

NFA to DFA

Subset Construction

- The *subset construction algorithm* converts an NFA into a DFA using:
 - $\epsilon\text{-closure}(s)$ = Set of NFA states reachable from NFA state s on ϵ transition alone (T)
 - $\epsilon\text{-closure}(T)$ = Set of NFA states reachable from NFA state s in T on ϵ transition alone
 - $\text{move}(T, a)$ = Set of NFA states to which there is a transition on input symbol a from some NFA state s in T
- The algorithm produces:
 - $Dstates$ is the set of states of the new DFA consisting of sets of states of the NFA
 - $Dtran$ is the transition table of the new DFA

ϵ - Closure

push all states in T onto *stack*;

initialize ϵ -closure(T) to T ;

while *stack* is not empty **do begin**

 pop t , the top element, off of *stack*;

for each state u with an edge from t to u labeled ϵ *do*

if u is not in ϵ -closure(T) **do begin**

 add u to ϵ -closure(T);

 push u onto *stack*;

end

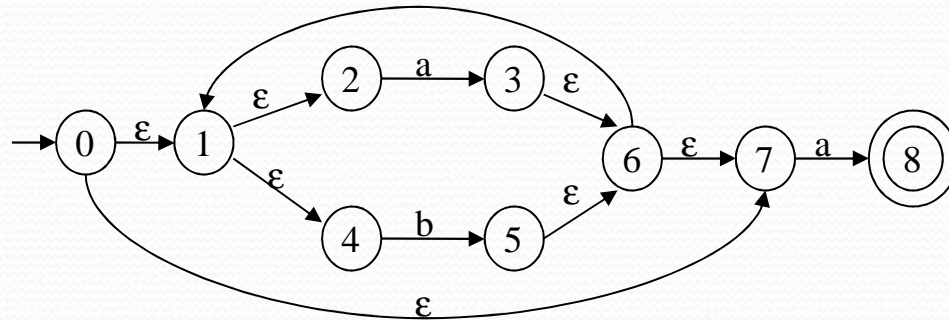
end

Subset Construction Algorithm

Initially, $\epsilon\text{-closure}(s_0)$ is the only state in $Dstates$ and it is unmarked

```
while there is an unmarked state  $T$  in  $Dstates$  do  
    mark  $T$   
    for each input symbol  $a \in \Sigma$  do  
         $U := \epsilon\text{-closure}(\text{move}(T, a))$   
        if  $U$  is not in  $Dstates$  then  
            add  $U$  as an unmarked state to  $Dstates$   
        end if  
         $Dtran[T, a] := U$   
    end do  
end do
```

Example



$$A = \epsilon\text{-closure}(\{0\}) = \{0, 1, 2, 4, 7\}$$

A into *Dstates* as an unmarked state

↓ mark A

$$\epsilon\text{-closure}(\text{move}(A, a)) = \epsilon\text{-closure}(\{3, 8\}) = \{1, 2, 3, 4, 6, 7, 8\} = B$$

B into *Dstates*

$$\epsilon\text{-closure}(\text{move}(A, b)) = \epsilon\text{-closure}(\{5\}) = \{1, 2, 4, 5, 6, 7\} = C$$

C into *Dstates*

$$\text{Dtran}[A, a] \leftarrow B \quad \text{Dtran}[A, b] \leftarrow C$$

↓ mark B

$$\epsilon\text{-closure}(\text{move}(B, a)) = \epsilon\text{-closure}(\{3, 8\}) = \{1, 2, 3, 4, 6, 7, 8\} = B$$

$$\epsilon\text{-closure}(\text{move}(B, b)) = \epsilon\text{-closure}(\{5\}) = \{1, 2, 4, 5, 6, 7\} = C$$

$$\text{Dtran}[B, a] \leftarrow B \quad \text{Dtran}[B, b] \leftarrow C$$

↓ mark C

$$\epsilon\text{-closure}(\text{move}(C, a)) = \epsilon\text{-closure}(\{3, 8\}) = \{1, 2, 3, 4, 6, 7, 8\} = B$$

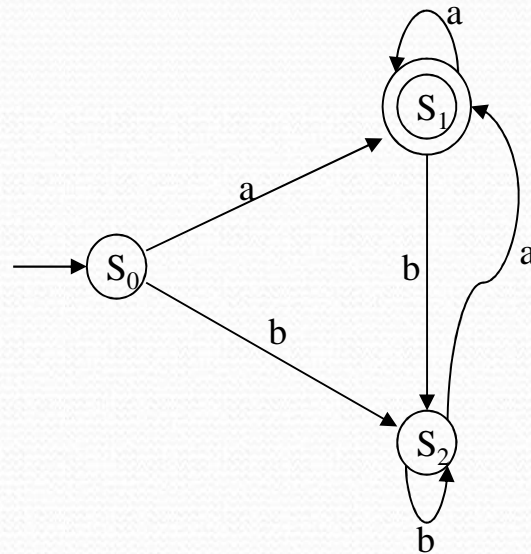
$$\epsilon\text{-closure}(\text{move}(C, b)) = \epsilon\text{-closure}(\{5\}) = \{1, 2, 4, 5, 6, 7\} = C$$

