

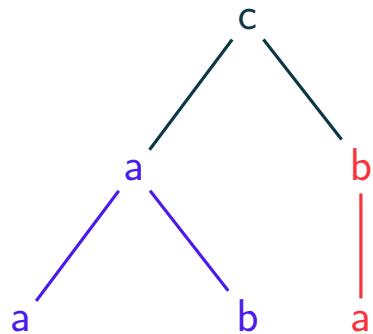
# Progress on weak validation of streamed trees

Corentin Barloy

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## Processing streamed trees

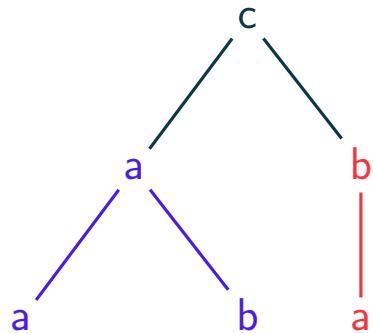
XML encoding of trees:



$\langle c \rangle$   
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 $\langle a \rangle \langle /a \rangle$   
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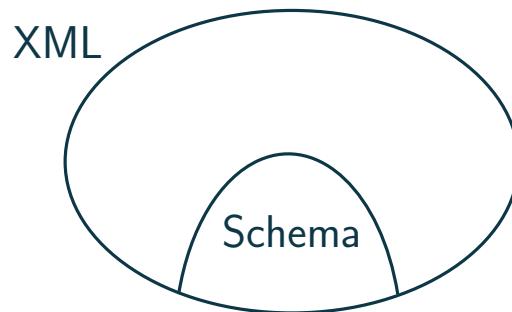
Validation: Does a XML document belongs to a given regular language of trees?

## Weak validation

Weak validation: the document is well-formed.

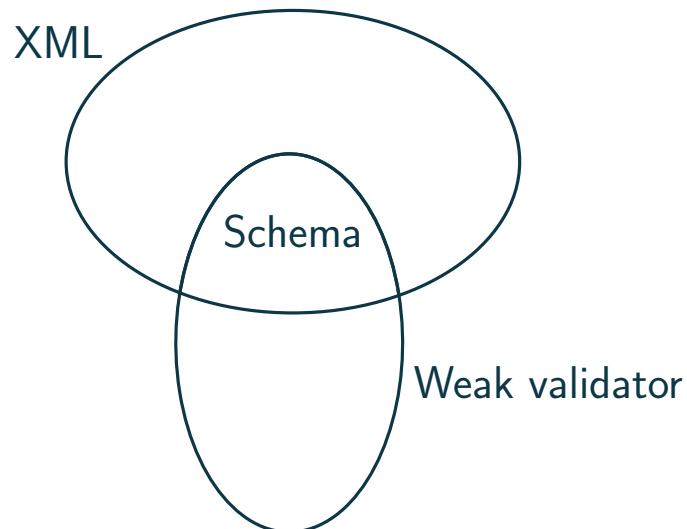
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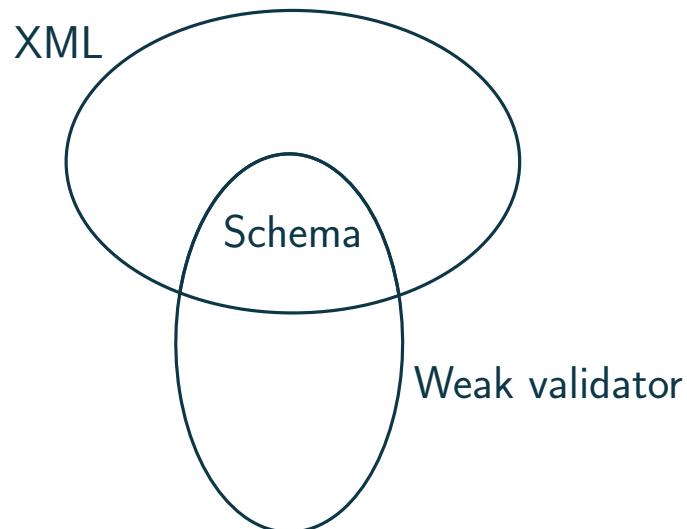
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### Question

What is the space complexity of weakly validating documents in streaming?

## Space complexity

Segoufin, Vianu:

What can be weakly validated in constant space?

