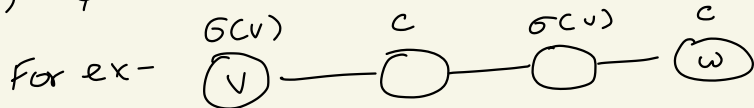


Flip Dynamics

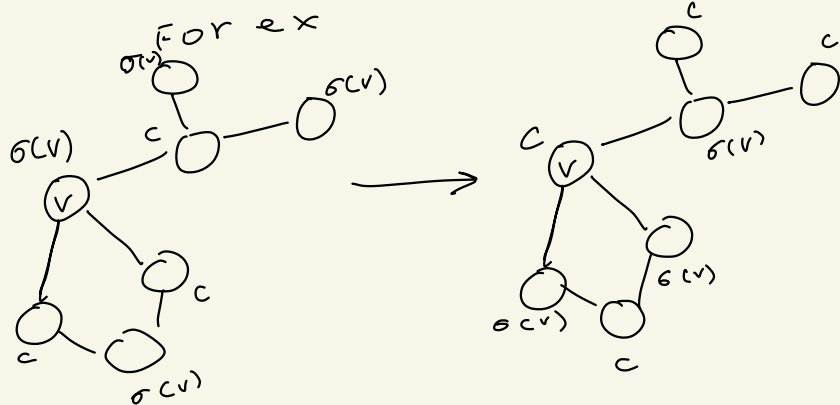
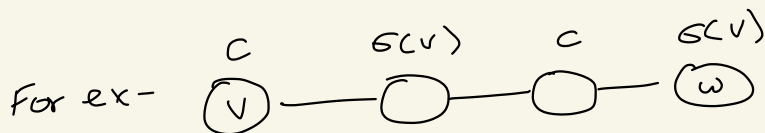
1. Choose $v \in V$, $c \in C$

2. Let $S_\sigma(v, c) = \{w \mid \text{there is an alternating path b/w } v, w \text{ using colors } c \text{ and } \sigma(v)\}$
 let $S_\sigma(v, \sigma(v)) = \emptyset$



Let $\alpha = |S_\sigma(v, c)|$

With prob $\frac{p_\alpha}{\alpha}$ flip $S_\sigma(v, c)$
 by interchanging c and $\sigma(v)$



Why $\frac{p_\alpha}{\alpha}$? Each $u \in S_\sigma(v, c)$ gives rise to the same cluster.
 \Rightarrow Each cluster is flipped with prob p_α .

Irreducible: When $c \notin \{\sigma(v) \mid v \in N(v)\}$
(when $|C| \geq \Delta + 2$) equivalent to Glauber Dynamics.
(which is irreducible).

Aperiodic: Self loops when
 $c = \sigma(v)$

Symmetric: Flipping $S_\sigma(v, c)$
results in coloring σ , Flipping
 $S_{\sigma'}(v, \sigma(v))$ recovers σ .

Both occur with prob P_α

where $\alpha = |S_\sigma(v, c)| = |S_{\sigma'}(v, \sigma(v))|$

Prob Values -

$$P_\alpha = \max\left(0, \frac{13}{42} - \frac{1}{7} \left[1 + \frac{1}{2} + \dots + \frac{1}{\alpha-2}\right]\right)$$

$$\Rightarrow P_1 = 1, P_2 = \frac{13}{42}, P_3 = \frac{1}{6}$$

$$P_4 = \frac{2}{21}, P_5 = \frac{1}{21}, P_6 = \frac{1}{84}$$

$$P_\alpha = 0, \alpha > 6.$$

Why these? : In the analysis we will need -

$$i) 2(i-1)p_i + p_{2i+1} \leq \frac{2}{3}$$

$$ii) (j-1)(p_j - p_{j+1}) + i(p_i - p_{i+1}) \leq \frac{5}{6}$$

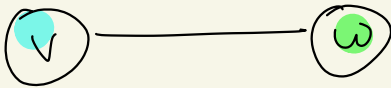
$$iii) ip_i \leq p_1 = 1$$

$$iv) (i-1)p_i \leq 2p_3 = \frac{1}{3} \quad v) (i-c)p_i \leq \frac{1}{4} \quad \forall c \geq 2$$

