Carousel: Scalable Traffic Shaping at End Hosts

Ahmed Saeed, Nandita Dukkipati, Vytautas Valancius, Vinh The Lam, Carlo Contavalli, and Amin Vahdat

















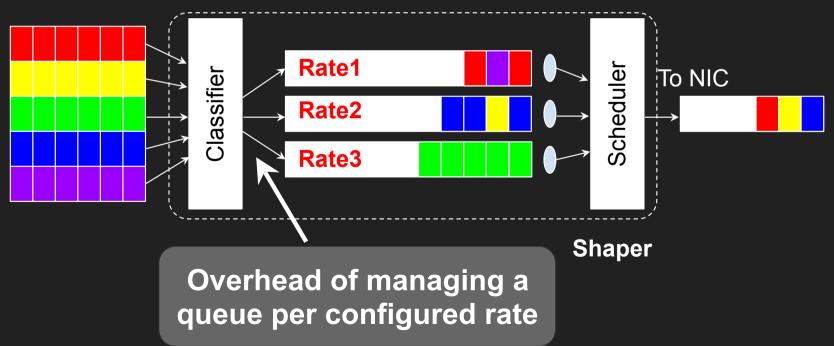


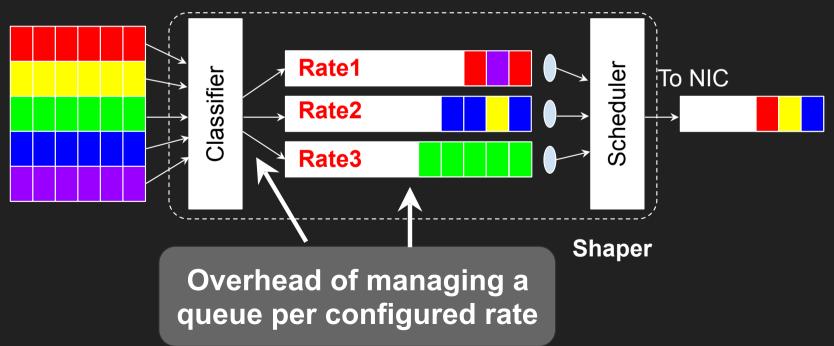
Packet sources

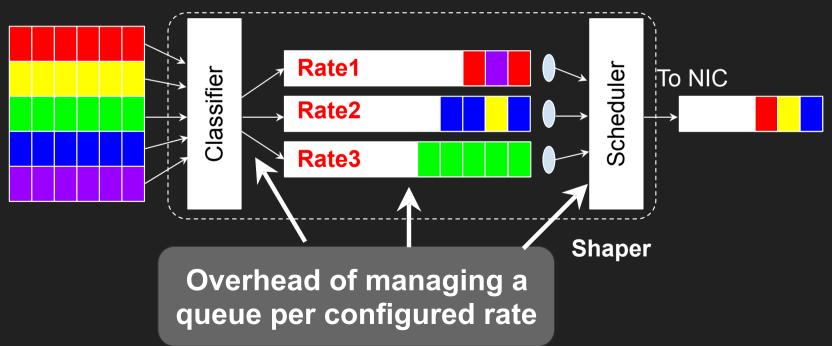


Overhead of managing a queue per configured rate

Shaper

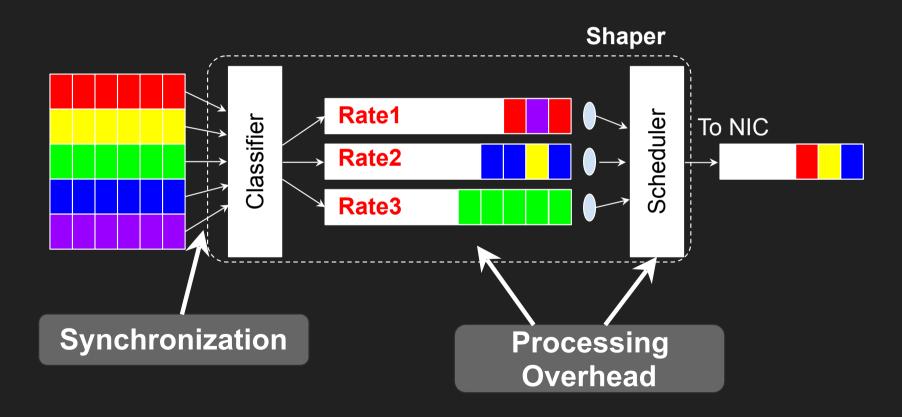


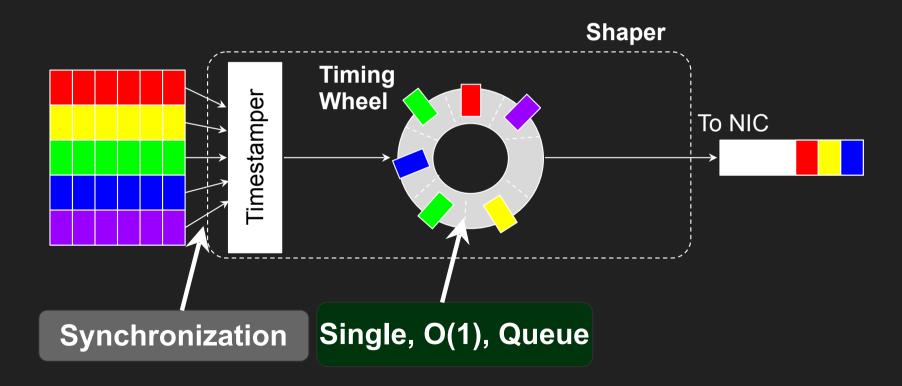


