COSC364 Assignment 2

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Abstract

The following pertains to COSC364-17S1 Assignment 2. This was a joint work by George Drummond and Ryan Cox in accordance with the course requirements of COSC364-17S1 and is a result of equal contribution (50%/50%).

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1 Problem formulation

We wish to formulate an optimization problem for generic values of X, Y and Z (with Y > 3) such that the load on all transit nodes is balanced. Here we outline the governing mathematical expressions concerning the objective function, the decision variables and all constraints.

Notation: $[X] = \{1, 2, 3, ..., X\}.$

1.1 Decision and auxiliary variables

Let u_{ik} be the amount of flow on the link between a given source node S_i and transit node T_k . Likewise let v_{kj} be the amount of flow on the link between a given transit node T_k and destination node D_i .

Therefore, letting x_{ikj} be the part of the demand volume between source node S_i and destination node D_j that is routed through transit node T_k , we achieve

$$u_{ik} + v_{kj} = x_{ikj} \tag{1}$$

 $\forall i \in [X], k \in [Y], j \in [Z].$

Also, the total traffic flow into a transit node is equal to the total traffic flow out of the node, achieving the following balance.

$$\sum_{i=1}^{X} u_{ik} = \sum_{j=1}^{Z} v_{kj} \tag{2}$$

 $\forall k \in [Y].$

1.2 Objective function

The goal of this linear program is to balance the load (total *incoming* traffic) across the transit nodes $T_1, T_2, ..., T_Y$. As the load on a given transit node l_k is simply the sum of the flows from all source nodes to T_k , we achieve

$$l_k = \sum_{i=1}^{X} u_{ik}, \quad \forall k \in [Y]$$

As demonstrated in the small example of figure 1.

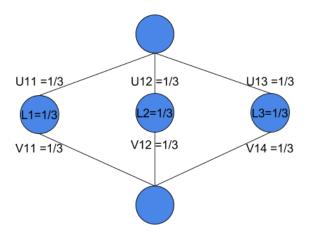


Figure 1: X=1,Y=4,Z=1

At this stage, it may be tempting to define the objective function l as equal to all l_k , $k \in [Y]$ in order to achieve a common minimum load. This however, would actually lead to an infeasible problem for certain cases such as shown in figure 2 below in which all l_k cannot be equal.

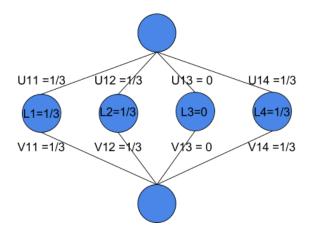


Figure 2: X=1,Y=4,Z=1

To balance this load across the transit nodes, we define our objective function l by

$$l = \max l_k = \max \sum_{i=1}^{X} u_{ik}, \quad \forall k \in [Y]$$

Proceeding in this fashion acts to minimise the greatest load on a transit node. The piecewise linear interpretation of this is as follows.

$$l \ge \sum_{i=1}^{X} u_{ik}, \quad \forall k \in [Y]$$
 (3)

1.3 Demand constraints

We now turn our attention to demand constraints and the global requirement that each demand volume shall be split over exactly three different paths, such that each path gets an equal share of the demand volume.

let w_{ikj} be the indicator variable for the path between source node S_i and destination node D_j through transit node T_k , taking the value 1 if path $S_i - > T_k - > Z_j$ carries data and 0 otherwise. With our 3 path restriction, we achieve

$$\sum_{k=1}^{Y} w_{ikj} = 3, \quad \forall i \in [X], j \in [Z]$$
 (4)

Now to consider the demand volume between source node S_i and destination node D_j . Firstly, we have that this demand volume is simply the sum of the parts of the demand volume through each transit router T_k , $k \in [Y]$.

$$\sum_{k=1}^{Y} x_{ikj} = h_{ij} = i + j$$

However, we also have that x_{ikj} will be non-zero if and only if there is flow on path $S_i - > T_k - > Z_j$ ($w_{ikj} = 1$). As this occurs for precisely 3 x_{ikj} (say $x_{ik_1j}, x_{ik_2j}, x_{ik_3j}$) and the demand h_{ij} is shared equally across the paths that do possess flows, we achieve the following.

$$x_{ikj} = \left\{ \begin{array}{l} \frac{(i+j)}{3}, & \text{if } w_{ikj} = 1\\ 0, & \text{if } w_{ikj} = 0 \end{array} \right\}$$

Or explicitly.

$$x_{ikj} = w_{ikj} \frac{h_{ij}}{3} = \frac{(i+j)}{3} w_{ikj}, \quad \forall i \in [X], k \in [Y], j \in [Z]$$
 (5)

1.4 Additional constraints

Capacity:

 $\forall i \in [X], k \in [Y], j \in [Z], \text{ we have,}$

$$u_{ik} \le c_{ik} \tag{6}$$

$$v_{kj} \le d_{kj} \tag{7}$$

Non-negativity:

 $\forall i \in [X], k \in [Y], j \in [Z], \text{ we have,}$

$$x_{ikj} \ge 0 \tag{8}$$

$$u_{ik} \ge 0 \tag{9}$$

$$v_{kj} \ge 0 \tag{10}$$

$$l \ge 0 \tag{11}$$

1.5 Model summary

The constraints generated by equations 1 through to 11 suffice to fully describe our optimisation problem.

