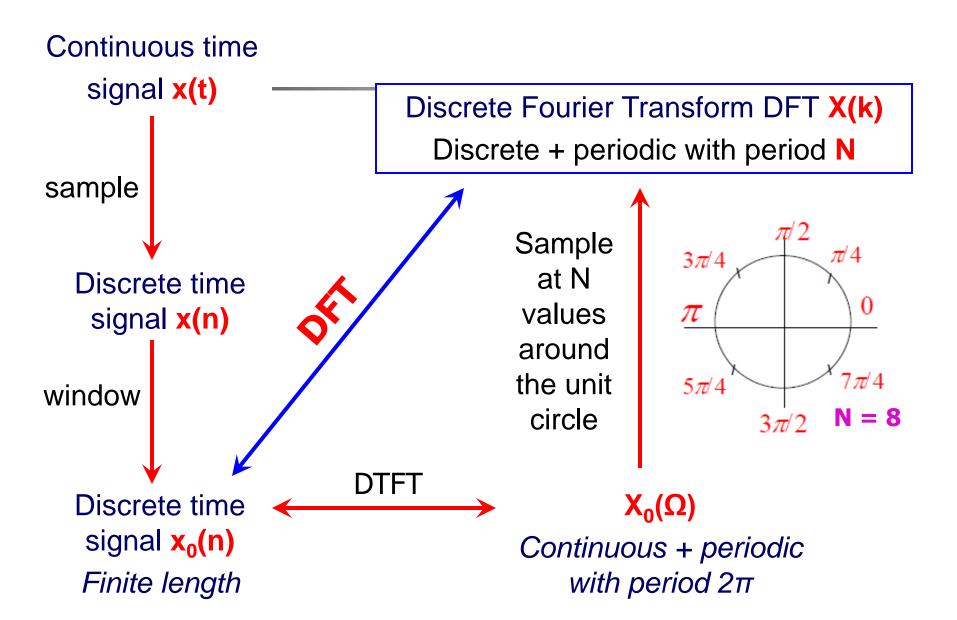
$$w_{R} = \begin{cases} 1 & n = 0, 1, \dots N-1 \\ 0, & otherwise \end{cases}$$

$$x_o[n]=x[n]$$
  $w_R[n]$ 

#### **Building the DFT formula (cont)**



#### **Building the DFT formula (cont)**



#### **DFT and inverse DFT formulas**

Notation: 
$$W_N = e^{-j\frac{2\pi}{N}}$$

$$X[k] = \sum_{n=0}^{N-1} x[n] W_N^{kn}$$

$$k = 0, 1, ..., N-1$$

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{-kn}$$

$$n = 0,1,...,N-1$$

Note that the DFT is a sequence of N numbers (in the frequency domain), just like x<sub>o</sub>[n] is a sequence of N numbers in the time domain

You only have to store N points

**Ex.1.** Find the DFT of x(n) = 1, n = 0, 1, 2, ..., (N-1)

$$X[k] = \sum_{n=0}^{N-1} x[n] W_N^{kn} = \sum_{n=0}^{N-1} W^{kn} = \frac{1 - W^{kN}}{1 - W^k}$$

$$\mathbf{k} = 0, 1, \dots, N-1$$

$$k = 0 \rightarrow X(k) = X(0) = N$$

$$k \neq 0 \rightarrow X(k) = 0$$

$$X[k] = N\delta[k]$$

**Ex.2.** Given  $y(n) = \delta(n-2)$  and N = 8, find Y(k)

$$Y[k] = \sum_{n=0}^{N-1} y[n] W_N^{kn} = \sum_{n=0}^7 \delta[n-2] e^{-j2\pi kn/N}$$

$$= e^{-j4\pi k/8} = e^{-j\pi k/2} = (-j)^k$$

$$= [1, -j, -1, j, 1, -j, -1, j]$$
Using N=8

**Ex.3.** Find the IDFT of X(k) = 1, k = 0, 1, ..., 7.

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{-kn}$$

$$x[n] = \frac{1}{8} \sum_{k=0}^{7} W_8^{-kn} = \frac{1}{8} N \delta[n] = \delta[n]$$