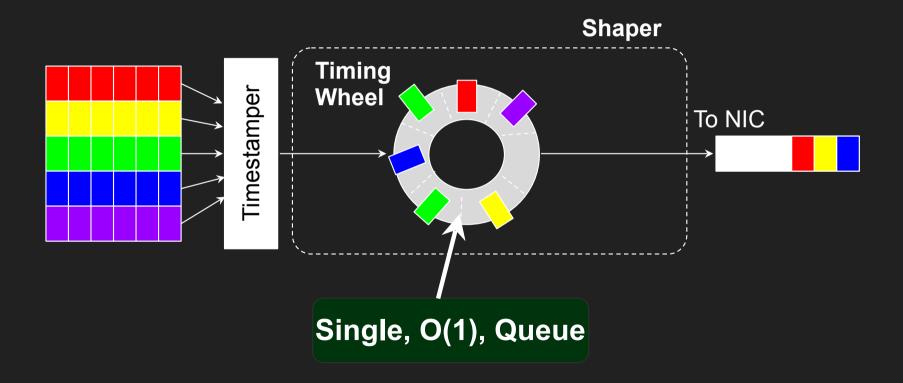


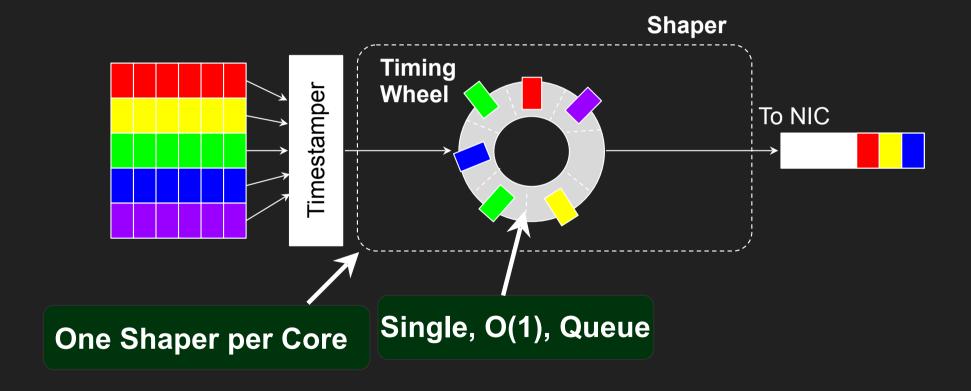
We need new traffic shapers that can handle tens of thousands of flows and rates

Main Idea Replace the many queues with a single low-overhead queue

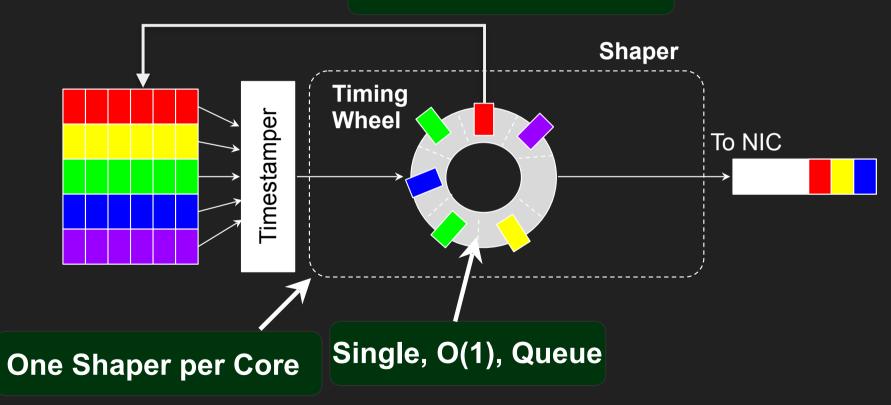








Apply backpressure



Outline

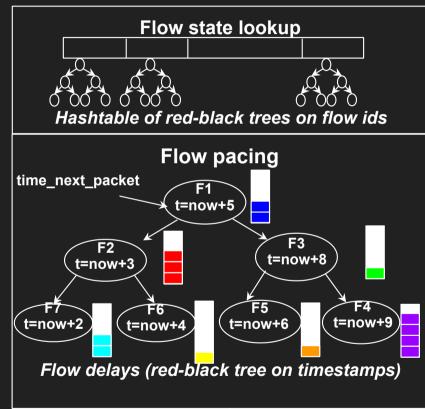
- Problems with Current Shapers
- Carousel Overview
- Single Queue Shaping
- Backpressure
- Evaluation

