



# A Tale of Two Topologies: Exploring Convertible Data Center Network Architectures with Flat-tree

Yiting Xia, Xiaoye Steven Sun, Simbarashe Dzinamarira, Dingming Wu, Xin Sunny Huang, T. S. Eugene Ng

Rice University

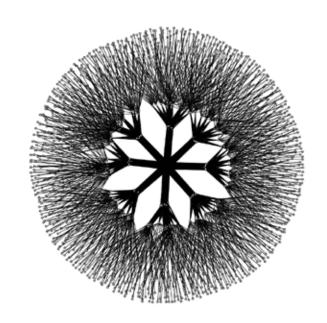
## Convertible Network

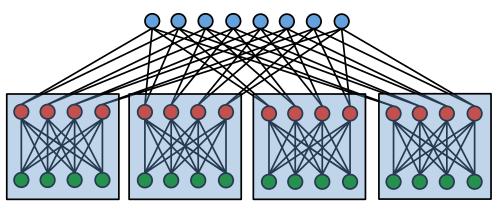
- Convertibility
  - A network's ability to change between multiple topologies with different characteristics
  - Managed by software, no human labor for rewiring
- Combine benefits of different worlds



# Clos Topology

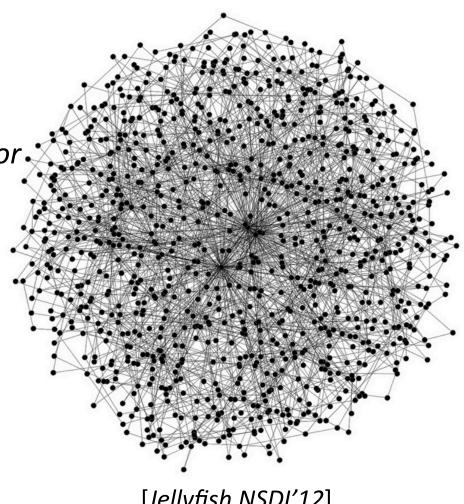
- Implementation friendly
  - Central wiring
  - Flexible scale and oversubscription
  - Pod modular design
- Suboptimal performance
  - Long paths
  - Congested network core





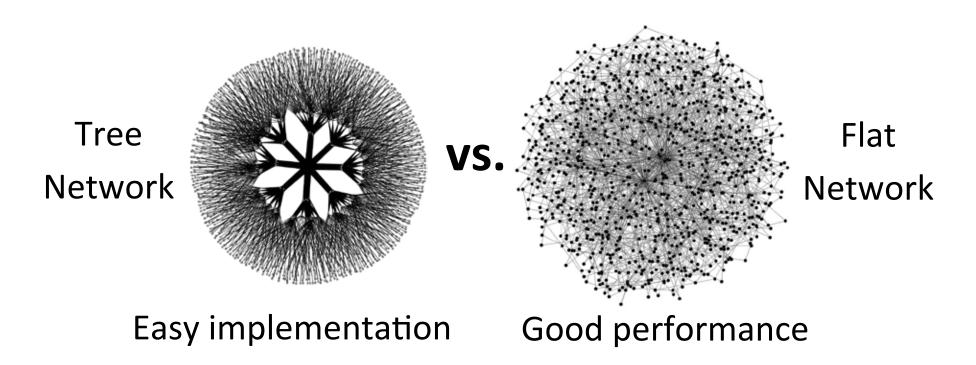
# Random Graph

- Good performance
  - Low average path length
  - Rich bandwidth
  - Near optimal throughput for uniform traffic
- Hard to implement
  - Neighbor-to-neighbor wiring complicated



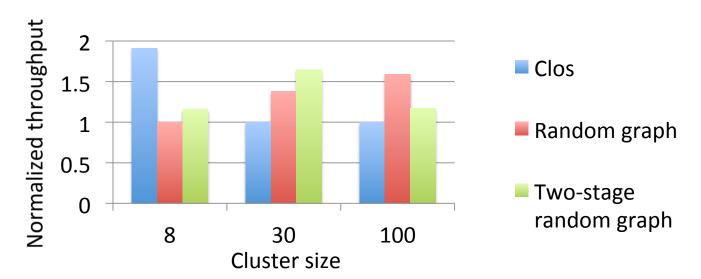
[Jellyfish NSDI'12]