**Ex.4.** Given  $x(n) = \delta(n) + 2\delta(n-1) + 3\delta(n-2) + \delta(n-3)$  and N = 4. Find X(k).

$$X[k] = 1 + 2e^{-j\pi k/2} + 3e^{-j\pi k} + e^{-j3\pi k/2}$$

$$= 1 + 2(-j)^k + 3(-1)^k + (j)^k$$

$$X[0] = 7$$

$$X[1] = 1 - 2j - 3 + j = -2 - j$$

$$X[2] = 1 - 2 + 3 - 1 = 1$$

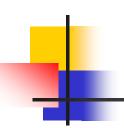
$$X[3] = 1 + 2j - 3 - j = -2 + j$$

## 4

- x = [1 2 3 1];
- $\rightarrow$  >> X = fft(x)

- 7.0000 1.7071 5.1213i -2.0000 1.0000i 0.2929 + 0.8787i
- Columns 5 through 8





## Lecture 9 DFT and Inverse DFT

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#### Frequency resolution of the DFT

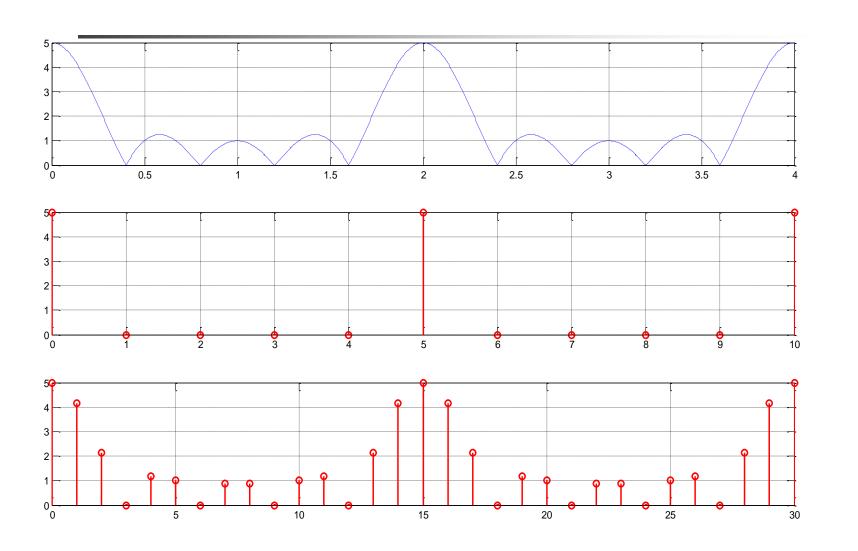
Discrete frequency spectrum computed from DFT has the spacing between frequency samples of:

