# CSCI 341 Problem Set 3

Determinization, the Structure of Fin, Regular Expressions, and Antimirov Derivatives

Due Monday, September 22 at 11:00 AM

Don't forget to check the webspace for hints and additional context for each problem!

## Determinization

**Problem 1** (Determinizing is Deterministic). Prove that for any automaton  $\mathcal{A} = (Q, A, \delta, F)$ ,  $\operatorname{Det}(\mathcal{A})$  is total deterministic.

Solution.  $\Box$ 

**Problem 2** (Determinized State, Completing the Proof). Let  $\mathcal{A} = (Q, A, \delta, F)$  be an automaton, and let  $\mathrm{Det}(\mathcal{A})$  be its determinization. Prove that for any state  $x \in Q$ ,

$$\mathcal{L}(\mathcal{A}, x) \supseteq \mathcal{L}(\mathrm{Det}(\mathcal{A}), \{x\})$$

Solution.

**Problem 3** (You Got Options). Find the smallest automaton (not necessarily total or deterministic) with a state that accepts the language

$$L = \{ab^n \mid n \in \mathbb{N}\} \cup \{ac^n \mid n \in \mathbb{N}\} \cup \{a(bc)^n \mid n \in \mathbb{N}\}\$$

over the alphabet  $A = \{a, b, c\}$ . Use determinization to find a deterministic automaton with a state that accepts the same language.

Solution.  $\Box$ 

#### The Structure of Fin

**Problem 4** (Finish Closed under Complement). Let  $\mathcal{A} = (Q, A, \delta, F)$  be a total deterministic automaton with a state x and let  $L = \mathcal{L}(\mathcal{A}, x)$ . Let  $\mathcal{A}' = (Q, A, \delta, Q \setminus F)$ . Prove that  $\mathcal{L}(\mathcal{A}', x) = A^* \setminus L$ .

Solution.  $\Box$ 

**Problem 5** (Intersection-product Construction). Let  $A_1 = (Q_1, A, \delta_1, F_1)$  and  $A_2 = (Q_2, A, \delta_2, F_2)$  be total deterministic automata, and let  $x \in Q_1$  and  $y \in Q_2$ . Let  $L_1 = \mathcal{L}(A_1, x)$  and  $L_2 = \mathcal{L}(A_2, x)$ . Change the accepting states in the union-product construction to obtain an automaton  $A_1 \otimes A_2 = (Q_1 \times Q_2, A, \delta^{\times}, F^{\otimes})$  such that  $\mathcal{L}(A_1 \otimes A_2, (x, y)) = L_1 \cap L_2$ . Explain how to obtain a proof of the Closure under Intersection Theorem from the proof of the Closure under Union Theorem (i.e., what would you change?).

Solution.

## **Regular Expressions**

**Problem 6** (Intersections and Complements). Show that the following two languages are regular over  $A = \{a, b\}$ .

(1) 
$$L_6 = \mathcal{L}(b^*a(a+b)^*) \cap \mathcal{L}(a^*b(a+b)^*)$$

(2) 
$$L_7 = A^* \backslash L_6$$

Solution.

### **Antimirov Derivatives**

Problem 7 (Some Nested Derivatives). Consider the regular expression  $r = (a(b+c^*)+b)^*$  over the alphabet  $A = \{a, b, c\}$ .

- (1) Name three different words  $w_0, w_1, w_1 \in A^*$  that are not in  $\mathcal{L}(r)$ , i.e.,  $w_0, w_1, w_2 \notin \mathcal{L}(r)$ .
- (2) Use the inequalities in the proof of the Linear Bound on Antimirov Derivatives Lemma to determine an upper bound on the number of states in the automaton  $\langle r \rangle_{\mathcal{A}_{Ant}}$  generated by r in  $\mathcal{A}_{Ant}$ , i.e., #(r).
- (3) Now draw a state diagram of  $\langle r \rangle_{\mathcal{A}_{Ant}}$ .
- (4) How many formation rules were used to form the regular expression r? How does this number of formation rules compare to the number of states in  $\langle r \rangle_{\mathcal{A}_{Ant}}$ ?

Solution.