

CPSC 312

Functional and Logic  
Programming

24 September 2015



# Assignment 1

- ❖ Due next Tuesday at the beginning of the class. I'll be collecting hard copies at the start of the class.
- ❖ You can also hand it in today (after the class), or slide it under my office door (ICCS 187) anytime before noon next Tuesday.



# Office Hours

- ❖ Instructor: (Sara Sagaii) [sarams@cs.ubc.ca](mailto:sarams@cs.ubc.ca)
  - ❖ Tuesdays 10-11am
  - ❖ Thursdays 10:30-11:30am
  - ❖ ICCS 187
- ❖ Rui Ge:
  - ❖ Wednesdays 10-12am
  - ❖ Table 1 at DLC
- ❖ Susanne Bradley:
  - ❖ Fridays 11:30-1:30
  - ❖ Table 1 at DLC
- ❖ Khurram Ali:
  - ❖ Mondays 3:30-5:30
  - ❖ Place: TBA



# Questions



# Recursion: review

- ❖ A recursive procedure consists of three parts:
- ❖ The base case or termination condition. Usually the first thing done upon entering a recursive procedure
- ❖ The reduction step -- the operation that moves the computation closer to the termination condition
- ❖ The recursive procedure calling itself



# Recursion: the prolog view

- ❖ The underlying logic is the same, what's different is the higher level thinking.
- ❖ remember: Prolog procedures don't return the value you're looking for, they just return yes or no.
- ❖ So instead of saying "compute the factorial of 4 and return it to me" ...
- ❖ you really want to write a procedure that will prove, e.g. that 24 is the factorial of 4. we talked about induction as a way to carry out this proof.
- ❖ ...then when you want to know what the factorial of 4 is, you just leave a variable in the query where you would have put the 24.



# Recursion: arithmetics

```
natural(0).
```

```
natural(s(X)) :- natural(X).
```

```
plus(0,X,X).
```

```
plus(s(X),Y,s(Z)) :-
```

```
plus(X,Y,Z).
```



# Recursion: arithmetics

```
natural(0).
```

```
natural(s(X)) :- natural(X).
```

```
plus(0,X,X).
```

```
plus(s(X),Y,s(Z)) :- plus(X,Y,Z).
```

```
times(0,X,0).
```

```
times(s(X),Y,Z) :-
```

```
times(X,Y,P),plus(P,Y,Z).
```



# Recursion: arithmetics

```
natural(0).
```

```
natural(s(X)) :- natural(X).
```

```
plus(0,X,X).
```

```
plus(s(X),Y,s(Z)) :- plus(X,Y,Z).
```

```
times(0,X,0).
```

```
times(s(X),Y,Z) :- times(X,Y,P),plus(P,Y,Z).
```

```
in Prolog: times(X,Y,Z) :-
```

```
succ(Xa,X),times(Xa,Y,P),plus(P,Y,Z).
```



# Recursion: arithmetics

$\text{exp}(X, N, Y) ?$



# Recursion: arithmetics

$\text{exp}(X, N, Y) ?$

Use Induction.

What is the base case?

$\text{exp}(X, 0, 1) .$

$\text{exp}(0, N, 0) .$



# Recursion: arithmetics

$\text{exp}(X, N, Y) ?$

Use Induction.

What is the base case?

$\text{exp}(X, 0, 1) .$

$\text{exp}(0, N, 0) .$

if  $\text{exp}(X, N, Y)$ , what can be said about  $\text{exp}$  with  $N+1$  or  $s(N)$ ?



# Recursion: arithmetics

`factorial(N, F)?`



# Recursion: arithmetics

`factorial(N, F)?` meaning?



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`factorial(0, s(0)).`



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Use induction.

what is the base case?

`factorial(0, s(0))`.

if `factorial(N, F)` is true, what is factorial of  $N+1$ ,  
or `s(N)`?



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`factorial(N, F)?` meaning?