$$\Delta f = \frac{f_s}{N}$$

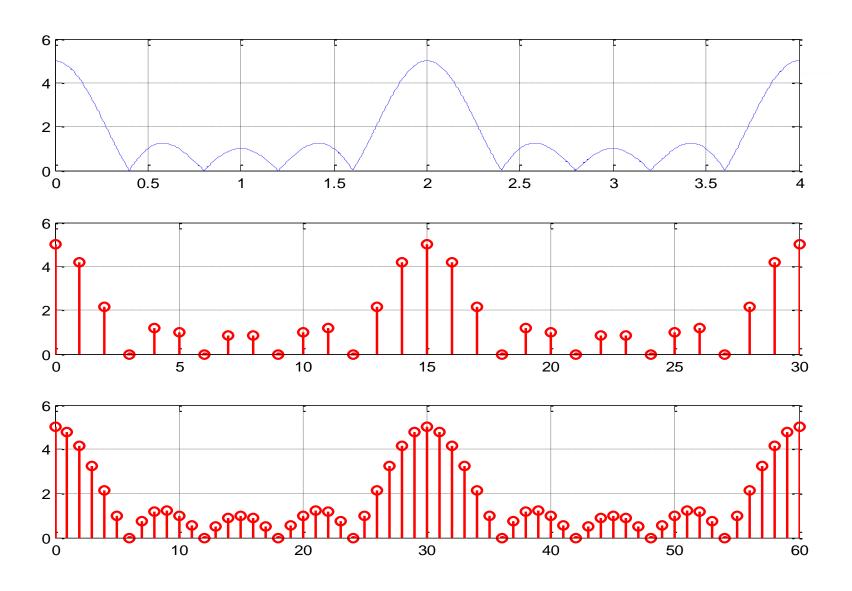
$$\Delta\Omega = \frac{2\pi}{N}$$

- → The choice of N determines the resolution of the frequency spectrum, or vice-versa
- → To obtain the adequate resolution, some zeros can be appended to the signal (*zero padding*)

#### Examples of N = 5 and N = 15



#### Examples of N = 15 and N = 30



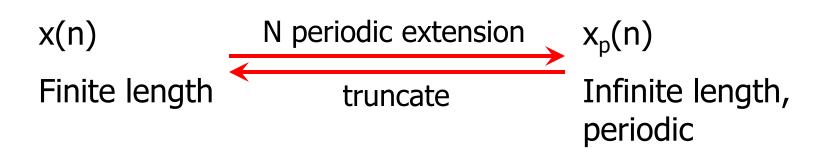




## Lecture 9 DFT and Inverse DFT

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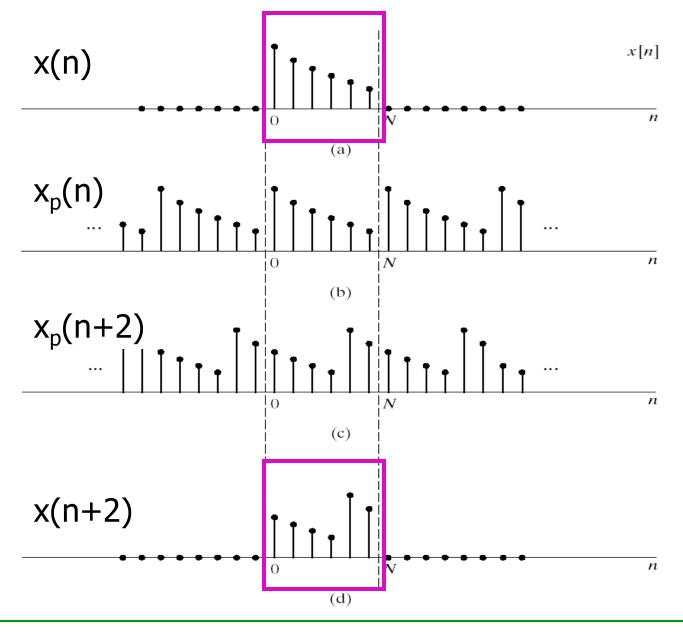
#### Circular shift property of the DFT



x(n) is one period of signal  $x_p(n)$ 

$$x[n-m] \xrightarrow{DFT} W^{km}X[k]$$

# Circular shift property of the DFT



Circular shift by *m* is the same as a shift by *m modulo N* 

#### **Recall linear convolution**

$$y[n] = x_1[n] * x_2[n] = \sum_{p=-\infty}^{\infty} x_1[p] x_2[n-p]$$

- N<sub>1</sub>: the non-zero length of  $x_1(n)$ ; N<sub>2</sub>: the non-zero length of  $x_2(n)$ ; N<sub>y</sub> = N<sub>1</sub> + N<sub>2</sub> -1
- The shift operation is the regular shift
- The flip operation is the regular flip

#### Circular convolution of the DFT

$$y[n] = x_1[n] \otimes x_2[n] = \sum_{p=0}^{N-1} x_1[p] x_2[n-p]_{\text{mod } N}$$

