

## Homework 2 - CSCI 181 - S20

For the following questions, when an algorithm is said to be  $O(f(n))$ , for simplicity assume that the time it takes to run this algorithm is  $k \cdot f(n)$  where  $k$  is a positive integer.

1. (20 pts)
  - (a) We use a multiplication algorithm that takes time  $O(\log^2(N))$  to multiply  $N_1 \cdot N_2$  if  $N_1 \approx N_2 \approx N$ . If it takes 3 nanoseconds to multiply two 1000 bit numbers then how long would it take to multiply two 5000 bit numbers?
  - (b) Let  $N$  be a large positive integer. The time it takes to factor  $N$  using trial division is  $O(\sqrt{N})$ . Assume that  $N'$  is a large positive integer and that the binary representation of  $N'$  has ten more bits than that of  $N$ . Assume that it takes 11 nanoseconds to factor  $N$  using trial division. Approximately how long would it take to factor  $N'$  using trial division?
2. (10 pts) Let's say we switch from 1024-bit RSA to 4096-bit RSA. How much longer does decryption take?
3. (20 pts)
  - (a) The sum of integers from 1 to  $N$  is  $N(N+1)/2$ , that is  $1+2+\dots+N = N(N+1)/2$ . Suppose a programmer does not know this formula and wants to compute the sum of integers from 1 to  $N$ . The programmer writes a program by adding 1 and 2 first and then getting the result and adding it by 3, and so on:  $((1+2)+3)+\dots+N$ . Find the running time this will take in terms of  $N$ .
  - (b) Find the running time in terms of  $N$  that it would take you to compute the sum using the formula  $N(N+1)/2$ . You'll see this is much faster.
4. (10 pts) Find the running time required to compute  $6^N$  in terms of  $N$ . The computer program would compute  $((((6 \cdot 6) \cdot 6) \dots) \cdot 6)$ .
5. (10 pts) Find the running time required to compute  $X^N$ , in terms of  $X$  and  $N$ . The computer program would compute  $((((X \cdot X) \cdot X) \dots) \cdot X)$ .
6. (10 pts) Let  $F_n$  denote the  $n$ th Fibonacci number. We have  $F_1 = F_2 = 1$  and for  $i \geq 3$ ,  $F_i = F_{i-1} + F_{i-2}$  (so  $F_3 = 2, F_4 = 3, F_5 = 5, F_6 = 8, \dots$ ). Recall  $F_n \approx \alpha^n$ , where  $\alpha = (1 + \sqrt{5})/2$ . Find the running time of an algorithm that exactly finds the integer  $\prod_{i=1}^n F_i$  using  $((F_1 \cdot F_2) \cdot F_3) \dots F_n$ . Your answer should be  $O()$  of a function of  $n$  and not have an  $F$  in it. Explain your answer. For simplicity, assume that you already have  $F_1, \dots, F_n$  in storage, so you don't have to worry about the time to compute them.