Given the Taylor polynomial expansions of two functions around
$$x = 0$$

$$\frac{1}{1-\Delta x} = 1 + \Delta x + \Delta x^2 + \Delta x^3 + \Theta(\Delta x^4)$$
and
$$\cosh \Delta x = 1 + \frac{\Delta x^2}{21} + \frac{\Delta x^4}{41} + \Theta(\Delta x^6)$$

Calculate their sum and product as well as the exponent that belongs in the O.

For small
$$f(\Delta x) = p(\Delta x) + O(\Delta x^n)$$

$$g(\Delta x) = q(\Delta x) + O(\Delta x^m)$$

$$r = \min(n, m)$$

$$f + g = p + q + O(\Delta x^r)$$

$$f \cdot g = p \cdot q + p \cdot O(\Delta x^m) + q \cdot O(\Delta x^n) + O(\Delta x^{n+m})$$

$$= p \cdot q + O(\Delta x^r)$$

product:
$$f \cdot g = p \cdot q + O(\Delta x^{r})$$

$$\frac{1}{1-\Delta x} \cdot \cosh \Delta x = (1 + \Delta x + \Delta x^{2} + \Delta x^{3}) \cdot (1 + \frac{\Delta x^{2}}{2!} + \frac{\Delta x^{4}}{4!}) + O(\Delta x^{4})$$

$$= 1 + \frac{\Delta x^{2}}{2!} + \frac{\Delta x^{4}}{4!} + \Delta x + \frac{\Delta x^{3}}{2!} + \frac{\Delta x^{5}}{4!} + \Delta x^{2} + \frac{\Delta x^{4}}{2!} + \frac{\Delta x^{6}}{4!}$$

$$\frac{1}{1-\Delta x} \cdot \cosh \Delta x = (1 + \Delta x + \Delta x^{2} + \Delta x^{3}) \cdot (1 + \frac{\Delta x^{2}}{2!} + \frac{\Delta x^{4}}{4!}) + O(\Delta x^{4})$$

$$= 1 + \frac{\Delta x^{2}}{2!} + \frac{\Delta x^{4}}{4!} + \Delta x + \frac{\Delta x^{3}}{2!} + \frac{\Delta x^{5}}{4!} + \Delta x^{2} + \frac{\Delta x^{6}}{4!} + \frac{\Delta x^{6}}{4!}$$

$$= 1 + \frac{3}{2} \Delta x^{2} + \frac{3}{2} \Delta x^{3} + \frac{13}{24} \Delta x^{4} + \frac{13}{24} \Delta x^{5} + \frac{1}{24} \Delta x^{6} + \frac{1}{24} \Delta x^{7} + O(\Delta x^{4})$$

$$= 1 + \frac{3}{2} \Delta x^{2} + \frac{3}{2} \Delta x^{3} + O(\Delta x^{4})$$

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The great Exp challenge:

write a function to calculate ex using its Taylor polynomial approximation expanded around $x_0 = 0$

$$e^{x} \approx T_{n}(x) = 1 + x + \frac{x^{2}}{2!} + \frac{x^{3}}{3!} + \cdots + \frac{x^{n}}{n!}$$

S.t. the relative error of $f=e^x$ and $F=T_n(x)$ is of order ε machine for $x \in [-50, 50]$

Discretization Error

Taylor's Theorem: Let $f(x) \in C^{N+1}[a,b]$ and $x_0 \in [a,b]$, then for all $x \in (a,b)$ there exists a number c = c(x) that lies between x_0 and x such that

$$f(x) = T_N(x) + R_N(x)$$

where $T_N(x)$ is the Taylor polynomial approximation

$$T_N(x) = \sum_{n=0}^{N} \frac{f^{(n)}(x_0) \cdot (x - x_0)^n}{n!}$$

and $R_N(x)$ is the residual (the part of the series we left off)

$$R_N(x) = \frac{f^{(N+1)}(c) \cdot (x - x_0)^{N+1}}{(N+1)!}$$

Start by replacing $x-x_0$ with Δx . The primary idea here is that the residual $R_N(x)$ becomes smaller as $\Delta x \to 0$ (at which point $T_N(x) = f(x_0)$).

