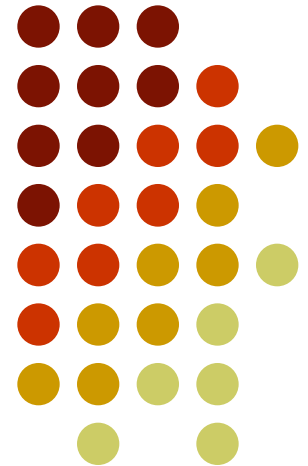


# SISTEME DE CALCUL DEDICATE

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Curs 4



# Outline



- SystemC
  - Dynamic processes
  - Basic channels
  - Evaluate-update channels
- Bibliography



# Dynamic processes

- SystemC 2.1 introduced the concept of ***dynamically spawned processes***
- useful in testbench scenarios
  - to track transaction completion
  - to spawn traffic generators dynamically

```
#define SC_INCLUDE_DYNAMIC_PROCESSES  
#include <systemc>
```



# Dynamic processes

- declare the functions to be spawned as processes
- unlike static processes
  - dynamic processes may have up to eight arguments and a return value
  - the return value will be provided via a reference variable in the actual spawn function

```
// Ordinary function declarations
void inject(void); // no args or return
int count_changes(sc_signal<int>& sig);

// Method function declarations
class TestChan : public sc_module {
    ...
    bool Track(sc_signal<packet>& pkt);
    void Errors(int maxwarn, int maxerr);
    void Speed(void);
    ...
};
```



# Dynamic processes

- define the implementation and register the function with the kernel
  - within an **SC\_THREAD**
  - with restrictions within an **SC\_METHOD**
- syntax to register dynamic processes with void return

```
sc_process_handle hname - // ordinary function
sc_spawn(
    sc_bind(&funcName, ARGS_)//no return value
    ,processName
    ,spawnOptions
);

sc_process_handle hname - // member function
sc_spawn(
    sc_bind(&methName, object, ARGS_)//no return
    ,processName
    ,spawnOptions
);
```



# Dynamic processes

- syntax to register dynamic processes with return values
- object is a reference to the calling module
  - normally just use the C++ keyword **this**

```
sc_process_handle hname - // ordinary function
sc_spawn(
    &returnVar
    ,sc_bind(&funcName, ARGS...)
    ,processName
    ,spawnOptions
);

sc_process_handle hname - // member function
sc_spawn(
    &returnVar
    ,sc_bind(&methodName, object, ARGS ...)
    ,processName
    ,spawnOptions
);
```



# Dynamic processes

- by default, arguments are passed by value
- to pass by reference or by constant reference, a special syntax is required

```
sc_ref(var)    // reference  
sc_cref(var)   // constant reference
```



# Dynamic processes

- spawn options are determined
  - by creating an **sc\_spawn\_option** object
  - invoking one of several methods that set the options

```
sc_spawn_option objname;  
objname.spawn_method();// register as SC_METHOD  
objname.dont_initialize();  
objname.set_sensitivity(event_ptr);  
objname.set_sensitivity(port_ptr);  
objname.set_sensitivity(interface_ptr);  
objname.set_sensitivity(event_finder_ptr);  
objname.set_stack_size(value); // experts only!
```





# Dynamic processes

```
#define SC_INCLUDE_DYNAMIC_PROCESSES
#include <systemc>

...
void spawned_thread() { // This will be spawned
    cout << "INFO: spawned_thread "
          << sc_get_current_process_handle().name()
          << " @ " << sc_time_stamp() << endl;
    wait(10, SC_NS);
    cout << "INFO: Exiting" << endl;
}

void simple_spawn::main_thread() {
    wait(15, SC_NS);
    // Unused handle discarded
    sc_spawn(sc_bind(&spawned_thread));
    cout << "INFO: main_thread " << name()
          << " @ " << sc_time_stamp() << endl;
    wait(15, SC_NS);
    cout << "INFO: main_thread stopping "
          << " @ " << sc_time_stamp() << endl;
}
```



