
How to build a (resolution) prover?

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Problem?

Dealing with millions of clauses ...

Efficient automated theorem prover by

good theory

- + efficient implementation
- + clever heuristics

References: [Voronkov, IJCAR, 2001]

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Proof Search by Saturation

Given: Set of clauses

+ Inferences (resolution, factorization, paramodulation)

⇒ Saturate clause set with all possible inferences

1. if empty clause in clause set, then terminate
2. select clauses
3. apply inferences to selected clauses
4. add result to clause set; goto 1

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Theory

Progress in the theory of resolution-based systems:

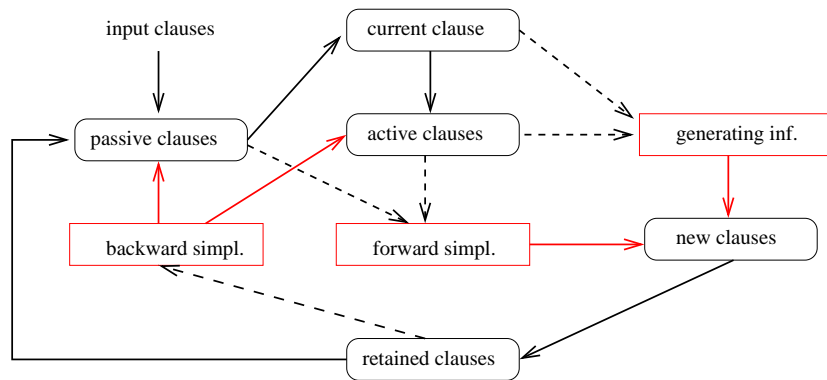
- reduction of possible inferences
without impair of completeness
(e.g., superposition calculus)
- decision procedures for some fragments
(e.g., realized in BLIKSEM)

References: [Bachmair+Ganzinger, Handbook of Aut. Reas. I]
[Nieuwenhuis+Rubio, Handbook of Aut. Reas. I]
[Bachmair+Ganzinger, Diverse CADE, LPAR]

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Implementation

Organization

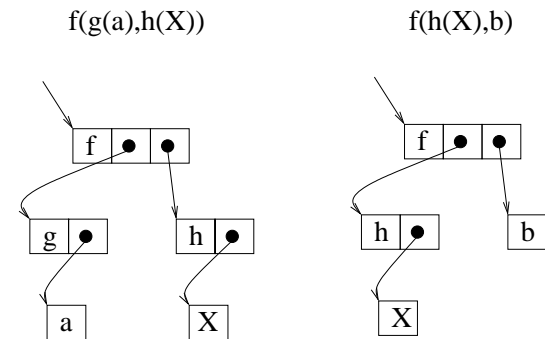


References: [McCune, OTTER 3.0 Manual, 1994]