$$T_N(x) = \sum_{n=0}^N \frac{f^{(n)}(x_0) \cdot \overrightarrow{\Delta x^n}}{n!}$$

and $R_N(x)$ is the residual (the part of the series we left off)

 $R_N(x) = \frac{f^{(n+1)}(c) \cdot \Delta x^{n+1}}{(n+1)!} \le M \Delta x^{n+1} = O(\Delta x^{n+1})$

where M is an upper bound on

Assume x > 0 and show that the upper bound on the relative error at term n $T_n = \underbrace{1e^x - T_n(x)1}_{1e^x} \underbrace{3P_n(x)}_{1e^x} \underbrace{(e^x = T_n(x) + R_n(x))}_{1e^x}$ a)

is given by
$$\Gamma_n \leq \frac{X^{n+1}}{(n+1)!}$$

s given by $\Gamma_n \leq \frac{\lambda}{(n+1)!}$ $\Gamma_n = \frac{|e^{x} - T_n(x)|}{|e^{x}|} \frac{\beta |R_N(x)|}{|e^{x}|} \leq \frac{|\Delta x^{n+1}|}{|(n+1)!||e^{x}|} \leq \frac{|\Delta x^{n+1}|}{|(n+1)!|}$ Lead as max

upper bound as max val. occurs

when $f^{(n+1)}(c) \geq |e^{x}|$ and we can't have funn) (c) >lext

as c is defined as: $\exists \text{ number } c = c(x) \text{ that lies blw}$ $ce(x,x_0)$ $x_0 \text{ and } x \text{ s.t. } f(x) = T_N(x) + R_N(x)$ $0 < c \le x \Rightarrow 0 < e^c \le e^x$

b) Show that for large x >> 1 and n, $r_n \leq \epsilon_{machine}$ implies that we need approximately $n > e \cdot x$ terms in the series (where $e = \exp(1)$)

We need to show that:

$$\frac{X^{n+1}}{(n+1)!} \leq \mathcal{E}_{\text{machine}}$$

We then have $X^{n+1} < Cn+1)!$

Taking log on both sides:

 $(n+1)\log x < (n+1)\log(n+1) - (n+1)$

$$\log x < \log(n+1)-1$$

 $\log x + 1 < \log(n+1)$

log (e·x) < log (n+1)

(for large n, $\log(n+1) \sim \log n$ $\Rightarrow \log(e \cdot x) < \log n$ $e \cdot x < n$

 $n > e \cdot x$

(4) Storage of #'s is base 4

assuming we have p=3 for the mantissa & the exponent $E \in [-3,3]$, & Sign chirality:

a) how many #'s we can represent w/ this floating point system:

$$N=4$$
, $\rho=3$, $E\in[-3,3]$
(normalized numbers $d_1\neq 0$)
 $d_1=[1,2,3]$ $d_2=d_3=[0,1,2,3]$

$$E = [-3, -2, -1, 0, 1, 2, 3]$$
 $|E| = 7$

$$2 \times 3 \times 4 \times 4 \times 7 + 1 = 673$$

underflow (imit: $1.00 \times 4^{-3} = 0.015625$

overflow (imit: $3.33 \times 4^3 = 213.12$ $ext{mach} = 4^{-1} = 0.25$

c) how many more #'s can we store in N base-pairs (base 4)

vs. N bits C base 2) where the mantissa and exponent are the same relative length (e.g. p=3, $E\in [-3,3]$)

In base 4, our # of representations $w/\rho=3$ and $E\in \mathcal{C}^{-3}$, 3) is

for base 2:

which gives us 2 × 1 × 2 × 2 × 7 + 1 = 57 as the # of representations

Thus the difference is 673 - 57 = 616 more numbers