Path Coupling Analysis

$\phi(\sigma, \tau) = \text{Hamming Distance b/w } \sigma \text{ and } \tau$

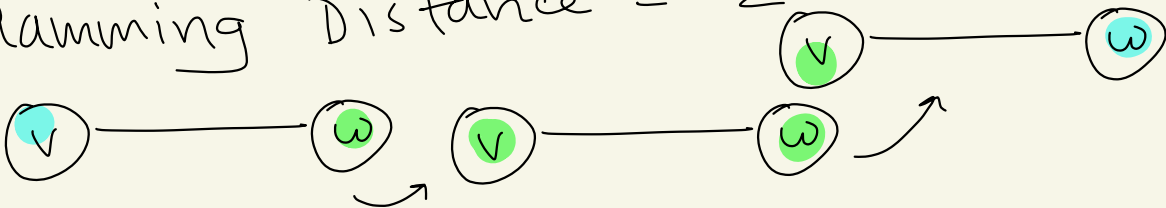
Consider two states that differs at v and w



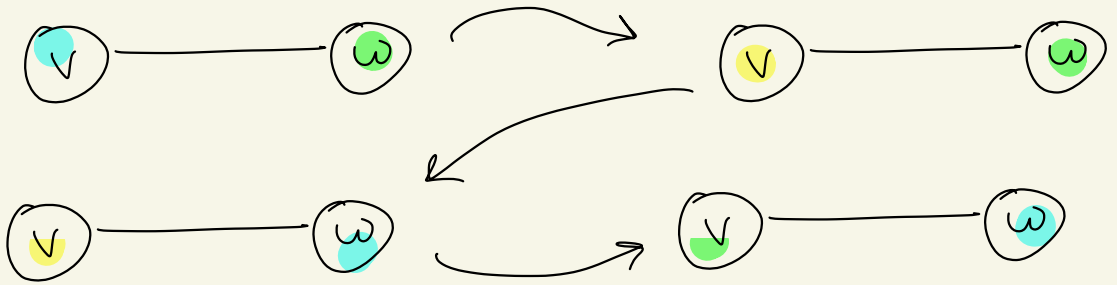
as -



Hamming Distance = 2



Distance in space of valid colorings = 3



Thus we define state space to be C^V .

- If the chain is started at a proper coloring it stays a proper coloring
 - If the chain starts from an improper, it'll eventually reach a proper
- (Doesn't this break irreducibility?)

Simplifying

→ Each $y \in S_G(x, c)$ gives rise to the same cluster.

Thus it is enough to reason about clusters instead of vertices.

When can clusters differ?

Suppose G, T differ only at v .

Let $w \in N(v)$

The clusters that are different

i) $S_G(v, c)$ and $S_T(v, c)$

↓
Has $G(v)$ and c

↓
Has $T(v)$ and c

ii) $S_G(w, \tau(v))$ and $S_\tau(w, \sigma(v))$

↙
 v "blocks" the cluster
 in σ , whereas
 it doesn't in τ

For all other clusters, they
 can either -

- i) Be re-indexed to one
 of the two above cases
- or ii) Be exactly the same in
 both G and τ .

For the 2nd case, we use
 the identity coupling and so
 distance does not change.

Analysis set up

Let

Γ_c = neighbours of v colored c

$$D_c = \{ S_\sigma(v, c),$$

↓

All differing
clusters
associated
with c .

$$\{ S_\tau(v, c), \\ S_\sigma(w, \tau(v)), \\ S_\tau(w, \sigma(v)) \}_{w \in \Gamma_c}$$

Note: $D_c \cap D_c = \emptyset$ except

for $D_\sigma(v) \cap D_\tau(v)$

(which may or may not)