# The Case for Convertibility



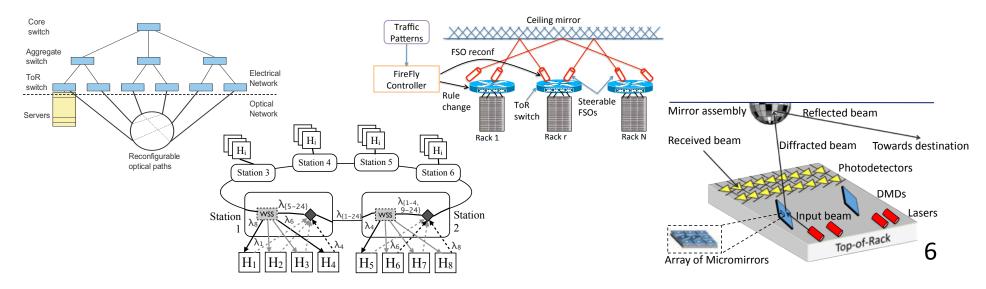
# The Case for Convertibility

- Data center traffic: different locality and cluster size
  - Random graph (global) for global traffic
  - Two-stage (local) random graph for in-Pod traffic
  - Clos for in-rack traffic
- Motivating example
  - 8 servers per rack, 64 servers per Pod, all-to-all traffic



# Topology can be configurable

- Helios, c-Through, Flyways, OSA, 3DBeam, Mordia, FireFly, Quartz, WaveCube, ProjecToR, etc
- Create ad-hoc links on the fly
- Technology available
  - 3D MEMS, WSS, WDM, DMD, free-space optics, 60GHz wireless, wireless beamforming



# Flat-tree Prototype Architecture

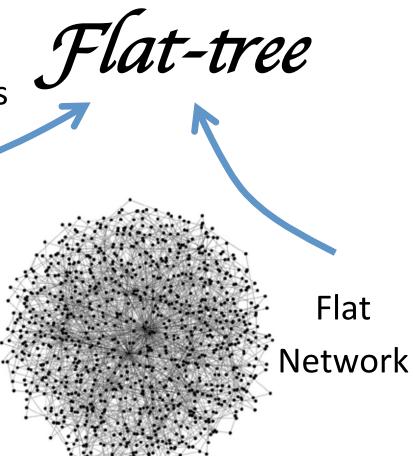
Start from Clos

Tree

Network

Flatten tree structure

Approximate random graphs

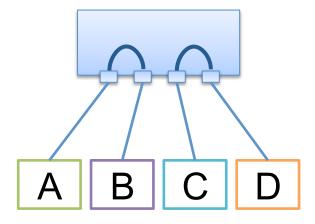


## Flatten the Tree

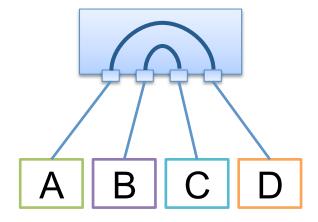
• How to flatten the tree structure?

Difference	Clos	Random graph	Solution
Server distribution	Edge switches	All switches	Relocate servers
Wiring	Central	Neighbor-to-neighbor	Diversify connections

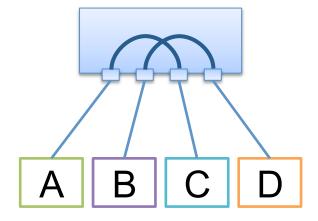
- Small port-count
- Low cost
  - Optical Fibers in data center
    - \* Small optical switch
    - \* \$10 per port
  - DAC in data center
    - \* Crosspoint switch
    - \* \$3 per port
- Physical layer device



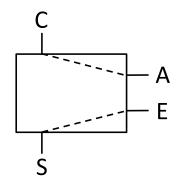
- Small port-count
- Low cost
  - Optical Fibers in data center
    - \* Small optical switch
    - \* \$10 per port
  - DAC in data center
    - \* Crosspoint switch
    - \* \$3 per port
- Physical layer device



- Small port-count
- Low cost
  - Optical Fibers in data center
    - \* Small optical switch
    - \* \$10 per port
  - DAC in data center
    - \* Crosspoint switch
    - \* \$3 per port
- Physical layer device



