CPSC 221 Basic Algorithms and Data Structures

May 22, 2015

Administrative stuff

Hw1 posted. Hw2 to come.

Office hours:

```
Monday - 9:30am to 10:30am | Shu Yang | Demco Table 3 (11:00 to noon later)
Tuesday - 9:00am to 10:00am | Kamil Khan | Demco Table 2
Tuesday - 10:30am to 12:00pm | Nasa Rouf | Demco Table 3 (really 10:45-11:45)
Tuesday - 12:00pm to 2:00pm | Henry Li | Demco Table 3
Tuesday - 3:00pm to 4:00pm | Kaitlyn Melton | Demco Table 3
Tuesday - 5:00pm to 6:00pm | Angad Kalra | Demco Table 1
Wednesday - 9:00am to 10:00am | Issam Laradji | Room ICCSX337
Thursday - 10:00am to 12:00pm | Devon | Room ICCSX141
Thursday, - 12:00pm to 2:00pm | Kurt Eiselt | Room ICCSX151
Thursday - 2:00pm to 3:00pm | Jason | Demco Table 3
```

Latest info always available at https://my.cs.ubc.ca/students/ta-hours

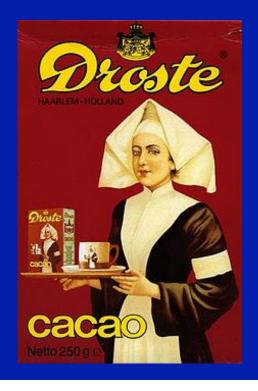
Having survived previous courses, you've had some experience with recursion.

Students aren't always well-versed in how recursion works, or what the resource implications may be.

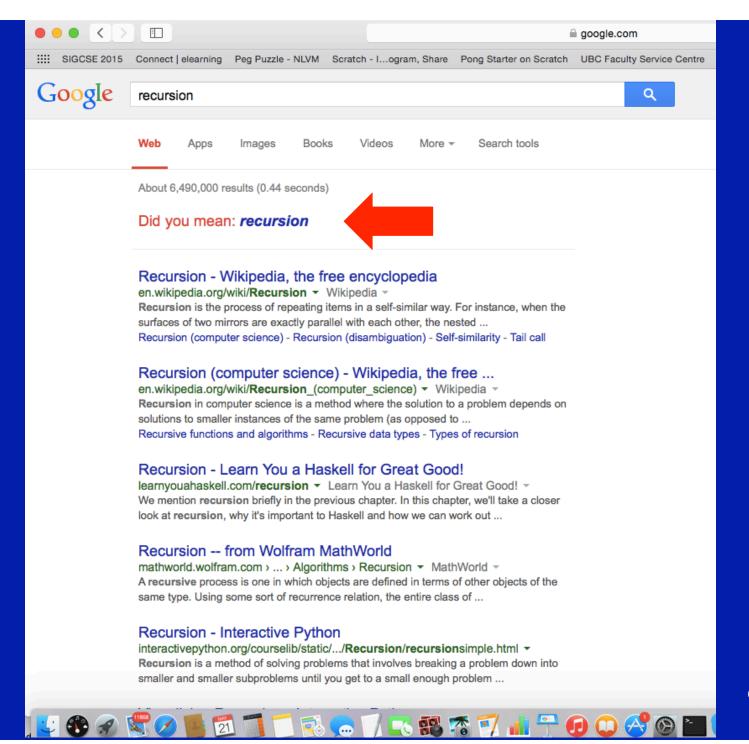
Recursion is a means of obtaining repetitive behaviour from a procedure that is defined in terms of itself. The procedure is self-referential. (That last part should sound familiar to many of you.)

Recursion is a means of obtaining repetitive behaviour from a procedure that is defined in terms of itself. The procedure is self-referential. (That last part should sound familiar to many of you.)