- The non-zero length of  $x_1(n)$ ,  $x_2(n)$  and y(n) can be no longer than N
- The shift operation is circular shift
- The flip operation is circular flip

$$x_1[n] \otimes x_2[n] \xrightarrow{DFT} X_1[k] X_2[k]$$

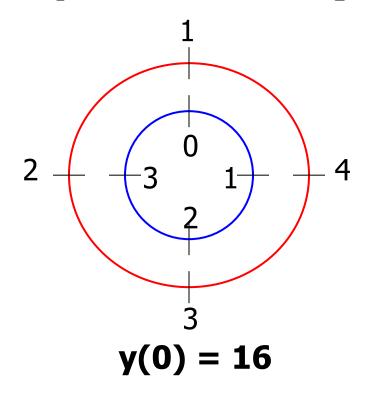
### Direct method to calculate circular convolution

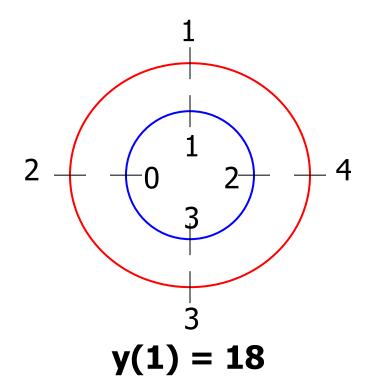
- **1.** Draw a circle with N values of x(n) with N equally spaced angles in a counterclockwise direction.
- **2.** Draw a smaller radius circle with N values of h(n) with equally spaced angles in a clockwise direction. Superimpose the centers of 2 circles, and have h(0) in front of x(0).
- **3.** Calculate y(0) by multiplying the corresponding values on each radial line, and then adding the products.
- **4.** Find succeeding values of y(n) in the same way after rotating the inner disk counterclockwise through the angle  $2\pi k/N$

#### Example to calculate circular convolution

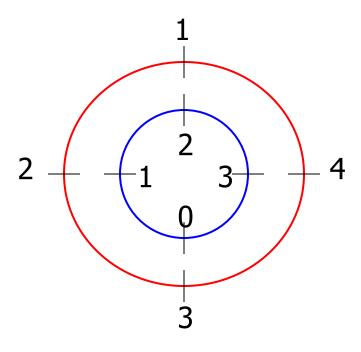
Evaluate the circular convolution, y(n) of 2 signals:

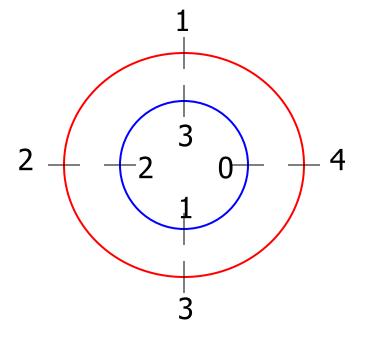
$$x_1(n) = [1234]; x_2(n) = [0123]$$





#### **Example (cont)**





$$y(2) = 16$$

$$y(3) = 10$$

## Another method to calculate circular convolution

$$x_1(n) \xrightarrow{DFT} X_1(k) \xrightarrow{IDFT} y(n)$$
 $x_2(n) \xrightarrow{DFT} X_2(k)$ 

Ex. 
$$x_1(n) = [1 2 3 4]; x_2(n) = [0 1 2 3]$$
  
 $X_1(k) = [10, -2+j2, -2, -2-j2];$   
 $X_2(k) = [6, -2+j2, -2, -2-j2];$   
 $Y(k) = X_1(k).X_2(k) = [60, -j8, 4, j8]$   
 $y(n) = [16, 18, 16, 10]$ 

#### Calculation of the linear convolution

The circular convolution of 2 sequences of length  $N_1$  and  $N_2$  can be made **equal to** the linear convolution of 2 sequences by zero padding both sequences so that they both consists of  $N_1+N_2-1$  samples.