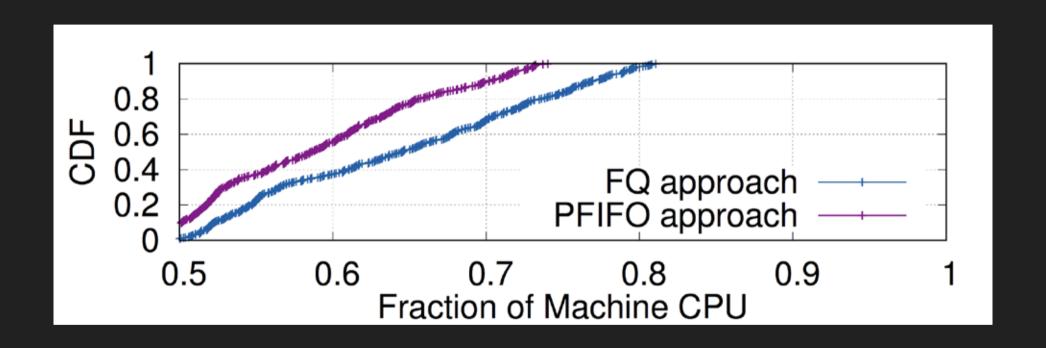
Problems with Current Shapers

FQ/Pacing

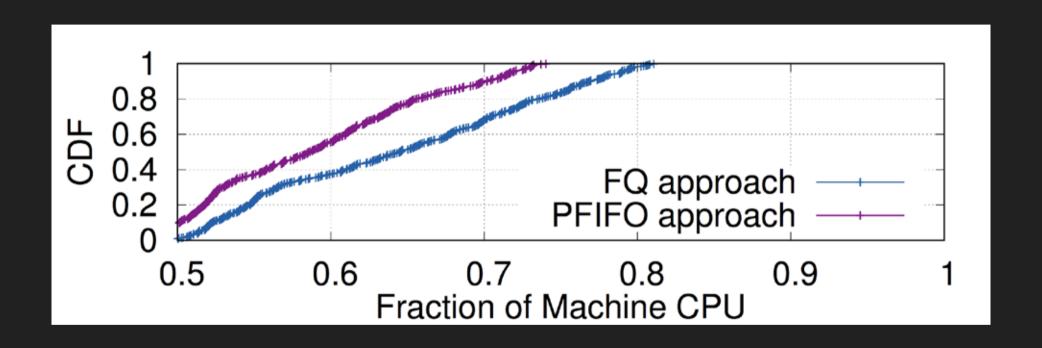
- Implements per TCP flow pacing
- Requires a queue per flow
 - Flows are kept in order of their scheduled transmission time
 - Flows are dequeued in order
- O(log n) operations per packet to operate on a sorted list of flows





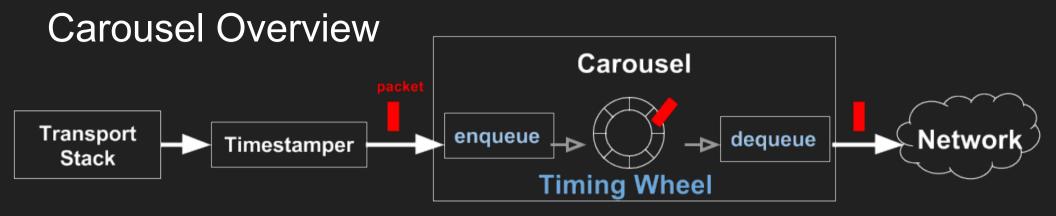


CPU utilization for FQ/pacing and a NOOP Qdisc for the same load



FQ/Pacing introduces 10% more CPU overhead

Carousel Overview



- Relies on a single queue for all packets from all flows
- Requires a high frequency timer or busy polling
- Pinned to a single core

Single Queue Shaping

Single Queue Shaping

- All packet are sorted by their transmission time in one data structure
- A single queue for all traffic will need to handle tens of thousands of packets

Challenge:

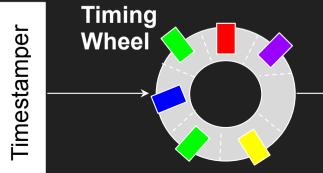
Enqueue and dequeue in a data structure of sorted elements at line rate

Single Queue Shaping

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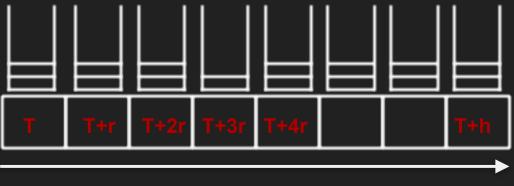
Enqueue and dequeue in a data structure of sorted elements at

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Timing Wheel [Varghese et al. SOSP '87]