Use induction.

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`factorial(0, s(0)).`

if `factorial(N, F)` is true, what is factorial of  $N+1$ , or `s(N)`?

`factorial(s(N), F) :-`

`factorial(N, Fn), times(s(N), Fn, F).`



# Recursion: arithmetics

`factorial(N, F)?` meaning?

Use induction.

what is the base case?

`factorial(0, s(0)).`

if `factorial(N, F)` is true, what is factorial of  $N+1$ , or `s(N)`?

`factorial(s(N), F) :-`

`factorial(N, Fn), times(s(N), Fn, F).`

rewrite in Prolog: `factorial(M, F) :-`

`succ(N, M), factorial(N, Fn), times(M, Fn, F).`



# List Processing

```
list([]).
```

```
list([_|T]) :- list(T).
```

```
member(X, [X|T]).
```

```
member(X, [_|T]) :- member(X, T).
```

```
last(X, [X]).
```

```
last(X, [_|T]) :- last(X, T).
```



# List Processing

adding a new item to a list?



# List Processing

adding a new item to a list?

as simple as `add(X,L,[X|L])`.



# List Processing

adding a new item to a list?

as simple as  $\text{add}(X, L, [X | L])$ .

or even:  $[X | L]$



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what's the meaning of this add function?



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what is the meaning of this add function?

what if we want to add X to the end?



# List Processing

adding a new item to a list?

as simple as  $\text{add}(X, L, [X|L])$ .

or even:  $[X|L]$

what if we want to add it to the end?

$\text{add}(X, [], [X])$ .

$\text{add}(X, [H|T], [H|Tx]) :-$

$\text{add}(X, T, Tx)$ .



# List Processing: delete



# List Processing: delete

`del(X, L, Lx)`



# List Processing: delete

`del(X, L, Lx)` meaning?



# List Processing: delete

$\text{del}(X, L, Lx)$  meaning?

base case?



# List Processing: delete

$\text{del}(X, L, Lx)$  meaning?

base case?  $\text{del}(X, [X | Xs], Xs) .$



# List Processing: delete

$\text{del}(X, L, Lx)$  meaning?

base case?  $\text{del}(X, [X | Xs], Xs) .$

what is the induction step?



# List Processing: delete

?- del(a,[b,c,a,e],X).



# List Processing: delete

?- del(a,[b,c,a,e],X).

X=[b,c,e].



# List Processing: delete

?- del(a,[b,c,a,e],X).

X=[b,c,e].

?- del(a,X,[b,c,e]).



# List Processing: delete

```
?- del(a,[b,c,a,e],X).
```

```
X=[b,c,e].
```

```
?- del(a,X,[b,c,e]).
```

```
X = [a, b, c, e] ;
```

```
X = [b, a, c, e] ;
```

```
X = [b, c, a, e] ;
```

```
X = [b, c, e, a] ;
```

```
false.
```



# List Processing: delete

```
?- del(a,[b,c,a,e],X).
```

```
X=[b,c,e].
```

```
?- del(a,X,[b,c,e]).
```

```
X = [a, b, c, e] ;
```

```
X = [b, a, c, e] ;
```

```
X = [b, c, a, e] ;
```

```
X = [b, c, e, a] ;
```

```
false.
```

what else does this result resemble?



# List Processing: delete

```
?- del(a,[b,c,a,e],X).
```

```
X=[b,c,e].
```

```
?- del(a,X,[b,c,e]).
```

```
X = [a, b, c, e] ;
```

```
X = [b, a, c, e] ;
```

```
X = [b, c, a, e] ;
```

```
X = [b, c, e, a] ;
```

```
false.
```

what else does this result resemble? how about insert a in any place in the list?



# List Processing:delete

$\text{add}(X, L, [X|L])$ .



# List Processing:delete

$\text{add}(X, L, [X | L])$ .

can we use add for deleting?