2 Results

3 Appendix

3.1 LP generation source file

```
1 import sys
2
  def print_aux(X,Y,Z):
4
        aux = ""
5
        for i in range (1,X+1):
6
              for k in range (1,Y+1):
                    for j in range (1,Z+1):
8
                         aux += u\{0\}\{1\} + v\{1\}\{2\} - x\{0\}\{1\}\{2\} = 0 , "ormat(i,
      k, j
10
        print (aux, end="")
11
        for k in range (1,Y+1):
13
              mysum = ""
14
              for i in range (1,X+1):
15
                   mysum += "+ u{}{}  ". format(i,k)
16
              for j in range (1,Z+1):
                   mysum += "- v\{\}\{\} ".format(k,j)
18
19
             mysum += "= 0"
20
              print (mysum [2:])
21
22
23
  def print_objective_constraints (X,Y,Z):
        for k in range (1,Y+1):
25
              con = "1 - "
26
              for i in range (1,X+1):
27
                   con += "u{}{} - ".format(i,k)
28
29
              con = con[:-2] + ">= 0"
30
31
              print (con)
32
33
  def print_demand(X,Y,Z):
34
        ''', prints demand constraints '''
35
        demand = ""
        for i in range (1,X+1):
              for j in range (1,Z+1):
38
                   mysum1 = mysum2 = mysum3 = ""
39
                    for k in range (1,Y+1):
40
                         demand += x\{0\}\{1\}\{2\} - \{3\}w\{0\}\{1\}\{2\} = 0 n. format (i, k)
41
      , j, (i+j)/3)
                         mysum1 += w\{\}\{\}\} + .format(i,k,j)
42
                         mysum2 \ += \ "w{0}{1}{2} \ x{0}{1}{2} \ + \ ".format(i,k,j)
43
                         mysum3 += x\{\}\{\}\} + format(i,k,j)
44
45
```

```
46
                   mysum1 = mysum1[:-2] + "="
47
                   mysum2 = mysum2[:-2] + "= "
48
                   mysum3 = mysum3[:-2] + "= "
49
50
                   mysum1 += "3"
51
                   mysum2 += str(i + j)
52
                   mysum3 += str(i + j)
53
54
                   demand += mysum1 + "\n" #+ mysum2 + "\n" + mysum3+ "\n"
55
56
57
        print (demand, end="")
59
60
61
  def print_capp(X,Y,Z):
62
        capp = ""
63
        for i in range (1,X+1):
64
              for k in range (1,Y+1):
65
                   capp += "u{0}{1} - c{0}{1} <= 0\n".format(i,k)
66
67
68
69
        for k in range (1,Y+1):
70
              for j in range (1,Z+1):
71
                   capp += "v\{0\}\{1\} - d\{0\}\{1\} \le 0 n". format(k, j)
72
73
74
        print (capp, end="")
75
76
  def print_integer (X,Y,Z):
78
        integer = ""
79
        for i in range (1,X+1):
80
              for k in range (1,Y+1):
                    for j in range (1,Z+1):
82
                          83
        print(integer, end="")
84
85
  def print_nonneg(X,Y,Z):
86
87
        nonneg = "
        for i in range (1,X+1):
88
              for k in range (1,Y+1):
89
                   nonneg += "0 \le u\{\}\{\} \setminus n" . format(i,k)
90
                   for j in range (1,Z+1):
91
92
                         nonneg += 0 <= x\{\}\{\}\{\} \setminus n \text{ . format } (i, k, j)
93
94
        for k in range (1,Y+1):
95
              for j in range (1,Z+1):
96
                   nonneg += "0 <= v\{\}\{\}\setminus n". format (k, j)
97
```

```
98
         nonneg += "0 <= 1 \ n"
99
100
         print (nonneg, end="")
101
102
103
   def main():
104
         (X,Y,Z) = sys.argv[1:4]
105
         X = int(X)
106
         Y = int(Y)
107
         Z = int(Z)
108
109
         print("X = {}, Y = {}, Z = {}".format(X,Y,Z))
110
         print ("Minimize")
print (" 1")
111
112
         print("Subject to")
113
         print_aux(X,Y,Z)
114
         print_objective\_constraints(X,Y,Z)
115
         print_demand(X, Y, Z)
116
         print_capp(X, Y, Z)
117
         print("Bounds")
118
         \texttt{print\_nonneg}\left(X,Y,Z\right)
119
          print("Integer")
120
         print_integer(X,Y,Z)
121
         print("End")
122
123
124
125
126 main()
```