- Bucket sort approach to Calendar Queue covering a time horizon
 - Relies on having a minimum rates
- Implemented as an array of buckets each a linked list of packets
 - Each bucket represents a certain time range

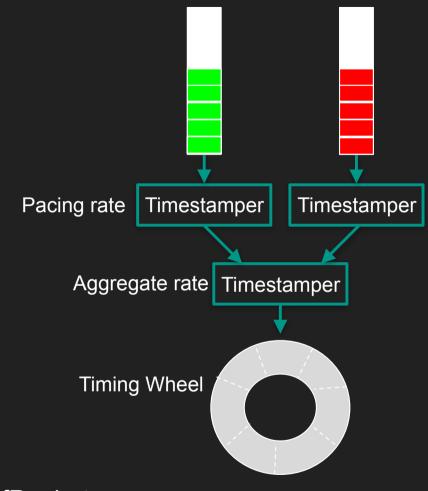


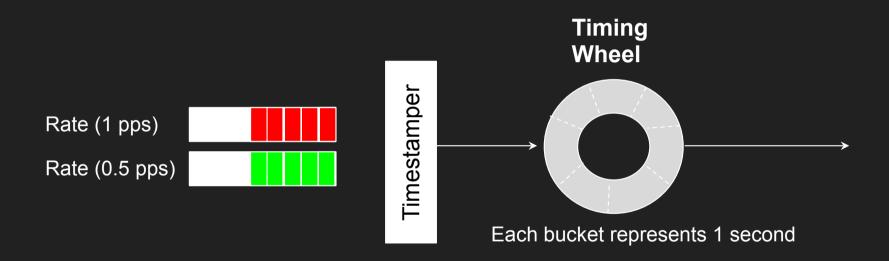
Timing Wheel Benchmark

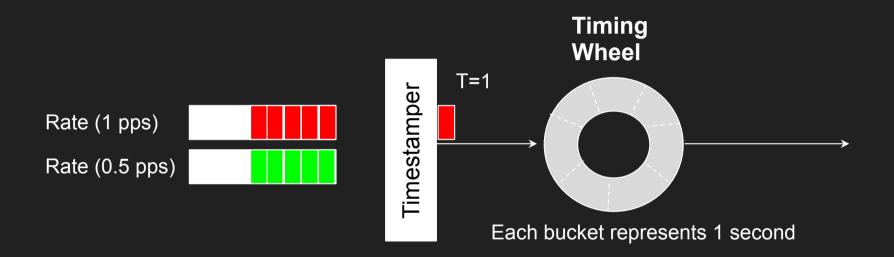
- Measured overhead per enqueue/dequeue pairs
- Overhead per element is between 21-22 nanoseconds
 - Fixed for 2000 to 2 million sorted elements
 - 21 nanoseconds per packet = 500 Gbps (for 1500 byte packets)

Timestampers

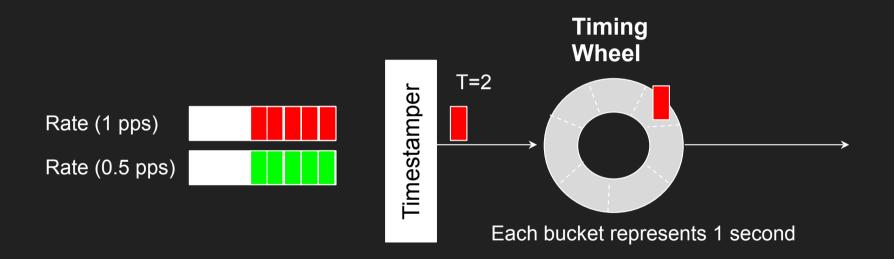
- Packets are timestamped by policy enforcers in their transmission path
 - TCP timestamps a packet based on its pacing rate
 - Bandwidth enforcer timestamps a packet based on its policy-based aggregate rate
- Carousel picks the largest timestamp
- NextTimestamp = LastTimestamp + SizeOfPacket ConfiguredRate