Zero padding 
$$X'_1(n)$$
 DFT  $N_1$  samples  $X'_1(n)$  DFT  $X_1(n)^*X_2(n)$   $X_2(n)$   $X_2$   $X_3$   $X_4$   $X_4$   $X_5$   $X_5$   $X_5$   $X_5$   $X_6$   $X_7$   $X_8$   $X_8$   $X_8$   $X_9$   $X$ 

#### Example of calculation the linear convolution

```
1 2 3 4 ]; x_2(n) = [0 1 2 3]
x'_{1}(n) = [1234000]; x'_{2}(n) = [0123000]
X'_{1}(k) = [10, -2.0245-j6.2240, 0.3460+j2.4791, 0.1784-
j2.4220, 0.1784+j2.4220, 0.3460-j2.4791, -2.0245-j6.2240 ];
X'_{2}(k) = [6, -2.5245-j4.0333, -0.1540+j2.2383, -0.3216-j4.0333, -0.1540+j2.2383, -0.3216-j4.0333, -0.1540+j2.2383, -0.3216-j4.0333, -0.3216-j4.0316-j4.0333, -0.3216-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316-j4.0316
j1.7950, -0.3216+j1.7950, -0.1540-j2.2383, -2.5245+j4.0333];
Y'(k) = [60, -19.9928 + j23.8775, -5.6024 + j0.3927, -5.8342 - j0.3927, -5.8427, -5.8427, -5.8427, -5.8427, -5.8427, -5.8427, -5.8427, -5.8427, -5.8427, -5.8427, -5.8427, -5.8427, -5
i0.8644, -4.4049+j0.4585, -5.6024-j0.3927, -19.9928
   +i23.8775 ]
  IDFT\{Y'(k)\} = y'(n) = [0 1 4 10 16 17 12]
```



#### HW



#### Prob.1

Compute the DFT with N time samples:

(a) 
$$x[n] = \delta[n]$$

(b) 
$$x[n] = a^{n} \{u[n] - u[n - N]\}$$

(c) 
$$x[n] = \begin{cases} 1 & \text{n even} \\ 0 & \text{n odd} \end{cases}$$
  $0 \le n \le N-1$ 







#### Prob.2

Given the two four-point sequences:

$$x(n) = [1 0.75 0.5 0.25]$$

$$y(n) = [0.75 \ 0.5 \ 0.25 \ 1]$$

Express the DFT Y(k) in terms of the DFT X(k)





#### Prob.3

Given signals below and their DFT-5

(a) 
$$x_1[n] = \delta[n-1] + 2\delta[n-2] + 3\delta[n-3] + 4\delta[n-4]$$

(b) 
$$x_2[n] = \delta[n-1]$$

(c) 
$$s[n] = \delta[n]$$

- 1. Find y[n] so that Y[k] =  $X_1[k].X_2[k]$
- 2. Does  $x_3[n]$  exist, if  $S[k] = X_1[k].X_3[k]$ ?





#### **Prob.4** Given x(n) and its 8-point DFT, X(k)

$$x[n] = \begin{cases} 1, & 0 \le n \le 3 \\ 0, & 4 \le n \le 7 \end{cases}$$

Express the DFTs of the signals below in terms of X(k).

(a) 
$$x_1[n] = \begin{cases} 1, & n = 0 \\ 0, & 1 \le n \le 4 \\ 1, & 5 \le n \le 7 \end{cases}$$

(b) 
$$x_2[n] = \begin{cases} 0, & 0 \le n \le 1 \\ 1, & 2 \le n \le 5 \\ 0, & 6 \le n \le 7 \end{cases}$$







**Prob.5** The 8-point DFTs of x(n) and h(n) are:

$$X(k) = [0, -j0.707, -j, -j0.707, 0, j0.707, j, j0.707]$$

and

$$H(k) = [3, 2.414, 1, -0.414, -1, -0.414, 1, 2.414]$$

Find the value of y(2), where y(n) is the circular convolution of x(n) and h(n).