Recursion exists outside of computing.







```
PHRASAL RULES:
S -> NP VP
NP -> DET N
NP -> PN
NP -> NP PP
VP -> VI
VP -> VT NP
VP -> VP PP
PP -> PREP NP
LEXICAL RULES:
PN
   -> John Mary
DET -> the
N -> man | dog | lake | house
VI -> runs
VT
   -> likes | sees
PREP -> with | at | on | by
John sees the man.
John sees the man with the dog.
John sees the man with the dog at the house.
John sees the man with the dog at the house by the lake.
. . .
```

```
PHRASAL RULES:
S -> NP VP
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NP -> PN
                                                           NP
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                                  John
                                                         the man
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```

```
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                                                VT
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                                                       NP
                                                                PP
VP -> VI
                                  John
                                                sees
VP -> VT NP
                                                               with the dog
                                                     the man
VP -> VP PP
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LEXICAL RULES:
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     -> John | Mary
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    -> man | dog | lake | house
   -> runs
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John sees the man.
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John sees the man with the dog at the house.
John sees the man with the dog at the house by the lake.
```

```
PHRASAL RULES:
S -> NP VP
NP -> DET N
NP -> PN
                                                 VT
                                                            NP.
NP -> NP PP
                                                        NP
                                                                 PP
VP -> VI
                                   John
                                                 sees
                                                                at the house
VP -> VT NP
VP -> VP PP
                                                   NP
                                                            PP
PP -> PREP NP
                                               the man
                                                          with the dog
LEXICAL RULES:
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John sees the man.
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John sees the man with the dog at the house.
John sees the man with the dog at the house by the lake.
. . .
```

```
PHRASAL RULES:
S -> NP VP
NP -> DET N
NP -> PN
                                                VT
                                                           NP
NP -> NP PP
                                                       NP
                                                                 PP
VP -> VI
                                   John
                                                sees
                                                                by the lake
VP -> VT NP
VP -> VP PP
                                                   NP
                                                           PP
PP -> PREP NP
                                                           at the house
LEXICAL RULES:
                                               NP
                                                       PP
PN
     -> John
               Mary
DET -> the
                                            the man
                                                      with the dog
     -> man | dog | lake | house
   -> runs
VI
VT
     -> likes
              sees
PREP -> with | at | on | by
John sees the man.
John sees the man with the dog.
John sees the man with the dog at the house.
John sees the man with the dog at the house by the lake.
```

Recursion is also an approach to problem solving where a problem is split into smaller instances of itself. The splitting continues until a trivial, easy-to-solve instance is reached, which can then be used to build up a solution to the original problem.

As we explore data structures, we'll see that recursion is a very useful and powerful tool.

Thinking recursively

Do not start with code. Write the *story* of the problem, including the data definition!

Define the problem: What should be done given a specific input?

Solve some example cases by hand.

Identify and solve the (usually simple) base case(s).

Figure out how to break the complex cases down in terms of any smaller case(s). For the smaller cases, call the function recursively and assume it works. Do not think about how!

Implementing recursion

Once you have all that, write out your solution in comments (i.e., make your own "template"). Then fill out the code and test.

Should be easy. If it's hard, maybe you're not assuming your recursive call works (i.e., thinking too much, not abstracting enough).

A recursive procedure consists of three parts:

- 1 The base case or termination condition. Usually the first thing done upon entering a recursive procedure
- 2 The reduction step -- the operation that moves the computation closer to the termination condition
- 3 The recursive procedure call itself

Finding the length of a list:

Data definition is just your basic list.

Return the number of elements in a list.

What's the length of the empty list?

What's the length of '(a b c)?

What's the base case?

How does it break down into smaller parts?

What's the reduction?

Where does the recursive call go?

Finding if some item is an element of a list:

```
Data definition is just your basic list.

Return the list that begins with the found element (other versions just return true)

What do you return if the list is empty?

What do you return for (element `b `(a b c))?

What do you return for (element `x `(a b c))?

What's the base case?

How does it break down into smaller parts?

What's the reduction?

Where does the recursive call go?
```

base case/termination

Sometimes the solution is handed to you on a platter.

The mathematical definition for factorial:

$$n! = \begin{cases} 1 \text{ for } n = 0 \\ \\ n * (n - 1)! \text{ for } n > 0 \end{cases}$$

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}
```

base case/termination

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}
```

Activation stack

Each function call (whether recursive or not) generates a stack frame (also known as an activation frame or activation record) holding local variables and the program point to return to, which is pushed on a stack (the activation stack or call stack) that tracks the current chain of function calls.

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl; // using 'fact' for brevity</pre>
```

How does it work? The activation stack model:

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

fact(4)

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(4)
4 * fact(3)
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(3)

fact(4)
4 * fact(3)
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

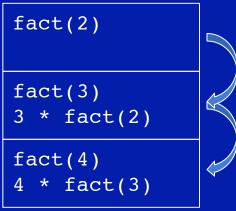
cout << fact(4) << endl;</pre>
```

```
fact(3)
3 * fact(2)

fact(4)
4 * fact(3)
```

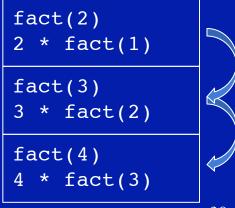
```
int factorial(int n)
{
   if (n == 0)
     return 1;
   else
     return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```



```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```