C: core switch



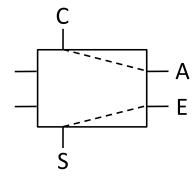
A: aggregation switch

E: edge switch

S: server

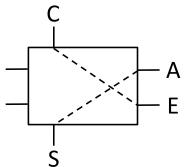
4-port Converter Switch

C: core switch



A: aggregation switch

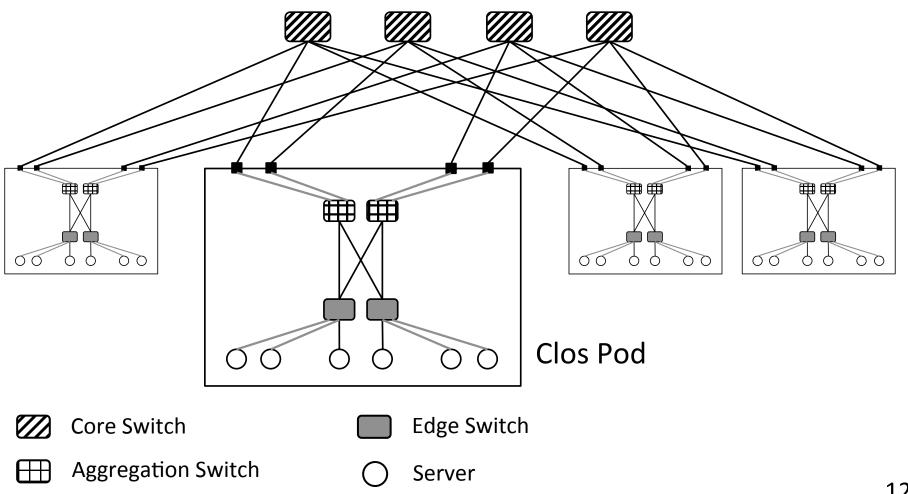
E: edge switch



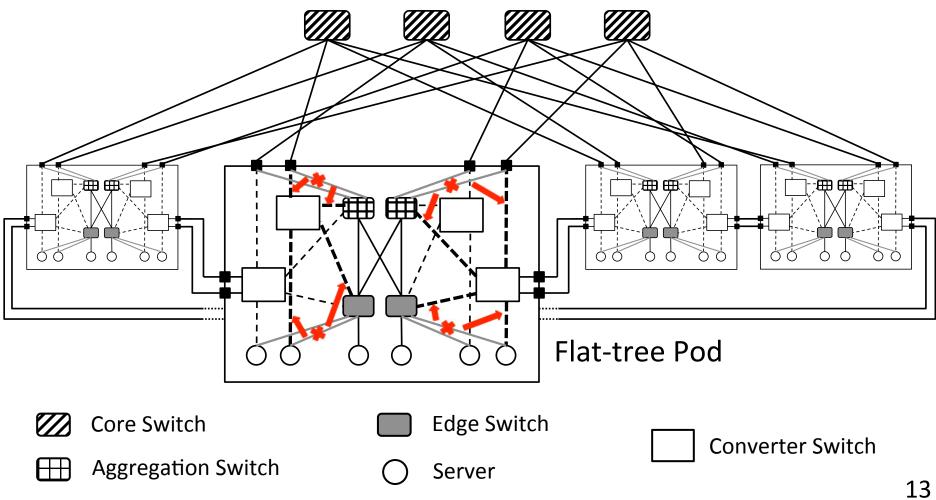
S: server