# Dynamic processes

```
// Add "& resume" to sensitivity while suspended
void sc_process_handle::suspend(descend);
void sc_process_handle::resume(descend);

// Ignore sensitivity while disabled
void sc_process_handle::disable(descend);
void sc_process_handle::enable(descend);

// Complete remove process
void sc_process_handle::kill(descend);

// Asynchronously restart a process
void sc_process_handle::reset(descend);

// Reset process on every resumption event
void sc_process_handle::sync_reset_on(descend);
void sc_process_handle::sync_reset_off(descend);

// Throw an exception in the specified process
template<typename T>
void sc_process_handle::throw_it(
    const T&, descend);
```



# Basic channels

- communication of information between concurrent processes is done using
  - events
    - require careful coding
    - use handshake variable
  - ordinary module member data
  - channels - encapsulate complex communications
    - primitive
    - hierarchical



# Basic channels

- primitive channels
  - inherit from the base class **sc\_prim\_channel**
  - also inherit from and implement one or more SystemC interface classes
  - types
    - **sc\_mutex**
    - **sc\_semaphore**
    - **sc\_fifo<T>**



# Basic channels

- mutex is short for *mutually exclusive text*
- **sc\_mutex** class
  - implements **sc\_mutex\_if** interface class
- blocking methods can only be used in **SC\_THREAD** processes

```
sc_mutex NAME;  
  
NAME.lock();    // Lock the mutex,  
                // wait until unlocked if in use  
int NAME.trylock() // Non-blocking, returns success  
  
NAME.unlock();  // Free a previously locked mutex
```



# Basic channels

```
class bus : public sc_module {
    sc_mutex bus_access;
    ...
    void write(int addr, int data) {
        bus_access.lock();
        // perform write
        bus_access.unlock();
    }
    ...
};
```

```
void grab_bus_method() {
    if (bus_access.trylock() == 0) {
        // access bus
        ...
        bus_access.unlock();
    }
}
```



# Basic channels

- **sc\_semaphore** class
  - inherits from and implements the **sc\_semaphore\_if** class

```
sc_semaphore NAME (COUNT);  
  
NAME.wait();           // Lock one semaphore  
                        // Wait until available if in use  
int NAME.trywait()     // Non-blocking, return success  
  
int NAME.get_value()   // Returns available semaphores  
  
NAME.post();           // Free one previously locked  
                        // semaphore
```



# Basic channels

```
class multiport_RAM {
    sc_semaphore read_ports(3);
    sc_semaphore write_ports(2);
    ...
    void read(int addr, int& data) {
        read_ports.wait();
        // perform read
        read_ports.post();
    }
    void write(int addr, const int& data) {
        write_ports.wait();
        // perform write
        write_ports.post();
    }
    ...
}; //endclass
```





# Basic channels

- the most popular channel for modeling at the architectural level is the **sc\_fifo<T>** channel
  - inherits from and implements two interface classes:
    - **sc\_fifo\_in\_if<T>**
    - **sc\_fifo\_out\_if<T>**



# Basic channels

```
sc_fifo<ELEMENT_TYPENAME> NAME(SIZE);

NAME.write(VALUE);
NAME.read(REFERENCE);
__ = NAME.read() /* function style */
if (NAME.nb_read(REFERENCE)) { // Non-blocking
                                // true if success
    ...
}
if (NAME.num_available() == 0)
    wait(NAME.data_written_event());
if (NAME.num_free() == 0)
    next_trigger(NAME.data_read_event());
```