## **Lecture 10 Fast Fourier Transform (FFT)**

- Duration: 2 hrs
- Outline:
  - 1. What is FFT?
  - 2. The decomposition-in-time Fast Fourier

Transform algorithm

### Recall DFT and IDFT definition

$$X[k] = \sum_{n=0}^{N-1} x[n]W_N^{nk} \quad 0 \le k \le N-1$$

$$X[k] = \sum_{n=0}^{N-1} x[n] W_N^{nk} \quad 0 \le k \le N-1$$

$$x[n] = \frac{1}{N} \sum_{n=0}^{N-1} X[k] W_N^{-nk} \quad 0 \le n \le N-1$$

$$W_N = e$$

$$W_N = e^{-j\frac{2\pi}{N}}$$

DFT plays an important role in the analysis, design and implementation of the DT signal processing algorithms and systems.

Major reason: existence of efficient algorithms for computing DFT called **FFT** 

## **Direct computation of the DFT**

$$X[k] = \sum_{n=0}^{N-1} x[n]W_N^{nk} \quad 0 \le k \le N-1$$

#### The direct computation of the DFT requires:

- 1. N complex multiplications for each of k
- 2.  $N^2$  complex multiplications for all N points of X(k)
- (N-1) complex summations for each of k
- 4. N(N-1) complex summations for all N points of X(k)
- → FFT optimize computational processes (1) & (2) in different algorithms

## **Decomposition in time FFT (DIT-FFT)**

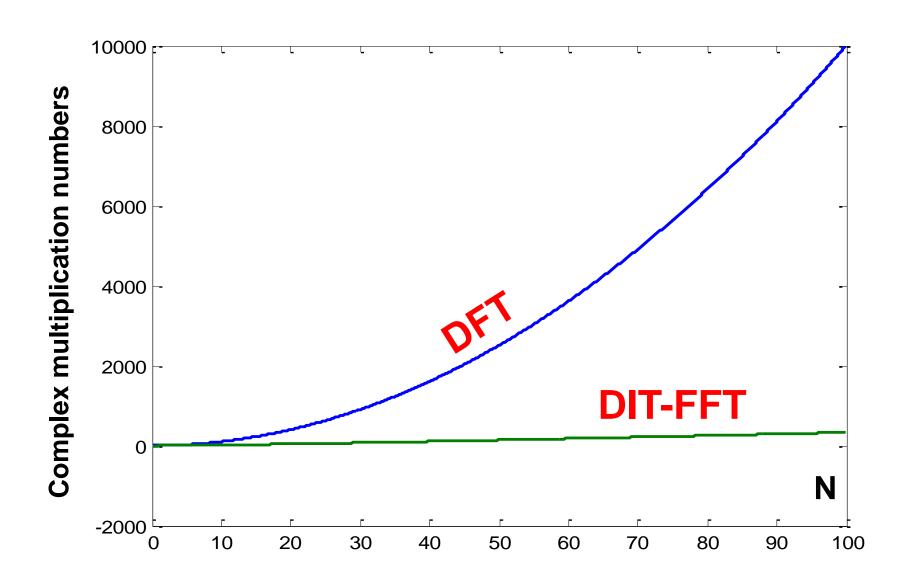
Breaking an N-point DFT into smaller DFTs

→ Fewer calculations with the same output

For example: N is radix-2 number

$$N^2 \Rightarrow \frac{N}{2} \log_2 N$$
 complex multiplications

## **Comparing DFT and FFT efficiency**



## **Computation of DFT**

• Some properties of  $\{W_N^{nk}\}$  can be exploited

$$W_N^{k(N-n)} = W_N^{-kn} = (W_N^{kn})^*$$
, complex conjugate symmetry  $W_N^{kn} = W_N^{k(n+N)} = W_N^{(k+N)n}$ , periodicity in n and k

