

MI3 DIALOG Future Research



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WP1: Interpretation of Informal Mathematical Input

Robust Sentence-Level Analysis

• Processing "ill-formed" input (syntactic errors, incompleteness, out-of-grammar) through combination of deep and shallow methods:

In diesem Fall: z.B. $K(A) = dem Begriff K(A \cup B)$

• Extension from written-only to simultaneous written and spoken input, accompanied with simple pointing/selection on screen

Discourse Representation of Informal Mathematical Input

Und wenn $B \subseteq K(A)$ sein soll, muss es auch Element von K(A) sein.

- Coreference of symbolic identifiers
- Anaphoric reference to parts of mathematical expressions
- Discourse structure as reflex of proof structure

Ontology-Based Domain-Specific Interpretation

• Informal and/or imprecise naming of domain concepts and relations:

A muss in B sein

... B vollstaendig ausserhalb von A liegen muss ...

...dann sind A und B vollkommen verschieden

• Semantically complex operators:

Wenn alle A in K(B) enthalten sind und dies auch umgekehrt gilt, ...

WP2: Proof Management and Proof Step Evaluation (PSE)

Abstract-level Proof Representation

 Required for PSE: cognitive oriented proof representation PSE example scenario

Assertions already introduced (A1) $A \wedge B$.

 $(A2) A \Rightarrow C.$ $(A3) C \Rightarrow D.$

 $(A4) F \Rightarrow B.$ $(G) D \lor E.$

Alternative proof step directives. (a) From the context follows D.

(b) B holds.

(c) It is sufficient to show D. (d) We show E.

PSE: Novel Theorem Proving Application

Criterion	Task (first approach)	Requirements for theorem prover		
Soundness	$E \vdash_C^? D \lor E$	'Yes' or 'No' answer; any theorem prover resp. calculus C		
Granularity	$proof ext{-steps}(E \vdash^?_C D \lor E)$	adequate abstract-level theorem prover resp. calculus C; measure 'shortest' proof; take tutorial constraints into account; proof planning or assertion level reasoning?		
Relevance	$A \wedge B$ $A \Rightarrow C$ $C \Rightarrow D \vdash_{C}^{?} E$ $F \Rightarrow B$	recognize detours; compare with other 'shorter' proofs; take tutorial constraints into account; forward case more challenging		

WP3: Domain Reasoning for Ambiguity Resolution

An example

Discourse:

- (1) From previous observations we know that A or B.
- (2) The former implies D by Lemma X.
- (3) Similarly, from the latter follows C.

Alternative user utterances with underspecification:

- (a) From this follows D since C implies D by Lemma Y.
- this may refer to (1)+(2)+(3), to (3), or even (2) with wrong justification
- (b) It holds D since C implies D by Lemma Y.
- no underspecified anaphoric reference but ambiguity at domain reasoning level
- Ambiguities may arise at linguistic and domain reasoning level.
- Ambiguities are resolved by a combination of linguistic processing and proof step evaluation.
- Remaining ambiguous readings are explicitly represented and ranked.
- The use of underspecification techniques (CHORUS) will be explored.

Further ambiguity examples

Example	Where does ambiguity arise?	Ambiguity resolution means
(1) $x \in B$ und somit $x \subseteq K(B)$ und $x \subseteq K(A)$ wegen Vorraussetzung	linguistic meaning level;	linguistic means;
(2) A enthaelt B	attachment, coordination	type checking in (2)
(3) $P((A \cup C) \cap (B \cup C)) = PC \cup (A \cap B)$	linguistic meaning level;	type checking for (3);
(4) $K((A \cup C) \cap (B \cup C)) = KC \cup (A \cap B)$	informal character of discourse	mathematical domain reasoning for (4)
(5) T1: Bitte zeigen Sie: $K((A \cup B) \cap (C \cup D)) = (K(A) \cap K(B)) \cup (K(C) \cap K(D))$!		
S1: nach deMorgan-Regel-2 ist $K((A \cup B) \cap (C \cup D)) = (K(A \cup B) \cup K(C \cup D))$	underspecified proof step	mathematical domain reasoning

