

ENRICHED LAWVERE THEORIES FOR OPERATIONAL SEMANTICS

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ABSTRACT. Enriched Lawvere theories are a generalization of Lawvere theories that allow us to describe operational semantics for formal systems. For example, a graph-enriched Lawvere theory describes structures that have a *graph* of operations of each arity, where the vertices are operations and the edges are rewrites going between operations. Enriched theories can be used to equip systems with operational semantics, and maps between enriching categories can serve to translate between different forms of operational and denotational semantics. We use a definition of Lucyshyn-Wright which allows for theories to be parameterized by a monoidal subcategory of arities, and show that presentation of enriched Lawvere theories is simplified by restricting to natural number arities. We illustrate these ideas with the *SKI* combinator calculus, a variable-free version of the lambda calculus, presented as a graph-enriched theory.

1. INTRODUCTION

Formal systems are sometimes defined without intrinsic connection to how they actually operate in practice. Lawvere theories [15] are an excellent formalism for describing algebraic structures obeying equational laws, but they do not specify how to compute in such a structure, for example taking a complex expression and simplifying it using rewrite rules. Recall that a Lawvere theory is a category with finite products \mathbf{T} generated by a single object t , for “type”, and morphisms $t^n \rightarrow t$ representing n -ary operations, with commutative diagrams specifying equations. There is a theory for groups, a theory for rings, and so on. We can specify algebraic structures of a given kind in some category \mathbf{C} with finite products by a power-preserving functor $\mu: \mathbf{T} \rightarrow \mathbf{C}$. This is a simple and

In a Lawvere theory the objects are types and the morphisms are terms; the problem is that there are no rewrites between terms, only equations. In operational semantics, program behavior is often specified by labelled transition systems, or labelled directed graphs [24]. The edges of such a graph can be seen as rewrites:

To bring rewrites into Lawvere theories we need a structure with types, terms, and also rewrites between terms. This suggests using an enhanced Lawvere theory where instead of merely a *set* of morphisms between objects one has a *graph* or perhaps a *category*. Enriched Lawvere theories are perfectly suited for this purpose.

sorts	: generating object t	
term constructors	: generating morphisms $t^n \rightarrow t$	
structural congruence	: commuting diagrams	
* rewrite rules	: generating hom-edges	*

- **Graphs:** Gph-theories represent “small-step” operational semantics
— a hom-graph edge represents a *single term* rewrite.
- **Categories:** Cat-theories represent “big-step” operational semantics*
— composition generates morphisms representing *big-step* rewrites
(*Conventionally, a big-step rewrite is one resulting in a “normal form”, i.e. an expression which has been computed completely. We will use the term more generally.)
- **Posets:** Pos-theories represent “full-step” operational semantics:
— a hom-poset boolean represents the *existence* of a big-step rewrite.
- **Sets:** Set-theories represent denotational semantics:
— a hom-set element represents an *equivalence class* of the symmetric closure of the big-step relation.

In section §2, we review Lawvere theories as a more explicit, but equivalent, presentation of finitary monads. In §3, we recall the basics of enrichment, and especially the theory of powers.

In §4 we give the central definition of V-theory, from Lucyshyn-Wright [17], which allows us to parametrize our theory by a monoidal subcategory of arities.

In §5 we discuss how functors between enriching categories induce change-of-base 2-functors between their 2-categories of enriched categories, and in §6 we show that functors preserving finite products induce *change-of-semantics*: that is, they map theories to theories and models to models. Our main examples arise from this chain of adjunctions:

$$\begin{array}{ccccc}
 & \xrightarrow{\text{FC}} & & \xrightarrow{\text{FP}} & & \xrightarrow{\text{FS}} \\
 \text{Gph} & \begin{array}{c} \perp \\ \hline \end{array} & \text{Cat} & \begin{array}{c} \perp \\ \hline \end{array} & \text{Pos} & \begin{array}{c} \perp \\ \hline \end{array} & \text{Set} \\
 & \xleftarrow{\text{UG}} & & \xleftarrow{\text{UC}} & & \xleftarrow{\text{UP}}
 \end{array}$$

The right adjoints here automatically preserve finite products, but the left adjoints do as well, and these are more important in applications:

- Change of base along FC: $\text{Gph} \rightarrow \text{Cat}$ maps small-step operational semantics to big-step operational semantics.
- Change of base along FP: $\text{Cat} \rightarrow \text{Pos}$ maps big-step operational semantics to full-step operational semantics.
- Change of base along FS: $\text{Pos} \rightarrow \text{Set}$ maps full-step operational semantics to denotational semantics.

In §7 we show that models of all possible theories with all possible enrichments can be assimilated into one category using the Grothendieck construction. Finally, in §8 we bring all the strands together by demonstrating these concepts with the *SKI*-combinator calculus, introducing the idea of developing actual programming languages with enriched Lawvere theories.

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2. LAWVERE THEORIES

Algebraic structures are traditionally treated as sets equipped with operations obeying equations, but we can generalize such structures to live in any category with finite products. For example, given any category \mathbf{C} with finite products, we can define a monoid internal to \mathbf{C} to consist of:

$$\begin{array}{ll}
 \text{an object} & M \\
 \text{an identity element} & e: 1 \rightarrow M \\
 \text{and multiplication} & m: M^2 \rightarrow M \\
 \text{obeying the associative law} & m \circ (m \times M) = m \circ (M \times m) \\
 \text{and the right and left unit laws} & m \circ (e \times \text{id}_M) = \text{id}_M = m \circ (\text{id}_M \times e).
 \end{array}$$

Lawvere theories formalize this idea. For example, there is a Lawvere theory $\text{Th}(\text{Mon})$, the category with finite products freely generated by an object t equipped with an identity element $e: 1 \rightarrow t$ and multiplication $m: t^2 \rightarrow t$ obeying the associative law and unit laws listed above. This captures the “Platonic idea” of a monoid internal to a category with finite products. A monoid internal to \mathbf{C} then corresponds to a functor $\mu: \text{T} \rightarrow \mathbf{C}$ that preserves finite products.

In more detail, let \mathbf{N} be any skeleton of the category of finite sets FinSet . Because \mathbf{N} is the free category with finite coproducts on 1, \mathbf{N}^{op} is the free category with finite products on 1. A **Lawvere theory** is a category with finite products T equipped with a functor $\tau: \mathbf{N}^{\text{op}} \rightarrow \text{T}$ that is bijective on

objects and preserves finite products. Thus, a Lawvere theory is essentially a category generated by one object $\tau(1) = t$ and n -ary operations $t^n \rightarrow t$, as well as the projection and diagonal morphisms of finite products.

For efficiency let us call a functor that preserves finite products **cartesian**. Lawvere theories are the objects of a category **Law** whose morphisms are cartesian functors $f: \mathbf{T} \rightarrow \mathbf{T}'$ that obey $f\tau = \tau'$. More generally, for any category with finite products \mathbf{C} , a **model** of the Lawvere theory \mathbf{T} in \mathbf{C} is a cartesian functor $\mu: \mathbf{T} \rightarrow \mathbf{C}$. The models of \mathbf{T} in \mathbf{C} are the objects of a category $\mathbf{Mod}(\mathbf{T}, \mathbf{C})$, in which the morphisms are natural transformations.

A theory can thus have models in many different contexts. For example, there is a Lawvere theory $\mathbf{Th}(\mathbf{Mon})$, the theory of monoids, described as above. Ordinary monoids are models of this theory in **Set**, while topological monoids are models of this theory in **Top**.

For completeness, it is worthwhile to mention the *presentation* of a Lawvere theory: after all, we are arguing their utility in everyday programming. How exactly does the above “sketch” of $\mathbf{Th}(\mathbf{Mon})$ produce a category with finite products? It is precisely analogous to the presentation of an algebra by generators and relations: we form the free category with finite products on the data given, and impose the required equations. The result is a category whose objects are powers of M , and whose morphisms are composites of products of the morphisms in $\mathbf{Th}(\mathbf{Mon})$, projections, deletions, symmetries and diagonals. A detailed account was given by Barr and Wells [4, Chap. 4]; for a more computer-science-oriented approach see Crole [9, Chap. 3].