Other useful properties:

$$\boldsymbol{W}_{N}^{(k+\frac{N}{2})n} = \boldsymbol{W}_{N}^{kn} \boldsymbol{W}_{N}^{\frac{nN}{2}} = \boldsymbol{W}_{N}^{kn} \boldsymbol{e}^{-jn\pi} = \begin{cases} \boldsymbol{W}_{N}^{kn}, & \text{if n even} \\ -\boldsymbol{W}_{N}^{kn}, & \text{if n odd} \end{cases}$$

$$\mathbf{W}_{N}^{2kn} = \mathbf{W}_{\frac{N}{2}}^{kn}$$





- Duration: 2 hrs
- Outline:
  - 1. What is FFT?
  - 2. The decomposition-in-time Fast Fourier Transform algorithm

#### DIT-FFT with N as a 2-radix number

- G(K) is N/2 points DFT of the even numbered data: x(0), x(2), x(4), ..., x(N-2).
- H(k) is the N/2 points DFT of the odd numbered data: x(1), x(3), ..., x(N-1).

$$X[k] = \sum_{n \text{ even}} x[n]W^{kn} + \sum_{n \text{ odd}} x[n]W^{kn}$$

## **DIT-FFT (cont)**

$$X[k] = \sum_{m=0}^{\frac{N}{2}-1} x[2m]W^{2mk} + \sum_{m=0}^{\frac{N}{2}-1} x[2m+1]W^{k(2m+1)}$$

$$= \sum_{m=0}^{\frac{N}{2}-1} x[2m](W^2)^{mk} + W^k \sum_{m=0}^{\frac{N}{2}-1} x[2m+1](W^2)^{mk} = W^2_{N} = (e^{-j2\pi/N})^2 = e^{-j2\pi/(N/2)} = W_{N/2}$$

$$X(k) = G(k) + W_N^k H(k), k = 0,1,..., N-1$$

G(k) and H(k) are of length N/2; X(k) is of length N/G(k)=G(k+N/2) and H(k)=H(k+N/2)

