

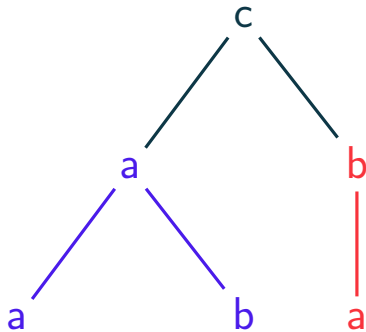
Progress on weak validation of streamed trees

Corentin Barloy

UNIVERSITY OF LILLE

Processing streamed trees

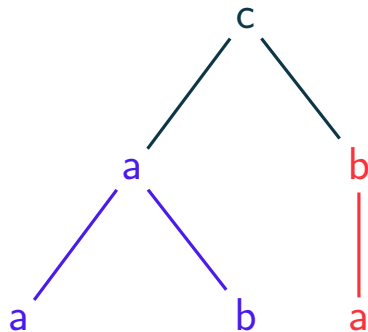
XML encoding of trees:



```
<c>  
  <a>  
    <a></a>  
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  </a>  
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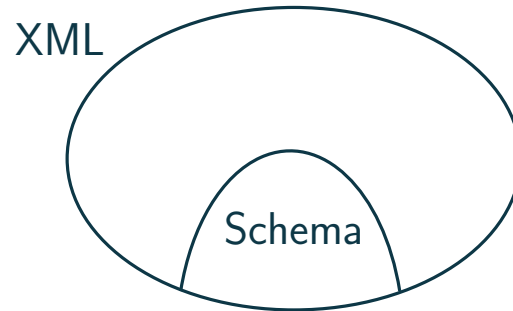
Validation: Does a XML document belongs to a given regular language of trees?

Weak validation

Weak validation: the document is well-formed.

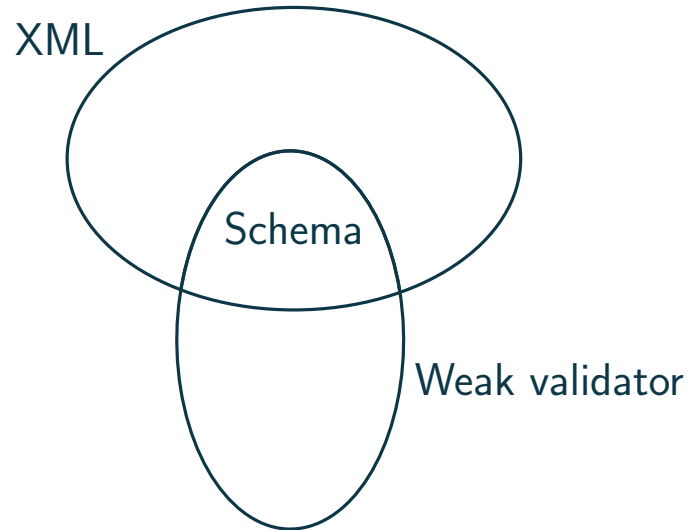
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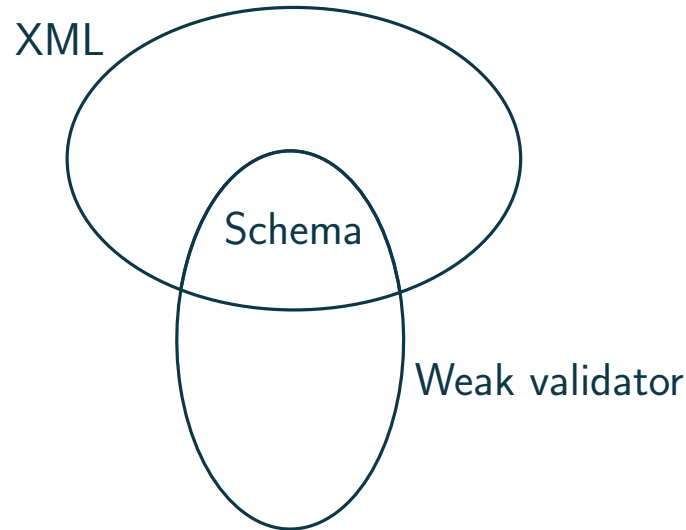
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Question

What is the **space complexity** of **weakly validating** documents in **streaming**?

Space complexity

Segoufin, Vianu:

What can be weakly validated in **constant space**?