# Basic channels

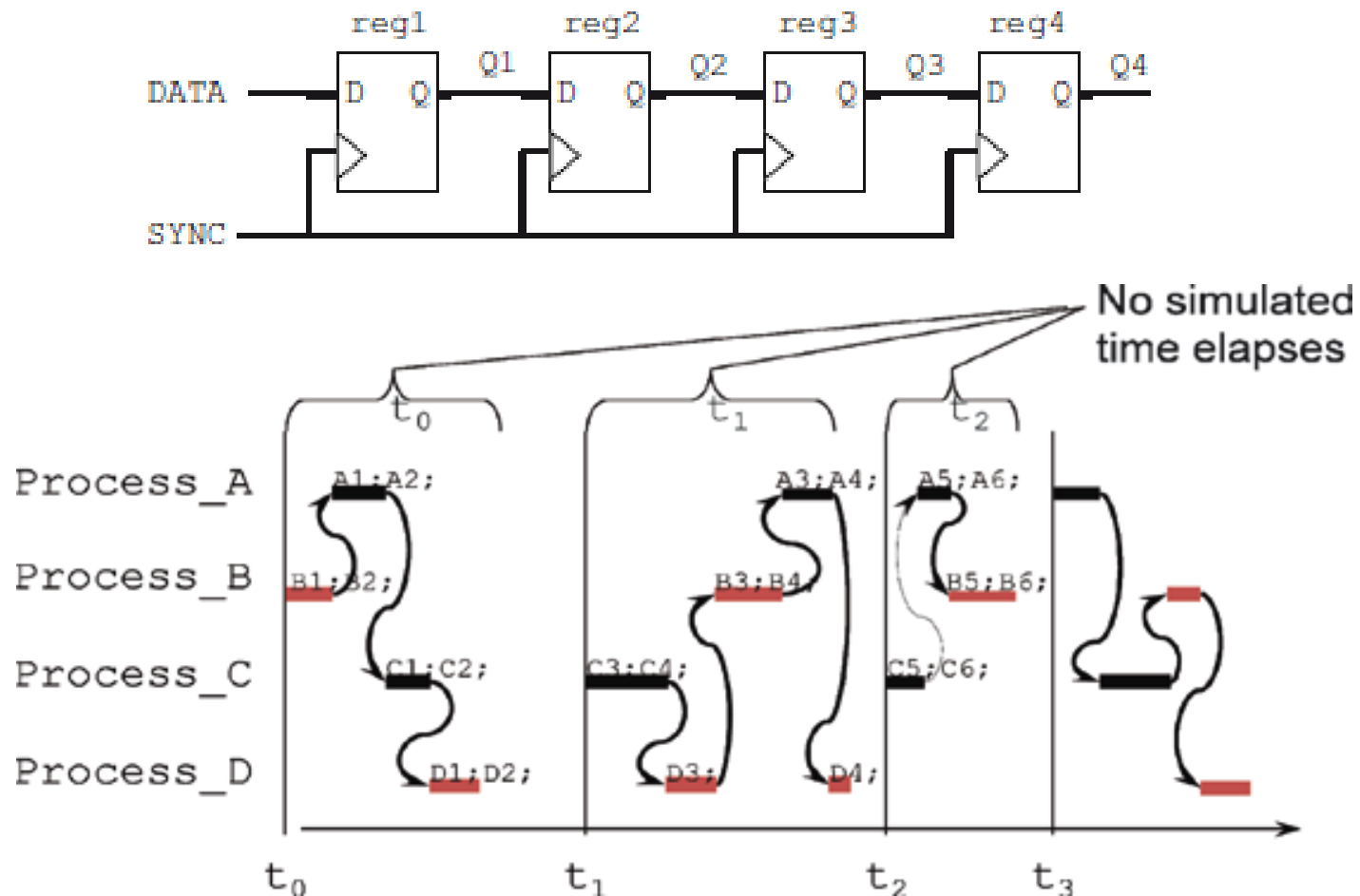
```
SC_MODULE(kahn_ex) {
    ...
    sc_fifo<double> a, b, y;
    ...
};
// Constructor
kahn_ex::kahn_ex() : a(24), b(24), y(48)
{
    ...
}
void kahn_ex::stim_thread() {
    for (int i=0; i!=1024; ++i) {
        a.write(double(rand())/1000);
        b.write(double(rand())/1000);
    }
}
void kahn_ex::addsub_thread() {
    while(true) {
        y.write(kA*a.read() + kB*b.read());
        y.write(kA*a.read() - kB*b.read());
    } //endforever
}
void kahn_ex::monitor_method() {
    cout << y.read() << endl;
}
```