6-port Converter Switch

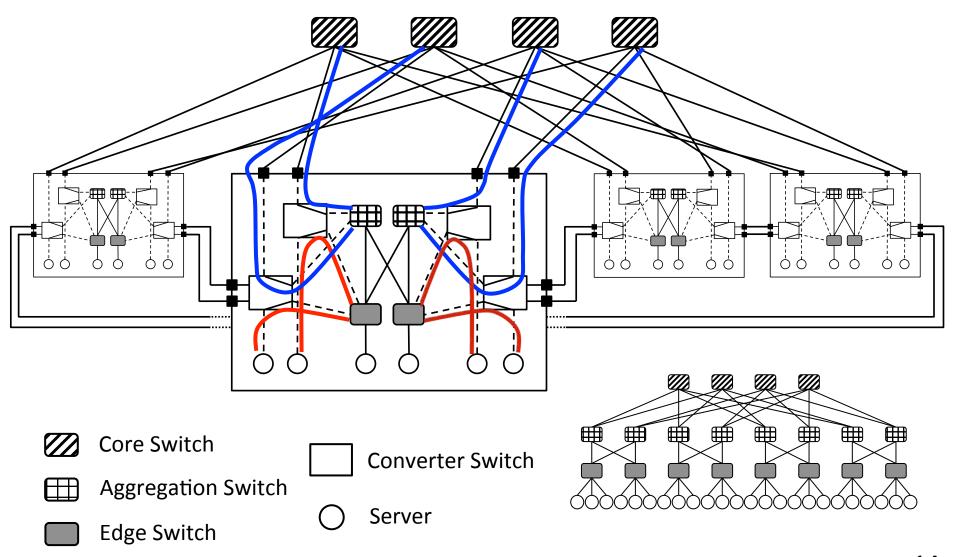
# Flat-tree Example



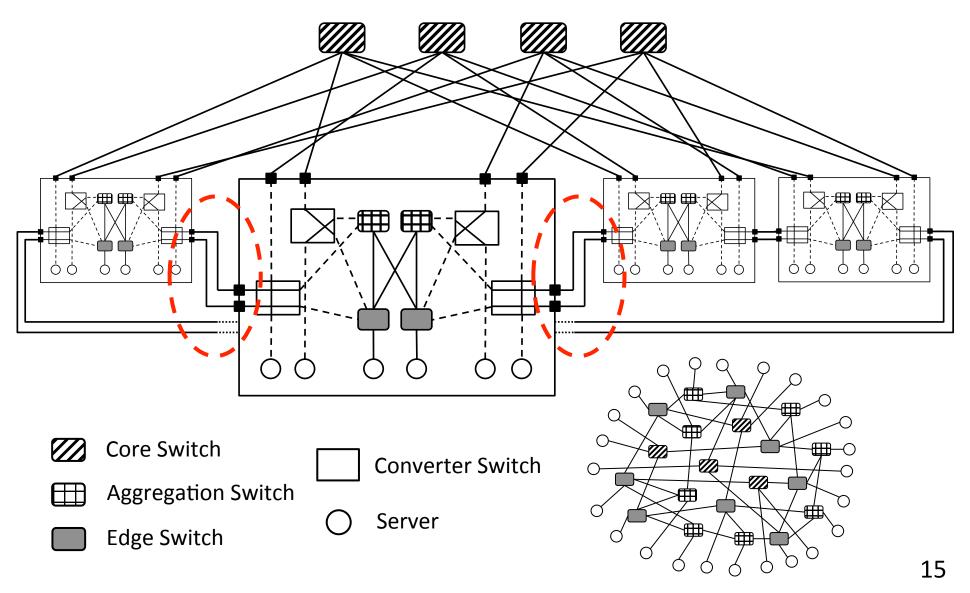
# Flat-tree Example



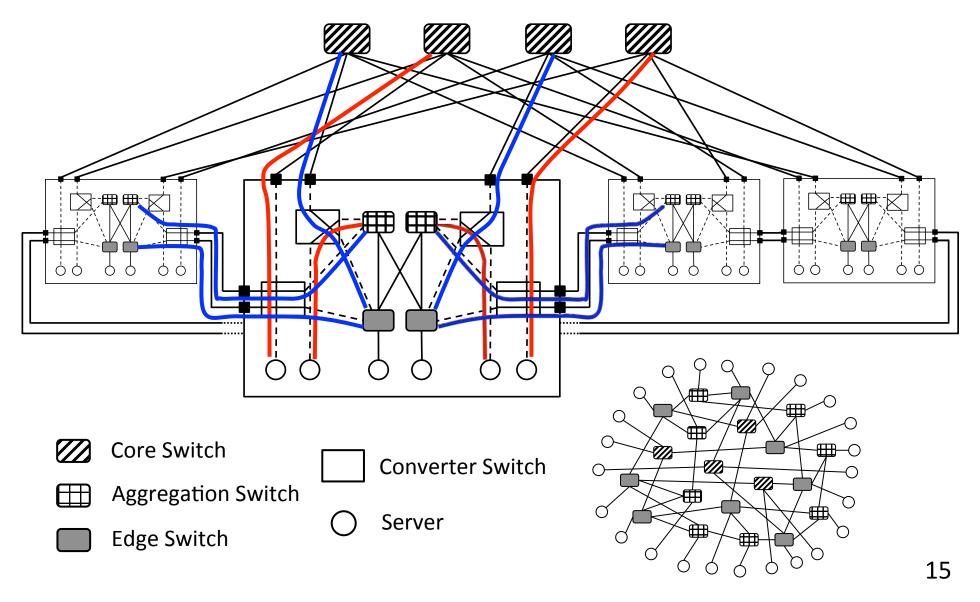
## Clos Network



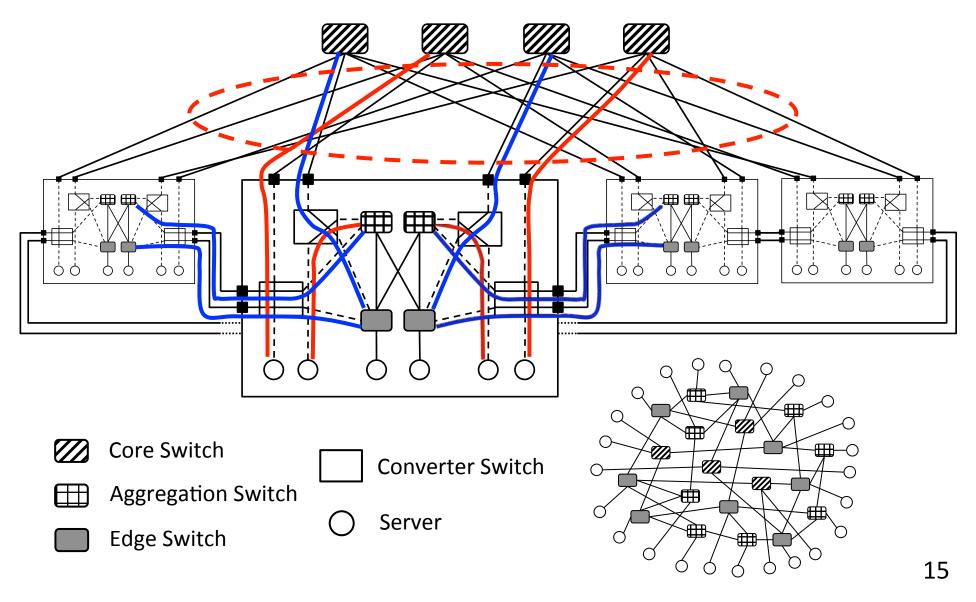
# Approximate Random Graph



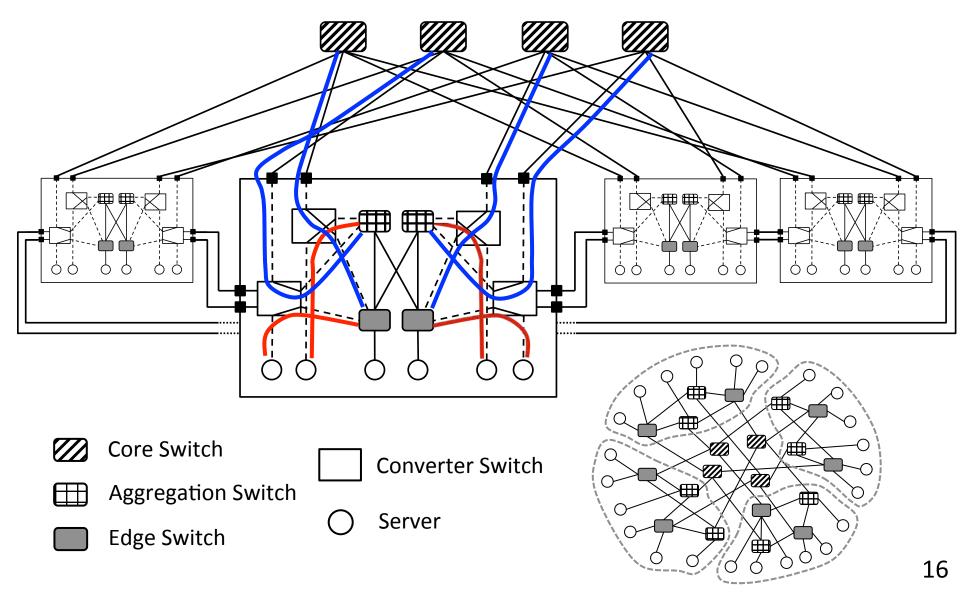