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- Space trichotomy between $\Theta(1)$, $\Theta(\log(h))$, $\Theta(h)$.

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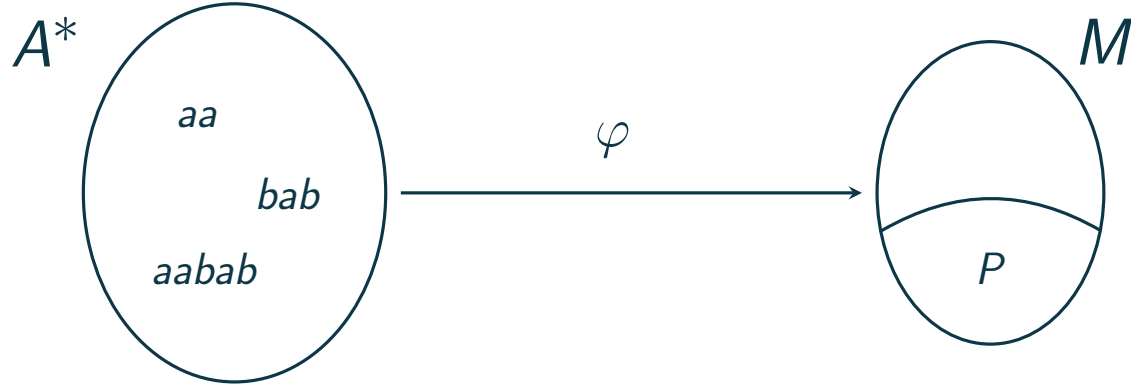
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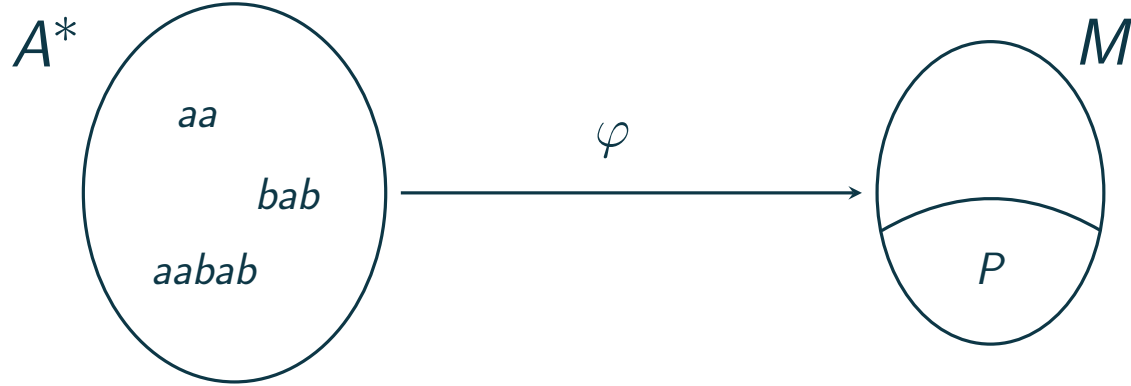
Does the trichotomy holds in general?

