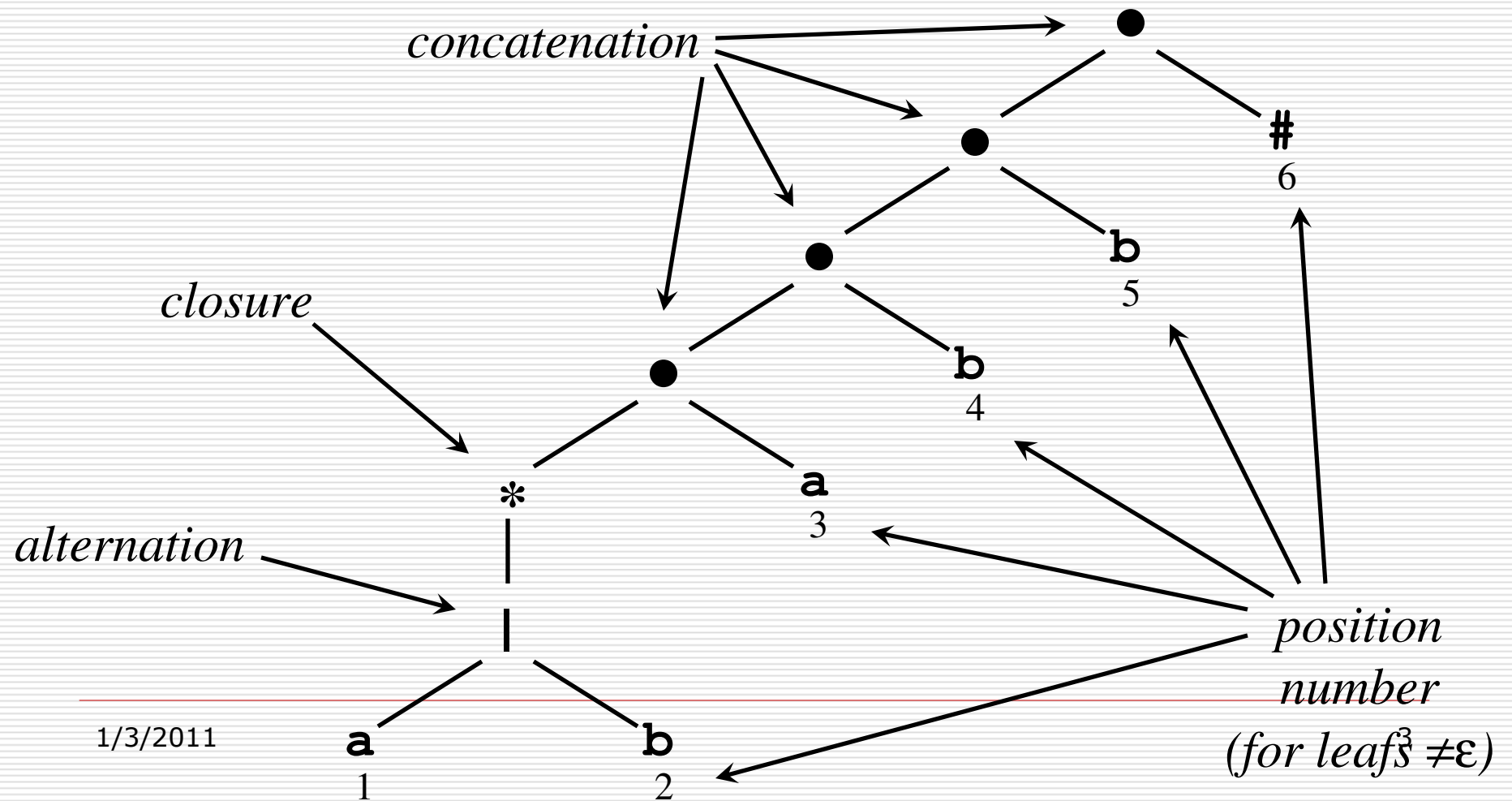


Direct Conversion of RE to DFA

From Regular Expression to DFA Directly (Algorithm)

- Augment the regular expression r with a special end symbol $\#$ to make accepting states important
 $r \rightarrow (r)\#$
- Construct a syntax tree for $(r)\#$
Alphabets \rightarrow leaf node
Operators \rightarrow inner node
- Number each alphabet including $\#$
- Traverse the tree to construct functions *nullable*, *firstpos*, *lastpos*, and *followpos*