# Approximate Random Graph



# Approximate Random Graph



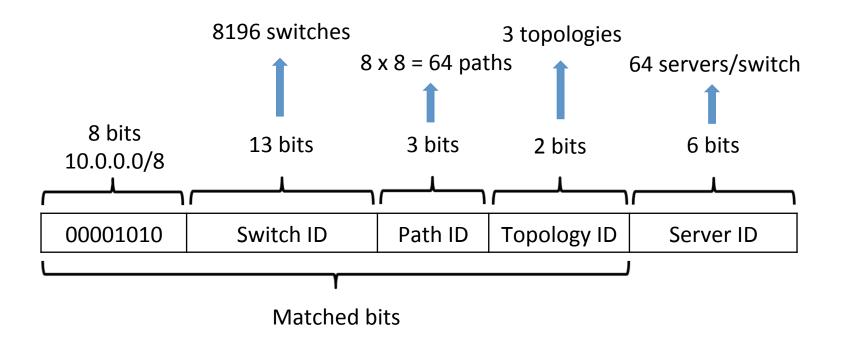
# Approximate Local Random Graph

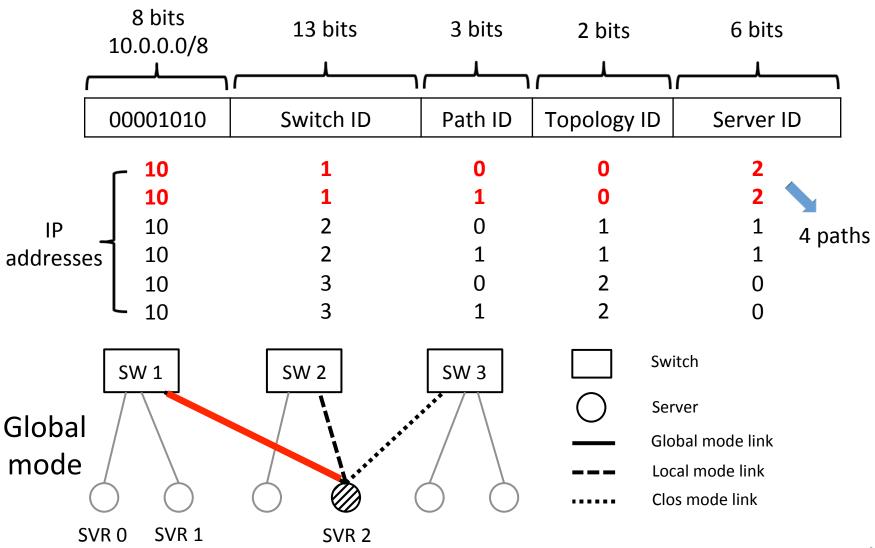


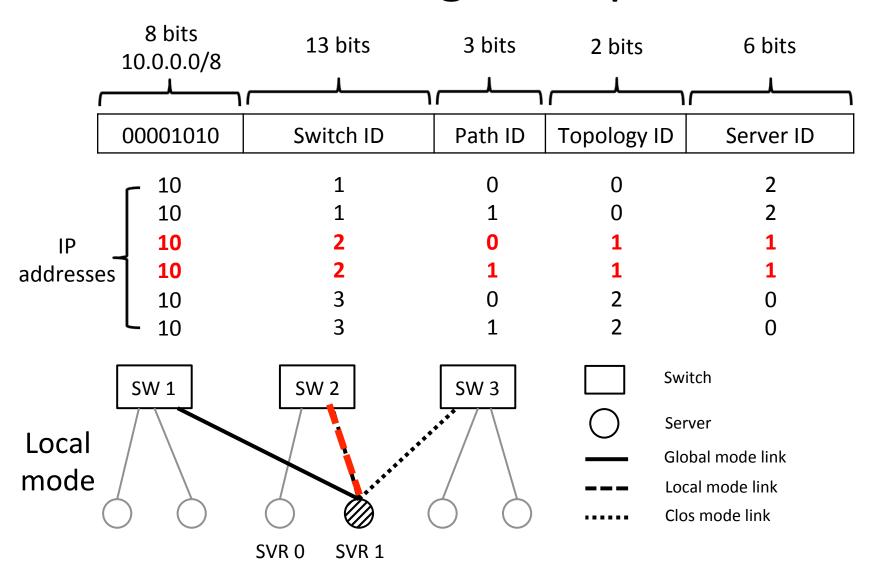
## **Control Plane**

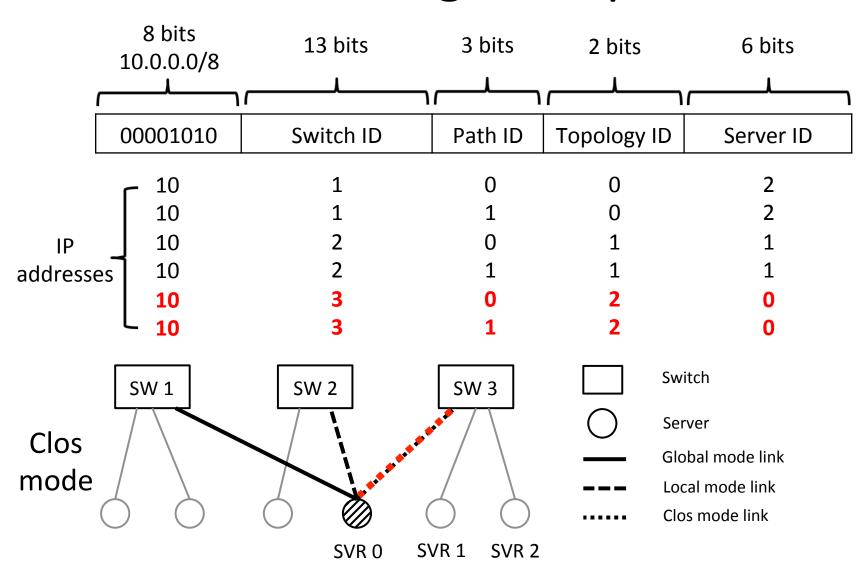
- k-shortest-path routing + MPTCP
  - k paths for every sever pairs
  - Enormous number of states  $\rightarrow$  exceed switch capacity
  - No solution from random graph networks
  - Scalability concern!!!
- Aggregation
  - Addressing: prefix matching of ingress/egress switch
  - Source routing
- Highly challenging in flat-tree
  - Server mobility to different switches