Currently, monads are more widely used in computer science than Lawvere theories. However, Hyland and Power have suggested that Lawvere theories could do much of the work that monads do today [10]. In 1965, Linton [16] proved that Lawvere theories correspond to “finitary monads” on the category of sets. For every Lawvere theory \mathbf{T} , there is an adjunction:

$$\begin{array}{ccc} & F & \\ \text{Set} & \xrightarrow{\quad} & \mathbf{Mod}(\mathbf{T}, \mathbf{Set}) \\ & U & \end{array} \quad \perp$$

The functor

$$U: \mathbf{Mod}(\mathbf{T}, \mathbf{Set}) \rightarrow \mathbf{Set}$$

sends each model μ to its underlying set, $X = \mu(\tau(1))$. Its left adjoint, the free model functor

$$F: \mathbf{Set} \rightarrow \mathbf{Mod}(\mathbf{T}, \mathbf{Set}),$$

sends each finite set $n \in \mathbf{N}$ to the representable functor $\mathbf{T}(\tau(n), -): \mathbf{T} \rightarrow \mathbf{Set}$, and in general any set X to the colimit of all such representables as n ranges over the poset of finite subsets of X . In rough terms, $F(X)$ is the set of all possible n -ary operations from \mathbf{T} on the set X .

If we momentarily abbreviate $\mathbf{Mod}(\mathbf{T}, \mathbf{Set})$ as \mathbf{Mod} , we obtain an adjunction

$$\mathbf{Mod}(F(n), \mu) = \mathbf{Mod}(\mathbf{T}(\tau(n), -), \mu) \cong \mu(\tau(n)) \cong \mu(\tau(1))^n = \mathbf{Set}(n, U(\mu))$$

where the left isomorphism arises from the Yoneda lemma, and the right isomorphism from the product preservation of μ .

This adjunction induces a monad T on **Set**:

$$(1) \quad T(X) = \int^{n \in \mathbf{N}} X^n \times \mathbf{T}(n, 1).$$

The integral here is a coend, essentially a coproduct quotiented by the equations of the theory and the equations induced by the cartesian structure of the category. The monad constructed this way is always **finitary**: that is, it preserves filtered colimits [1].

Conversely, for a monad T on **Set**, its Kleisli category $\mathbf{Kl}(T)$ is the category of all free algebras of the monad, which has all coproducts. There is a functor $k: \mathbf{Set} \rightarrow \mathbf{Kl}(T)$ that is the identity on objects and preserves coproducts (because it is a right adjoint). Thus,

$$k^{\text{op}}: \mathbf{Set}^{\text{op}} \rightarrow \mathbf{Kl}(T)^{\text{op}}$$

is a cartesian functor, and restricting its domain to \mathbf{N}^{op} is a Lawvere theory. When T is finitary, the monad arising from this Lawvere theory is naturally isomorphic to T itself. For more details see [4, 15, 21].

This correspondence sets up an equivalence between the category **Law** of Lawvere theories and the category of finitary monads on **Set**. There is also an equivalence of between the category $\mathbf{Mod}(T)$ of models of any given Lawvere theory and the category of algebra of the corresponding finitary monad T . Furthermore, all this generalizes with **Set** replaced by any “locally finitely presentable” category [1].

3. ENRICHMENT

To allow more general semantics, we now turn to Lawvere theories that have hom-*objects* rather than mere hom-*sets*. To do this we use enriched category theory [12], where we replace sets with objects of a monoidal category (V, \otimes, I) , the “enriching” category. A **V-category** or **V-enriched category** C is:

$$\begin{array}{ll} \text{a collection of objects} & \text{Ob}(C) \\ \text{a hom-object function} & C(-, -): \text{Ob}(C) \times \text{Ob}(C) \rightarrow \text{Ob}(V) \\ \text{composition morphisms} & \circ_{a,b,c}: C(b, c) \otimes C(a, b) \rightarrow C(a, c) \quad \forall a, b, c \in \text{Ob}(C) \\ \text{identity-assigning morphisms} & i_a: I \rightarrow C(a, a) \quad \forall a \in \text{Ob}(C) \end{array}$$

such that composition is associative and unital. A **V-functor** $F: C \rightarrow D$ is:

$$\begin{array}{ll} \text{a function} & F_0: \text{Ob}(C) \rightarrow \text{Ob}(D) \\ \text{a collection of morphisms} & F_{ab}: C(a, b) \rightarrow D(F_0(a), F_0(b)) \quad \forall a, b \in C \end{array}$$

such that F is compatible with composition and identity. A **V-natural transformation** $\alpha: F \Rightarrow G$ is:

$$\text{a family } \alpha_a: I \rightarrow D(F_0(a), G_0(a)) \quad \forall a \in \text{Ob}(C)$$

such that α is “natural” in a . There a 2-category \mathbf{VCat} of **V-categories**, **V-functors**, and **V-natural transformations**.

To go further, let V be a symmetric monoidal closed category, providing

$$\begin{array}{ll} \text{internal hom} & [-, -]: V^{\text{op}} \times V \rightarrow V \\ \text{tensor-hom adjunction} & V(u \otimes v, w) \cong V(u, [v, w]) \quad \forall u, v, w \in \text{Ob}(V) \\ \text{braiding natural isomorphism} & \tau_{u,v}: u \otimes v \cong v \otimes u \quad \forall u, v \in \text{Ob}(V) \end{array}$$

obeying the hexagon and triangle equations and the symmetry axiom saying $\tau_{u,v}$ is the inverse of $\tau_{v,u}$ [19]. Then we can construct new **V-categories** from old by taking tensor products and opposites in an obvious way. Moreover V is itself a **V-category**, denoted \tilde{V} , with the hom-object function providing the internal hom. The tensor-hom adjunction is called “currying” in computer science; the counit is “evaluation”.

Generalizing how the symmetric monoidal closed category \mathbf{V} has a tensor product and internal hom, a \mathbf{V} -category can be “ \mathbf{V} -tensored” and “ \mathbf{V} -powered”. Suppose \mathbf{C} is any \mathbf{V} -category. Given $a \in \mathbf{C}$ and $v \in \mathbf{V}$, we say an object $v \odot a \in \mathbf{C}$ is a **tensor product** of a by v if it is equipped with isomorphisms

$$(2) \quad \mathbf{C}(v \odot a, b) \cong [v, \mathbf{C}(a, b)]$$

\mathbf{V} -natural in b . In the special case $\mathbf{V} = \mathbf{Set}$ this forces $v \cdot a$ to be the v -fold coproduct of copies of a :

$$v \odot a = \sum_{i \in v} a.$$

In general, if we can choose a tensor product $v \odot a$ for all $v \in \mathbf{V}$, $a \in \mathbf{C}$ we say \mathbf{C} is **\mathbf{V} -tensored**, and we obtain a \mathbf{V} -functor

$$\odot : \tilde{\mathbf{V}} \otimes \mathbf{C} \rightarrow \mathbf{C}.$$

For example, $\tilde{\mathbf{V}}$ is always \mathbf{V} -tensored via the tensor product $\otimes : \tilde{\mathbf{V}} \otimes \tilde{\mathbf{V}} \rightarrow \tilde{\mathbf{V}}$.

Similarly, given $b \in \mathbf{C}$ and $v \in \mathbf{V}$, we say an object $v \pitchfork b \in \mathbf{C}$ is a **power** of b by v if it is equipped with isomorphisms

$$(3) \quad [v, \mathbf{C}(a, b)] \cong \mathbf{C}(a, v \pitchfork b).$$

\mathbf{V} -natural in a . In the special case $\mathbf{V} = \mathbf{Set}$ this forces $v \pitchfork b$ to be the v -fold product of copies of b :

$$v \pitchfork b = \prod_{i \in v} b.$$

In general, if we can choose a power $v \pitchfork b$ for all $v \in \mathbf{V}$, $b \in \mathbf{C}$ we say \mathbf{C} is **\mathbf{V} -powered**, and we obtain a \mathbf{V} -functor

$$\pitchfork : \tilde{\mathbf{V}}^{\text{op}} \otimes \mathbf{C} \rightarrow \mathbf{C}.$$

For example, $\tilde{\mathbf{V}}$ is always \mathbf{V} -powered via the internal hom $[-, -] : \mathbf{V}^{\text{op}} \times \mathbf{V} \rightarrow \mathbf{V}$.