For ex -

$$D_{\delta(v)} = S_{\delta(v)}(v, \delta(v))$$

$$S_{\tau}(v, \delta(v))$$

$$S_{\delta}(w_i, \tau(v))$$

$$S_{\tau}(w_i, \sigma(v))$$

$$\phi \leftarrow w_i \in \delta(v)$$

$$D_{\tau(v)} = S_{\tau(v)}(v, \tau(v))$$

$$S_{\tau}(v, \tau(v))$$

$$S_{\delta}(w_i, \tau(v))$$

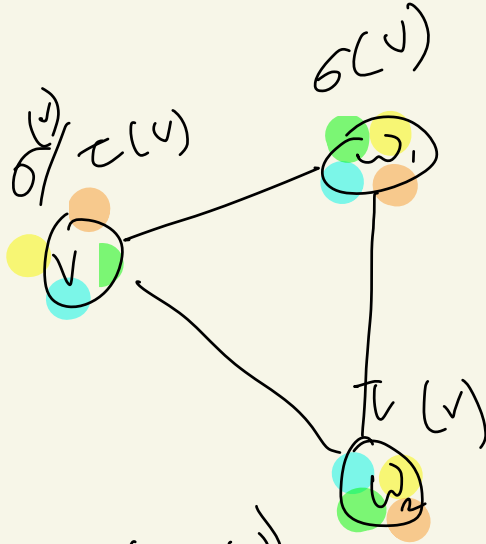
$$S_{\tau}(w_i, \sigma(v))$$

$$w_i \in \delta_{\tau(v)}$$

$$\rightarrow S_{\delta}(w_1, \tau(v)) = S_{\tau}(w_2, \sigma(v))$$

$$\text{cyan circle} \rightarrow S_{\delta}(w_1, \tau(v))$$

$$\text{yellow circle} \rightarrow S_{\tau}(w_2, \sigma(v))$$



$$\text{green circle} \rightarrow S_{\delta}(v, \tau(v))$$

$$\text{orange circle} \rightarrow S_{\tau}(w_2, \sigma(v))$$

This case will be dealt with separately.

Our coupling f is such that

$$(i) \text{ If } S \notin \bigcup_c D_c, \quad f(S) = S$$

\downarrow
Identity.

$$(ii) \text{ If } S \in D_c, \quad f(S) \in D_c$$

$$\text{Thus } \mathbb{E}(\Delta \Phi) =$$

$$\sum_c \frac{1}{k} \mathbb{E}[\Delta_c \Phi]$$

\downarrow

expected change in Φ given both chains
Flip a cluster in D_c .

equiv to one chain flips
a cluster in D_c by (ii)

Let $\delta_c = |\Gamma_c|$ = no. of c colored neighbours.

Main Lemma

a) if $\delta_c = 0$, $\mathbb{E}[\Delta_{\Phi_c} \Phi] \leq -1$

b) if $\delta_c > 0$, $\mathbb{E}[\Delta_{\Phi_c} \Phi] \leq \frac{11}{6} \delta_c - 1$

Main Thm: Flip Dynamics mixes in $O(nk \log n)$ given $k > \frac{11\Delta}{6}$

Proof (given the lemma):

colors c , $\delta_c = 0$ (i.e. doesn't appear in neighborhood of v)

$$= k - \left| \{ c' \mid \delta_{c'} > 0 \} \right|$$

↓
colors in at least
one neighbour.

$$= k - \sum_{\substack{c' \\ \delta_{c'} > 0}} 1$$

Now $\sum_c \mathbb{E}(\Delta_{0_c} \phi)$

$$\leq \sum_{\substack{c' \\ \delta_{c'} = 0}} (-1) + \sum_{\substack{c' \\ \delta_{c'} > 0}} \frac{11}{6} \delta_{c'} - 1$$

$$= \left(k - \sum_{\substack{c' \\ \delta_{c'} > 0}} 1 \right) (-1) + \sum_{\substack{c' \\ \delta_{c'} > 0}} \frac{11}{6} \delta_{c'} - 1$$

$$= -k + \sum_{\substack{c' \\ \delta_{c'} > 0}} \frac{11}{6} \delta_{c'}$$

$$= -k + \frac{11}{6} \delta$$

\therefore If $k > \frac{11}{6} \Delta$, this will be negative.

Proof of Main Lemma

a) If $\delta_c = 0$, $\Rightarrow D_c = \{ S_\sigma(v, c) \}$
 $\{ S_\tau(v, c) \}$
 Also $S_\sigma(v, c) = S_\tau(v, c)$
 $= \{v\}$.