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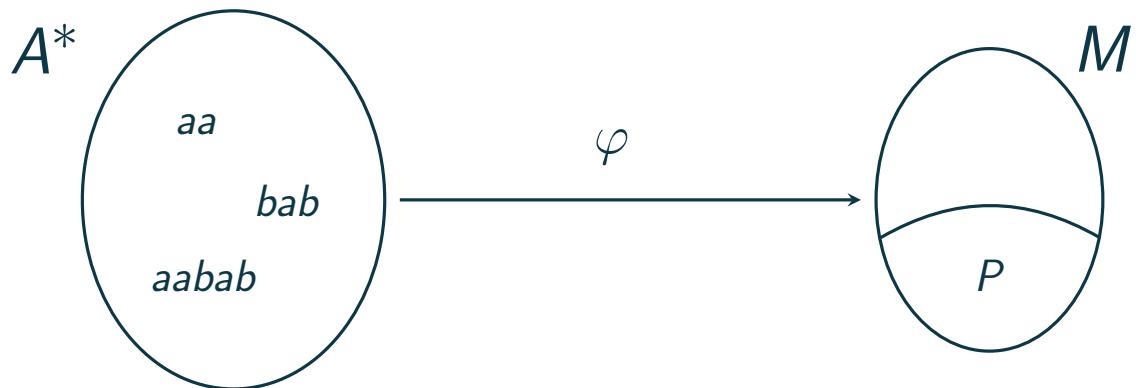
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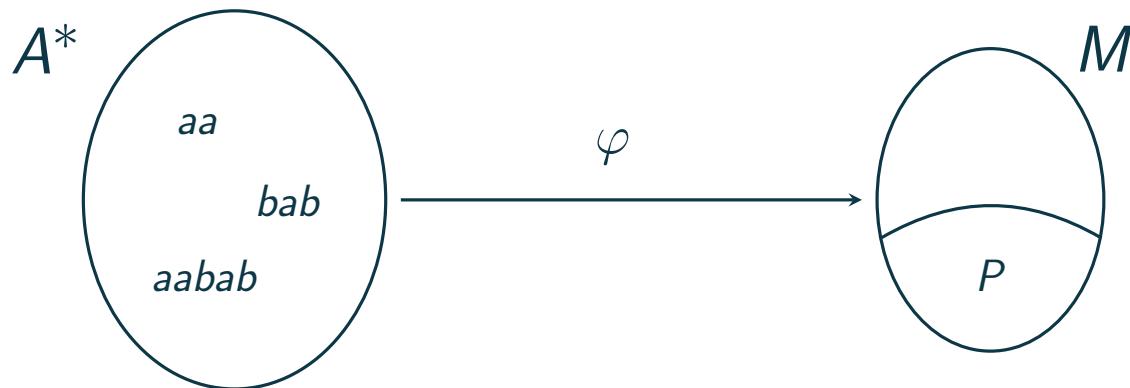
Does the trichotomy holds in general?

## Monoïds



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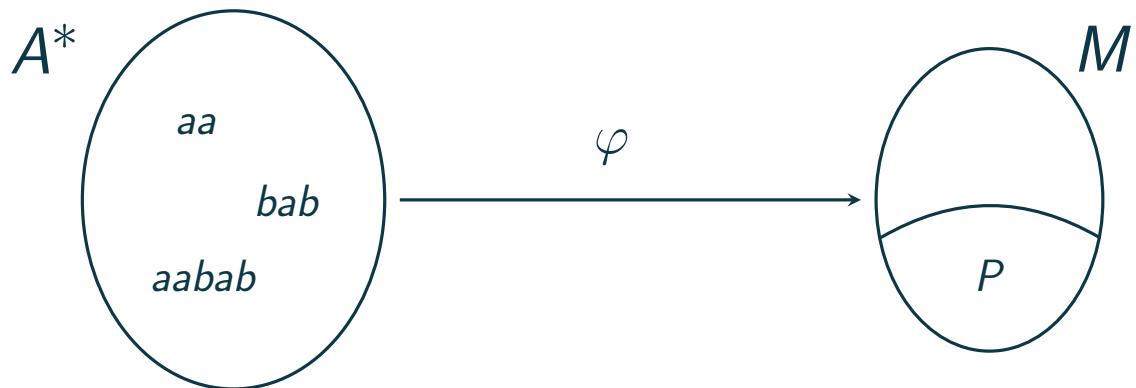