When \mathbf{C} is both \mathbf{V} -tensored and \mathbf{V} -powered, we have isomorphisms

$$(4) \quad \mathbf{C}(v \odot a, b) \cong [v, \mathbf{C}(a, b)] \cong \mathbf{C}(a, v \pitchfork b).$$

that are \mathbf{V} -natural in $a, b \in \mathbf{C}$ and $v \in \mathbf{V}$. When $\mathbf{V} = \mathbf{Set}$ this occurs if \mathbf{C} has all products and coproducts. In what follows we use exponential notation $b^v := v \pitchfork b$ for \mathbf{V} -powers, and denote the unit I by 1 , because the enriching categories we consider are cartesian. We should note that tensors are sometimes called “copowers”, and powers are sometimes called “cotensors”.

There are just a few more technicalities. A category is **locally finitely presentable** if it is the category of models for a finite limits theory, and an object is **finite** if its representable functor is **finitary**: that is, it preserves filtered colimits [1]. A \mathbf{V} -category \mathbf{C} is **locally finitely presentable** if its underlying category \mathbf{C}_0 is locally finitely presentable, \mathbf{C} has finite powers, and $(-)^x : \mathbf{C}_0 \rightarrow \mathbf{C}_0$ is finitary for all finitely presentable x . The details are not crucial here: all categories to be considered are locally finitely presentable. We will use \mathbf{V}_f to denote the full subcategory of \mathbf{V} of finite objects: in \mathbf{Gph} , these are simply graphs with finitely many vertices and edges.

Even though the definition of Lawvere theory seems to be all about products, it is actually about *powers*, because these constitute the arities of the operations. These are greatly generalized in the enriched case, because whereas the only finite objects in \mathbf{Set} are *finite sets*, there are more complex finite objects in any other enriching \mathbf{V} . In the next section we discuss some difficulties of this generality and how to circumvent them.

4. ENRICHED LAWVERE THEORIES

Power introduced the notion of enriched Lawvere theory about twenty years ago, “in seeking a general account of what have been called notions of computation” [25]. The original definition is as follows: for a symmetric monoidal closed category $(V, \otimes, 1)$, a “ V -enriched Lawvere theory” is a V -category T that has powers by objects in V_f , equipped with an identity-on-objects V -functor

$$\tau: V_f^{\text{op}} \rightarrow T$$

that preserves these powers. A “model” of a V -theory is a V -functor $\mu: T \rightarrow V$ which preserves powers by finite objects of V . There is a category $\text{Mod}(T, V)$ whose objects are models and whose morphisms are V -natural transformations. The monadic adjunction and equivalence of §2 generalize to this context.

However, this sort of V -enriched Lawvere theory has arities for every finite object of V . These *generalized arities* may be very powerful—rather than only inputting n -tuples of terms, we can input any finite object of terms. Despite its potential, this generalization remains largely unexploited in computer science. Power [13] introduced “enriched sketches” as a way of presenting enriched Lawvere theories, and “monads with arities” have been explored by Berger, Melliés and Weber [6], but they are still not yet widely understood or used. What does it really mean for an operation to take in a finite graph of terms? How can we learn to use this generality? We hope that someone answers these questions, so that we can use more general arities in computer programming.

In this paper, however, we only consider *natural number* arities, while still retaining enrichment. To do this we use the work of Lucyshyn-Wright [17], who along with Power [23] has generalized Power’s original ideas to allow a more flexible choice of arities. For simplicity, we also limit ourselves to the case where the tensor product of V is cartesian.

Thus, in all that follows, we let $(V, \times, 1_V)$ be a cartesian closed category equipped with chosen finite coproducts of the terminal object 1_V , say

$$n_V := \sum_{i \in n} 1_V.$$

These objects come with natural isomorphisms

$$(5) \quad V(n_V, a) \cong V(1_V, a)^n.$$

Define N_V to be the full subcategory of V containing just these objects n_V , and let

$$j: N_V \hookrightarrow V$$

be the inclusion. There is also a V -category \tilde{N}_V whose objects are those of N_V and whose hom-objects are given as in V :

$$\text{hom}(a, b) = [j(a), j(b)]$$

for all $a, b \in N_V$.

Definition 1. We define a **V -theory** (T, τ) to be a V -category T equipped with a V -functor

$$\tau: \tilde{N}_V^{\text{op}} \rightarrow T$$

that is bijective on objects and **N_V -power-preserving**: that is, it preserves powers by objects of (the subcategory image of) N_V .

This definition is adapted from that of Lucyshyn-Wright [17]. While he demands that τ be the identity on objects, we allow it to be bijective on objects, to more easily handle change of base.

Definition 2. A **model** of \mathbb{T} in a \mathbf{V} -category \mathbf{C} is a $\mathbf{N}_{\mathbf{V}}$ -power-preserving \mathbf{V} -functor

$$\mu: \mathbb{T} \rightarrow \mathbf{C}.$$

In the same way that all the objects of a Lawvere theory are \mathbf{N} -powers of a single object, we can see that all the objects of a \mathbf{V} -theory are $\mathbf{N}_{\mathbf{V}}$ -powers $t^{n_{\mathbf{V}}}$ of the object

$$t = \tau(1_{\mathbf{V}}).$$

Note that t itself is $t^{1_{\mathbf{V}}}$. Moreover, just as every object of \mathbf{N}^{op} is a power of 1, every object of $\mathbf{N}_{\mathbf{V}}^{\text{op}}$ is a power of $1_{\mathbf{V}}$. This fact comes from the isomorphism $n_{\mathbf{V}} \cong n_{\mathbf{V}}^1$ in $\mathbf{N}_{\mathbf{V}}$ by taking opposites. Then, since τ preserves $\mathbf{N}_{\mathbf{V}}$ -powers and is bijective on objects, it follows that every object of \mathbb{T} is an $\mathbf{N}_{\mathbf{V}}$ -power of t .

Thus, the full sub- \mathbf{V} -category $\tilde{\mathbf{N}}_{\mathbf{V}}$ has hom-objects which behave as follows:

$$[n_{\mathbf{V}}, m_{\mathbf{V}}] \cong (m \cdot 1_{\mathbf{V}})^n \cong (m^n)_{\mathbf{V}}.$$

CHECK!!!

Definition 3. We define the **category of \mathbf{V} -theories**, $\mathbf{V}\text{Law}$, to be the category for which an object is a \mathbf{V} -theory and a morphism from (\mathbb{T}, τ) to (\mathbb{T}', τ') is an $\mathbf{N}_{\mathbf{V}}$ -power-preserving \mathbf{V} -functor $f: \mathbb{T} \rightarrow \mathbb{T}'$ such that $f\tau = \tau'$.

Definition 4. For every \mathbf{V} -theory (\mathbb{T}, τ) and every \mathbf{V} -category \mathbf{C} with $\mathbf{N}_{\mathbf{V}}$ -powers, we define the **category of models** of \mathbb{T} in \mathbf{C} , $\text{Mod}(\mathbb{T}, \mathbf{C})$, to consist of $\mathbf{N}_{\mathbf{V}}$ -power preserving \mathbf{V} -functors $\mathbb{T} \rightarrow \mathbf{C}$ and \mathbf{V} -natural transformations between these.

If \mathbf{V} is complete and cocomplete, then $\mathbf{V}\text{Law}$ and $\text{Mod}(\mathbb{T}, \mathbf{C})$ naturally become \mathbf{V} -categories. This is potentially very useful, and the enriching categories $\mathbf{V} = \mathbf{Gph}, \mathbf{Cat}, \mathbf{Pos}, \mathbf{Set}$ that are the focus of this paper are indeed complete and cocomplete.