Thus the only possible coupling is identity. That reduces distance by 1.

$$\Rightarrow \mathbb{E}[\Delta_{D_c} \phi] = -1$$

b) $S_\sigma(v, c) = \bigcup_i S_\tau(w_i, \overset{\text{c colored neighbors}}{\sigma(v)})$
 $\bigcup \{v\}$
 \downarrow
 v "blocks" the clusters.

Similarly -

$$S_{\tau}(v, c) = \bigcup_j S_{\sigma}(\omega_j, \tau(v)) \cup \{v\}.$$

Let

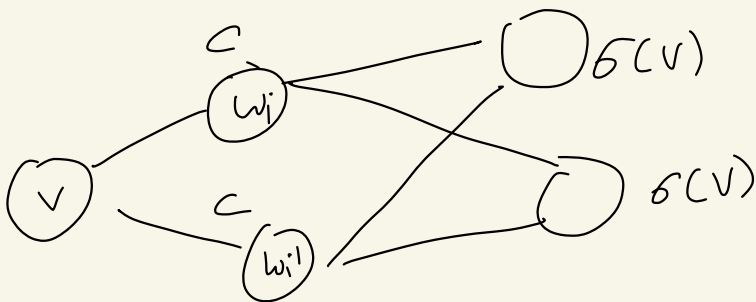
$$a_i = |S_{\tau}(\omega_i, \sigma(v))|$$

$$A = |S_{\sigma}(v, c)|$$

$$b_j = |S_{\sigma}(\omega_j, \tau(v))|$$

$$B = |S_{\tau}(v, c)|$$

Technicality: It may be that
 $S_{\tau}(\omega_i, \sigma(v)) = S_{\tau}(\omega_{i+1}, \sigma(v))$



For such cases just set $a_i = 0$. i.e. each such cluster is indexed by exactly 1 neighbor.

Coupling—

Let

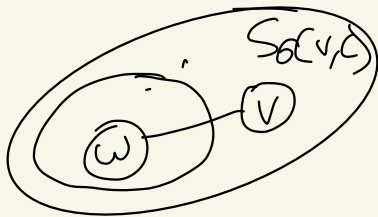
$$a_{\max} = \max_i a_i$$

$$i_{\max} = \underset{i}{\operatorname{argmax}} a_i$$

$$b_{\max} = \max_j b_j$$

$$j_{\max} = \underset{j}{\operatorname{argmax}} b_j$$

Couple:



1. With prob P_A flip $S_\delta(v, c)$
and $S_\tau(\omega_{i_{\max}}, \delta v)$

2. With prob P_B flip $S_\tau(v, c)$
and $S_\delta(\omega_{j_{\max}}, \tau(v))$

3. Let

$$q_l = \begin{cases} P_{al} - P_A, & \text{if } l = i_{\max} \\ P_{al}, & \text{o/w} \end{cases}$$

$$q'_l = \begin{cases} P_{bl} - P_B, & \text{if } l = j_{\max} \\ P_{bl}, & \text{o/w} \end{cases}$$

For each w_x ,

3a. With prob $\min(q_x, q'_x)$

flip $S_\tau(w_x, \sigma(v))$ and

$S_\sigma(w_x, \tau(v))$

3b. With prob $q_x - \min(q_x, q'_x)$

flip $S_\tau(w_x, \sigma(v))$

3c. With prob $q'_x - \min(q_x, q'_x)$

flip $S_\sigma(w_x, \tau(v))$

Marginal : a) $S_{\sigma}(v, c)$ is
flipped with prob P_A
 $S_{\tau}(v, c)$ is
flipped with prob P_B
in step 1/2

b) $S_{\sigma}(w_{i\max}, \tau(v))$
WLOG assume $q_e \leq q_e^i$

flipped with prob P_A in
step 1.

Then with prob $P_{A_e} - P_A$ in
step 2 a

\therefore totally with prob P_{A_e} .

$$c) S_T(w_{\max}, G_{CV})$$

WLOG assume $q_L < q_L'$

\Rightarrow flipped with prob P_B in
step 2.

with prob q_L in step
3a.

with prob $P_{bL} - P_B - q_L$

\therefore totally with prob

$$\begin{aligned} P_B + q_L + P_{bL} - P_B - q_L \\ = P_{bL}. \end{aligned}$$

Similarly for the others.
Thus each cluster is flipped

with prob P_d .