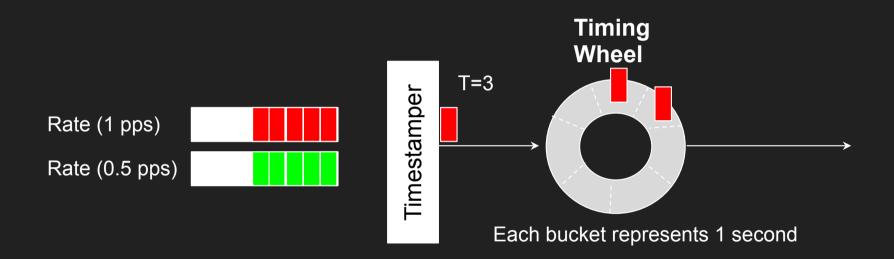




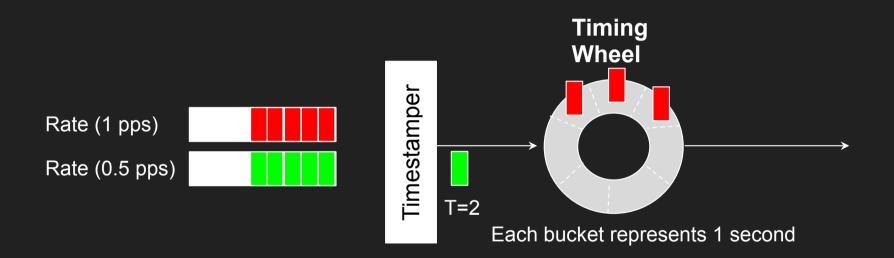
A time step 0

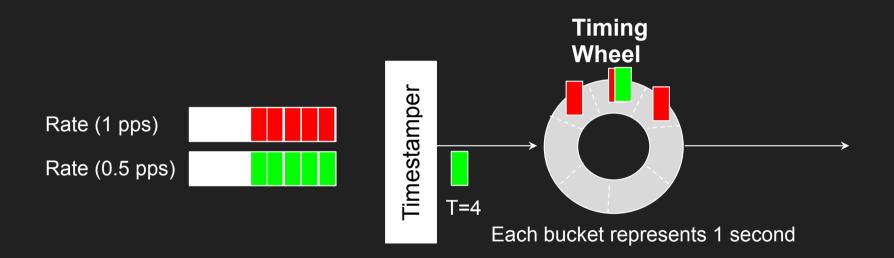


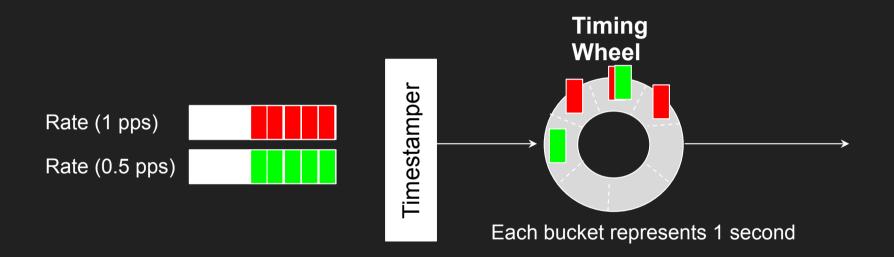
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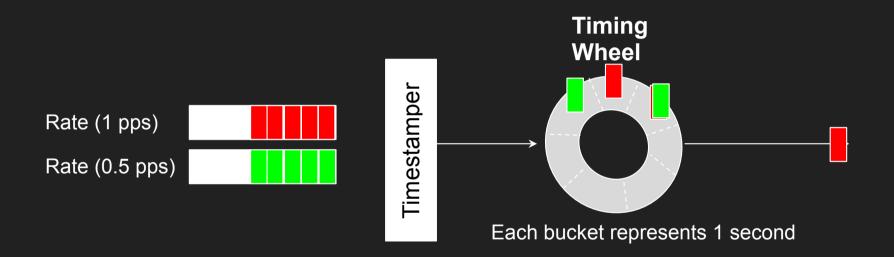


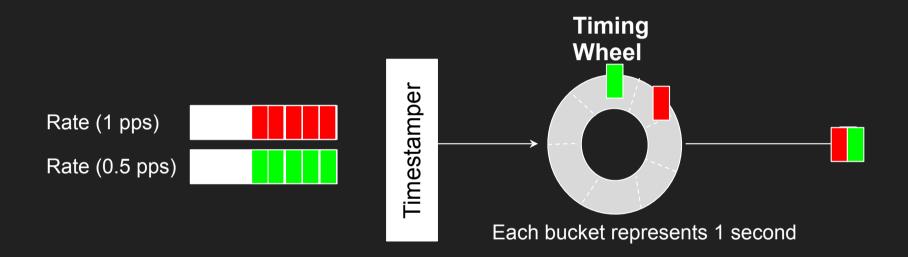
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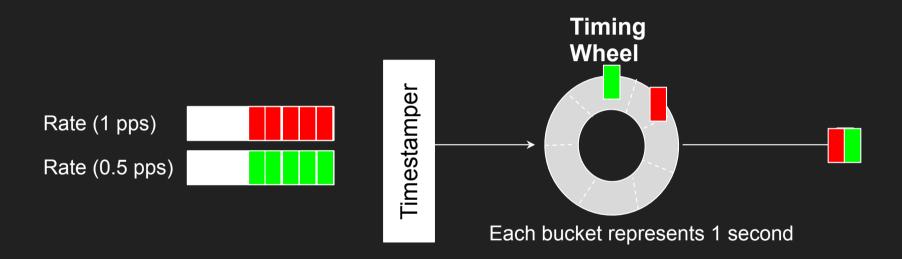












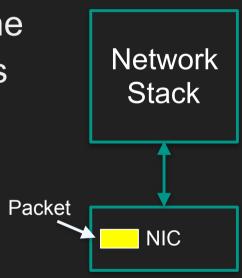
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 - Unnecessary drops (when the queue is full the queue tail drops)