The basic monadicity results for Lawvere theories generalize to \mathbf{V} -theories. It follows from the work of Lucyshyn-Wright [17, Sec. 8] that we have a \mathbf{V} -enriched adjunction

$$\begin{array}{ccc} & F & \\ \mathbf{V} & \xrightarrow{\quad} & \text{Mod}(\mathbb{T}, \mathbf{V}) \\ & U & \end{array} \quad \perp$$

The forgetful \mathbf{V} -functor $U: \text{Mod}(\mathbb{T}, \mathbf{V}) \rightarrow \mathbf{V}$ is evaluation at $t \in \mathbb{T}$. Its left adjoint F relies on the existence of the **free model** on an object $n_{\mathbf{V}} \in \mathbf{V}$. This is the enriched generalization of the free model described in §2: it is the composite of $\tau^{\text{op}}: \tilde{\mathbf{N}}_{\mathbf{V}}^{\text{op}} \rightarrow \mathbb{T}^{\text{op}}$ with the enriched Yoneda embedding $y: \mathbb{T}^{\text{op}} \rightarrow [\mathbb{T}, \mathbf{V}]$:

$$\begin{array}{ccccc} \mathbf{N}_{\mathbf{V}}^{\text{op}} & \xrightarrow{\tau^{\text{op}}} & \mathbb{T}^{\text{op}} & \xrightarrow{y} & [\mathbb{T}, \mathbf{V}] \\ n_{\mathbf{V}} & \mapsto & t^{n_{\mathbf{V}}} & \mapsto & \mathbb{T}(t^{n_{\mathbf{V}}}, -) \end{array}$$

Since an object of \mathbf{V} does not necessarily have a “poset of finite subobjects” over which to take a filtered colimit (as in \mathbf{Set}), the extension of this “free model” functor $y\tau^{\text{op}}$ to all of \mathbf{V} is specified

by a somewhat higher-powered generalization: it is the left Kan extension of $y\tau$ along j :

$$\begin{array}{ccc} \mathbf{N}_V & \xrightarrow{y\tau^{\text{op}}} & [\mathbf{T}, \mathbf{V}] \\ & \searrow j & \downarrow \eta \\ & & \mathbf{V} \end{array} \quad \begin{array}{c} \nearrow F := \text{Lan}_j y\tau \end{array}$$

This is essentially the universal “best solution” to the problem of making the triangle commute up to a natural transformation. That is, for any functor $G: \mathbf{V} \rightarrow [\mathbf{T}, \mathbf{V}]$ and \mathbf{V} -natural transformation $\theta: y\tau \Rightarrow Gj$, the latter factors uniquely through η .

From the adjunction between \mathbf{V} and the category of models $\text{Mod}(\mathbf{T}, \mathbf{V})$ we obtain a \mathbf{V} -enriched monad

$$T = UF: \mathbf{V} \rightarrow \mathbf{V},$$

and this has a more concrete formula as an enriched coend:

$$(6) \quad T(V) = \int^{n_V \in \mathbf{N}_V} V^{n_V} \times \mathbf{T}(t^{n_V}, t)$$

Finally, it follows from Thm. 8.9 and Example 7.5.2 of Lucyshyn-Wright [17] that $\text{Mod}(\mathbf{T}, \mathbf{V})$ is \mathbf{V} -monadic over \mathbf{V} via this \mathbf{V} -enriched monad T .

Example 5. When $\mathbf{V} = \text{Cat}$, a \mathbf{V} -category is a 2-category, so a \mathbf{V} -theory deserves to be called a **2-theory**. For example, let $\mathbf{T} = \text{Th}(\text{PsMon})$ be the 2-theory of pseudomonoids [14]. A pseudomonoid is a weakened version of a monoid: rather than associativity and unitality *equations*, it has 2-isomorphisms called the associator and unitors, which we can treat as *rewrite rules*. To equate various possible rewrite sequences, these 2-isomorphisms must obey equations called “coherence laws”. Here is a presentation of the 2-theory for pseudomonoids:

Th(PsMon)

sorts	M	pseudomonoid
operations	$e: 1 \rightarrow M$ $m: M^2 \rightarrow M$	identity multiplication
rewrites	$\alpha: m \circ (m \times \text{id}_M) \xRightarrow{\sim} m \circ (\text{id}_M \times m)$ $\lambda: m \circ (e \times \text{id}_M) \xRightarrow{\sim} \text{id}_M$ $\rho: m \circ (\text{id}_M \times e) \xRightarrow{\sim} \text{id}_M$	associator left unitor right unitor
coherence laws		

$$\begin{array}{ccc} M^4 & \xrightarrow{1 \times 1 \times m} & M^3 \\ \downarrow m \times 1 \times 1 & \searrow 1 \times m \times 1 & \downarrow 1 \times m \\ & M^3 & \xrightarrow{1 \times m} M^2 \\ & \searrow \alpha \times 1 & \downarrow m \\ & M^3 & \xrightarrow{m \times 1} M^2 \\ & \searrow m \times 1 & \downarrow m \\ & M^2 & \xrightarrow{m} M \end{array} \quad \begin{array}{c} \xRightarrow{1 \times \alpha} \\ \xRightarrow{\alpha} \end{array} \quad \begin{array}{ccc} M^4 & \xrightarrow{1 \times 1 \times m} & M^3 \\ \downarrow m \times 1 \times 1 & \searrow 1 \times m & \downarrow m \times 1 \\ & M^3 & \xrightarrow{1 \times m} M^2 \\ & \searrow m \times 1 & \downarrow m \\ & M^2 & \xrightarrow{m} M \end{array}$$

$$\begin{array}{ccc}
& M^2 & \\
1 \times e \times 1 \swarrow & \xrightarrow{1 \times \lambda} & \searrow 1 \\
M^3 & \xrightarrow{1 \times m} & M^2 \\
m \times 1 \downarrow & \xrightarrow{\alpha} & \downarrow m \\
M^2 & \xrightarrow{m} & M
\end{array}
=
\begin{array}{ccc}
& M^2 & \\
1 \times e \times 1 \swarrow & & \searrow 1 \\
M^3 & \xrightarrow{\rho \times 1} & M^2 \\
m \times 1 \downarrow & & \downarrow m \\
M^2 & \xrightarrow{m} & M
\end{array}$$

The top diagram expresses the pentagon identity for the associator, while the bottom one expresses the usual coherence law involving the left and right unitors. CHECK - SHOULDN'T THERE BE TWO UNITOR LAWS???

Models of $\mathbf{T} = \mathbf{Th}(\mathbf{PsMon})$ in \mathbf{Cat} are monoidal categories. A \mathbf{Cat} -enriched monad is called a 2-monad, and \mathbf{T} gives rise to the “free monoidal category” 2-monad on \mathbf{Cat} [7]. Let us explore this example in more detail. A model of \mathbf{T} is a finite-product-preserving 2-functor $\mu: \mathbf{T} \rightarrow \mathbf{Cat}$, which sends

$$\begin{array}{lll}
t & \mapsto & \mathbf{C} \\
m & \mapsto & \otimes: \mathbf{C}^2 \rightarrow \mathbf{C} \\
e & \mapsto & I: 1 \rightarrow \mathbf{C} \\
\alpha & \mapsto & \otimes \circ (\otimes \times 1_{\mathbf{C}}) \Rightarrow \otimes \circ (1_{\mathbf{C}} \times \otimes) \\
\lambda & \mapsto & I \circ 1_{\mathbf{C}} \Rightarrow 1_{\mathbf{C}} \\
\rho & \mapsto & 1_{\mathbf{C}} \circ I \Rightarrow 1_{\mathbf{C}}
\end{array}$$

such that the coherence laws of the rewrites are preserved. Thus, a model is a category equipped with a tensor product \otimes and unit object I such that these operations are unital and associative up to natural isomorphism; so these models are precisely monoidal categories. In this way, 2-theories generalize equipping *set*-like objects with operations obeying equations to equipping *category*-like objects with operations obeying equations up to transformations that obey certain equations of their own.

To form the free model on a category $\mathbf{C} \in \mathbf{Cat}$, we follow the above method: the formula for left Kan extension (writing n instead of $n_{\mathbf{Cat}}$ for simplicity) gives $F(\mathbf{C}): \mathbf{T} \rightarrow \mathbf{V}$ by

$$F(\mathbf{C}) = \int^{n \in \mathbf{N}_{\mathbf{Cat}}} \mathbf{T}(t^n, t^{(-)}) \times \mathbf{C}^n$$

which is constructed by pairing n -ary morphisms in \mathbf{T} with n -tuples of objects in \mathbf{C} for all $n \in \mathbf{N}_{\mathbf{Cat}}$, then quotienting the coproduct of these pairs by the equations of \mathbf{T} and \mathbf{C} .

This functor is not very intuitive, but composing with the left adjoint, i.e., evaluating $F(\mathbf{C})$ at 1, gives the *free monoidal category* on \mathbf{C} : in the same way that the (underlying set of the) free monoid on a set X consists of all finite strings of elements of X , $F(\mathbf{C})(1)$ consists of all finite tensors of objects and morphisms of \mathbf{C} , and all composites of these morphisms, up to the relations induced by the (composites and tensors of the) images of the associator and unitors.