Coupling Analysis

Step 1: The nodes shared
b/w $S_0(v, c)$ and $S_T(w_{i_{\max}}, a_{\max})$
are identical.

\Rightarrow Distance increases by
at most $A - a_{\max} - 1$
 $\hookrightarrow V$ was
already matched

Step 2: Similarly,
 $B - b_{\max} - 1$.

Step 3a: $S_\tau(w_u, \sigma(v))$

and $S_\sigma(w_u, \tau(v))$ are disjoint
apart from w_u . \therefore Dist

increases by $a_u + b_u - 1$

\downarrow
 w_u counted
twice.

Step 3b: Dist increases by a_u

3c: Dist increases by b_u

$$\text{let } f(w_u) = \min(q_u, q'_u) \cdot (a_u + b_u - 1) \\ + (q_u - \min) a_u + (q'_u - \min) b_u$$

$$= q_1 a_1 + q'_1 b_1 - \min$$

$$\text{Now } \mathbb{E}[\Delta_{D_c} \Phi] \leq p_A (A - a_{\max} - 1) + p_B (B - b_{\max} - 1) + \sum_{\omega_1 \in \Gamma_c} f(\omega_1)$$

Cases

$$\underline{S_c = 1}$$

$$\text{Here } A = a_1 + 1, B = b_1 + 1$$

$$\text{wlog: } q_1 \geq q'_1$$

$$\Rightarrow f(\omega_1) = q_1 a_1 + q'_1 b_1 - q'_1$$

$$\Rightarrow \mathbb{E}(\Delta_{D_c} \Phi) \leq p_A (a_1 + 1 - a_1 - 1) + p_B (b_1 + 1 - b_1 - 1) + q_1 a_1 + q'_1 b_1 - q'_1$$

substituting $q_i = p_{a_i} - p_A$

$$q'_i = p_{b_i} - p_B$$

$$\Rightarrow \mathbb{E}(\Delta_{PC}\phi) \leq (p_{a_i} - p_A) a_i + (p_{b_i} - p_B) b_i - (p_{b_i} - p_B)$$

$$\leq a_i (p_{a_i} - p_A) + (b_i - 1) (p_{b_i} - p_B)$$

$$\leq a_i (p_{a_i} - p_{a_{i+1}}) + (b_i - 1) (p_{b_i} - p_{b_{i+1}})$$

Recall "key property"

$$i(p_i - p_{i+1}) + (j-1)(p_j - p_{j+1}) \leq \frac{5}{6}$$

$$\Rightarrow \mathbb{E}(\Delta_{D_c} \Phi) \leq \frac{5}{6}$$

$$\underline{\underline{\delta_c = 2}}$$

Claim: When $\delta_c = 2$, $\mathbb{E}[\Delta_{D_c} \Phi]$ is maximized (worst case) for $a_1 = a_2 = a \leq 3$ and $b_1 = b_2 = b = 1$

Given the claim -

$$f(\omega_1) = a p_a + b p_b - p_a$$

$$f(\omega_2) = a(p_a - p_A) + b(p_b - p_B) - (p_a - p_A)$$

$$\begin{aligned} \mathbb{E}[\Delta_{D_c}\phi] &\leq P_A(A - a - 1) + P_B(B - b - 1) \\ &\quad + (a-1)P_a + bP_b \\ &\quad + (a-1)(P_a - P_A) + b(P_b - P_B) \end{aligned}$$

Substituting $A = 2a + 1$
 $B = 2b + 1 = 3$

$$\begin{aligned} &\leq P_{2a+1}(2a+1 - a - 1) + P_{2b+1}(2b - b) \\ &\quad + (a-1)(2P_a - P_{2a+1}) + b(2P_b - P_B) \end{aligned}$$

$$\leq 2(a-1)P_a + P_{2a+1} + 2bP_b$$

$$\leq 2(a-1)P_a + P_{2a+1} + 2 \left[\because P_i = 1 \right]$$

First "key property"