3.2 LP File for X = 3, Y = 2 and Z = 4

```
1 X = 3, Y = 2, Z = 4
<sup>2</sup> Minimize
з 1
4 Subject to
5 u11 + v11 - x111 = 0
u11 + v12 - x112 = 0
v_1 = v_1 + v_1 + v_1 + v_1 + v_1 + v_1 = 0
u11 + v14 - x114 = 0
u12 + v21 - x121 = 0
u12 + v22 - x122 = 0
u12 + v23 - x123 = 0
u12 + v24 - x124 = 0
u21 + v11 - x211 = 0
u21 + v12 - x212 = 0
u21 + v13 - x213 = 0
u21 + v14 - x214 = 0
u22 + v21 - x221 = 0
u22 + v22 - x222 = 0
u22 + v23 - x223 = 0
u22 + v24 - x224 = 0
u31 + v11 - x311 = 0
u31 + v12 - x312 = 0
u31 + v13 - x313 = 0
u31 + v14 - x314 = 0
u32 + v21 - x321 = 0
u32 + v22 - x322 = 0
u32 + v23 - x323 = 0
u32 + v24 - x324 = 0
  u11 + u21 + u31 - v11 - v12 - v13 - v14 = 0
30 u12 + u22 + u32 - v21 - v22 - v23 - v24 = 0
u_{11} - u_{11} - u_{21} - u_{31} >= 0
u_{12} 1 - u_{12} - u_{22} - u_{32} > = 0
x111 - 0.666666666666666668 w 111 = 0
x121 - 0.666666666666666668 w121 = 0
35 \text{ w}111 + \text{w}121 = 3
36 \times 112 - 1.0 \times 112 = 0
x122 - 1.0 w122 = 0
38 \text{ w}112 + \text{w}122 = 3
x113 - 1.3333333333333333333333 = 0
40 x123 - 1.333333333333333333333 = 0
w113 + w123 = 3
x114 - 1.666666666666667 w114 = 0
x124 - 1.666666666666667 w124 = 0
44 \text{ w}114 + \text{w}124 = 3
45 \times 211 - 1.0 \times 211 = 0
46 \times 221 - 1.0 \times 221 = 0
47 \text{ w} 211 + \text{w} 221 = 3
x212 - 1.33333333333333333333333 = 0
49 \quad x222 \quad - \quad 1.333333333333333333333222 \quad = \quad 0
50 \text{ w} 212 + \text{w} 222 = 3
```

```
x213 - 1.6666666666666667 w213 = 0
x^{223} - 1.6666666666666667 w^{223} = 0
53 \text{ w} 213 + \text{w} 223 = 3
54 \times 214 - 2.0 \times 214 = 0
55 \text{ x} 224 - 2.0 \text{ w} 224 = 0
56 \text{ w} 214 + \text{w} 224 = 3
w311 + w321 = 3
x312 - 1.666666666666667 w312 = 0
x322 - 1.6666666666666667 \times 322 = 0
62 \text{ w} 312 + \text{w} 322 = 3
63 \times 313 - 2.0 \times 313 = 0
64 \times 323 - 2.0 \times 323 = 0
65 \text{ w}313 + \text{w}323 = 3
x314 - 2.3333333333333333333333333 = 0
  x324 - 2.33333333333333335 w324 = 0
  w314 + w324 = 3
u11 - c11 <= 0
u12 - c12 <= 0
u21 - c21 <= 0
u22 - c22 <= 0
u31 - c31 <= 0
u32 - c32 <= 0
v11 - d11 <= 0
v12 - d12 <= 0
v13 - d13 <= 0
v14 - d14 <= 0
v21 - d21 <= 0
v22 - d22 <= 0
v23 - d23 <= 0
v24 - d24 <= 0
83 Bounds
84 \ 0 <= u11
85 \ 0 <= x111
86\ 0 <= x112
87 \ 0 <= x113
88\ 0 <= x114
0 <= u12
90 0 <= x121
91\ 0 <= x122
92\ 0 <= x123
93 0 <= x124
94 0 <= u21
95 \ 0 <= x211
96\ 0 <= x212
97 \ 0 <= x213
98 0 <= x214
99 0 \le u22
100 \ 0 <= x221
101 0 <= x222
102 \ 0 <= x223
```

```
103 \ 0 <= x224
104 0 <= u31
105 \ 0 <= x311
_{106} 0 <= _{\rm x312}
107 0 <= x313
108 \ 0 <= x314
109 0 <= u32
110 \ 0 <= x321
111 0 <= x322
112 0 <= x323
113 \ 0 <= x324
114 0 <= v11
115 0 <= v12
116 0 <= v13
117 0 <= v14
118 \ 0 <= v21
119 0 <= v22
120 \ 0 <= v23
121 0 <= v24
122 0 <= 1
123 Integer
124\ w111
125 \text{ w} 112
126 \text{ w} 113
_{127}\ w114
128 \text{ } \text{w} 121
129 \text{ } \text{w} 122
130 \text{ w} 123
131 \text{ w} 124
132 w211
133 w212
134 \text{ } \text{w}213
135 w214
136 \text{ } \text{w} 221
138 \text{ } \text{w}223
139 \text{ } \text{w} 224
_{140}\ w311
_{141}\ w312
142 \text{ } \text{w}313
143 \text{ } \text{w}314
144 \text{ } \text{w}321
145 \text{ } \text{w}322
146\ w323
147 \text{ } \text{w}324
148 End
```