

3.2.4. Utilizăm metoda tabelului semantic, verificând dacă
 are loc ?

$$2. \vdash (\exists x)(\forall y) p(x, y) \leftrightarrow (\forall y)(\exists x) p(x, y)$$

$$(A \leftrightarrow B) \quad (A \rightarrow B) \wedge (B \rightarrow A)$$

$$\neg \left(\left((\exists x)(\forall y) p(x, y) \rightarrow (\forall y)(\exists x) p(x, y) \right) \wedge \left((\forall y)(\exists x) p(x, y) \rightarrow (\exists x)(\forall y) p(x, y) \right) \right)$$

$$\neg \left((\exists x)(\forall y) p(x, y) \rightarrow (\forall y)(\exists x) p(x, y) \right)$$

$$\neg \left((\forall y)(\exists x) p(x, y) \rightarrow (\exists x)(\forall y) p(x, y) \right)$$

$$\neg \alpha(1)$$

$$(\forall y)(\exists x) p(x, y) (1)$$

$$\neg \left((\exists x)(\forall y) p(x, y) \right) (5)$$

$$\neg \delta(1), \text{ a-cant. exist.}$$

$$(\exists x) p(x, a) (6) \checkmark$$

$$(\forall y)(\exists x) p(x, y) (4)$$

$$\delta(1), \text{ b-cant. exist.}$$

$$p(b, a)$$

$$\neg \delta(1), \text{ b-cant. exist.}$$

$$\neg (\forall y) p(a, y) (1) \checkmark$$

$$\neg (\forall y) p(b, y) (3) \checkmark$$

$$\neg \left((\exists x)(\forall y) p(x, y) \right) (5)$$

$$\delta(2), \text{ c-cant. exist.}$$

$$\neg p(a, c)$$

$$\delta(1), \text{ d-cant. exist.}$$

$$\neg p(b, d)$$

$$\delta(1), \text{ c, d-cant. exist.}$$

$$\neg (\forall y) p(c, y) (5)$$

$$\neg (\forall y) p(d, y) (6) \checkmark$$

$$\neg (\exists x)(\forall y) p(x, y) (5')$$

$$\delta(2), \text{ e-cant. exist.}$$

$$\neg p(c, e)$$

$$\delta(10), \text{ f-cant. exist.}$$

$$\neg p(d, f)$$

\Rightarrow am un ciclu infinit \Rightarrow nu putem deschide tipul finit

5. 2. 8. 2. Sunt unificabili termenii din paranteze urm. ? Dacă da, identificați cel mai general unificator

$$A_1 = P(a, x, f(g(y)))$$

$$A_2 = P(y, f(b), f(z))$$

$$A_3 = P(x, g(f(a)), f(b))$$

$$A_4 = P(f(y), z, z)$$

$$A_5 = P(a, x, f(g(y)))$$

$$A_6 = P(z, h(z, a), f(b), z)$$

P_1 : Skolepi simbol const? $\Delta A(P_1)$

P_2 : Skolepi aritate $\Delta A(3)$

P_3 : \neq
 \hookrightarrow \exists char \neq
 subst. (var)

P_1 : $\Delta A, P$
 P_2 : $\Delta A, 3$

P_1 : $\Delta A(P)$

P_6 : $3 \neq 4 \Rightarrow A_5$ și A_6

P_3 : $\Theta_1 = [x \leftarrow d(y)]$

$$\Theta_1(A_3) = P(d(y), g(d(a)), d(b))$$

$$\Theta_1(A_4) = P(d(y), z, z)$$

$$\Theta_2 = [z \leftarrow g(d(a))]$$

$$\Theta_2(\Theta_1(A_3)) = P(d(y), g(d(a)), d(y))$$

$$\Theta_2(\Theta_1(A_4)) = P(d(y), g(d(a)), g(d(a)))$$

$d \neq g$ - nu \exists substituție
 între termenii simbolului de funcție
 $\Rightarrow A_3, A_4$ - nu sunt unificabili

nu sunt unificabile

$$\lambda_1 = [x \leftarrow a]$$

$$\lambda_{1(A_1)} = P(a, x, f(g(a)))$$

$$\lambda_{1(A_2)} = P(a, f(z), f(z))$$

$$\lambda_2 = [x \leftarrow f(z)]$$

$$\lambda_{2(A_1)} = P(a, f(z), f(g(a)))$$

$$\lambda_2 \circ \lambda_{1(A_2)} = P(a, f(z), f(z))$$

$$\lambda_3 = [z \leftarrow g(a)]$$

$$\lambda_3 \circ \lambda_2 \circ \lambda_{1(A_1)} = P(a, f(g(a)), f(g(a)))$$

$$\lambda_3(\lambda_2(\lambda_{1(A_2)})) = P(a, f(g(a)), f(g(a)))$$



A_1, A_2 unificabile,

$$\text{mgu}(A_1, A_2) = [x \leftarrow a, x \leftarrow f(g(a)), z \leftarrow g(a)]$$

9.2.16 Verificați dacă următoarele formule sunt teoreme utilizând
2. rezoluție generală.

$$U = (\forall x)(\forall y) (P(x,y) \rightarrow P(y,x)) \rightarrow (\forall x)P(x,x)$$

Mod. Res. expr. imp.

$$\neg U \equiv \neg (\forall x)(\forall y) (P(x,y) \rightarrow P(y,x)) \rightarrow (\forall x)P(x,x)$$

$$P_1 \neg \rightarrow \neg U \equiv \neg (\forall x)(\forall y) (\neg P(x,y) \vee P(y,x)) \vee (\forall x)P(x,x)$$

$$P_2 \neg \rightarrow \neg U \equiv (\exists x)(\exists y) (\neg P(x,y) \vee P(y,x)) \wedge (\exists x) \neg P(x,x)$$

P3 Reduc. var. legate a 2 rădăcine

$$\neg U \equiv (\exists x)(\exists y) (\neg P(x,y) \vee P(y,x)) \wedge (\exists z) \neg P(z,z)$$

P4 $\neg \rightarrow \exists x$

$$\neg U \equiv (\exists x) \neg U^x \equiv (\exists z)(\exists y) (\neg P(x,y) \vee P(y,x)) \wedge \neg P(z,z)$$

P5 $\neg \rightarrow \exists y$

$$(\neg U)^y \equiv (\exists x)(\exists y) (\neg P(x,y) \vee P(y,x)) \wedge \neg P(a,a)$$

P6 $\neg \rightarrow \exists x$

$$(\neg U)^z \equiv (\neg P(x,y) \vee P(y,x)) \wedge \neg P(a,a) = (\neg U)^c$$

P7 —

$$\Rightarrow S = \left\{ \underbrace{\neg P(x,y) \vee P(y,x)}_{C_1}, \underbrace{\neg P(a,a)}_{C_2} \right\}$$

$$C_1 = \neg P(x,y) \vee P(y,x)$$

$$C_2 = \neg P(a,a)$$

$$\frac{C_1 \quad C_2}{\neg P(a,a) \wedge P(a,a)} \quad C_3 = \neg P(a,a) \wedge P(a,a) \wedge \neg P(a,a) \wedge P(a,a)$$

$\Rightarrow U$ nu este teoremă