```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(1)

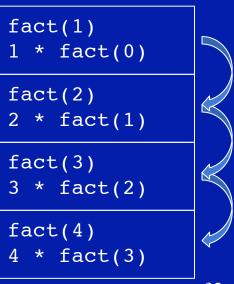
fact(2)
2 * fact(1)

fact(3)
3 * fact(2)

fact(4)
4 * fact(3)
```

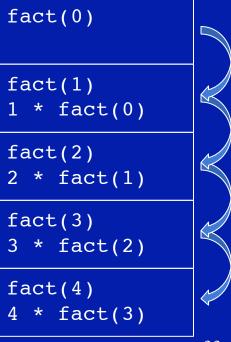
```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```



```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```



```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(0)
1

fact(1)
1 * fact(0)

fact(2)
2 * fact(1)

fact(3)
3 * fact(2)

fact(4)
4 * fact(3)
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(1)
1 * 1

fact(2)
2 * fact(1)

fact(3)
3 * fact(2)

fact(4)
4 * fact(3)
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(1)
1

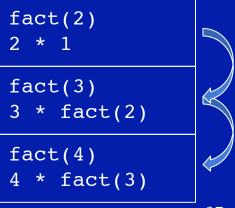
fact(2)
2 * fact(1)

fact(3)
3 * fact(2)

fact(4)
4 * fact(3)
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```



```
int factorial(int n)
{
   if (n == 0)
     return 1;
   else
     return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(2)
2

fact(3)
3 * fact(2)

fact(4)
4 * fact(3)
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(3)
3 * 2

fact(4)
4 * fact(3)
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(3)
6

fact(4)
4 * fact(3)
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(4)
4 * 6
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
fact(4)
24
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;
24</pre>
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;
  4 * fact(3)</pre>
```

```
int factorial(int n)
{
   if (n == 0)
     return 1;
   else
     return n * factorial(n - 1);
}

cout << fact(4) << endl;
   4 * fact(3)
   4 * 3 * fact(2)</pre>
```

```
int factorial(int n)
{
   if (n == 0)
     return 1;
   else
     return n * factorial(n - 1);
}

cout << fact(4) << endl;
   4 * fact(3)
   4 * 3 * fact(2)
   4 * 3 * 2 * fact(1)</pre>
```

```
int factorial(int n)
{
   if (n == 0)
     return 1;
   else
     return n * factorial(n - 1);
}

cout << fact(4) << endl;
   4 * fact(3)
   4 * 3 * fact(2)
   4 * 3 * 2 * fact(1)
   4 * 3 * 2 * 1 * fact(0)</pre>
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;
  4 * fact(3)
  4 * 3 * fact(2)
  4 * 3 * 2 * fact(1)
  4 * 3 * 2 * 1 * fact(0)
  4 * 3 * 2 * 1 * 1</pre>
```

```
int factorial(int n)
 if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
cout << fact(4) << endl;</pre>
  4 * fact(3)
  4 * 3 * fact(2)
  4 * 3 * 2 * fact(1)
  4 * 3 * 2 * 1 * fact(0)
  4 * 3 * 2 * 1 * 1
  4 * 3 * 2 * 1
```

```
int factorial(int n)
 if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
cout << fact(4) << endl;</pre>
  4 * fact(3)
  4 * 3 * fact(2)
  4 * 3 * 2 * fact(1)
  4 * 3 * 2 * 1 * fact(0)
  4 * 3 * 2 * 1 * 1
  4 * 3 * 2 * 1
  4 * 3 * 2
```

```
int factorial(int n)
 if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
cout << fact(4) << endl;</pre>
  4 * fact(3)
  4 * 3 * fact(2)
  4 * 3 * 2 * fact(1)
  4 * 3 * 2 * 1 * fact(0)
  4 * 3 * 2 * 1 * 1
  4 * 3 * 2 * 1
  4 * 3 * 2
  4 * 6
```

```
int factorial(int n)
 if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
cout << fact(4) << endl;</pre>
  4 * fact(3)
  4 * 3 * fact(2)
  4 * 3 * 2 * fact(1)
  4 * 3 * 2 * 1 * fact(0)
  4 * 3 * 2 * 1 * 1
  4 * 3 * 2 * 1
  4 * 3 * 2
  4 * 6
  24
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;
24</pre>
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

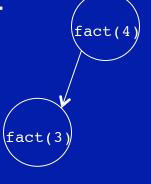
```
fact(4)
```

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

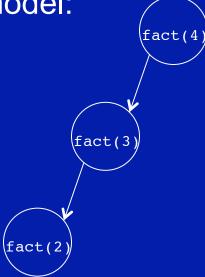
```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```