$$\text{Dtran}[C, a] \leftarrow B \quad \text{Dtran}[C, b] \leftarrow C$$

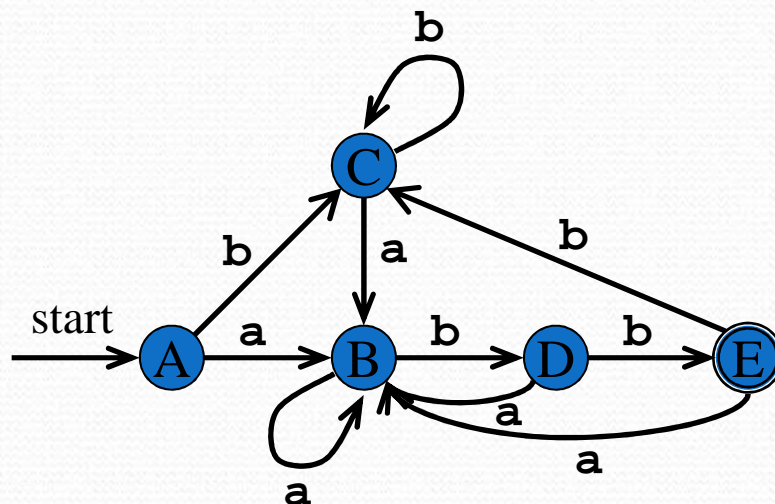
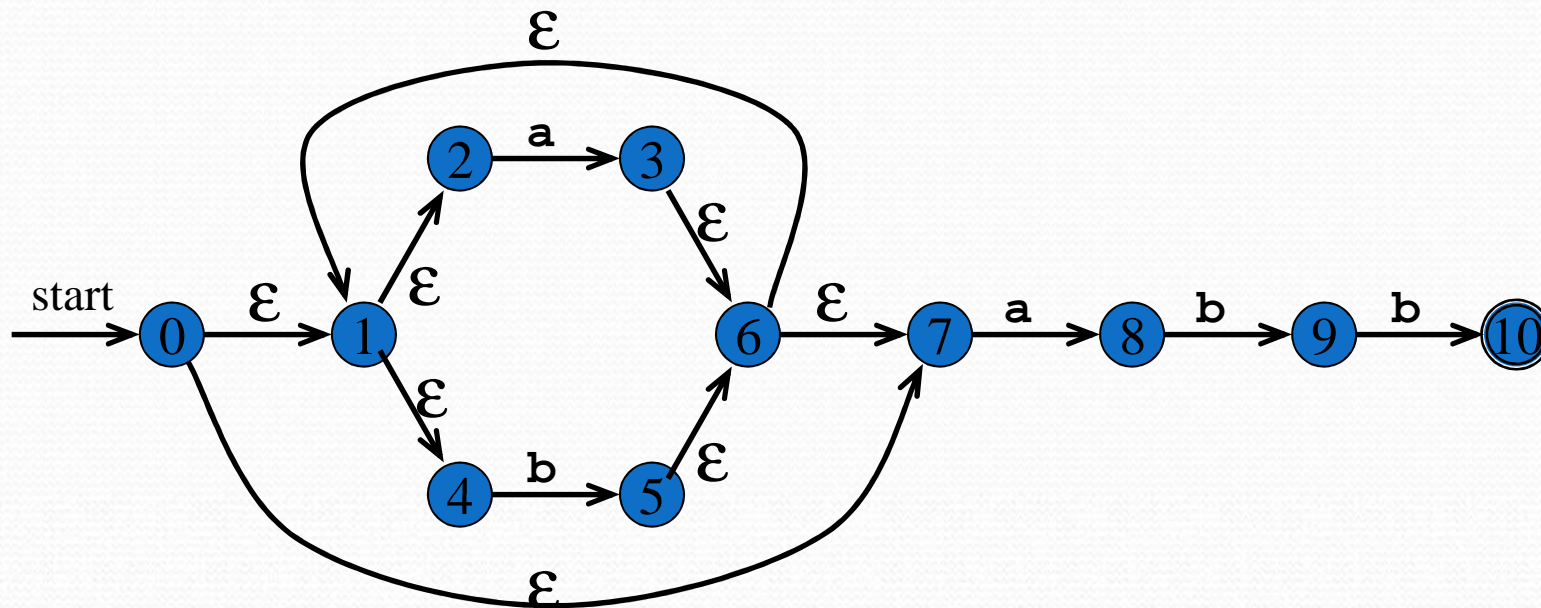
Example Cont...

A is the start state of DFA since 0 is a member of $A = \{0, 1, 2, 4, 7\}$

B is an accepting state of DFA since 8 is a member of $B = \{1, 2, 3, 4, 6, 7, 8\}$



Example 2



Dstates

$A = \{0, 1, 2, 4, 7\}$

$B = \{1, 2, 3, 4, 6, 7, 8\}$

$C = \{1, 2, 4, 5, 6, 7\}$

$D = \{1, 2, 4, 5, 6, 7, 9\}$

$E = \{1, 2, 4, 5, 6, 7, 10\}$

Simulation of an NFA

$S := \epsilon\text{-closure}(\{s_0\})$

$a := \text{nextchar}()$

while $S \neq \emptyset$ **do**

$S := \epsilon\text{-closure}(\text{move}(S, a))$

$a := \text{nextchar}()$

end do

if $S \cap F \neq \emptyset$ **then**

return “yes”

else return “no”