# Evaluate-Update Channels



- electronic hardware
  - behave in a manner approaching instantaneous activity
- electronic signals have
  - a single source (producer)
  - multiple sinks (consumer)
    - it is quite important that all sinks “see” a signal update at the same time

# Evaluate-Update Channels



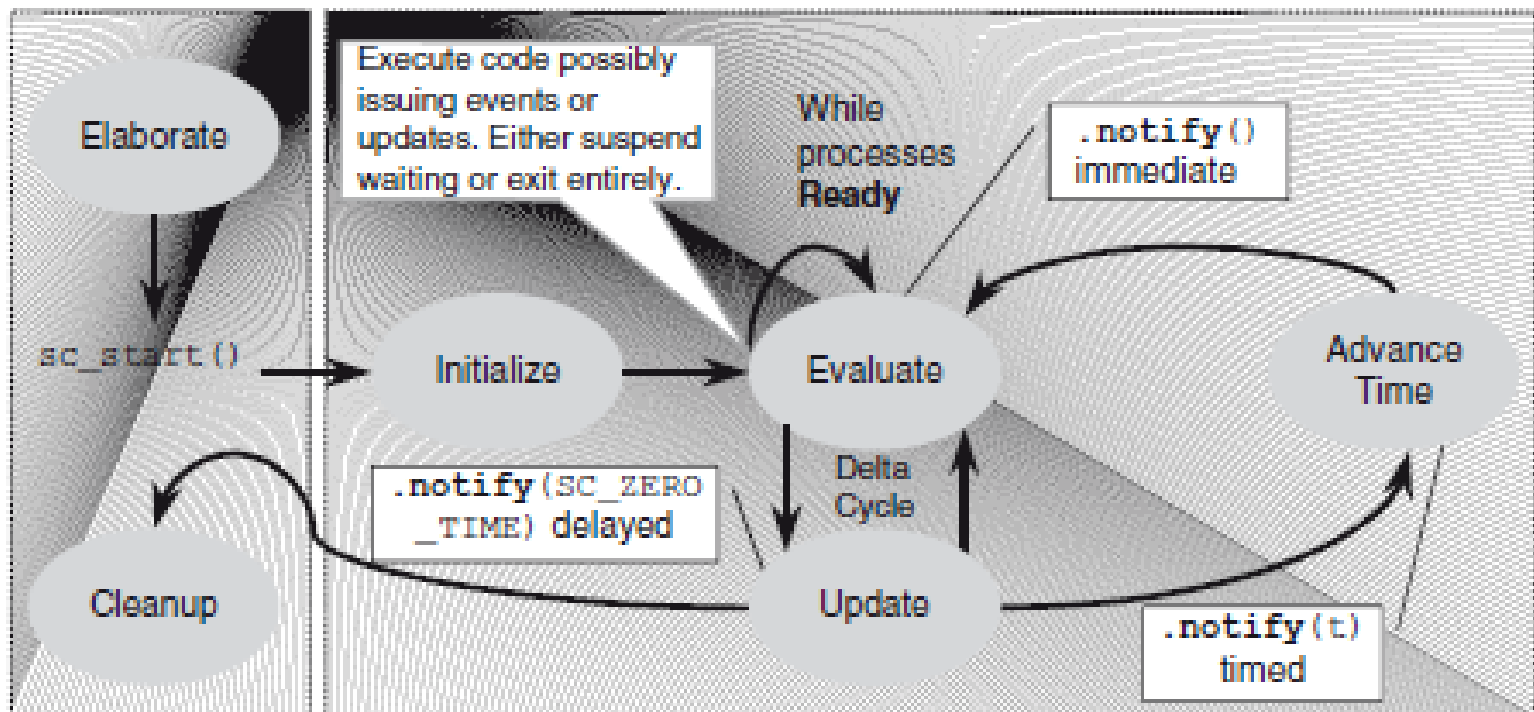


# Evaluate-Update Channels

- evaluate-update paradigm
  - delta-cycle

`sc_main()`

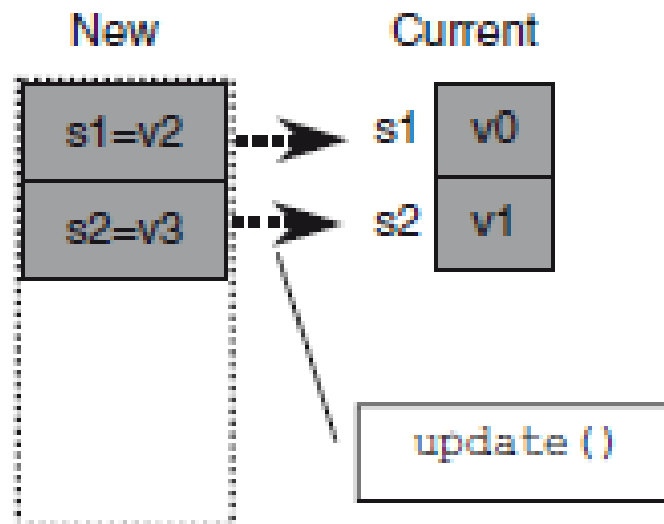
SystemC Simulation Kernel





# Evaluate-Update Channels

- signal channels, use update phase as a point of data synchronization
- to accomplish this synchronization, every channel has two storage locations:
  - the current value and the new value





# Evaluate-Update Channels

- **sc\_signal** $\langle T \rangle$  primitive channel and its close relative, **sc\_buffer** $\langle T \rangle$  both use the evaluate-update paradigm

```
sc_signal<datatype> signame[, signame_i]...; //define
...
signame.write(newvalue);
varname = signame.read();
wait(signame.value_changed_event() | ...);
wait(signame.default_event() | ...);
if (signame.event() == true) {
    // occurred in previous delta-cycle
```





```
// Declare variables
int          count;
string       message_temp;
sc_signal<int> count_sig;
sc_signal<string> message_sig;