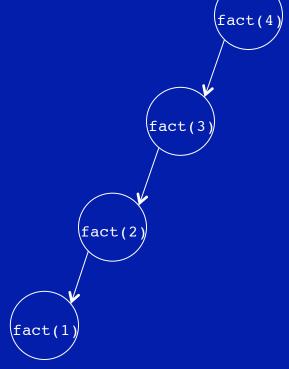
```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```



```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```



How does it work? The recursion tree model: fact(4 int factorial(int n) if (n == 0)return 1; fact(3) else return n * factorial(n - 1);fact(2) cout << fact(4) << endl;</pre> fact(1

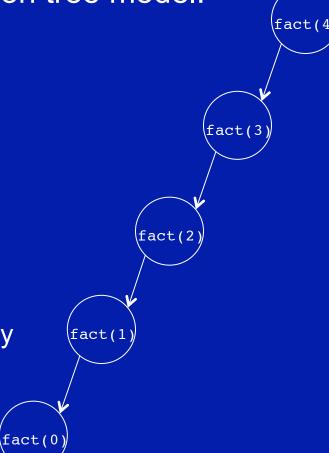
fact(0

How does it work? The recursion tree model:

```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}

cout << fact(4) << endl;</pre>
```

Seriously, this visualization isn't nearly as useful as the others in the case of factorial, but it will give a good idea of how much work has to be done with the next example of recursion.

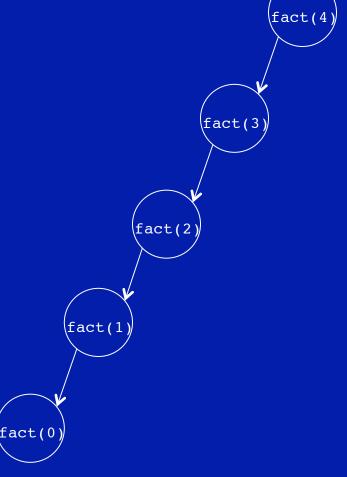


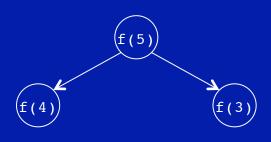
How does it work? The recursion tree model:

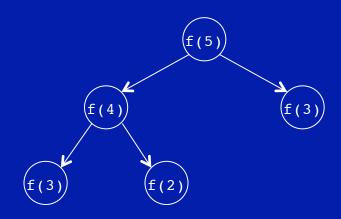
```
int factorial(int n)
{
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}
```

```
cout << fact(4) << endl;</pre>
```

What's the time complexity for this approach to computing factorials?







How does it work? The recursion tree model:

How efficient is that?

Can you use the visualization to help you estimate the growth rate of the number of function calls (= activation frames to be processed)?

How does it work? The recursion tree model:

How efficient is that?

Not very. Your book confirms that T(n) for fibonacci(n) increases exponentially with n, because of all the duplicated function calls.

How does it work? The recursion tree model:

How efficient is that?

Your book says that fib(100) requires about 2^{100} activation frames. If your computer can process 1,000,000 activation frames per second, it'll take 3 x 10^{16} years.

How does it work? The recursion tree model:

How efficient is that?

But computers have become much faster since the book was written. With gigahertz speeds, we might process 1,000,000,000 frames/sec. Now we're down to a mere $_{75}$ 3 x 10^{13} years. I feel much better now.