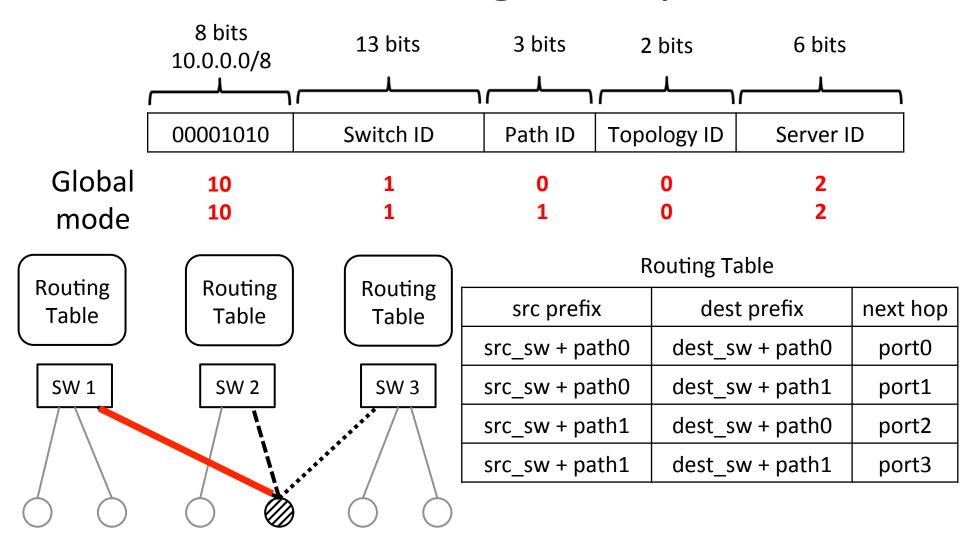
# Addressing



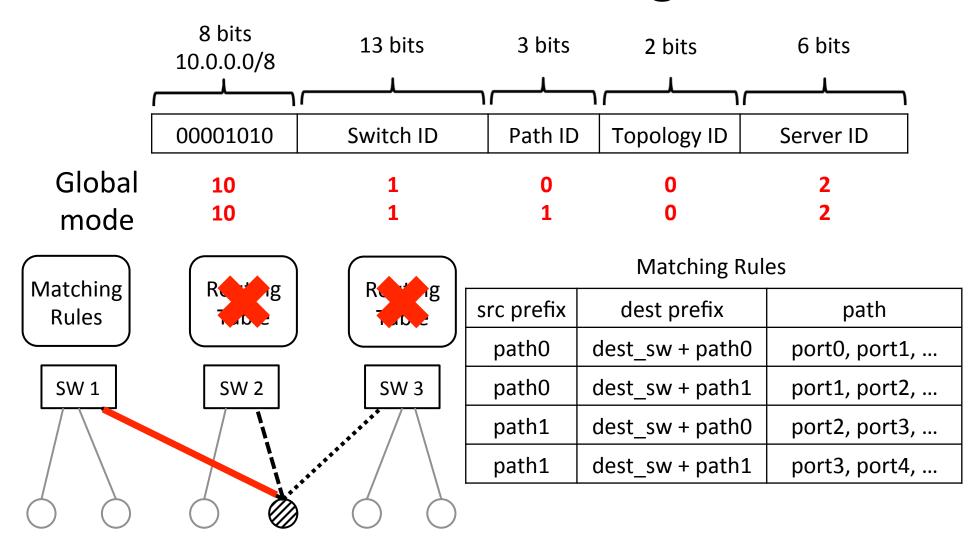








## Source Routing

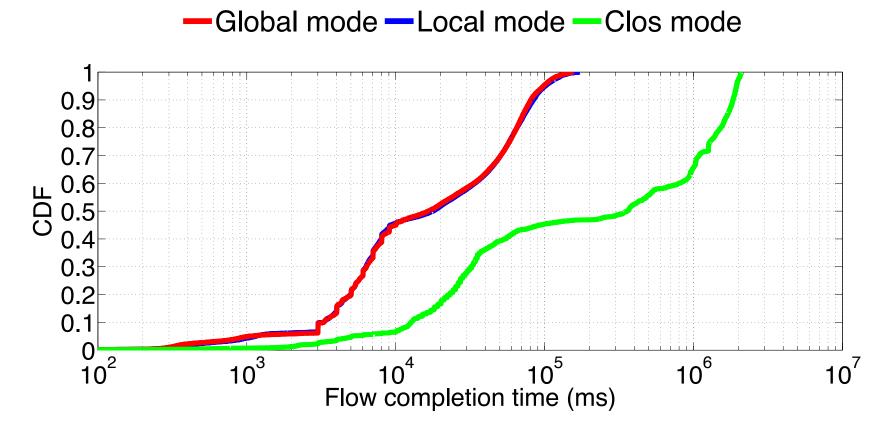


## **Control Plane**

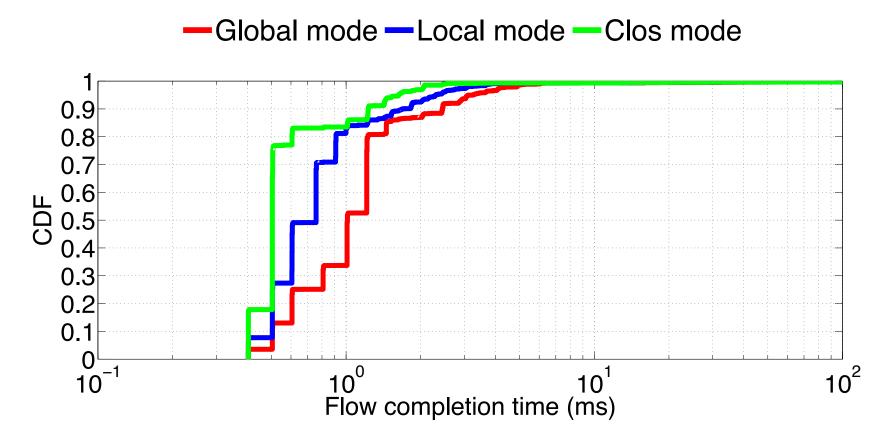
- Addressing
  - Server-level -> switch-level k-shortest path routing
  - k paths per server pair  $\rightarrow k$  paths per switch pair
- Source routing
  - Transit switches: no states
  - Ingress switch: k paths per egress switch
- Applicable to static random graph networks