cout << "Initialize during 1st delta cycle" << endl;
count_sig.write(10);
message_sig.write("Hello");
count = 11;
message_temp = "Whoa";
cout << "count is " << count << " "
    << "count_sig is " << count_sig << endl
    << "message_temp is '" << message_temp << "' "
    << "message_sig is '" << message_sig << "'"
    << endl << "Waiting" << endl << endl;
wait(SC_ZERO_TIME);

cout << "2nd delta cycle" << endl;
count = 20;
count_sig.write(count);
cout << "count is " << count << ", "
    << "count_sig is " << count_sig << endl
    << "message_temp is '" << message_temp << "', "
    << "message_sig is '" << message_sig << "'"
    << endl << "Waiting" << endl << endl;
wait(SC_ZERO_TIME);

cout << "3rd delta cycle" << endl;
message_sig.write(message_temp = "Rev engines");
cout << "count is " << count << ", "
    << "count_sig is " << count_sig << endl
    << "message_temp is '" << message_temp << "', "
    << "message_sig is '" << message_sig << "'"
    << endl << endl << "Done" << endl;
```

```
Initialize during 1st delta cycle
count is 11, count_sig is 0
message_temp is 'Whoa', message_sig is ''
Waiting

2nd delta cycle
count is 20, count_sig is 10
message_temp is 'Whoa', message_sig is 'Hello'
Waiting

3rd delta cycle
count is 20, count_sig is 20
message_temp is 'Rev engines', message_sig is
'Hello'

Done
```