How does it work? The recursion tree model:

How efficient is that? We'll let you do the formal proof on your own.

```
int fibonacci(int n)
 if (n <= 2)
   return 1;
 else
   return fibonacci(n - 1)
        + fibonacci(n - 2);
cout << fib(5) << endl;</pre>
           fib2
                fib1
                                   fib2
        fib3 fib3 fib3 fib3 fib3
                        fib2
     fib4 fib4 fib4 fib4 fib4 fib4 fib4 fib4
                                fib3 fib3 fib3
```

```
int fibonacci(int n)
 if (n <= 2)
   return 1;
 else
   return fibonacci(n - 1)
        + fibonacci(n - 2);
cout << fib(5) << endl;</pre>
           fib2
                fib1
        fib3 fib3 fib3 fib3 fib3
                        fib2
                                   fib2
                                         fib1
     fib4 fib4 fib4 fib4 fib4 fib4 fib4 fib4
                                fib3 fib3 fib3 fib3
```

```
int fibonacci(int n)
 if (n <= 2)
   return 1;
 else
   return fibonacci(n - 1)
        + fibonacci(n - 2);
cout << fib(5) << endl;</pre>
          fib2
                fib1
        fib3 fib3 fib3 fib3 fib3
                        fib2
                                   fib2
                                        fib1
     fib4 fib4 fib4 fib4 fib4 fib4 fib4 fib4
                                fib3 fib3 fib3 fib3 fib3
```

```
int fibonacci(int n)
 if (n <= 2)
   return 1;
 else
   return fibonacci(n - 1)
        + fibonacci(n - 2);
cout << fib(5) << endl;</pre>
          fib2
               fib1
                                  fib2
        fib3 fib3 fib3 fib3 fib3
                       fib2
                                       fib1
     fib4 fib4 fib4 fib4 fib4 fib4 fib4 fib4
                               fib3 fib3 fib3 fib3 fib3
```

```
int fibonacci(int n)
 if (n <= 2)
   return 1;
 else
   return fibonacci(n - 1)
        + fibonacci(n - 2);
cout << fib(5) << endl;</pre>
          fib2
               fib1
                                 fib2
       fib3 fib3 fib3 fib3 fib3
                       fib2
                                       fib1
                               fib3 fib3 fib3 fib3 fib3
     fib4 fib4 fib4 fib4 fib4 fib4 fib4 fib4
```

```
int fibonacci(int n)
 if (n <= 2)
   return 1;
 else
   return fibonacci(n - 1)
        + fibonacci(n - 2);
cout << fib(5) << endl;</pre>
           fib2
                fib1
        fib3 fib3 fib3 fib3 fib3
                        fib2
                                   fib2
                                         fib1
     fib4 fib4 fib4 fib4 fib4 fib4 fib4 fib4
                                fib3 fib3 fib3 fib3 fib3
  fib(3) is evaluated 2 times
```

```
int fibonacci(int n)
 if (n <= 2)
   return 1;
 else
   return fibonacci(n - 1)
         + fibonacci(n - 2);
cout << fib(5) << endl;</pre>
           fib2
                 fib1
        fib3 fib3 fib3 fib3 fib3
                         fib2
                                     fib2
                                           fib1
     fib4 fib4 fib4 fib4 fib4 fib4 fib4 fib4
                                  fib3 fib3 fib3 fib3 fib3
  fib(3) is evaluated 2 times
fib(2) is evaluated 3 times
```

```
int fibonacci(int n)
 if (n \le 2)
   return 1;
 else
   return fibonacci(n - 1)
          + fibonacci(n - 2);
cout << fib(5) << endl;</pre>
                  fib1
            fib2
         fib3 fib3 fib3 fib3 fib3
                                        fib2
                                               fib1
      fib4 fib4 fib4 fib4 fib4 fib4 fib4 fib4
                                     fib3 fib3 fib3 fib3 fib3
   fib(3) is evaluated 2 times
fib(2) is evaluated 3 times
fib(1) is evaluated 2 times (that's just a little consistency check)
                                                      98
```

How does it work? The activation stack model:

```
int fibonacci(int n)
 if (n <= 2)
   return 1;
 else
   return fibonacci(n - 1)
        + fibonacci(n - 2);
cout << fib(5) << endl;</pre>
               fib1
          fib2
       fib3 fib3 fib3 fib3 fib3
                       fib2
                                  fib2
                                       fib1
     fib4 fib4 fib4 fib4 fib4 fib4 fib4 fib4
                               fib3 fib3 fib3 fib3 fib3
```

What are your thoughts about memory usage (space complexity)? Exponential too?