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- Congruence on  $A^*$
- Computed from the minimal automaton
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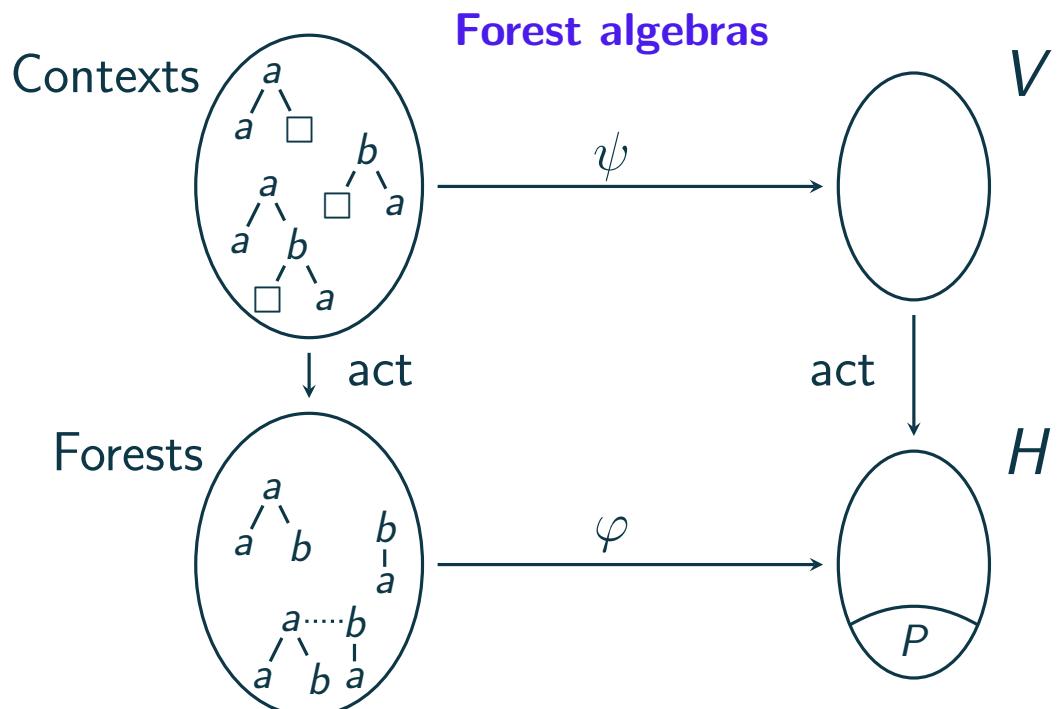
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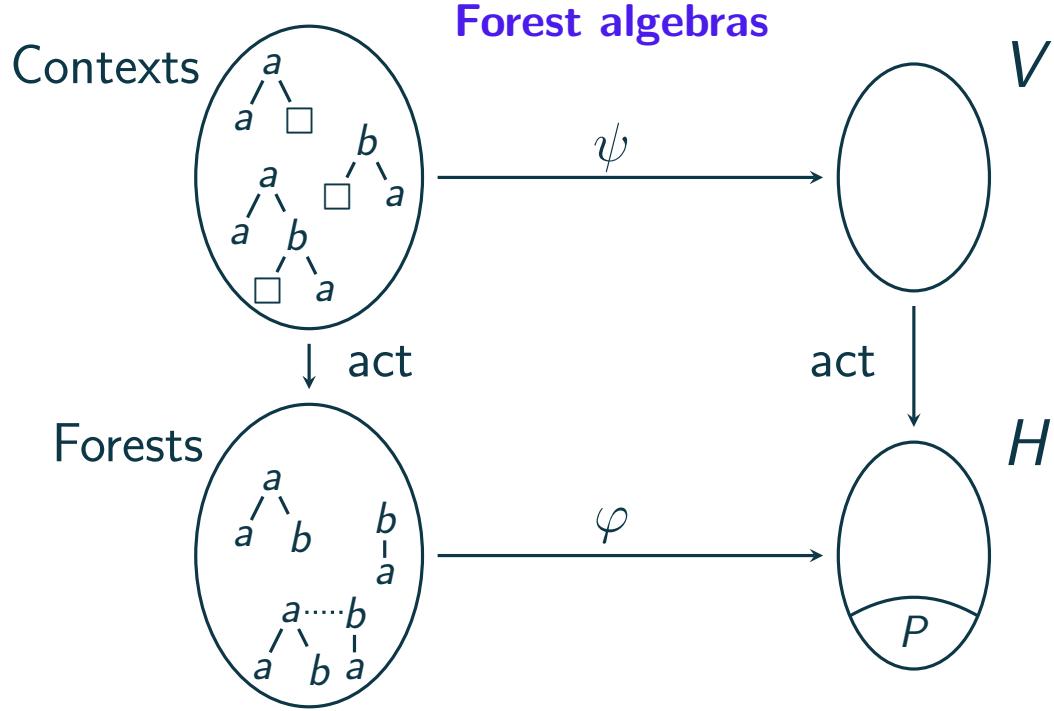
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Varieties: membership to a class can be read on the syntactic monoïd.

