

Structs

Create our own data types via a `struct`. Syntax:

```
typedef struct
{
    string name;
    string number;
} person;
```

typedefin islenmesi:

```
// Implements a phone book with structs

#include <cs50.h>
#include <stdio.h>
#include <string.h>

typedef struct
{
    string name;
    string number;
} person;

int main(void)
{
    person people[3];

    people[0].name = "Yuliia";
    people[0].number = "+1-617-495-1000";

    people[1].name = "David";
    people[1].number = "+1-617-495-1000";

    people[2].name = "John";
    people[2].number = "+1-949-468-2750";

    // Search for name
    string name = get_string("Name: ");
    for (int i = 0; i < 3; i++)
    {
        if (strcmp(people[i].name, name) == 0)
        {
            printf("Found %s\n", people[i].number);
            return 0;
        }
    }
    printf("Not found\n");
    return 1;
}
```

```
}
```

Sort

Selection sort – burada ilk olaraq bütün array gezilib en kiçik axtarılır sonra 0 ci (i-ci) indexe getirilir. Sonra 1 den (novbeti i-den) başlayaraq yenə bütün array gezilir (amma bu dəfə n-1 step olur). Və balaca axtarılıb, 1-e getirilir.

```
For i from 0 to n-1
  Find smallest number between numbers[i] and numbers[n-1]
  Swap smallest number with numbers[i]
```

- Summarizing those steps, the first time iterating through the list took $n - 1$ steps. The second time, it took $n - 2$ steps. Carrying this logic forward, the steps required could be represented as follows:

$$(n - 1) + (n - 2) + (n - 3) + \dots + 1$$

- This could be simplified to $n(n-1)/2$ or, more simply, $O(n^2)$. In the worst-case or upper-bound, selection sort is in the order of $O(n^2)$. In the best-case, or lower-bound, selection sort is in the order of $\Omega(n^2)$.

Bu alqoritmde $O(n^2)$ olmuş olacaq (yuxarıdakı şəkildə bax, orada $/2$ kimi və $n/2$ kimi kiçik fərq yaradan ədədləri ciddiyyə almiriq), və best scenario yenə $\Omega(n^2)$ olacaq çünki, no matter what axıra qədər check edir bütün arrayi. Deməli $\Theta(n^2)$ olacaq. Çünki O ilə maksimum operation göstərilir, Ω ilə ən yaxşı hal göstərilir, və eger O və Ω eyni olsa o zaman Θ da olur o ədəd.

Bubble sort – təkrar-təkrar 2-2 müqayisə edərək sort edir.

```
Repeat n-1 times
  For i from 0 to n-2
    If numbers[i] and numbers[i+1] out of order
      Swap them
  If no swaps
    Quit
```

Eger son conditionu anlamadınsa,

<https://www.youtube.com/live/iCx3zwK8Ms8?si=ACVVRM0aNj1Oi8c9J&t=5466>

This way it still is $O(n^2)$, but the lower bound can be $\Omega(n)$. therefore, it is more efficient than Selection sort.

Bubble Sort

- In bubble sort, the idea of the algorithm is to move higher valued elements generally towards the right and lower value elements generally towards the left.

In pseudocode:

- Set swap counter to a non-zero value
- Repeat until the swap counter is 0:
 - Reset swap counter to 0
 - Look at each adjacent pair
 - If two adjacent elements are not in order, swap them and add one to the swap counter

Merge sort

```
If only one number
  Quit
Else
  Sort left half of number
  Sort right half of number
  Merge sorted halves
```

$O(n\log(n))$ olmus olacaq burada, 2 hisseye bole bole sort olunur deye. Ω -de O ile bir olacaq. Yeni Θ da o eded olmus olur. N dene elementi $\log(n)$ defe divide and conquer edirik, ona gore de $n\log(n)$ olmus scenariolar.

Runtime Analysis

| Algorithm | O | Ω |
|----------------|---------------|--------------------|
| Merge Sort | $O(N\log(N))$ | $\Omega(N\log(N))$ |
| Selection Sort | $O(N^2)$ | $\Omega(N^2)$ |
| Bubble Sort | $O(N^2)$ | $\Omega(N)$ |

For example:

```
answers.txt X
1 sort1 uses: Bubble Sort
2
3 How do you know?: on the reversed list, sort1 took 5s; on the sorted list, sort1 took .3s.
4
5 sort2 uses: Merge Sort
6
7 How do you know?: on the reversed list, sort2 took 0.275s; on the sorted list, 0.586s.
8
9 sort3 uses: Selection Sort
10
11 How do you know?: on the reversed list, sort3 took 2.4s; on the sorted list, 2.1s;
```

Recursion

Recursion – when a function calls itself, then it is recursive. **Base case** – ise recursive functionun sonudur. Meselen, binary searchde middle ile muqayise funksiyasi tekrar-tekrar devam edir, ve sonu start=end olanda olur, hansi ki base case dir.

Meselen 1ci de mario piramidinin evvel yazdigimiz recursion olmayan usuludur. 2ci de ise recursion usulu

```
// Draws a pyramid using iteration

#include <cs50.h>
#include <stdio.h>

void draw(int n);

int main(void)
{
    // Get height of pyramid
    int height = get_int("Height: ");

    // Draw pyramid
    draw(height);
}

void draw(int n)
{
    // Draw pyramid of height n
    for (int i = 0; i < n; i++)
    {
        for (int j = 0; j < i + 1; j++)
        {
            printf("#");
        }
    }
}
```

```
    }  
    printf("\n");  
}  
}
```

```
// Draws a pyramid using recursion  
  
#include <cs50.h>  
#include <stdio.h>  
  
void draw(int n);  
  
int main(void)  
{  
    // Get height of pyramid  
    int height = get_int("Height: ");  
  
    // Draw pyramid  
    draw(height);  
}  
  
void draw(int n)  
{  
    // If nothing to draw  
    if (n <= 0)  
    {  
        return;  
    }  
  
    // Draw pyramid of height n - 1  
    draw(n - 1);  
  
    // Draw one more row of width n  
    for (int i = 0; i < n; i++)  
    {  
        printf("#");  
    }  
    printf("\n");  
}
```

Let's walk through your example with $n = 3$:

1. `draw(3)` is called. It calls `draw(2)`.
2. `draw(2)` is called. It calls `draw(1)`.
3. `draw(1)` is called. It calls `draw(0)`.
4. `draw(0)` is called. Since $n \leq 0$, it returns immediately.

Now, the function starts to "come back" or "unwind" from the recursive calls:

5. `draw(1)` resumes after `draw(0)` returns. It draws one row of width 1.
6. `draw(2)` resumes after `draw(1)` returns. It draws one row of width 2.
7. `draw(3)` resumes after `draw(2)` returns. It draws one row of width 3.

Recursion

- The Collatz conjecture applies to positive integers and speculates that it is always possible to get "back to 1" if you follow these steps:
 - If n is 1, stop.
 - Otherwise, if n is even, repeat this process on $n/2$.
 - Otherwise, if n is odd, repeat this process on $3n + 1$.
- Write a recursive function `collatz(n)` that calculates how many steps it takes to get to 1 if you start from n and recurse as indicated above.

```
int collatz(int n)
{
    // base case
    if (n == 1)
        return 0;
    // even numbers
    else if ((n % 2) == 0)
```

```
return 1 + collatz(n/2);  
// odd numbers  
else  
return 1 + collatz(3*n + 1);  
}
```

Yuxarıdaki kodu anla sonra.