

# CS120 ProblemSet1

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## 1 Schema Refinement

(a)

No, because the primary key is  $\{EMP\_ID, REC\_ID\}$ . For example, in a functional dependency like  $EMP\_ID \rightarrow ROLE$ ,  $EMP\_ID$  is not a super key of this table.

(b)

- Functional dependencies:
  - $EMP\_ID \rightarrow EMP\_NAME, ROLE$
  - $REC\_ID \rightarrow REC\_NAME, REC\_ARTIST, PRICE$
- Step 1 (using  $EMP\_ID \rightarrow EMP\_NAME, ROLE$ ):
  - $R1(EMP\_ID, EMP\_NAME, ROLE)$
  - $R2(EMP\_ID, REC\_ID, REC\_NAME, REC\_ARTIST, PRICE)$
- Step 2 (in  $R2$ , using  $REC\_ID \rightarrow REC\_NAME, REC\_ARTIST, PRICE$ ):
  - $R3(REC\_ID, REC\_NAME, REC\_ARTIST, PRICE)$
  - $R4(EMP\_ID, REC\_ID)$
- Final BCNF relations:
  - $Employee(EMP\_ID, EMP\_NAME, ROLE)$
  - $Record(REC\_ID, REC\_NAME, REC\_ARTIST, PRICE)$
  - $Sales(EMP\_ID, REC\_ID)$

123 emp_id	AZ emp_name	AZ role
1	Bob	Cashier
2	Lupe	Manager
3	Sara	Asst. Manager
4	Geoff	Security
5	Dakota	Cashier

Figure 1: Employee

123 rec_id	AZ rec_name	AZ rec_artist	123 price
1	Rocky Mountain High	John Denver	10
5	Back to the Country	Loretta Lynn	15
3	Jolene	Dolly Parton	12
2	Rocky Mountain High	John Denver	10
4	22	Taylor Swift	10

Figure 2: Record

123 emp_id	123 rec_id
1	1
2	5
3	3
4	2
5	4

Figure 3: Sales

## 2 Schema Design

(a)

### Venue

- Attributes: `venue_id`, `name`, `address`
- Primary Key: `venue_id`
- Constraints: `name` and `address` must be non-null
- Justification: Each venue has a unique identifier and stores basic information about the location.

### Section

- Attributes: `section_id`, `venue_id`, `name`, `capacity`
- Primary Key: `section_id`
- Constraints: `capacity` must be strictly positive; `venue_id` is a foreign key referencing Venue
- Justification: Each section belongs to exactly one venue, and its seating capacity must be valid.

### Performance

- Attributes: `performance_id`, `artist`, `show_date`, `venue_id`
- Primary Key: `performance_id`
- Constraints: `show_date` must be non-null; `venue_id` is a foreign key referencing Venue
- Justification: Each performance is uniquely identified and must take place at exactly one venue.

### TicketPrice

- Attributes: `performance_id`, `section_id`, `price`
- Primary Key: Composite key (`performance_id`, `section_id`)
- Constraints: `price` must be non-negative; `performance_id` references Performance; `section_id` references Section
- Justification: The price of a ticket depends on both the performance and the section within the venue.

(b)

- If the ticket price for a section is not fixed but rather dynamic (for example, early-bird tickets versus regular tickets), this schema would not be able to support it.
- If a performance involves multiple artists performing together, and each artist has a separate ticket price, this schema would also not be compatible.

### 3 Relational Algebra and SQL

(a) Find the wids of workers who made  $\geq 1$  type of toys in category 'LEGO'

**Relational Algebra:**

$$\pi_{wid}(\sigma_{category='LEGO'}(Catalog))$$

**SQL:**

```
SELECT DISTINCT wid
FROM Catalog
WHERE category = 'LEGO';
```

(b) Find the wids of workers who produced only 'LEGO'

**Relational Algebra:**

$$\pi_{wid}(Catalog) - \pi_{wid}(\sigma_{category \neq 'LEGO'}(Catalog))$$

**SQL:**

```
SELECT wid
FROM Catalog
GROUP BY wid
HAVING COUNT(DISTINCT category) = 1
      AND MIN(category) = 'LEGO';
```

(c) Find the names of workers and the customers their toys are sold to

**Relational Algebra:**

$$\pi_{wname, cname}(Worker \bowtie Catalog \bowtie SalesRecord \bowtie Customer)$$

**SQL:**

```
SELECT DISTINCT w.wname, c.cname
FROM Worker w
JOIN Catalog ca ON w.wid = ca.wid
JOIN SalesRecord s ON ca.iid = s.iid
JOIN Customer c ON s.cid = c.cid;
```

(d) Find each customer that bought at least 10 different toys

**Relational Algebra:**

Not expressible in relational algebra without aggregation

**SQL:**

```
SELECT c.cname
FROM Customer c
JOIN SalesRecord s ON c.cid = s.cid
GROUP BY c.cid, c.cname
HAVING COUNT(DISTINCT s.iid) >= 10;
```

(e) Find the name(s) of the most expensive toy(s) bought by Mrs Claus

**Relational Algebra:**

Not expressible in relational algebra without the MAX operator.

**SQL:**

```
-- All toys she bought before
WITH ClausToys AS (
    SELECT ca.iname, s.unit_price
    FROM SalesRecord s
    JOIN Catalog ca ON s.iid = ca.iid
    JOIN Customer c ON s.cid = c.cid
    WHERE c.cname = 'Mrs Claus'
)

-- Find the most expensive ones
SELECT iname
FROM ClausToys
WHERE unit_price = (SELECT MAX(unit_price) FROM ClausToys);
```