PROGRAMMING

Lecture 17

Dept. of Computer Engineering Hanbat National University

OUTLINE

Divide and conquer Binary search Merge sort

DIVIDE AND CONQUER

"Divide and conquer" is a strategy to solve a problem.

Given a (hard) problem, it **divides** the problem into **sub-problems**, and **solves** the **sub-problems independently**, and finally **combines** their **solutions** to obtain a solution of the original problem.

Julius Caesar(BC 100 - 44): "divide and impera" (or divide and rule) to govern the European continent.

The British practiced this idea to control the Indian subcontinent(1858 - 1947).

Benjamin Franklin(1706 - 1790): "We must all **hang together**, or assuredly we shall all **hang separately**."

Sequential Search

```
L = [("John", 20), ("Mary", 19), ("David", 22), ("James", 21), ("Young", 19), ("Susan",22)]
```

Given the name of a **person** and an **unsorted list**, report his/her age.

```
def seq_search(L, qname):
    for i in range(len(L)):
        name, age = L[i]
        if qname == name:
            return age
    return -1
seq_search(L, "David")
```

How many comparisons? At most n comparisons, where n = len(L)!

What if L is sorted?

Try "Jeff." You may stop after comparing it with "John". Why?

```
L = [(David, 22), (James, 21), (John, 20), (Mary, 19), (Susan, 22)], (Young, 19)]
```

However, we need n comparisons in the worst case, anyway, where n = len(L). Why?

BINARY SEARCH

Basic idea

Assumption: the list L is sorted.

With a single comparison, we can reduce the search space by half. **Divide and conquer!**

23

(low = high = mid)

Observations

1. When **low** == **high**, one **additional comparison** is required to make the final decision. At this time, the **length of the sub-list** is reduced to **one**.

base case

2. As long as **high > low**, the **length of the sub-list** reduced **by half** with a single comparison.

recursive case

With a **single comparison**, the **search space** is **reduced by half**. In order to make the **final decision**, the **search space** should be reduced to **one**.

How many comparisons are required to make the final decision? At most $log_2 n$. Why?

```
n/2^k = 1 \implies k = \log_2 n, where n = len(L).

log_2 n + 1

to make the final decision to reduce the search space
```

```
def b_search(L, q, low, high):
      if low == high:
       return L[low] == q
                                                     Base case
      mid = (low + high) / 2
      if L[mid] == q:
                                                     Recursive case
         return True
      elif L[mid] < q:
          return b_search(L, q, mid+1, high)
      e se:
                                                  Why this?
          if low != mid:
                                                  Try L = [5, 7] and q = 3!
             return b_search(L, q, low, mid-1)
          else:
             return False
```

MERGE SORT

Basic Idea

1. If the list is of length 0 or 1, it is already sorted.

the base case

2. Otherwise, divide the list into two sub-lists of (almost) equal size and sort each of them recursively by using merge sort.

the recursive case

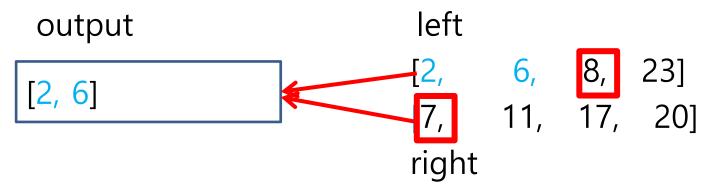
3. Merge the results.

Divide and conquer

[23,	2,	8,	6, 17,	11,	20,	7]	
[23,	2,	8,	6] [17,	11,	20,	7]	divide
[23,	2]	[8,	6] [17,	11]	[20,	7]	arvide
[23]	[2]	[8]	[6] [17]	[11]	[20]	[7]	
[2,	23]	[6,	8] [11,	17]	[7,	20]	
[2,	6,	8,	23] [7,	11,	17,	20]	merge
[2,	6,	7,	8, 11,	17,	20,	23]	

```
def merge_sort(L):
   if len(L) < 2:
return L[:] base case
   mid = len(L) / 2
   left = merge_sort(L[:mid ])
                                 recursive case
   right = merge_sort(L[mid:])
   return merge(left, right)
L = [23, 2, 8, 6, 17, 11, 20, 7]
L = merge_sort(L)
print L
```

How to merge to sorted lists



How many comparisons?

len(left) + len(right)
$$- 1 = n - 1$$
, where $n = len(L)$. Why?

At most **n comparisons**.

```
def merge(left, right):
  result = []
  i, j = 0, 0
  while i < len(left) and j < len(right):
     if left[i] < right[j]:
       result.append(left[i])
       i += 1
     else:
       result.append(right[j])
       i += 1
  while i < len(left):
     result.append(left[i])
                                No more elements in right
     i += 1
  while j < len(right):
     result.append(right[j])
                                No more elements in left
     i += 1
  return result
```

```
def merge_sort(L):
   if len(L) < 2:
      return L[:]
   mid = len(L) / 2
   left = merge_sort(L[:mid])
                                  How many comparisons?
   right = merge_sort(L[mid:])
   return merge(left, right)
L = [23, 2, 8, 6, 17, 11, 20, 7]
L = merge_sort(L)
```

print L

```
[23, 2, 8, 6, 17, 11, 20,
                               7]
[23,
    2, 8, 6] [17, 11,
                          20,
                                     divide
         [8, 6] [17, 11] [20, 7]
[23,
    2]
[23]
    [2] [8] [6]
                [17]
                          [20]
     23]
         [6,
             8] [11,
                      17]
                               20]
                          [7,
     6, 8, 23] [7, 11,
                         17,
                               20]
                                     merge
     6, 7, 8, 11, 17, 20,
                               23]
[2,
```

How many rounds?
How many comparisons per round?

How many rounds?

$$n/ 2^k = 1$$

 $2^k = n \implies k = \log_2 n$

How many comparisons at each round?

At most n comparisons

Therefore, at most $n \log_2 n$ comparisons

 $O(n log_2 n)$

Behavior of the merge sort.

The **dividing** and merging **phases** are **mixed** unlike the figure in the previous slide!

```
def merge_sort(L):
   global lv
   lv += 1
   if len(L) < 2:
      display(L)
      <u>lv -= 1</u>
      return L[:]
   display(L)
   mid = len(L) / 2
   left = merge_sort(L[:mid])
   right = merge_sort(L[mid:])
   result = merge(left, right)
   display(result)
   lv -= 1
   return result
                           1_{V} = -1
def display(L):
   if |v| = 0:
                           L = [23,2,6,8,17,11,20,7]
     print L
                           merge_sort(L)
   else:
     print " " * 4 * lv, L
```

```
[23, 2, 6, 8, 17, 11, 20, 7]
     [23, 2, 8, 6]
          [2, 23]
              [23]
              [2]
         [2, 23]
         [8 6]
              [8]
              [6]
         [6, 8]
     [2, 6, 8, 23]
     [17, 11, 20, 7]
         [17, 11]
              [17]
              [11]
         [11, 17]
         [20, 7]
              [20]
              [7]
         [7, 20]
     [7, 11, 17, 20]
[2, 6, 7, 8, 11, 17, 20, 23]
```