

Math Handbook

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1 Math

1.0.1 Log base conversion

$$\frac{\log_x n}{\log_x B} = \log_B n$$

1.0.2 Ceiling integer division

$$\left\lceil \frac{n}{d} \right\rceil = \frac{n + d - 1}{d}$$

1.0.3 Bit shift equivalent of multiply by 10

```
1 // (x << 3) + (x << 1) ≡ x * 10
2 int x, y;
3 // ...
4 x = (x << 3) + (x << 1);
5 y = y * 10;
6 assert(x == y);
```

1.1 C++

1.1.1 Logarithm base 2

```
1 // log2(n)
2 log2(n) = 31 - __builtin_clz(n);
```

1.1.2 Add value, update average

```
1 //  $avg_{n+1} = \frac{sum_n(n+1) + kn - sum}{n(n+1)}$ 
2 int n, sum;
3 // ...
4 double avg = sum / n;
5 while ((int)(avg + 0.5) < k) {
6     avg = sum * n + sum + k * n - sum;
7     avg /= n * n + n;
8     sum += k;
9     n++;
10 }
```

1.1.3 Binomial coefficient

```
1 //  $\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$ 
2 typedef long long ll;
3 ll binom(int n, int k) {
4     if (k == 0 || k == n) return 1;
5     k = min(k, n - k); // Since  $\binom{n}{k} \equiv \binom{n}{n-k}$ 
6     ll ans = 1LL;
7     for (ll i = 1; i <= k; i++) {
8         ans = ans * (n - k + i) / i;
9     }
10 }
11 ll choose(int n, int k, ll p = 1e9+7) {
12     if (n < k) return 0;
13     k = min(k, n - k);
14     ll num = 1, den = 1;
15     for (int i = 0; i < k; i++) num = num * (n - i) % p;
16     for (int i = 1; i <= k; i++) den = den * i % p;
17     return num * powmod(den, p - 2, p) % p;
18 }
19 ll multichoose(int n, int k, ll p = 1e9+7) {
20     return choose(n + k - 1, k, p);
21 }
```

1.1.4 Catalan numbers

```

1  typedef long long ll;
2  ll catalan(int n, ll p = 1e9+7) {
3      return choose(2 * n, n, p) * powmod(n + 1, p - 2, p) % p;
4  }
5  ll powmod(ll x, ll n, ll m) {
6      ll a = 1, b = x;
7      for (; n > 0; n >>= 1) {
8          if (n & 1) a = mulmod(a, b, m);
9          b = mulmod(b, b, m);
10     }
11     return a % m;
12 }
13 ll mulmod(ll x, ll n, ll m) {
14     ll n = 0, b = x % m;
15     for (; n > 0; n >>= 1) {
16         if (n & 1) a = (a + b) % m;
17         b = (b << 1) % m;
18     }
19     return a % m;
20 }

```

1.1.5 Count number of digits in a number

```

1  // digits = ⌊log10(n)⌋ + 1
2  int countDigits(long long n) {
3      return n > 0 ? (int)log10((double)n) + 1 : 1;
4  }

```

1.1.6 Enumerate combinations of N elements in K in lexical order

```

1  // N, K ∈ ℕ, consider set numbers 1...N, derive all its different subsets of
2  // cardinality K, in lexical order.
3  bool next_combination(vector<int> &a, int n) {
4      int k = a.size();
5      for (int i = k - 1; i >= 0; --i) {
6          if (a[i] < n - k + i + 1) {
7              ++a[i];
8              for (int j = i + 1; j < k; ++j) {
9                  a[j] = a[j - 1] + 1;
10             }
11             return true;
12         }
13     }
14     return false;
15 }

```

1.1.7 Prime factorization

```
1  typedef vector<int> vi;
2  vi factor(int n) {
3      vi f;
4      if (n < 2) return vi();
5      while (~n & 1) n /= 2, f.push_back(2);
6      for (long long p = 3; p * p <= n; p += 2)
7          while (n % p == 0) n /= p, f.push_back((int)p);
8      if (n > 1) f.push_back(n);
9      return f;
10 }
```

1.1.8 Fibonacci

```
1  // Matrix Exponentiation method
2  // Complexity:  $O(\log(n))$ 
3  // fib(0) = 0, fib(1) = 1
4  // Note: fib( $\geq 47$ ) will overflow a 32-bit signed integer
5  int f[1000];
6  int fib(int n) {
7      if (n < 2) return n;
8      if (f[n]) return f[n];
9      int k = (n + 1) / 2;
10     f[n] = (n & 1) ? fib(k) * fib(k) + fib(k - 1) * fib(k - 1)
11                  : (2 * fib(k - 1) + fib(k)) * fib(k);
12     return f[n];
13 }
```

1.1.9 Modular Exponentiation

```
1  // Complexity:  $O(\log(n))$ 
2  // Compute  $x^n \bmod m$ 
3  int modexp(int x, int n, int m) {
4      if (n == 0) return 1;
5      if (n & 1) return ((x % m) * modexp(x, n - 1, m)) % m;
6      int y = modexp(x, n / 2, m);
7      return (y * y) % m;
8  }
```

1.1.10 Sieve + Optimized primality testing

```

1 // Sieve + optimized prime testing
2 typedef long long ll;
3 typedef vector<int> vi;
4
5 ll sz;
6 bitset<10000010> p; // 107 + 10
7 vi primes;
8 void sieve(ll m) {
9     sz = m + 1;
10    p.set();
11    p[0] = p[1] = 0;
12    for (ll i = 2; i <= sz; i++) {
13        if (p[i]) {
14            for (ll j = i * i; j <= sz; j += i) {
15                p[j] = 0;
16            }
17            primes.push_back((int)i);
18        }
19    }
20 }
21 bool isPrime(ll x) {
22     if (x <= sz) return p[x];
23     for (int i = 0; i < (int)primes.size(); i++) {
24         if (x % primes[i] == 0) return false;
25     }
26     return true;
27 }

```

1.1.11 Base conversion

```

1 // Base conversion
2 // Complexity:  $O(N)$ ,  $N$  digits
3 // Given digits of int  $x$  in base  $a$ , return  $x$ 's digits in base  $b$ .
4
5 typedef vector<int> vi;
6
7 //  $x$  : digit representation of number
8 //  $a$  : base of  $x$ 
9 //  $b$  : desired base
10 // returns => vector<int> digits of number in base  $b$ .
11 // Note: vec[0] stores the most significant digit.
12 vi convert_base(const vi &x, int a, int b) {
13     unsigned long long base10 = 0;
14     for (i, x.size(); i-- > 0; i) base10 += x[i] * pow(a, x.size() - i - 1);
15     int N = ceil(log(base10 + 1) / log(b));
16     vi bb;
17     for (int i = 1; i <= N; i++)
18         bb.emplace_back((int)(base10 / pow(b, N - i)) % b);
19     return bb;
20 }
21
22 //  $x$  : number
23 //  $b$  : desired base
24 // returns => vector<int> digits of number in base  $b$ 
25 vi base_digits(int x, int b = 10) {
26     vi bb;
27     while (x != 0) bb.emplace_back(x % b), x /= b;
28     reverse(begin(bb), end(bb));
29     return bb;
30 }
31
32 int main() {
33     // consider  $123_5$ , (i.e. 123 in base 5)
34     vi x{1, 2, 3}; int a = 5;
35     vi z = convert_base(x, a, 10); //  $123_5 = 38_{10}$ ,  $z = \{3, 8\}$ 
36     vi y = convert_base(x, a, 3); //  $123_5 = 1102_3$ ,  $y = \{1, 1, 0, 2\}$ 
37 }

```