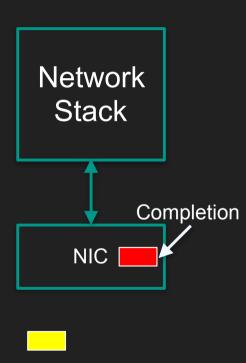
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 - Head of Line Blocking
- Backpressure allows shapers to control sender rate and avoid overwhelming the shaper

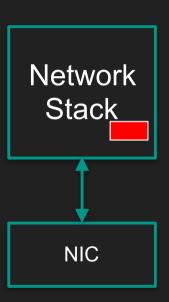
 Completions are signals from the NIC to the network stack to inform it that a packet has been transmitted



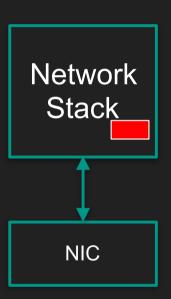
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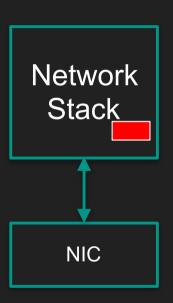
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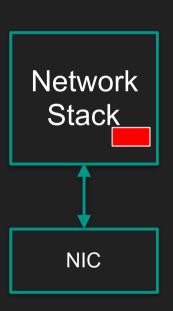
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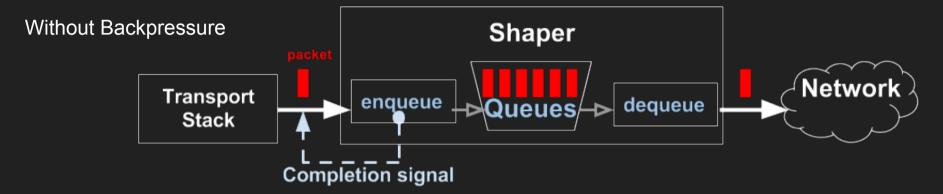


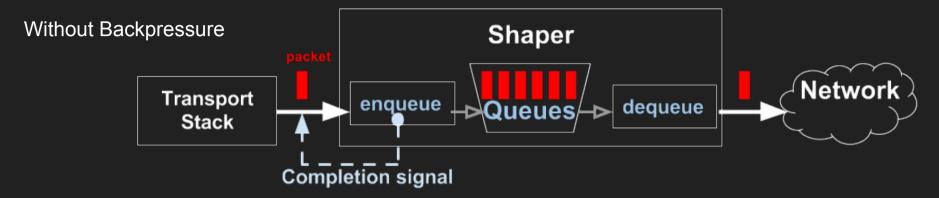
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 - Completion should be controlled by the hypervisor not the virtual NIC

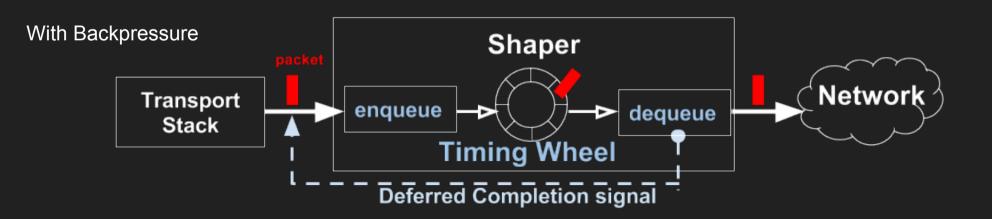


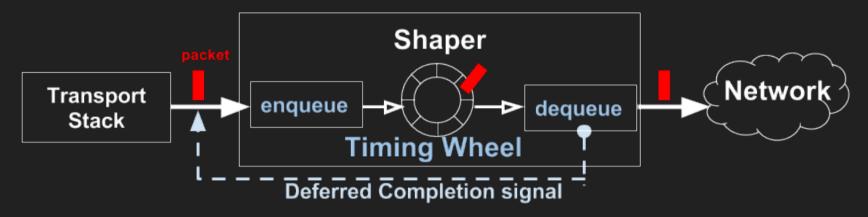
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 - Completion should be controlled by the hypervisor not the virtual NIC
- Completions should be delivered out of order and completely controlled by Shapers