In general for each $m \in \mathbf{N}_{\mathbf{Cat}}$, $F(\mathbf{C})(m)$ gives the category of all m -tuples of elements of $F(\mathbf{C})(1)$, forming the “free monoidal category on \mathbf{C} with m variables”. Finally, an algebra of the monad

$$T = UF: \mathbf{C} \mapsto \int^n \mathbf{C}^n \times \mathbf{T}(n, 1)$$

is a category A equipped with a functor $\otimes_A: F(A)(1) \rightarrow A$ such that it is compatible with the multiplication and unit of the monad, which are the “free” tensor bifunctor and monoidal unit on

A. Hence, (A, \otimes_A, I_A) is precisely a monoidal category, and we have the equivalence:

$$\mathbf{Mod}(\mathbf{T}, \mathbf{Cat}) \simeq \mathbf{Alg}(T).$$

5. CHANGE OF BASE

We now have the tools to formulate the main idea: a choice of enrichment for Lawvere theories corresponds to a choice of *semantics*, and changing enrichments corresponds to a *change of semantics*. We propose a general framework in which one can translate between different forms of semantics: small-step, big-step, full-step operational semantics, and denotational semantics.

Suppose that \mathbf{V} and \mathbf{W} are enriching categories of the sort we are considering: cartesian closed categories equipped with chosen finite coproducts of the terminal object. Suppose $F: \mathbf{V} \rightarrow \mathbf{W}$ preserves finite products. This induces a **change of base** functor $F_*: \mathbf{VCat} \rightarrow \mathbf{WCat}$ [8] which takes any \mathbf{V} -category \mathbf{C} and produces a \mathbf{W} -category $F_*(\mathbf{C})$ with the same objects but with

$$F_*(\mathbf{C})(a, b) := F(\mathbf{C}(a, b))$$

for all objects a, b . Composition in $F_*(\mathbf{C})$ is defined by

$$F(\mathbf{C}(b, c)) \times F(\mathbf{C}(a, b)) \xrightarrow{\sim} F(\mathbf{C}(b, c) \times \mathbf{C}(a, b)) \xrightarrow{F(\circ_{a,b,c})} F(\mathbf{C}(a, b)).$$

The identity-assigning morphisms are given by

$$1 \xrightarrow{\sim} F(1) \xrightarrow{F(i_a)} F(\mathbf{C}(a, b)).$$

Moreover, if $f: \mathbf{C} \rightarrow \mathbf{D} \in \mathbf{VCat}$ is a \mathbf{V} -functor, there is a \mathbf{W} -functor $F_*(f): F_*(\mathbf{C}) \rightarrow F_*(\mathbf{D})$ that on objects equals f and on hom-objects equals $F(f)$.

We can also define F_* on enriched natural transformations. If $\alpha: f \Rightarrow g$ is a \mathbf{V} -natural transformation and $c \in \mathbf{Ob}(\mathbf{C})$, then we define $F_*(\alpha)_c$ to be the composite

$$1 \xrightarrow{\sim} F(1) \xrightarrow{F(i_a)} F(\mathbf{C}(a, b)).$$

Thus, change of base actually gives a 2-functor from the 2-category of \mathbf{V} -categories, \mathbf{V} -functors and \mathbf{V} -natural transformations to the corresponding 2-category for \mathbf{W} .

We now study how change of base affects theories and their models. In Theorem 9 we see that everything works as one would hope, but we need some preliminary work to prove this. We begin by showing that $n_{\mathbf{V}}$ -powers in the enriched category $\tilde{\mathbf{V}}$ are the same as ordinary cartesian powers in \mathbf{V} . More precisely:

Lemma 6. The functors $[n_{\mathbf{V}}, -]: \mathbf{V} \rightarrow \mathbf{V}$ and $(-)^n: \mathbf{V} \rightarrow \mathbf{V}$ are naturally isomorphic.

Proof. If $a, b \in \mathbf{V}$, then

$$\begin{aligned} \mathbf{V}(a, [n_{\mathbf{V}}, b]) &\cong \mathbf{V}(a \times n_{\mathbf{V}}, b) && \text{hom-tensor adjunction} \\ &= \mathbf{V}(a \times (n \cdot 1_{\mathbf{V}}), b) && \text{definition of } n_{\mathbf{V}} \\ &\cong \mathbf{V}(n \cdot (a \times 1_{\mathbf{V}}), b) && \text{distributivity} \\ &\cong \mathbf{V}(n \cdot a, b) && \text{unitality} \\ &\cong \mathbf{V}(a, b)^n && \text{continuity of hom in its first argument} \\ &\cong \mathbf{V}(a, b^n) && \text{continuity of hom in its second argument.} \end{aligned}$$

Each of these isomorphisms is natural in a and b , so by the Yoneda lemma $[n_{\mathbf{V}}, -] \cong (-)^n$. \square

Now let us consider a \mathbf{V} -theory (\mathbf{T}, τ) as in Definition 1. All the objects of \mathbf{T} are of the form $t^{n\mathbf{V}}$. That is, they are $\mathbf{N}_{\mathbf{V}}$ -powers of the object t . However, since \mathbf{V} is cartesian, we can also define cartesian powers in \mathbf{T} , and we shall show that $t^{n\mathbf{V}}$ is the cartesian power t^n .

In any \mathbf{V} -category \mathbf{C} , given an object $a \in \text{Ob}(\mathbf{C})$ and a natural number n , we say $b \in \text{Ob}(\mathbf{C})$ is an **n th cartesian power** of a if there is a \mathbf{V} -natural isomorphism of \mathbf{V} -functors

$$\mathbf{C}(b, -) \cong \mathbf{C}(a, -)^n.$$

In \mathbf{T} there is just one object in each isomorphism class, so we can speak quite strictly of “the” n th cartesian power of t , and denote it by t^n , if it exists. In fact it does:

Lemma 7. Let (\mathbf{T}, τ) be a \mathbf{V} -theory. Then n th cartesian power t^n exists and equals $t^{n\mathbf{V}}$.

Proof. If $a \in V$ then

$$\begin{aligned} \mathbf{T}(a, t^{n\mathbf{V}}) &\cong [n_{\mathbf{V}}, \mathbf{T}(a, t)] && \text{definition of power} \\ &\cong [n \cdot 1_{\mathbf{V}}, \mathbf{T}(a, t)] && \text{definition of } n_{\mathbf{V}} \\ &\cong \mathbf{T}(a, t)^n && \text{continuity of hom in its first argument} \end{aligned}$$

Each of these isomorphisms is natural in a , so $t^{n\mathbf{V}}$ is the n th cartesian power of t . \square

We want to know when the functor $F: \mathbf{V} \rightarrow \mathbf{W}$ induces a change of base $F_*: \mathbf{VCat} \rightarrow \mathbf{WCat}$ that “preserves enriched theories”. That is, given a \mathbf{V} -theory

$$\tau_{\mathbf{V}}: \tilde{\mathbf{N}}_{\mathbf{V}}^{\text{op}} \rightarrow \mathbf{T}$$

we want to determine a minimal condition for the base-changed functor

$$F_*(\tau_{\mathbf{V}}): F_*(\tilde{\mathbf{N}}_{\mathbf{V}}^{\text{op}}) \rightarrow F_*(\mathbf{T})$$

to induce a \mathbf{W} -theory in a canonical way. So, assuming there is a clear identification of $F_*(\tilde{\mathbf{N}}_{\mathbf{V}}^{\text{op}})$ and $\tilde{\mathbf{N}}_{\mathbf{W}}^{\text{op}}$, it suffices to require that $F_*(\tau_{\mathbf{V}})$ preserves $\mathbf{N}_{\mathbf{W}}$ -powers.

Because $F_*(-)_{\text{hom}}$ is defined as

$$F_*(\mathbf{T})(a, s^{n\mathbf{V}}) = F(\mathbf{T}(a, s^{n\mathbf{V}})),$$

combined with the previous lemmas, the preservation of “ $\mathbf{N}_{(-)}$ -power preserving functors” by F_* is implied by the preservation of finite products by F : and since our enriching categories and base-change functors are cartesian, this is automatic.

