

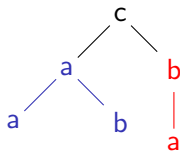
Stackless Processing of Streamed Trees

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Highlights 2021

Processing streamed trees

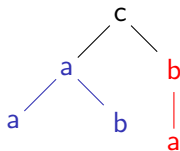
XML encoding of trees:



```
<c>  
  <a>  
    <a></a>  
    <b></b>  
  </a>  
  <b>  
    <a></a>  
  </b>  
</c>
```

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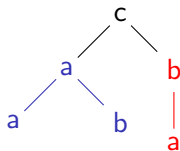


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RPQs: the path from the root belongs to a given regular language.

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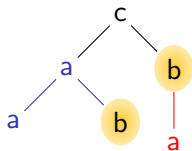


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For instance, the RPQ associated to ca^*b .

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XML encoding of trees:



$\langle c \rangle$
 $\langle a \rangle$
 $\langle a \rangle \langle /a \rangle$
 $\langle b \rangle \langle /b \rangle$
 $\langle /a \rangle$
 $\langle b \rangle$
 $\langle a \rangle \langle /a \rangle$
 $\langle /b \rangle$
 $\langle /c \rangle$

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Evaluation in constant memory VS linear memory

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- ▶ It is very limited: $//a/b$ and $//a//b$ are not doable.
- ▶ All RPQs can be evaluated with a **stack**, but this is costly.

Evaluation in logarithmic memory (Stackless automata)

- Main ingredients:
- a finite **state machine**,
 - a **counter** that stores the current depth in the tree,
 - a finite number of **registers** where the counter values can be stored,
 - can compare register values with the current depth.

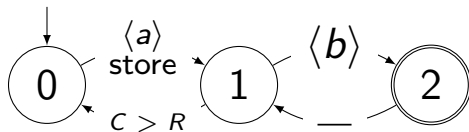
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Evaluating:

$$(a + b + c)^* a (a + b + c)^* b$$

//a//b



Main result

Theorem

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($//a//b$ is doable but still not $//a/b$)

Conclusion

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- ▶ Ongoing work on **vectorization**.