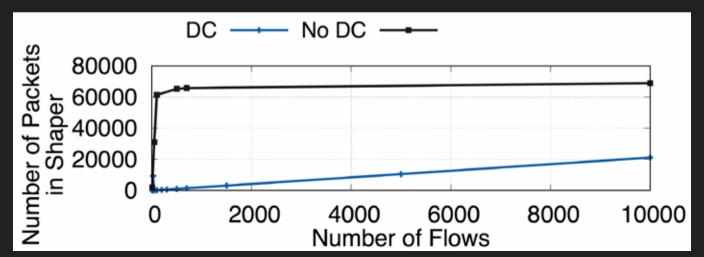


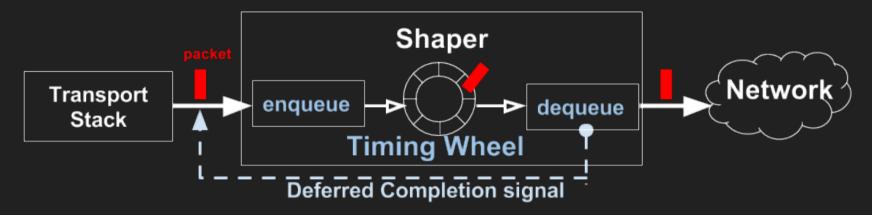


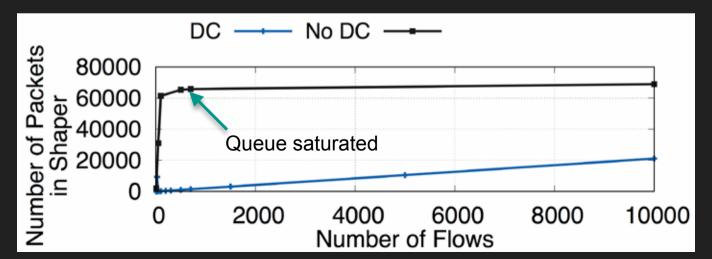


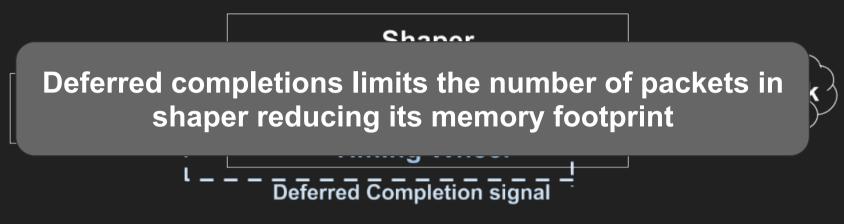


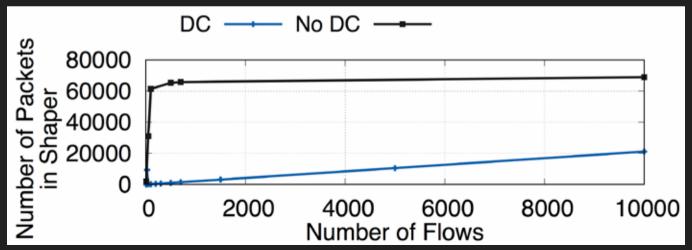








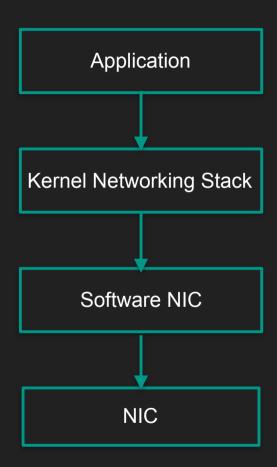




Evaluation

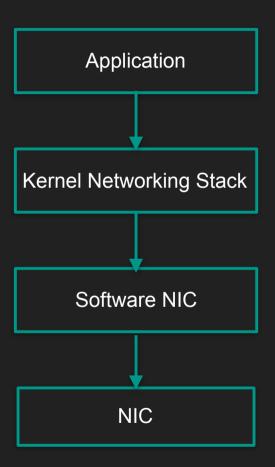
Evaluation Setup

Carousel deployed within a Software NIC



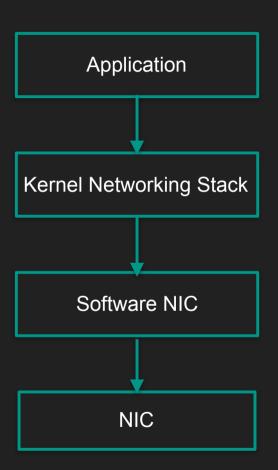
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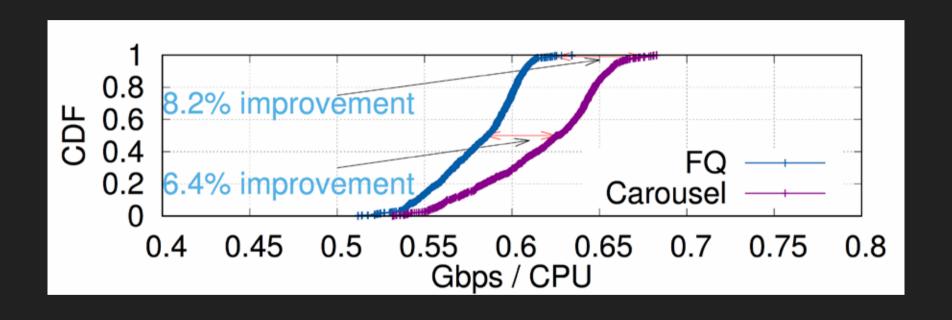
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- Each server handles up to 50k sessions concurrently