Lemma 8. Let $F: \mathbf{V} \rightarrow \mathbf{W}$ be a cartesian functor, and let $\mathbf{N}_{\mathbf{V}}, \mathbf{N}_{\mathbf{W}}$ be defined as above. If $f: \mathbf{C} \rightarrow \mathbf{D}$ is a \mathbf{V} -functor that preserves $\mathbf{N}_{\mathbf{V}}$ -powers, then $F_*(f): F_*(\mathbf{C}) \rightarrow F_*(\mathbf{D})$ is a \mathbf{W} -functor that preserves $\mathbf{N}_{\mathbf{W}}$ -powers.

Proof.

$$\begin{aligned} F_*(\mathbf{D})(F_*(f)(a), F_*(f)(s^{n\mathbf{V}})) &= F(\mathbf{D}(f(a), f(s^{n\mathbf{V}}))) && \text{definition of base change} \\ &\cong F(\mathbf{D}(f(a), f(s)^{n\mathbf{V}})) && f \text{ preserves } \mathbf{N}_{\mathbf{V}}\text{-powers} \\ &\cong F(\mathbf{D}(f(a), f(s))^n) && \text{Lemma 6 for } \mathbf{V} \\ &\cong F(\mathbf{D}(f(a), f(s)))^n && F \text{ cartesian} \\ &= F_*(\mathbf{D})(f(a), f(s))^n && \text{definition of base change} \\ &\cong F_*(\mathbf{D})(f(a), f(s)^{n\mathbf{W}}) && \text{Lemma 6 for } \mathbf{W} \end{aligned}$$

\square

Finally, let $\tilde{n}: \tilde{\mathbf{N}}_{\mathbf{W}} \rightarrow F_*(\tilde{\mathbf{N}}_{\mathbf{V}})$ be the isomorphism which sends $n_{\mathbf{W}} \mapsto n_{\mathbf{V}}$ and is the identity on morphisms. We can then construct a \mathbf{W} -functor which precisely fits the definition of a \mathbf{W} -theory:

Theorem 9. Let \mathbf{V}, \mathbf{W} be cartesian closed categories with finite coproducts of their unit objects, and let $F: \mathbf{V} \rightarrow \mathbf{W}$ be a cartesian functor. Then F is a **change of semantics**; i.e. for every \mathbf{V} -theory $\tau_{\mathbf{V}}: \tilde{\mathbf{N}}_{\mathbf{V}}^{\text{op}} \rightarrow \mathbf{T}$, the \mathbf{W} -functor

$$\tau_{\mathbf{W}} := F_*(\tau_{\mathbf{V}}) \circ \tilde{n}^{\text{op}}: \tilde{\mathbf{N}}_{\mathbf{W}}^{\text{op}} \rightarrow F_*(\mathbf{T})$$

is a \mathbf{W} -theory. Moreover, F preserves *models*, i.e. for every $\mathbf{N}_{\mathbf{V}}$ -power preserving \mathbf{V} -functor $\mu: \mathbf{T} \rightarrow \mathbf{C}$, the \mathbf{W} -functor $F_*(\mu)$ preserves $\mathbf{N}_{\mathbf{W}}$ -powers.

Proof. The \mathbf{W} -functor $\tau_{\mathbf{W}}$ is bijective on objects because $\tau_{\mathbf{V}}$ and \tilde{n} are; and it preserves $\mathbf{N}_{\mathbf{W}}$ -powers because \tilde{N} does and $F_*(\tau_{\mathbf{V}})$ does by the previous lemma. This preservation is strict because $F_*(\mathbf{T})$ has the same objects as \mathbf{T} , so the isomorphism implies that $\tau_{\mathbf{W}}(1_{\mathbf{W}}^{n_{\mathbf{W}}}) = \tau_{\mathbf{W}}(1_{\mathbf{W}})^{n_{\mathbf{W}}}$. The preservation of models follows from the previous lemma. \square

Hence, any cartesian functor constitutes a “change of semantics” — this is a simple, ubiquitous condition, which provides for a method of translating formal languages between various “modes of operation”.

Moreover, this reasoning generalizes to **multisorted** \mathbf{V} -theories, enriched theories which have multiple sorts: given any $n \in \mathbb{N}$, the monoidal subcategory $(\mathbf{N}_{\mathbf{V}})^n$ is also an eleutheric system of arities. At the end of §8, we give an example demonstrating why this is a very useful generalization.

Before exploring applications, we introduce two more useful kinds of translations, and demonstrate how all of this information be encapsulated in one categorical notion.

6. THE CATEGORY OF ALL \mathbf{V} -THEORIES

In addition to change-of-base, there are two other natural and useful translations for these theories. Let $\mathbf{V}\mathbf{Law}$ be the category of \mathbf{V} -theories, and let $f: \mathbf{T} \rightarrow \mathbf{T}'$ be a morphism of theories; this induces a “change-of-theory” functor between the respective categories of models

$$f^*: \mathbf{VMod}(\mathbf{T}', \mathbf{C}) \rightarrow \mathbf{VMod}(\mathbf{T}, \mathbf{C})$$

defined as precomposition with f . Similarly, given a cartesian functor $g: \mathbf{C} \rightarrow \mathbf{C}'$, this induces a “change-of-model” functor

$$g_*: \mathbf{VMod}(\mathbf{T}, \mathbf{C}) \rightarrow \mathbf{VMod}(\mathbf{T}, \mathbf{C}')$$

defined as postcomposition with g .

These translations, as well as change-of-base, can all be packed up nicely using the **Grothendieck construction**: given a (pseudo)functor $F: \mathbf{D} \rightarrow \mathbf{Cat}$, there is a category $\int F$ that encapsulates all of the categories in the image of F , defined as follows:

$$\begin{array}{ll} \text{objects} & (d, x): d \in \mathbf{D}, x \in F(d) \\ \text{morphisms} & (f: d \rightarrow d', a: F(f)(x) \rightarrow x') \\ \text{composition} & (f, a) \circ (f', a') = (f \circ f', a \circ F(f)(a')). \end{array}$$

Moreover there is a functor $\overline{F}: \int F \rightarrow \mathbf{D}$ given as follows:

$$\begin{array}{ll} \text{on objects} & \overline{F}: (d, x) \mapsto d \\ \text{on morphisms} & \overline{F}: (f, a) \mapsto f. \end{array}$$

For more details see [8, 11]. We noted in §4 that $\mathbf{V}\mathbf{Law}$ and $\mathbf{Mod}(\mathbf{T}, \mathbf{C})$ are \mathbf{V} -categories when \mathbf{V} is complete and cocomplete: this and further conditions imply we can use the *enriched* Grothendieck construction [5]; but we will focus on the **Set**-enriched case for simplicity.

This idea allows us to bring together all of the different enrichments, theories, and models into one big category. For every enriching category \mathbf{V} , let \mathbf{VCat}_{np} be the subcategory of \mathbf{VCat} of \mathbf{V} -categories with \mathbf{N} -powers and \mathbf{N} -power preserving functors; then there is a functor

$$\mathbf{VMod}: \mathbf{VLaw}^{\text{op}} \times \mathbf{VCat}_{np} \rightarrow \mathbf{Cat}$$

which sends (\mathbf{T}, \mathbf{C}) to $\mathbf{VMod}(\mathbf{T}, \mathbf{C})$. The (bi)functoriality of \mathbf{VMod} gives the contravariant change-of-theory and the covariant change-of-model above.

Using the Grothendieck construction, we obtain a category $\int \mathbf{VMod}$, with a morphism

$$((f, g), \alpha): ((\mathbf{T}, \mathbf{C}), \mu) \rightarrow ((\mathbf{T}', \mathbf{C}'), \mu')$$

being finite power-preserving \mathbf{V} -functors $f: \mathbf{T} \rightarrow \mathbf{T}'$, $g: \mathbf{C} \rightarrow \mathbf{C}'$, and \mathbf{V} -natural transformation $\alpha: \mathbf{VMod}(f, g)(\mu) \rightarrow \mu'$.

Lemma 10. There is a functor

$$\text{thy}: \mathbf{CCC} \rightarrow \mathbf{Cat}$$

which assigns \mathbf{V} to $\int \mathbf{VMod}$ and $(F: \mathbf{V} \rightarrow \mathbf{W})$ to a functor $(F_*: \int \mathbf{VMod} \rightarrow \int \mathbf{WMod})$.