- Packet-level simulation
- Traffic traces from 4 Facebook data centers
  - Hadoop-1: no locality
  - Haddop-2: rack-level locality
  - Web: Pod-level locality
  - Cache: Pod-level locality

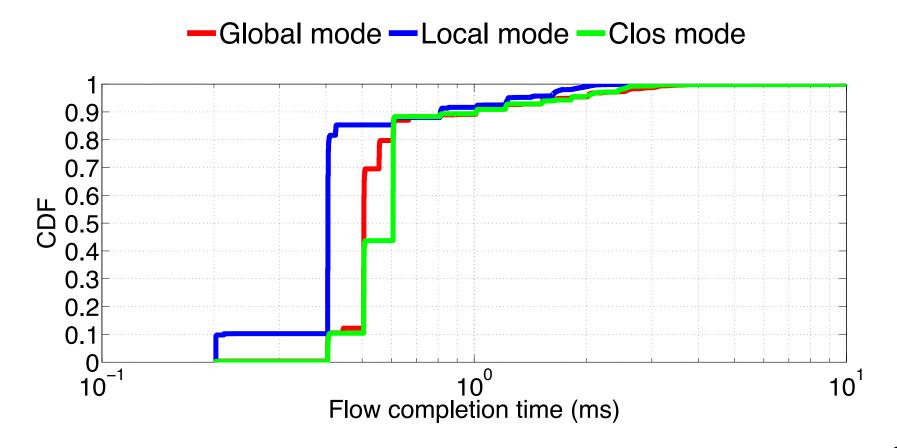
Hadoop-1: no locality



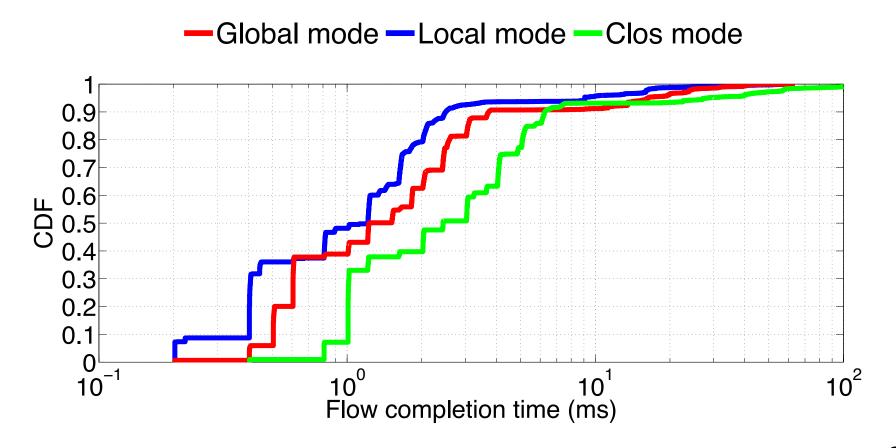
Hadoop-2: rack-level locality



Web: Pod-level locality

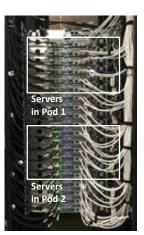


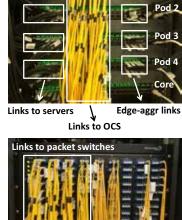
Cache: Pod-level locality

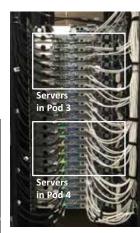


- Convertibility!!!
- Different topology for different workload
- Convert topology as workload changes
- Partition network into zones

- Theoretical performance
  - Average path length
  - Throughput from Linear Programming solver
- Effectiveness of k-shortest-path routing and MPTCP
  - Throughput from simulation
    - ≈ LP solver
- Testbed implementation
  - Hadoop & Spark
  - 27.6% more bandwidth
  - 10% less data read time







# Configurability vs. Convertibility

- Helios, c-Through, Flyway, OSA, 3DBeam, Mordia, FireFly, Quartz, WaveCube, ProjecToR, etc
- Different design philosophy

	Configurable network	Convertible network	
Traffic to service	Instantaneous flows	Long-lasting workloads	
Capacity	Add bandwidth	Better use of bandwidth	
Topology change	Incremental & frequent	Network-wide & infrequent	

## Conclusion

- Convertible data center network architecture
- Flat-tree converts between Clos topology and approximate random graphs of different scales
- Complete architecture and control plane design
  - Inexpensive converter switches
  - Distributed converter switches  $\rightarrow$  scalability of architecture
  - Addressing + source routing  $\rightarrow$  scalability of control plane
- Extensive performance evaluation
  - LP simulations
  - Packet-level simulations
  - Testbed implementation