How does it work? The activation stack model:

```
int fibonacci(int n)
 if (n \le 2)
   return 1;
 else
   return fibonacci(n - 1)
        + fibonacci(n - 2);
cout << fib(5) << endl;</pre>
               fib1
          fib2
       fib3 fib3 fib3 fib3 fib3
                                 fib2
                                      fib1
     fib4 fib4 fib4 fib4 fib4 fib4 fib4 fib4
                               fib3 fib3 fib3 fib3 fib3
```

The height of the recursion tree (plus 1) from several slides ago is the maximum number of activation frames on the stack at any given time – in this case it's 5. Memory for the stack is limited. If we attempt to push more frames than can be stored on the stack, a stack overflow occurs and the program will crash.

```
int fibonacci(int n)
                                 fib(5)
                                 fib(4) + fib(3)
 if (n <= 2)
                                 fib(3) + fib(2) + fib(3)
    return 1;
                                 fib(2) + fib(1) + fib(2) + fib(3)
  else
                                 1 + fib(1) + fib(2) + fib(3)
    return fibonacci(n - 1)
                                 1 + 1 + fib(2) + fib(3)
           + fibonacci(n - 2);
                                2 + fib(2) + fib(3)
                                 2 + 1 + fib(3)
                                 3 + fib(3)
cout << fib(5) << endl;</pre>
```

```
int fibonacci(int n)
                                 fib(5)
                                 fib(4) + fib(3)
 if (n <= 2)
                                 fib(3) + fib(2) + fib(3)
    return 1;
                                 fib(2) + fib(1) + fib(2) + fib(3)
  else
                                 1 + fib(1) + fib(2) + fib(3)
    return fibonacci(n - 1)
                                 1 + 1 + fib(2) + fib(3)
           + fibonacci(n - 2);
                                2 + fib(2) + fib(3)
                                 2 + 1 + fib(3)
                                 3 + fib(3)
cout << fib(5) << endl;</pre>
                                 3 + fib(2) + fib(1)
```

```
int fibonacci(int n)
                                 fib(5)
                                 fib(4) + fib(3)
  if (n \le 2)
                                 fib(3) + fib(2) + fib(3)
    return 1;
                                 fib(2) + fib(1) + fib(2) + fib(3)
  else
                                 1 + fib(1) + fib(2) + fib(3)
    return fibonacci(n - 1)
                                 1 + 1 + fib(2) + fib(3)
           + fibonacci(n - 2);
                                2 + fib(2) + fib(3)
                                 2 + 1 + fib(3)
                                 3 + fib(3)
cout << fib(5) << endl;</pre>
                                 3 + fib(2) + fib(1)
                                 3 + 1 + fib(1)
```

```
int fibonacci(int n)
                                 fib(5)
                                 fib(4) + fib(3)
  if (n \le 2)
                                 fib(3) + fib(2) + fib(3)
    return 1;
                                 fib(2) + fib(1) + fib(2) + fib(3)
  else
                                 1 + fib(1) + fib(2) + fib(3)
    return fibonacci(n - 1)
                                 1 + 1 + fib(2) + fib(3)
           + fibonacci(n - 2);
                                2 + fib(2) + fib(3)
                                 2 + 1 + fib(3)
                                 3 + fib(3)
cout << fib(5) << endl;</pre>
                                 3 + fib(2) + fib(1)
                                 3 + 1 + fib(1)
                                 4 + fib(1)
```

```
int fibonacci(int n)
                                 fib(5)
                                 fib(4) + fib(3)
  if (n \le 2)
                                 fib(3) + fib(2) + fib(3)
    return 1;
                                 fib(2) + fib(1) + fib(2) + fib(3)
  else
                                 1 + fib(1) + fib(2) + fib(3)
    return fibonacci(n - 1)
                                 1 + 1 + fib(2) + fib(3)
           + fibonacci(n - 2);
                                2 + fib(2) + fib(3)
                                 2 + 1 + fib(3)
                                 3 + fib(3)
cout << fib(5) << endl;</pre>
                                 3 + fib(2) + fib(1)
                                 3 + 1 + fib(1)
                                 4 + fib(1)
                                 4 + 1
```

```
int fibonacci(int n)
                                 fib(5)
                                 fib(4) + fib(3)
  if (n \le 2)
                                 fib(3) + fib(2) + fib(3)
    return 1;
                                 fib(2) + fib(1) + fib(2) + fib(3)
  else
                                 1 + fib(1) + fib(2) + fib(3)
    return fibonacci(n - 1)
                                 1 + 1 + fib(2) + fib(3)
           + fibonacci(n - 2);
                                2 + fib(2) + fib(3)
                                 2 + 1 + fib(3)
                                 3 + fib(3)
cout << fib(5) << endl;</pre>
                                 3 + fib(2) + fib(1)
                                 3 + 1 + fib(1)
                                 4 + fib(1)
                                 4 + 1
```