Proof. Given $F: \mathbf{V} \rightarrow \mathbf{W}$, base change $F_*: \mathbf{VCat} \rightarrow \mathbf{WCat}$ is a 2-functor, thereby inducing the functor $F_*: \mathbf{VMod} \rightarrow \mathbf{WMod}$ which sends a morphism $((f, g), \alpha)$ to $((F_*(f), F_*(g)), F_*(\alpha))$. Checking functoriality is left to the reader. \square

Thus, we can use the Grothendieck construction once more to encapsulate even the enrichment:

Theorem 11. There is a category $\mathbf{Thy} := \int \text{thy}$ with a morphism

$$(F, ((f, g), \alpha)): (\mathbf{V}, ((\mathbf{T}, \mathbf{C}), \mu)) \rightarrow (\mathbf{W}, ((\mathbf{T}', \mathbf{C}'), \mu'))$$

being a cartesian functor F and a morphism $(f, g, \alpha): F_*((\mathbf{T}, \mathbf{C}), \mu) \rightarrow ((\mathbf{T}', \mathbf{C}'), \mu')$ in \mathbf{WMod} .

This category assimilates a whole lot of useful information. Most importantly, there are morphisms between objects of “different kinds”, something we consider often but is normally not possible in category theory. For example, in \mathbf{Thy} there is a morphism:

$$(\mathbf{Set}, ((u_{\mathbf{Grp}}, \text{Disc}), \text{exp})): (\mathbf{Set}, ((\mathbf{T}_{\mathbf{Grp}}, \mathbf{Set}), (\mathbb{R}, +, 0))) \rightarrow (\mathbf{Set}, ((\mathbf{T}_{\mathbf{Grp}}, \mathbf{Top}), (\mathbb{R}, \times, 1)))$$

There are many unexplored questions about the large, heterogeneous categories which arise from the Grothendieck construction, regarding what unusual structure may be gained, such as limits and colimits with objects of different types, or identifying “processes” in which the kinds of objects change in an essential way. This is just a remark; for our purposes we need only recognize that enriched Lawvere theories can be assimilated into one category, which provides a unified context for change-of-base, change-of-theory, and change-of-modelling.

7. APPLICATIONS

In theoretical computer science literature, enriched algebraic theories have primarily been studied in the context of “computational effects”. Mike Stay and Greg Meredith have recognized that Lawvere theories can actually be utilized for the design of *programming languages* [31]. This idea comes from an important but underappreciated subject in foundations — combinatory logic.

7.1. The *SKI*-combinator calculus. The λ -calculus is an elegant formal language which is the foundation of functional computation, the model of intuitionistic logic, and the internal logic of cartesian closed categories: this is the Curry–Howard–Lambek correspondence [2].

Terms are constructed recursively by *variables*, *application*, and *abstraction*, and the basic rewrite is *beta reduction*:

$$M, N := x \mid (M N) \mid \lambda x. M$$

$$(\lambda x. M N) \Rightarrow M[N/x]$$

Despite its apparent simplicity, there are complications regarding *substitution*, or evaluation of functions. Consider the term $M = \lambda x. (\lambda y. (xy))$: if this is applied to the variable y , then $(M y) \Rightarrow \lambda y. (y y)$ — but this is not intended, because the y in M is just a placeholder, it is “bound” by whatever will be plugged in, while the y being substituted is “free”, meaning it can refer to some other value or function in the program. Hence whenever a free variable is to be substituted for a bound variable, we need to rename the bound variable to prevent “variable capture” (e.g. $(My) \Rightarrow \lambda z. (y z)$).

This problem was noticed early in the history of mathematical foundations, even before the λ -calculus, and so Moses Schönfinkel invented **combinatory logic** [26], a basic form of logic without the red tape of variable binding, hence without functions in the usual sense. The *SKI*-calculus is the “variable-free” representation of the λ -calculus; λ -terms are translated via “abstraction elimination” into strings of combinators and applications. This is an important method for programming languages to minimize the subtleties of variables. A great introduction into the world of strange and powerful combinators can be found in [29].

The key insight of Stay and Meredith [30] is that Lawvere theories are by definition free of variables, and it is precisely through abstraction elimination that a programming language can be made an algebraic object. When representing a computational calculus as an **Gph**-theory, the general rewrite rules are simply edges in the hom-graphs $t^n \rightarrow t$, with the object t serving in place of the variable. Below is the theory of the *SKI*-calculus:

Th(*SKI*)

	sort	t
term constructors	S :	$1 \rightarrow t$
	K :	$1 \rightarrow t$
	I :	$1 \rightarrow t$
	$(- -)$:	$t^2 \rightarrow t$
structural congruence	n/a	
rewrites	σ :	$((S -) =) \equiv \Rightarrow ((- \equiv) (= \equiv))$
	κ :	$((K -) =) \Rightarrow -$
	ι :	$(I -) \Rightarrow -$

These rewrites are implicitly universally quantified; i.e. they apply to arbitrary subterms $-$, $=$, \equiv without any variable binding involved, by using the cartesian structure of the category. They are

simply edges with vertices:

$$\begin{array}{ccc}
(((S -) =) \equiv): & t^3 \xrightarrow{l^{-1} \times t^3} 1 \times t^3 \xrightarrow{S \times t^3} t^4 \xrightarrow{(-) \times t^2} t^3 \xrightarrow{(-) \times t} t^2 \xrightarrow{(-)} t & \\
\sigma \Downarrow & \Downarrow & \\
((- \equiv) (= \equiv)): & t^3 \xrightarrow{t^2 \times \Delta} t^4 \xrightarrow{t \times \tau \times t} t^4 \xrightarrow{(-) \times (-)} t^2 \xrightarrow{(-)} t & \\
\\
((K -) =): & t^2 \xrightarrow{l^{-1} \times t^2} 1 \times t^2 \xrightarrow{K \times t^2} t^3 \xrightarrow{(-) \times t} t^2 \xrightarrow{(-)} t & \\
\kappa \Downarrow & \Downarrow & \\
-: & t^2 \xrightarrow{t \times !} t \times 1 \xrightarrow{r} t & \\
\\
(I -): & t \xrightarrow{l^{-1}} 1 \times t \xrightarrow{I \times t} t^2 \xrightarrow{(-)} t & \\
\iota \Downarrow & \Downarrow & \\
-: & t \xrightarrow{t} t &
\end{array}$$

These abstract rules are evaluated on concrete terms by “plugging in” via precomposition:

$$\begin{array}{ccc}
((KS)I): & 1 \xrightarrow{S \times I} t^2 \xrightarrow{((K -) =)} t & \\
\kappa \circ (S \times I) \Downarrow & \Downarrow & \\
S: & 1 \xrightarrow{S \times I} t^2 \xrightarrow{-} t &
\end{array}$$

(Morphisms $1 \rightarrow t$ are the “closed” terms, meaning they have no free variables; in general morphisms $t^n \rightarrow t$ are terms with n free variables, and the same reasoning applies.)

A model of this theory is a power-preserving **Gph**-functor $\mu: \mathbf{Th}(SKI) \rightarrow \mathbf{Gph}$. This gives a graph $\mu(t)$ of all terms and rewrites in the SKI -calculus as follows:

$$1 \cong \mu(1) \xrightarrow{\mu(S)} \mu(t) \xrightarrow{\mu((-))} \mu(t^2) \cong \mu(t)^2$$

The images of the nullary operations S, K, I are distinguished vertices of the graph $\mu(t)$, because μ preserves the terminal object which “points out” vertices. The image of the binary operation $(- -)$ gives for every pair of vertices $(u, v) \in \mu(t)^2$, through the isomorphism $\mu(t)^2 \cong \mu(t^2)$, a vertex $(u \ v)$ in $\mu(t)$ which is their application. In this way we get all possible terms (writing $\mu(S), \mu(K), \mu(I)$ as S, K, I for simplicity):

$$(((S (K (I I))) S) \dots$$

The rewrites are transferred by the enrichment of the functor: rather than functions between hom-sets, the morphism component of μ consists of graph homomorphisms between hom-graphs. So,

$$\mu_{1,t}: \mathbf{Th}(SKI)(1, t) \rightarrow \mathbf{Gph}(1, \mu(t))$$

maps the “syntactic” graph of all closed terms and rewrites coherently into the “semantic” graph, meaning a rewrite in the theory $a \Rightarrow b$ is sent to a rewrite in the model $\mu(a) \Rightarrow \mu(b)$.