Monoïds



\mathcal{L} recognized by M : $\mathcal{L} = \varphi^{-1}(P)$

Monoïds

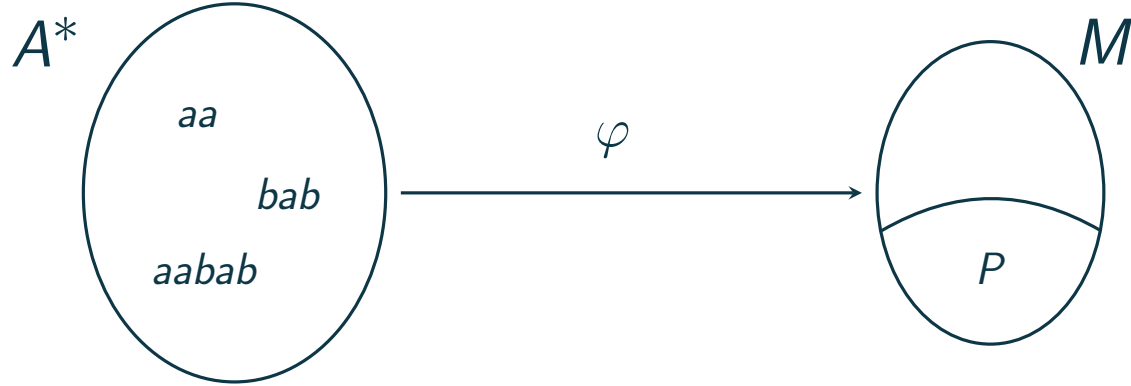


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Syntactic monoïd:

- Congruence on A^*
- Computed from the minimal automaton
- Smallest monoid recognizing \mathcal{L}

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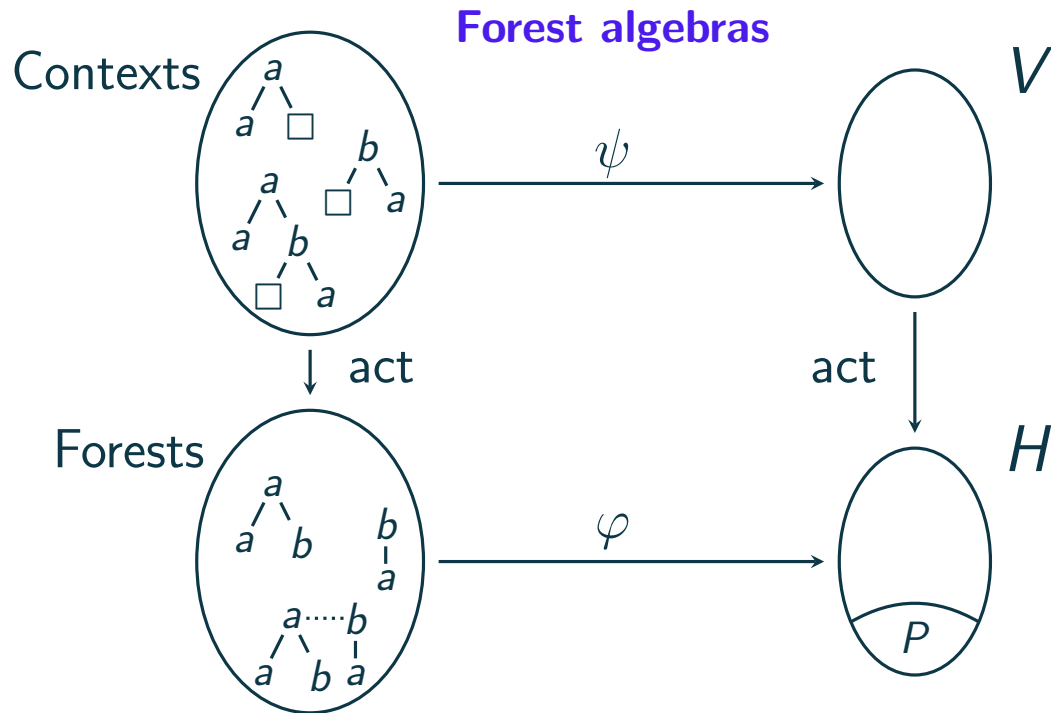
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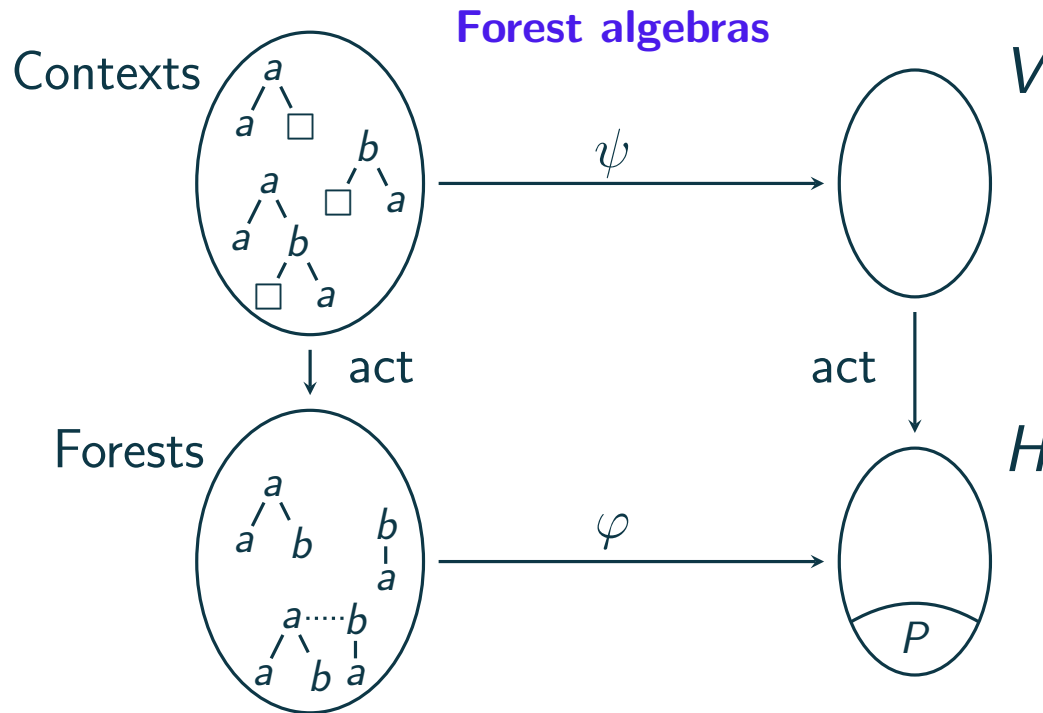
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Varieties: membership to a class can be read on the syntactic monoïd.