# List Processing:delete

$\text{add}(X, L, [X | L]) .$

can we use add for deleting?

?- $\text{add}(a, X, [a, b, c]) .$



# List Processing:delete

$\text{add}(X, L, [X | L]) .$

can we use add for deleting?

?- $\text{add}(a, X, [a, b, c]) .$

?- $\text{add}(a, X, [b, a, c]) .$



# List Processing:delete

$\text{add}(X, L, [X | L]) .$

can we use add for deleting?

?- $\text{add}(a, X, [a, b, c]) .$

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yes, but only from the head...



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yes, but only from the head...and that's fine  
because the meaning of add is also add to the  
head.



# List Processing:delete

`add(X,L,[X|L]).`

can we use add for deleting?

`?-add(a,X,[a,b,c]).`

`?-add(a,X,[b,a,c]).`

yes, but only from the head...and that's fine because the meaning of add is also add to the head.

from textbook (on how to blend declarative and procedural thinking p. 65): "Construct a program with a given use in mind; then consider if the alternative uses make declarative sense".



# List Processing:delete

what about the meaning of delete? Is the delete function we wrote declaratively sound?



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?- del(a,[b,a,c,a],X).



# List Processing:delete

what about the meaning of delete? Is the delete function we wrote declaratively sound?

```
?- del(a,[b,a,c,a],X).
```

```
X = [b, c, a] ;
```

```
X = [b, a, c] ;
```

```
false.
```



# List Processing:delete

what about the meaning of delete? Is the delete function we wrote declaratively sound?

```
?- del(a,[b,a,c,a],X).
```

```
X = [b, c, a] ;
```

```
X = [b, a, c] ;
```

```
false.
```

How can we fix that?



# List Processing

How can we rewrite this to make it remove all occurrences?

```
del(X, [X | Xs], Xs) .
```

```
del(X, [H | Xs], [H | Y]) :-  
del(X, Xs, Y) .
```



# List Processing

How can we rewrite this to make it remove all occurrences?

```
del(X, [X|Xs], Y) :- del(X, Xs, Y).
```

```
del(X, [H|Xs], [H|Y]) :-  
del(X, Xs, Y).
```



# List Processing

How can we rewrite this to make it remove all occurrences?

```
del(X, [], []).
```

```
del(X, [X|Xs], Y) :- del(X, Xs, Y).
```

```
del(X, [H|Xs], [H|Y]) :-  
del(X, Xs, Y).
```



# List Processing

The two versions of del are semantically or declaratively different.



# List Processing

The two versions of del are semantically or declaratively different. Here is an example how:

```
member(X,L) :-
```

```
    del(X,L,_).
```

```
%only the first version of del
```



# List Processing

The two versions of del are semantically or declaratively different. Here is an example how:

```
member(X,L) :-
```

```
    del(X,L,_).
```

```
%only the first version of del
```

which version gives a more *natural* definition of delete?



# List Processing

There is another problem with the second version of del as well..

```
?- del(a,[b,a,c,a],X).
```

```
X = [b, c]
```



# List Processing

There is another problem with the second version of del as well..

```
?- del2(a,[b,a,c,a],X).
```

```
X = [b, c] ;
```

```
X = [b, c, a] ;
```

```
X = [b, a, c] ;
```

```
X = [b, a, c, a] ;
```

```
false.
```



# List Processing

How can we fix that?

```
del(X, [], []).
```

```
del(X, [X|Xs], Y) :- del(X, Xs, Y).
```

```
del(X, [H|Xs], [H|Y]) :-  
del(X, Xs, Y).
```



# List Processing

How can we rewrite this to make it remove all occurrences?

```
del(X, [], []).
```

```
del(X, [X|Xs], Y) :- del(X, Xs, Y).
```

```
del(X, [H|Xs], [H|Y]) :- X \= H  
del(X, Xs, Y).
```

this is an issue of avoiding backtracking..we will talk about it more when we discuss the *cut* (!)