Syntax Tree of $(a|b)^*abb\#$



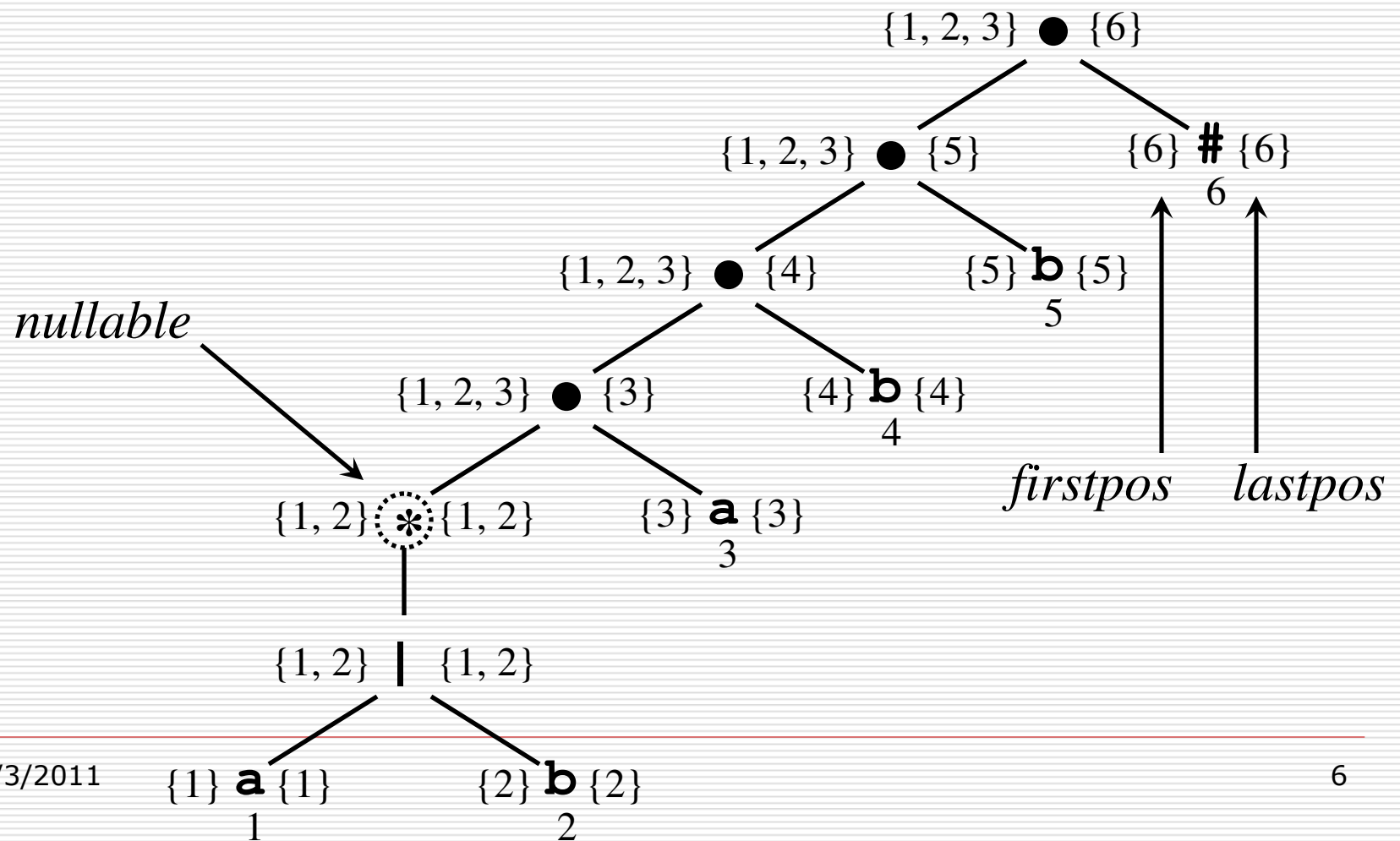
Annotating the Tree

- ***nullable(n)***: the subtree at node n generates languages including the empty string
- ***firstpos(n)***: set of positions that can match the first symbol of a string generated by the subtree at node n
- ***lastpos(n)***: the set of positions that can match the last symbol of a string generated by the subtree at node n
- ***followpos(i)***: the set of positions that can follow position i in the tree

Annotating the Tree Cont...

| Node n | $nullable(n)$ | $firstpos(n)$ | $lastpos(n)$ |
|---|---|---|--|
| Leaf ε | true | \emptyset | \emptyset |
| Leaf i | false | $\{i\}$ | $\{i\}$ |
| $\begin{array}{c} \\ / \quad \backslash \\ c_1 \quad c_2 \end{array}$ | $nullable(c_1)$ or $nullable(c_2)$ | $firstpos(c_1)$ \cup $firstpos(c_2)$ | $lastpos(c_1)$ \cup $lastpos(c_2)$ |
| $\begin{array}{c} \bullet \\ / \quad \backslash \\ c_1 \quad c_2 \end{array}$ | $nullable(c_1)$ and $nullable(c_2)$ | if $nullable(c_1)$ then $firstpos(c_1)$ \cup $firstpos(c_2)$ else $firstpos(c_1)$ | if $nullable(c_2)$ then $lastpos(c_1)$ \cup $lastpos(c_2)$ else $lastpos(c_2)$ |
| $\begin{array}{c} * \\ \\ 1/3/2011 \quad c_1 \end{array}$ | true | $firstpos(c_1)$ | $lastpos(c_1)$ |

Syntax Tree of $(a|b)^*abb\#$



followpos

```

for each node  $n$  in the tree do
    if  $n$  is a cat-node with left child  $c_1$  and right child  $c_2$  then
        for each  $i$  in  $lastpos(c_1)$  do
             $followpos(i) := followpos(i) \cup firstpos(c_2)$ 
        end do
    else if  $n$  is a star-node
        for each  $i$  in  $lastpos(n)$  do
             $followpos(i) := followpos(i) \cup firstpos(n)$ 
        end do
    end if
end do

```

Algorithm

$s_0 := \text{firstpos}(\text{root})$ where root is the root of the syntax tree

$Dstates := \{s_0\}$ and is unmarked

while there is an unmarked state T in $Dstates$ **do**

 mark T

for each input symbol $a \in \Sigma$ **do**

 let U be the set of positions that are in $\text{followpos}(p)$

 for some position p in T ,

 such that the symbol at position p is a

if U is not empty and not in $Dstates$ **then**

 add U as an unmarked state to $Dstates$

end if

$Dtran[T, a] := U$

end do

end do

Example