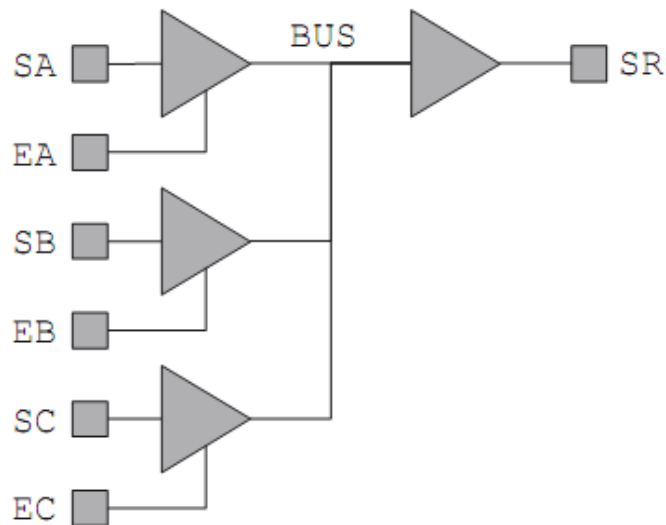


# Evaluate-Update Channels

- multiple writers

```
sc_signal_resolved name;  
sc_signal_rv<WIDTH> name;
```

Multiple Drivers on a Bus



**Table 9.1** Resolution functionality for `sc_signal_resolved`

| <i>A \ B</i> | '0' | '1' | 'X' | 'Z' |
|--------------|-----|-----|-----|-----|
| '0'          | '0' | 'X' | 'X' | '0' |
| '1'          | 'X' | '1' | 'X' | '1' |
| 'X'          | 'X' | 'X' | 'X' | 'X' |
| 'Z'          | '0' | '1' | 'X' | 'Z' |

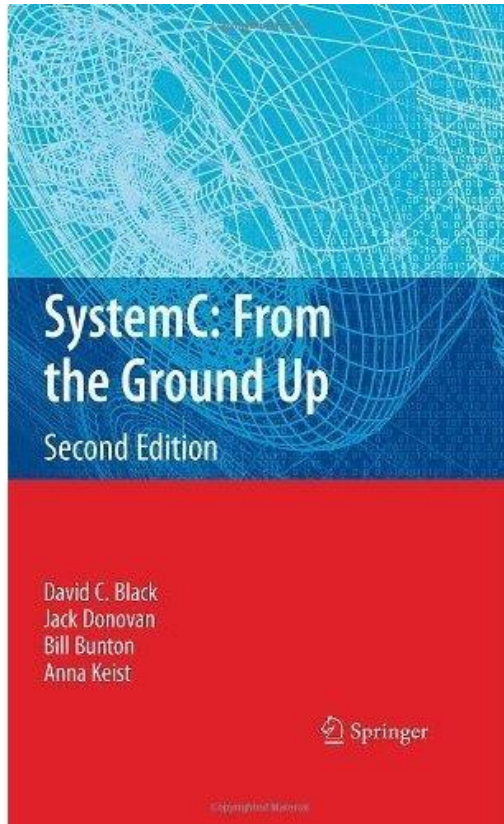


# Evaluate-Update Channels

- template specializations
  - a template specialization occurs when a definition is provided for a specific template value
  - specialized templates
    - `sc_signal<bool>`
    - `sc_signal<sc_logic>`

```
sensitive << signame.posedge_event()  
           << signame.negedge_event();  
wait(signame.posedge_event()  
    | signame.negedge_event());  
if (signame.posedge_event()  
   | signame.negedge_event()) {
```

# Bibliography



- David C. Black, Jack Donovan, Bill Bunton, Anna Keist, ***SystemC: From the Ground Up***, Springer Science+Business Media, LLC 2010
  - “The authors designed this book primarily for the student or engineer new to SystemC. This book’s structure is best appreciated by reading sequentially from beginning to end.”