Equational theory.



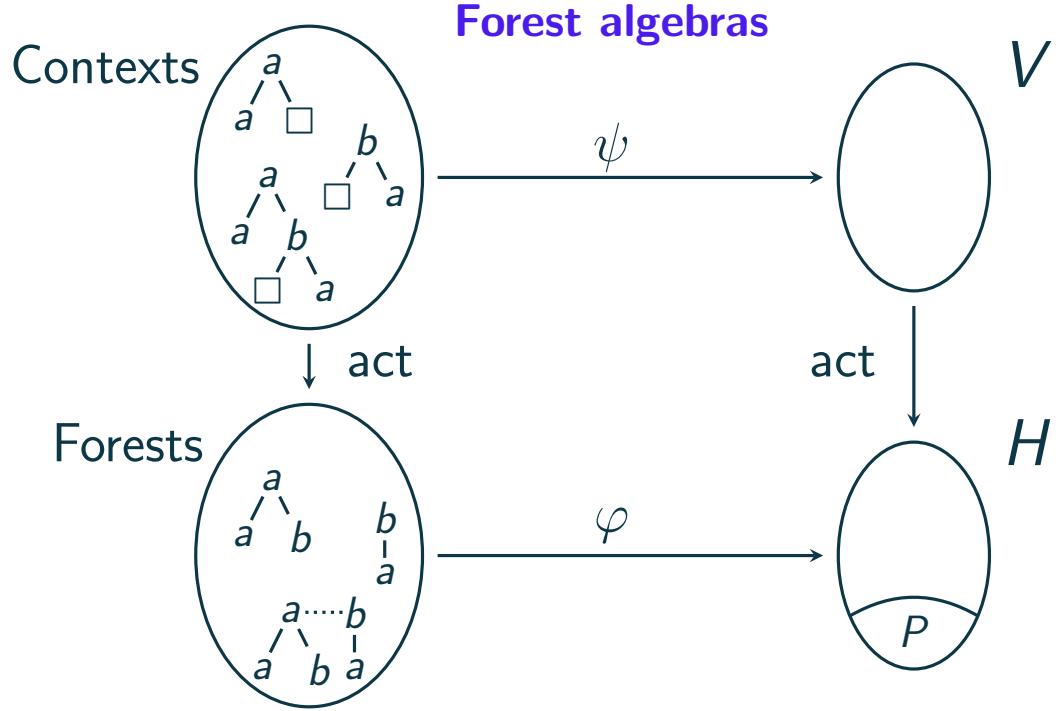
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Syntactic algebra:

- Congruences on contexts and trees Varieties
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## Monoïdal forest algebras

A forest algebra  $(V, H)$  is monoïdal iff:

For all  $v \in V$ , there exists  $x, y \in H$  such that for all  $h \in H$

$$v \cdot h = x + h + y$$

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The class of monoïdal algebras is not a variety.

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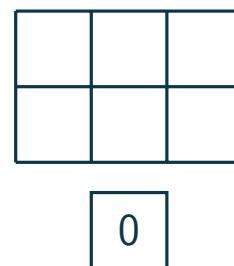
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Solved for algebras with  $H$  that is 0-simple :



## A candidate set of equations

$\mathcal{V}$  is the variety of all forests satisfying:

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Do we need more equations?

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Thanks!