How does it work? The substitution model:

```
int fibonacci(int n)
                                 fib(5)
                                 fib(4) + fib(3)
  if (n \le 2)
                                 fib(3) + fib(2) + fib(3)
    return 1;
                                 fib(2) + fib(1) + fib(2) + fib(3)
  else
                                 1 + fib(1) + fib(2) + fib(3)
    return fibonacci(n - 1)
                                 1 + 1 + fib(2) + fib(3)
           + fibonacci(n - 2);
                                2 + fib(2) + fib(3)
                                 2 + 1 + fib(3)
                                 3 + fib(3)
cout << fib(5) << endl;</pre>
                                 3 + fib(2) + fib(1)
                                 3 + 1 + fib(1)
                                 4 + fib(1)
                                 4 + 1
```

fib(3) is evaluated 2 times

How does it work? The substitution model:

```
int fibonacci(int n)
                                 fib(5)
                                 fib(4) + fib(3)
  if (n \le 2)
                                 fib(3) + fib(2) + fib(3)
    return 1;
                                 fib(2) + fib(1) + fib(2) + fib(3)
  else
                                 1 + fib(1) + fib(2) + fib(3)
    return fibonacci(n - 1)
                                 1 + 1 + fib(2) + fib(3)
           + fibonacci(n - 2);
                                2 + fib(2) + fib(3)
                                 2 + 1 + fib(3)
                                 3 + fib(3)
cout << fib(5) << endl;</pre>
                                 3 + fib(2) + fib(1)
                                 3 + 1 + fib(1)
                                 4 + fib(1)
                                  4 + 1
```

fib(3) is evaluated 2 times fib(2) is evaluated 3 times

How does it work? The substitution model:

```
int fibonacci(int n)
                                 fib(5)
                                 fib(4) + fib(3)
  if (n \le 2)
                                 fib(3) + fib(2) + fib(3)
    return 1;
                                 fib(2) + fib(1) + fib(2) + fib(3)
  else
                                 1 + fib(1) + fib(2) + fib(3)
    return fibonacci(n - 1)
                                 1 + 1 + fib(2) + fib(3)
           + fibonacci(n - 2); 2 + fib(2) + fib(3)
                                 2 + 1 + fib(3)
                                 3 + fib(3)
cout << fib(5) << endl;</pre>
                                 3 + fib(2) + fib(1)
                                 3 + 1 + fib(1)
                                 4 + fib(1)
                                 4 + 1
```

fib(3) is evaluated 2 times fib(2) is evaluated 3 times fib(1) is evaluated 2 times (one more consistency check)

A common complaint about recursion, as we've just seen, is the resource consumption in terms of memory (activation stack space) as well as time (handling the function calls and putting frames on/taking frames off the stack).

The culprit here is the repeated postponement of computations by pushing those computations on the stack.

Consider factorial. What's the space complexity? How many frames go on the stack for factorial(10)? factorial(100)?

But what if there were a type of recursion that worked without postponing computations? If this were so, we could have factorial using O(1) stack space instead of O(n) stack space.

But what if there were a type of recursion that worked without postponing computations? If this were so, we could have factorial using O(1) stack space instead of O(n) stack space.

This type of recursion exists, and it's inspired by the fact that if we can pass a partially-completed computation through the recursive function calls, we don't have to postpone any computations...we just do them as we go.

If we don't postpone any computations, we don't really need to push anything on the stack.

In other words, I could start with an accumulator "variable" to hold the product,

product:

In other words, I could start with an accumulator "variable" to hold the product, initialize it to 1,

product: 1

In other words, I could start with an accumulator "variable" to hold the product, initialize it to 1, multiply it by 4,

product:
$$4 f(4) = 4 * f(3)$$

In other words, I could start with an accumulator "variable" to hold the product, initialize it to 1, multiply it by 4, then multiply that value by 3,

product: 12 f(4) = 4 * 3 * f(2)

In other words, I could start with an accumulator "variable" to hold the product, initialize it to 1, multiply it by 4, then multiply that value by 3, and then multiply that value by 2,

product: 24 f(4) = 4 * 3 * 2 * f(1)

In other words, I could start with an accumulator "variable" to hold the product, initialize it to 1, multiply it by 4, then multiply that value by 3, and then multiply that value by 2, and finally multiply that value by 1 to give the result of the function call factorial(4):

product: 24 f(4) = 4 * 3 * 2 * 1

In other words, I could start with an accumulator "variable" to hold the product, initialize it to 1, multiply it by 4, then multiply that value by 3, and then multiply that value by 2, and finally multiply that value by 1 to give the result of the function call factorial(4):