## 8-point FFT

$$X[k]_{8} = G[k]_{4} + W_{8}^{k}H[k]_{4}$$

$$X[0] = G[0] + W_8^0 H[0]$$

$$X[1] = G[1] + W_8^1 H[1]$$

$$X[2] = G[2] + W_8^2 H[2]$$

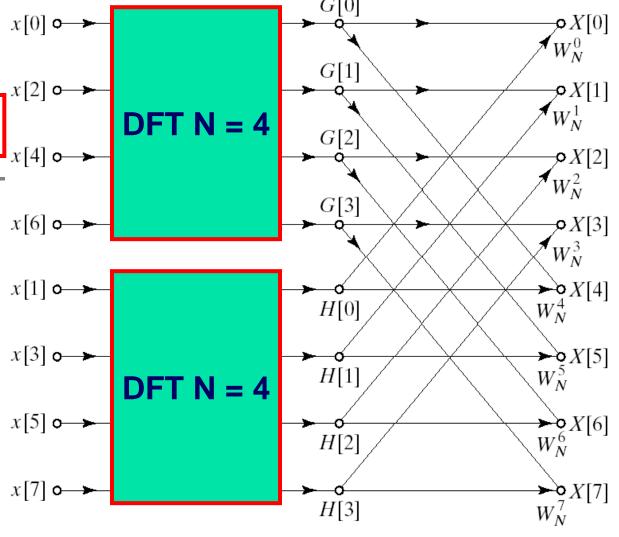
$$X[3] = G[3] + W_8^3 H[3]$$

$$X[4] = G[0] + W_8^4 H[0]$$

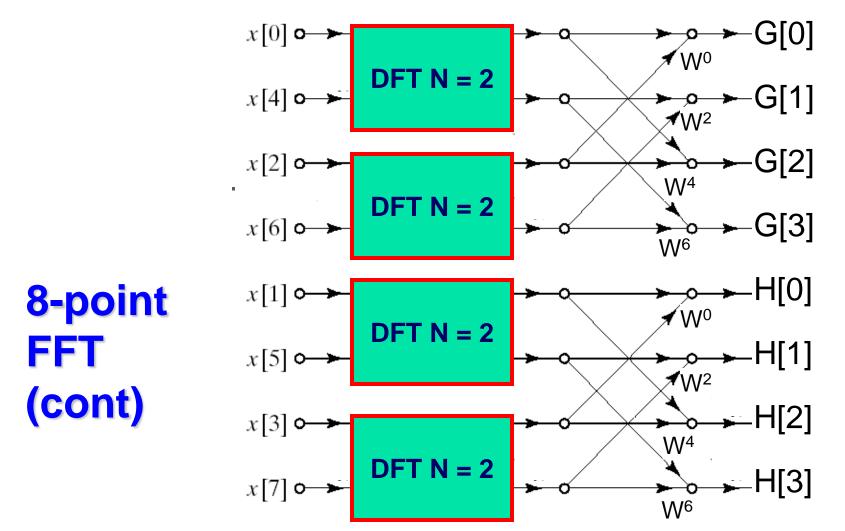
$$X[5] = G[1] + W_8^5 H[1]$$

$$X[6] = G[2] + W_8^6 H[2]$$

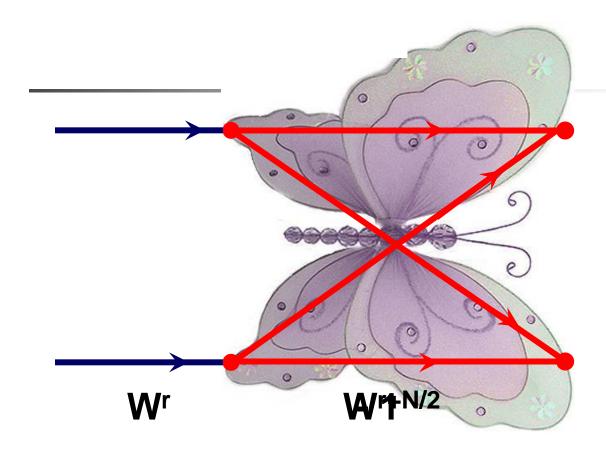
$$X[7] = G[3] + W_8^7 H[3]$$



The new computation counts are reduced

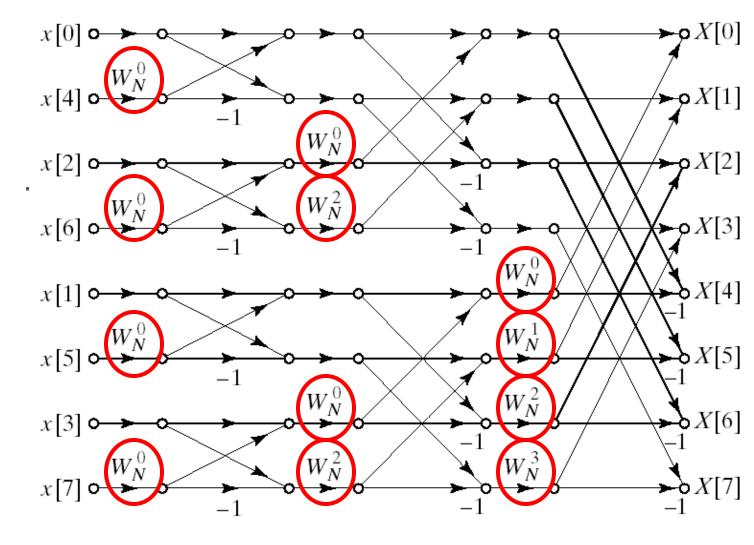


## 2-point FFT



 $W^{r+N/2} = -W^r$ 

# 8-point FFT



Now the overall computation is reduced to:

$$N^2 \Rightarrow \frac{N}{2} \log_2 N$$
 complex multiplications





#### Prob.6

(a) Draw an eight-point DIT FFT signal-flow diagram, and use it to solve for the DFT of the sequence x(n)

$$x(n) = (0.5)^{n} [u(n) - u(n-8)]$$

(b) Use Matlab (fft) to confirm the result of part (a)