# List Processing

On your own:

- ❖ length of list? `length([a,b,c],3)` or `length([a,b,c],s(s(s(0))))`.
- ❖ `equal_length(L1,L2)` ? (without using `length`)
- ❖ insert at a given point? `insert(X,L,N,Lx)`.
- ❖ `reverse([a,b,c],[c,b,a])`.



# List Processing: append

- ❖ Something we do frequently in list processing is join one list to another, giving one larger list. This is often called append or conc (for concatenate) or sometimes +.
- ❖ declarative perspective: we need to write a procedure which proves that a given list is the result of appending one list to another. For example, we want our procedure to prove

```
append([a,b,c],[d,e],[a,b,c,d,e])
```



# List Processing: append

❖ So how do you prove this?

`append([a,b,c],[d,e],[a,b,c,d,e])`

Induction.

Prove `append([b,c],[d,e],[b,c,d,e])`

Prove `append([c],[d,e],[c,d,e])`

Prove `append([], [d,e], [d,e])`. That's easy.

`append([], X, X)`.

What's the induction step?



# List Processing: append

- ❖ The previous slides give us these relationships:
- ❖ `append( [a,b,c], [d,e], [a,b,c,d,e] )`
- ❖ `append( [b,c], [d,e], [b,c,d,e] )`
- ❖ `append( [c], [d,e], [c,d,e] )`
- ❖ `append( [], [d,e], [d,e] )`



# List Processing: append

- ❖ The previous slides give us these relationships:
- ❖ `append( [a,b,c], [d,e], [a,b,c,d,e] )`
- ❖ `append( [b,c], [d,e], [b,c,d,e] )`
- ❖ `append( [c], [d,e], [c,d,e] )`
- ❖ `append( [], [d,e], [d,e] )`
- ❖ Can you generalize this? Look for patterns



# List Processing: append

- ❖ The previous slides give us these relationships:
- ❖ `append( [a,b,c], [d,e], [a,b,c,d,e] )`
- ❖ `append( [b,c], [d,e], [b,c,d,e] )`
- ❖ `append( [c], [d,e], [c,d,e] )`
- ❖ `append( [], [d,e], [d,e] )`
- ❖ Can you generalize this? Look for patterns
- ❖ `append( [H1|T1], X , [H2|T2] ) <- append( T1 , X , T2 )`



# List Processing: append

- ❖ The previous slides give us these relationships:
- ❖ `append( [a,b,c], [d,e], [a,b,c,d,e] )`
- ❖ `append( [b,c], [d,e], [b,c,d,e] )`
- ❖ `append( [c], [d,e], [c,d,e] )`
- ❖ `append( [], [d,e], [d,e] )`
- ❖ Can you generalize this? Look for patterns
- ❖ `append( [H|T1], X, [H|T2] ) <-  
append( T1, X, T2 )`



# List Processing: append

```
append( [ ], X, X ).
```

```
append( [ H | T1 ], X, [ H | T2 ] ) :- append( T1, X, T2 )
```



# List Processing: append

```
append( [ ], X, X ).
```

```
append( [ H | T1 ], X, [ H | T2 ] ) :- append( T1, X, T2 )
```

Append turns out to be a powerful tool whose utility extends beyond just joining lists together. Append can be used to split lists into component pieces by partially specifying the nature of the components.