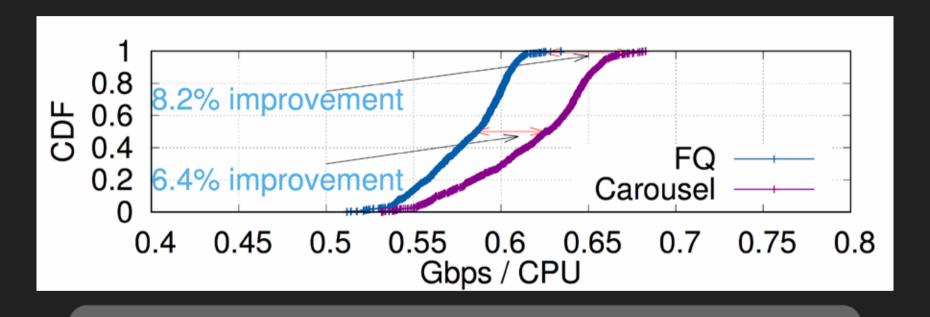
Evaluation Metric

- Measures Gbps served per CPU utilization
 - Metric used is Gbps/CPU (higher is better)
 - Compare machines with similar CPU utilization
 - Measurements performed during peak 12-hours per day
- Evaluation is performed for:
 - Overall CPU utilization
 - Software NIC utilization

Overall CPU Utilization

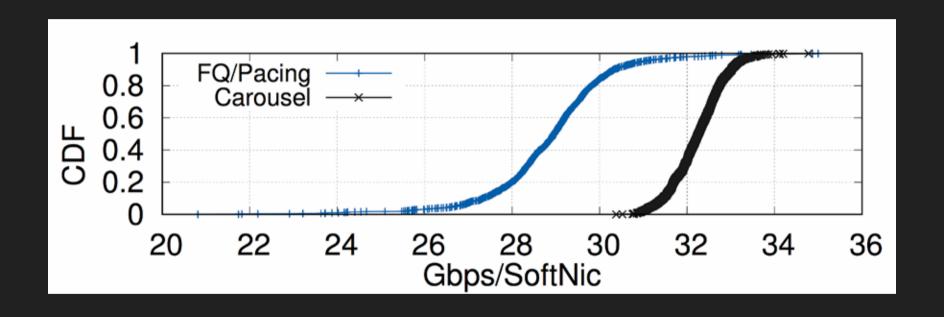


Overall CPU Utilization

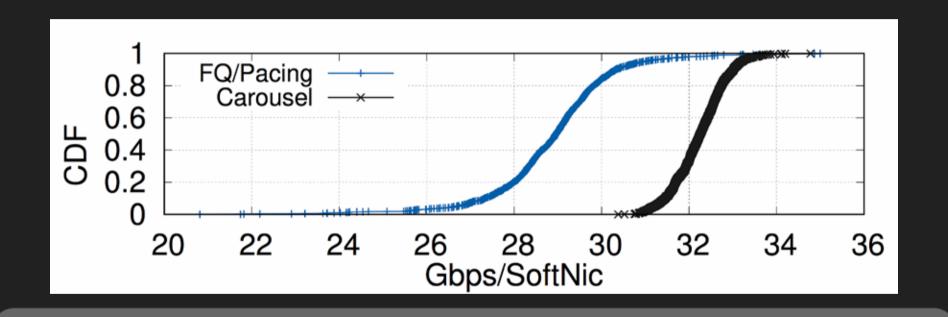


Carousel saves up to 8.2% of overall CPU utilization (5.9 cores on a 72 core machine)

SoftNIC Utilization

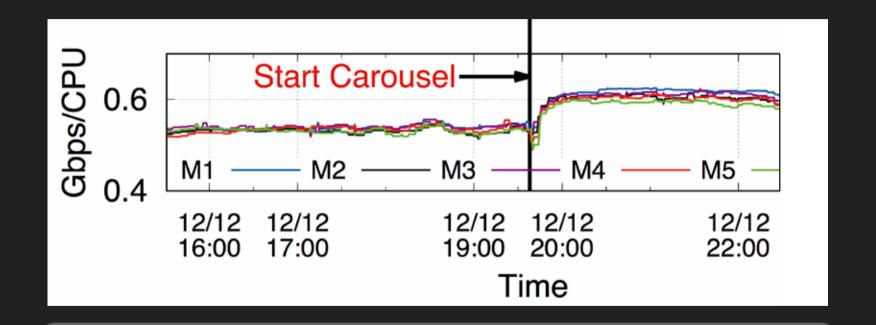


SoftNIC Utilization



Carousel improves even Software NIC utilization by 12% by increasing size of batches of packets enqueue in the Software NIC

Evaluation Summary



Performance improvement when Carousel starts on 5 different machines

Conclusion

- Carousel allows networks operators for the first time to shape tens of thousands of flows individually
- Carousel advantages make a strong case for providing single-queue shaping and backpressure in kernel, userspace stacks, hypervisors, and hardware

Questions?