$$2(i-1)p_i + p_{2i+1} \leq \frac{2}{3}$$

$$\Rightarrow \mathbb{E}[\Delta_{DC}\Phi] \leq \frac{2}{3} + 2$$

$$= \frac{2}{3} + 3 - 1$$

$$= \frac{4+18}{6} - 1$$

$$= 2\left(\frac{2+9}{6}\right) - 1$$

$$= 8_c\left(\frac{11}{6}\right) - 1$$

$$\underline{\delta_c > 2}$$

$$\text{let } g(w_d) = a_d p_{a_d} + b_d p_{b_d} - \min(p_{a_d}, p_{b_d})$$

$$\text{Note: } g(w_d) = f(w_d)$$

$$\text{for } d \neq i_{\max} \neq j_{\max}.$$

$$\text{If } d = i_{\max} = j_{\max}$$

$$f(w_d) = a_{\max}(p_{a_{\max}} - p_A) + b_{\max}(p_{b_{\max}} - p_B) - \min(p_{a_{\max}} - p_A, p_{b_{\max}} - p_B)$$

$$\leq a_{\max}(p_{\max} - p_A) + b_{\max}(p_{\max} - p_B) \\ - \min(p_{\max}, p_{\max}) \\ + p_A + p_B$$

$$= g(w_e) + p_A(1 - a_{\max}) \\ + p_B(1 - b_{\max})$$

If $i_{\max} \neq j_{\max}$

$$f(w_{i_{\max}}) + f(w_{j_{\max}}) =$$

$$a_{\max}(p_{\max} - p_A) + b_i p_{b_i}$$

$$- \min(p_{\max} - p_A, p_{b_i}) + a_j p_{a_j}$$

$$+ b_{\max}(p_{b_{\max}} - p_B) - \min(p_{a_j}, p_{b_{\max}} - p_B)$$

$$\begin{aligned}
&\leq a_{\max} p_{\max} + b_i p_{oi} \\
&\quad - \min(p_{\max}, p_{oi}) + p_A \\
&\quad + b_{\max} p_{\max} + a_j p_{aj} \\
&\quad - \min(p_{\max}, p_{aj}) + p_B \\
&= a_{\max} p_A - b_{\max} p_B
\end{aligned}$$

$$\begin{aligned}
&\leq g(\omega_{i\max}) + g(\omega_{j\max}) \\
&\quad + p_A(1 - a_{\max}) \\
&\quad + p_B(1 - b_{\max})
\end{aligned}$$

$$\Rightarrow \sum_l f(\omega_l) \leq \sum_l g(\omega_l) +$$

$$P_A(1 - a_{\max})$$

$$+ P_B(1 - b_{\max})$$

$$\Rightarrow \mathbb{E}(\Delta_{D_c} \Phi) \leq P_A(A - a_{\max} - 1)$$

$$+ P_B(B - b_{\max} - 1) + \sum_l g(\omega_l)$$

$$+ P_A(1 - a_{\max}) + P_B(1 - b_{\max})$$

$$\leq P_A(A - 2a_{\max}) + P_B(B - 2b_{\max})$$

$$+ \sum_l g(\omega_l)$$

Recall "key property"

$$(i-m) p_i < \frac{1}{4} \quad \text{for } m \geq 2$$

$$\text{here } 2a_{\max} \geq 2$$

$$\text{hence } p_A(A - 2a_{\max}) < \frac{1}{4}$$

$$p_B(B - 2b_{\max}) < \frac{1}{4}$$

$$g(w_e) = a_e p_{ae} + b_e p_{be} - \min(p_{ae}, p_{be})$$

wlog: $a_e \leq b_e$ and hence

$$p_{ae} \geq p_{be}$$

$$\Rightarrow g(w_e) = a_e p_{ae} + (b_e - 1) p_{be}$$

$$\left. \begin{aligned} a_l p_l &\leq \overbrace{1 \cdot p_1}^1 \\ (b_l - 1) p_{bl} &\leq \underbrace{2 p_3}_{1/3} \end{aligned} \right\} \text{"key properties"}$$

$$\Rightarrow g(\omega_l) \leq \frac{4}{3}$$

$$\Rightarrow \mathbb{E}(\Delta_{D_c} \phi) \leq \frac{1}{4} + \frac{1}{4} + \sum_l \frac{4}{3}$$

$$\leq \frac{1}{2} + \delta_c \cdot \frac{4}{3}$$

$$\leq \frac{11}{6} \delta_c - 1, \quad \delta_c > 2$$