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Implementation

Efficient Data Structures

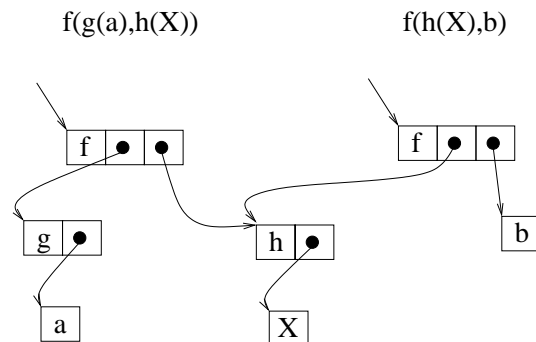


tree representation

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Implementation

Efficient Data Structures

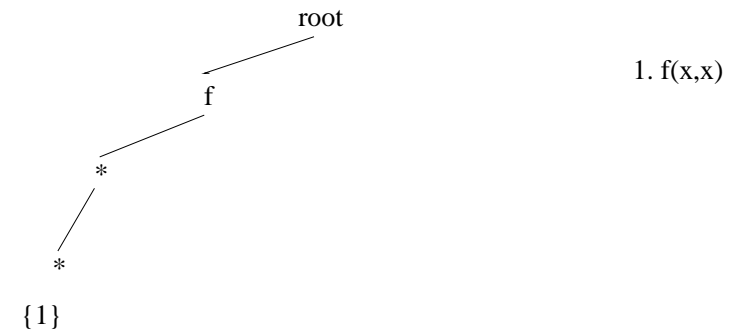


shared tree-like representation

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Implementation

Efficient Data Structures

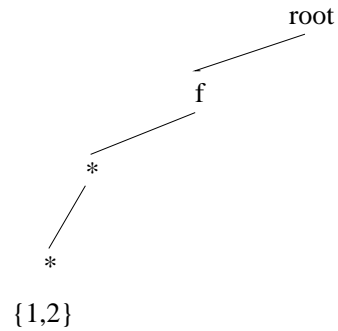


discrimination tree index

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Implementation

Efficient Data Structures



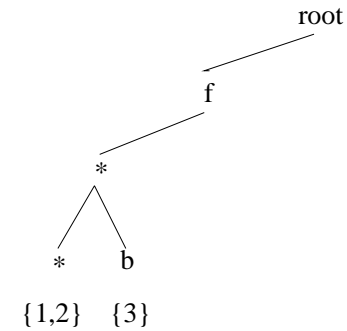
1. $f(x,x)$
2. $f(x,y)$

discrimination tree index

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Implementation

Efficient Data Structures



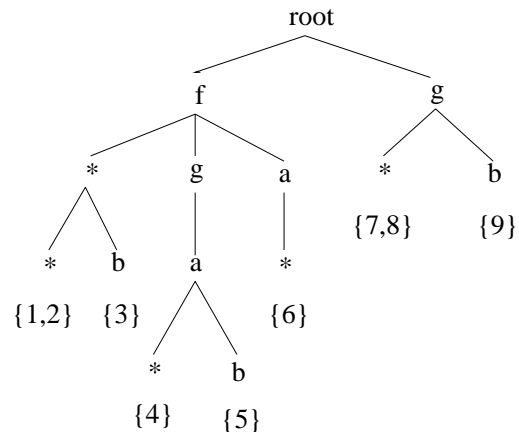
1. $f(x,x)$
2. $f(x,y)$
3. $f(x,b)$

discrimination tree index

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Implementation

Efficient Data Structures



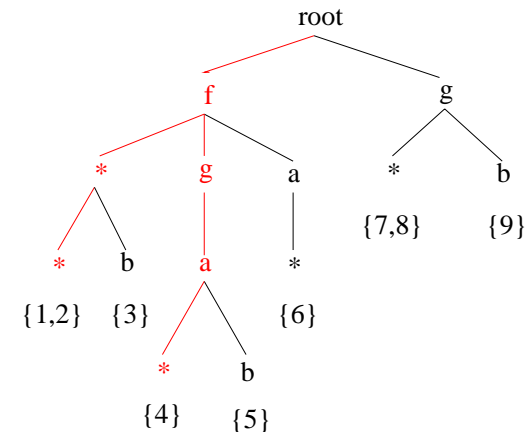
1. $f(x,x)$
2. $f(x,y)$
3. $f(x,b)$
4. $f(g(a),x)$
5. $f(g(a),b)$
6. $f(a,y)$
7. $g(x)$
8. $g(z)$
9. $g(b)$

discrimination tree index

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Implementation

Efficient Data Structures



1. $f(x,x)$
2. $f(x,y)$
3. $f(x,b)$
4. $f(g(a),x)$
5. $f(g(a),b)$
6. $f(a,y)$
7. $g(x)$
8. $g(z)$
9. $g(b)$

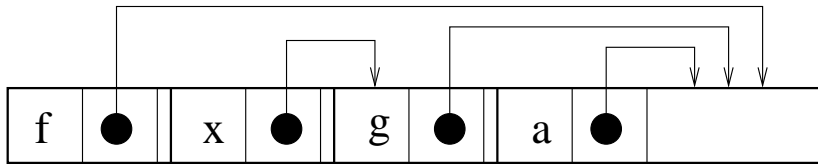
retrieving variants of: $f(g(x), a)$

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Implementation

Efficient Data Structures

$f(x, g(a))$



array-based flat-term representation

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Control

parameterized algorithms

(e.g., 32 parameters to determine OTTER's main algorithm)

⇒ selection of the “right” parameterization is crucial

either by the user

or automatically by the system

For instance:

- OTTER: auto mode
- VAMPIRE: preprocessor
- WALDMEISTER: self-adaption component

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Implementation

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References: [Graf, LNCS 1053, 1996]

[McCune, JAR, 1992]

[Ramakrishnan+Sekar+Voronkov,

Handbook of Aut. Reas. II]

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Control

parameterized algorithms

(e.g., 32 parameters to determine OTTER's main algorithm)

selection of the “right” parameterization is crucial

OR: try different instances competitively

- RCTHEO randomized decision points in SETHEO [Ertel, LPAR, 1992]
- SICOTHEO pre-defined strategies in SETHEO [Schumann, 1995]

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