How do we make this happen?

product: 24

This type of recursion is called tail recursion, and it often involves the introduction of an additional parameter used as a "variable" to hold the partially-computed result instead of storing postponed computations on the stack.

```
int fact_aux(int n, int result)
{
  if (n <= 1)
    return result;
  else
    return fact_aux(n - 1, n * result);
}
int fact(n)
{
  return fact_aux(n, 1);
}</pre>
```

fact_aux is an auxiliary or helper function to keep the syntax of fact(n) the same as before. The recursive function is really fact_aux, not fact. Note that fact_aux has no postponed or pending computations on return from recursive calls. All the work is done "inside" the recursive call in the tail call position. No work is done "outside" the recursive call in that tail call position.

```
int fact_aux(int n, int result)
{
  if (n <= 1)
    return result;
  else
    return fact_aux(n - 1, n * result);
}
int fact(n)
    fact(4)
{
  return fact_aux(n, 1);
}</pre>
```

Using the substitution model, we don't see any growth to the right, indicating there's no postponed computations. But does that mean there's no stack growth? We're still making function calls and pushing stack frames, aren't we?

A function is tail recursive if for any recursive call in the function, that call is the absolute last thing the function needs to do before returning.

In that case, we don't really need to push a new activation frame, do we? There's nothing new to remember. We could just re-use the old frame. (Or we could pop the old frame and push the new frame. Either way, there's no stack growth.)

Many compilers will do exactly that: reuse the frame. By definition, Scheme (and therefore Racket) must do tail call optimization. Some versions of Java will, but others won't. The gcc C++ compiler does, if you switch on optimization.

For our purposes in CPSC 221, don't use tail recursion because you think it will get you a more efficient C++ program. Do use tail recursion if it makes more sense to you than other approaches.

"Recursion isn't useful very often, but when used judiciously it produces exceptionally elegant solutions.... In general, recursion leads to small code and slow execution and chews up stack space. For a small group of problems, recursion can produce simple, elegant solutions. For a slightly larger group of problems, it can produce simple, elegant, hard-to-understand solutions. For most problems, it produces massively complicated solutions -- in those cases, simple iteration is usually more understandable. Use recursion selectively."

Steve McConnell in Code Complete

"Recursion isn't useful very often, but when used judiciously it produces exceptionally elegant solutions.... In general, recursion leads to small code and slow execution and chews up stack space."

As you've just seen, optimizing compilers and tail recursion can make this problem go away. You need to have a better understanding or recursion than just the reflexive response of "Eek! Recursion! I don't get it."

More from McConnell:

If a programmer who worked for me used recursion to compute a factorial, I'd hire someone else. Here's the recursive version of the factorial routine... (in Pascal)

```
Function Factorial( Number: integer ): integer;
begin
   if ( Number = 1 ) then
      Factorial := 1
   else
      Factorial := Number * Factorial( Number - 1);
end;
```

More from McConnell:

In addition to being slow and making the use of run-time memory unpredictable, the recursive version of this routine is harder to understand than the iterative version. Here's the iterative version:

Things to consider:

- 1. You now know that the recursive version isn't necessarily slow and doesn't necessarily chew up memory.
- 2. There may be reasons for employing recursion that are more important than efficiency. Like ease of implementation and understandability.
- 3. Understandability, like beauty, is in the eye of the beholder.

The authors of your textbook are much more accepting:

Generally, if it is easier to conceptualize an algorithm using recursion, then you should code it as a recursive function, because the reduction in efficiency does not outweigh the advantage of readable code that is easy to debug.

Koffman and Wolfgang, p. 416

Though they may be accepting, even they don't fully understand...

From Objects, Abstraction, Data Structures and Design Using C++ by Koffman and Wolfgang (also on p. 416):

In [tail-recursive] algorithms, there is a single recursive call and it is the last line of the function. An example...

```
int factorial(int n) {
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}
```

Can you see what's wrong here?

From Objects, Abstraction, Data Structures and Design Using C++ by Koffman and Wolfgang (also on p. 416):

In [tail-recursive] algorithms, there is a single recursive call and it is the last line of the function. An example...

```
int factorial(int n) {
  if (n == 0)
    return 1;
  else
    return n * factorial(n - 1);
}
```

Can you see what's wrong here? This is the function from our first recursion example, and it's not tail recursive.