# List Processing: append

```
append( [ ], X, X ).
```

```
append( [ H | T1 ], X, [ H | T2 ] ) :- append( T1, X, T2 )
```

Append turns out to be a powerful tool whose utility extends beyond just joining lists together. Append can be used to split lists into component pieces by partially specifying the nature of the components.

```
?- append( X, Y, [ a, b, c ] ).
```



# List Processing: append

```
append( [ ], X, X ).
```

```
append( [H|T1], X, [H|T2] ) :- append(T1, X, T2)
```

Append turns out to be a powerful tool whose utility extends beyond just joining lists together. Append can be used to split lists into component pieces by partially specifying the nature of the components.

```
?- append(X, Y, [a, b, c]).
```

```
X = [ ]
```

```
Y = [a, b, c] ;
```

```
Y = [b, c] ;
```

```
X = [a, b] Y = [c] ;
```

```
X = [a, b, c] Y = [ ] ;
```

```
false.
```



# List Processing: append

```
append([ ], X, X) .
```

```
append([H|T1], X, [H|T2]) :- append(T1, X, T2)
```

Append turns out to be a powerful tool whose utility extends beyond just joining lists together. Append can be used to split lists into component pieces by partially specifying the nature of the components.

```
?- append(X, Y, [a, b, c]) .
```

```
X = [ ]
```

```
Y = [a, b, c] ;
```

```
Y = [b, c] ;
```

```
X = [a, b] Y = [c] ;
```

```
X = [a, b, c] Y = [ ] ;
```

```
false.
```

All the legal ways to split a list..



# List Processing

?- prefix([a,b],[a,b,c]).



# List Processing

?- prefix([a,b],[a,b,c]).

Try this: list X is a prefix of list Y if Y can be split into two other lists, where X is the first list of those two other lists.



# List Processing

```
?- prefix([a,b],[a,b,c]).
```

Try this: list X is a prefix of list Y if **Y can be split into two other lists**, where X is the first list of those two other lists.

```
prefix(X,Y) :- append( , ,Y).
```



# List Processing

?- prefix([a,b],[a,b,c]).

Try this: list X is a prefix of list Y if Y can be split into two other lists, **where X is the first list of those two other lists.**

prefix(X,Y) :- append(**X**, ,Y).



# List Processing

?- prefix([a,b],[a,b,c]).

Try this: list X is a prefix of list Y if Y can be split into two other lists, where X is the first list of those two other lists.

prefix(X,Y) :- append(X,\_,Y).

The other list?



# List Processing

```
?- prefix([a,b],[a,b,c]).
```

Try this: list X is a prefix of list Y if Y can be split into two other lists, where X is the first list of those two other lists.

```
prefix(X,Y) :- append(X,_,Y).
```

The other list? It's just any other list that's not X or Y.



# List Processing

```
?- prefix([a,b],[a,b,c]).
```

Try this: list X is a prefix of list Y if Y can be split into two other lists, where X is the first list of those two other lists.

```
prefix(X,Y) :- append(X,  ,Y).
```

The other list? It's just any other list that's not X or Y.



# Exercise

?- suffix([c,d],[a,b,c,d]).

?- sublist([a,b],[c,d,a,b,e]).



# Exercise

1. How do you write a program using `append` that splits this list into a list of months before September and months after and including September?

```
[jan, feb, mar, apr, may, jun, jul, aug, sep,  
oct, nov, dec]
```

2. divide a list into two equal length lists?

```
dividelist([a,b,c,d,e],[a,b,c],[d,e])
```

3. Rewrite `member` with `append`.



# Assignment 2

- ❖ will be announced later today or tomorrow..



# infinite lists

- ❖ lists can be potentially infinite. what's the use of that? well, firstly, it's good for representing incomplete data.
- ❖ remember `member(b, X)`? `X` is an incomplete list. All we know is that it has at least `b` as a member. it could have other members too, we just don't know.
- ❖ Remember we said before that prolog operates under the *Closed World Assumption*; meaning if something has not been expressed and cannot be deduced then it's definitely wrong, i.e. false. The answer is not maybe. Incomplete list is a way to circumvent that rule and add data on the fly.
- ❖ `try ?- member(a, [b,c|T]), last([b,c|T], X).`



# Questions



# Next Class

- ❖ We have covered all chapters from 1 to 7 (with the exception of 5). you can finish reading them now if you haven't.
- ❖ Next Week: Arithmetic, Cuts and Negation (Ch. 8 & 11)