Equational theory.



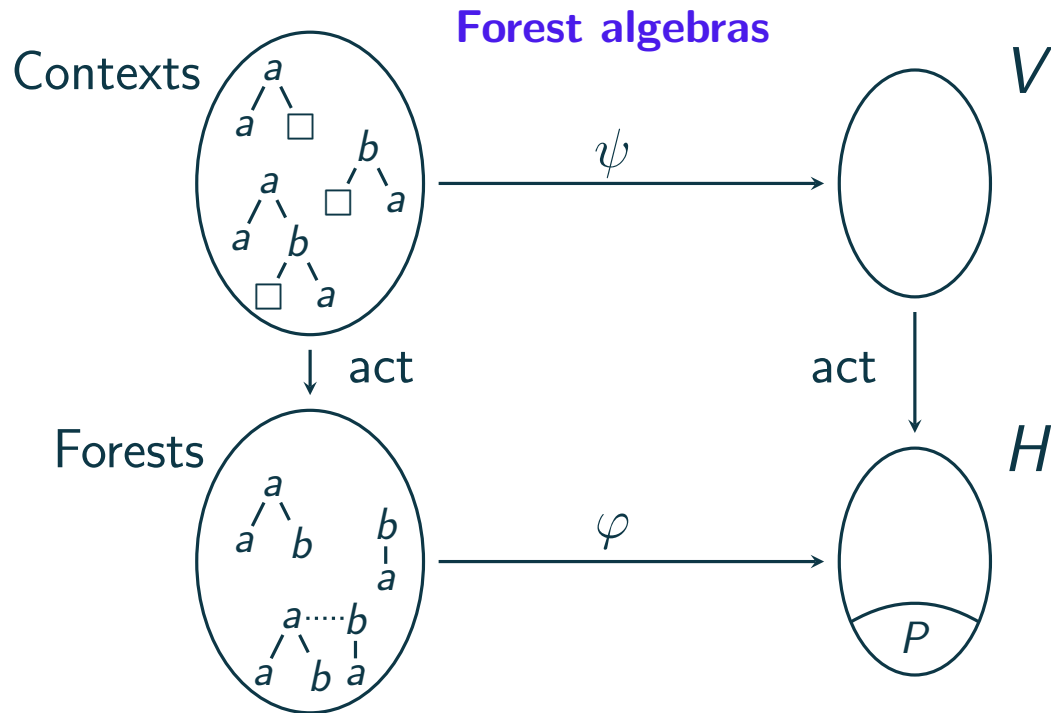
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- **Congruences** on contexts and trees
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Varieties

Equations

Monoïdal forest algebras

A forest algebra (V, H) is monoïdal iff:

For all $v \in V$, there exists $x, y \in H$ such that for all $h \in H$

$$v \cdot h = x + h + y$$

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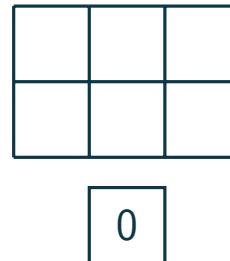
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Solved for algebras with H that is **0-simple** :



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Mimics a **pumping** argument.

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Do we need more equations?

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Thanks!