These rewrites in the image of μ are *graph transformations*, which are just like natural transformations of functors, without the commuting diagram: given two graph homomorphisms $f, g: G \rightarrow H$,

a graph transformation $\alpha: f \Rightarrow g$ is a function $G_0 \rightarrow H_1$ which sends a vertex $v \in G$ to an edge $\alpha(v)$ with source $f(v)$ and target $g(v)$.

This is how the model realizes the **Gph**-theory as an actual graph of terms and rewrites: in the same way that a transformation between two constant functors $a \Rightarrow b: 1 \rightarrow \mathbb{C}$ is just a morphism $a(1) \rightarrow b(1)$ in \mathbb{C} , a rewrite of closed terms $a \Rightarrow b: 1 \rightarrow \mu(t)$ corresponds to an edge in $\mu(t)$:

$$\mu((I S)) \bullet \xrightarrow{\mu(\iota)} \bullet \mu(S)$$

Finally, the fact that $\mu((-))$ is not just a function but a graph homomorphism means that pairs of edges (rewrites) $(a \rightarrow b, c \rightarrow d)$ are sent to rewrites $(a b) \rightarrow (c d)$. This gives the full complexity of the theory: given a large term (program), there are many different ways it can be computed — and some are better than others:

$$\begin{array}{ccc}
 ((K S) (((S K) I) (I K))) & \xrightarrow{((K S) \sigma)} & ((K S) ((K (I K)) (I (I K)))) \\
 \downarrow \kappa & & \downarrow (((K S) \iota) (I (I K))) \\
 & & ((K S) ((K K) (I (I K)))) \\
 & & \downarrow ((K S) ((K K) (I \iota))) \\
 & & ((K S) ((K K) (I K))) \\
 & & \downarrow ((K S) ((K K) \iota)) \\
 & & ((K S) ((K K) K)) \\
 & & \downarrow ((K S) \kappa) \\
 S & \xleftarrow{\kappa} & ((K S) K)
 \end{array}$$

This process is intuitive, but how do we actually define the model, as a functor, to pick out a specific graph? There are many models of $\text{Th}(SKI)$, but in particular we care about the canonical *free* model, which means that $\mu(t)$ is simply the graph of all closed terms and rewrites in the *SKI*-calculus. This utilizes the enriched adjunction of §4:

$$\begin{array}{ccc}
 \text{Gph} & \xrightleftharpoons[\mathcal{U}_{\text{Gph}}]{f_{\text{Gph}}} & \text{Mod}(\text{Th}(SKI), \text{Gph}) \\
 & \perp &
 \end{array}$$

Then the canonical model of closed terms and rewrites is simply the free model on the empty graph, $f_{\text{Gph}}(\emptyset)$, i.e. the \mathbf{V} -functor $\mathbf{T}(1, -): \mathbf{T} \rightarrow \mathbf{V}$. Hence for us, the syntax and semantics of the *SKI* combinator calculus are unified in the model

$$\mu_{SKI}^{\text{Gph}} := \text{Th}(SKI)(1, -): \text{Th}(SKI) \rightarrow \text{Gph}$$

Here we reap the benefits of the abstract construction: the graph $\mu_{SKI}^{\text{Gph}}(t)$ is the *transition system* which represents the **small-step operational semantics** of the *SKI*-calculus:

$$(\mu(a) \rightarrow \mu(b) \in \mu_{SKI}^{\text{Gph}}(t)) \iff (a \Rightarrow b \in \text{Th}(SKI)(1, t))$$

Interestingly, in the free model on a nonempty graph, the vertices represent designated “ground variables”, and edges represent rewrites of one variable into another. This is potentially useful for “building in” a language with other basic features not intrinsic to the theory.

7.2. Change-of-base. Now we can succinctly characterize the transformation from small-step to **big-step**, which is found throughout the operational semantics literature. The “free category” functor $\text{FC}: \mathbf{Gph} \rightarrow \mathbf{Cat}$ gives for every graph G the category $\text{FC}(G)$ whose objects are the vertices of G , and whose morphisms are freely generated by the edges of G , i.e. sequences

objects	vertices of G
morphisms	finite sequences of vertices and edges $(v_1, e_1, v_2, e_2, \dots, v_n)$
composition	$(v_1, e_1, v_2, e_2, \dots, v_n) \circ (v'_1, e'_1, v'_2, e'_2, \dots, v'_n) = (v_1, e_1, \dots, v_n = v'_1, e'_1, \dots, v'_n)$

This functor is cartesian, because the definition of graphical product and categorical product are identical except for composition: vertices/objects are pairs of vertices/objects from each component, and same for edges/morphisms; hence the above operation fulfills the preservation isomorphism:

$$\text{FC}(G \times H) \cong \text{FC}(G) \times \text{FC}(H)$$

because they have the same objects, and a morphism of the former is a sequence of pairs, while that of the latter is the corresponding pair of sequences.

Thus FC is the change-of-semantics which induces the transitive closure of the rewrite relation, hence

$$\mu_{SKI}^{\text{Cat}} := \text{FC}_*(\mu_{SKI}^{\text{Gph}})$$

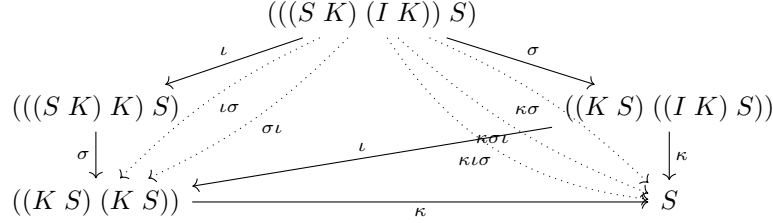
is the category which represents the big-step operational semantics of the SKI -calculus. To correspond with the conventional meaning of big-step semantics, as noted in §1, we can quotient by Curry’s equations to identify SKI -terms which are extensionally equivalent [3].

The same reasoning applies to the “free poset” functor $\text{FP}: \mathbf{Cat} \rightarrow \mathbf{Pos}$; it is a change-of-semantics because the product of posets is defined in the same way. This induces the lesser-known **full-step semantics**, which collapses hom-sets to truth values, simply asserting the existence of a rewrite sequence between terms, without distinguishing between different paths. Since there was no real algebraic information in the free category, this is simply adding the property that all the distinct paths between two terms are equal, while retaining transitivity.

Finally, we can pass to the purely abstract realm where all computation is already complete. One rarely speaks of the “free set on a poset”, but there is a left adjoint $\text{FS}: \mathbf{Pos} \rightarrow \mathbf{Set}$ of the functor $\text{UP}: \mathbf{Set} \rightarrow \mathbf{Pos}$ sending any set to the discrete poset on that set. The functor FP collapses every connected component of the poset to a point; equating every formal expression to its final value (this is cartesian because a pair of connected components corresponds to a connected component of pairs).

Assuming that the language is **terminating**, meaning every term has a finite sequence of possible rewrites, and **confluent**, meaning every pair of paths which branch from a term eventually rejoin, then this functor gives the denotational semantics of the language. (One more note here - the untyped SKI calculus is actually not terminating, and so we need to consider the simply typed SKI -calculus [3].)

Modulo a couple technicalities, from this simple sequence of functors, we can translate between the main kinds of semantics for the *SKI*-calculus. For example, we have the following computation:



The solid arrows are the one-step rewrites of the initial **Gph**-theory; applying FC_* gives the dotted composites, and FP_* asserts that all composites between any two objects are equal. Finally, FS_* collapses the whole diagram to S . This is a simple demonstration of the basic stages of computation: small-step, big-step, full-step, and denotational semantics.

Of course, most interesting languages are not always terminating, confluent, nor deterministic; the “spectrum” of semantics being presented here is simply an initial proof-of-concept. We expect that there are more interesting change-of-base functors which handle these subtleties — they have likely been studied in other contexts.

7.3. Change-of-theory: reduction contexts. We can equip term calculi with *reduction contexts*, which determine when rewrites are valid, thus giving the language a certain **evaluation strategy**. For example, the “weak head normal form” is given by only allowing rewrites on the left-hand side of the term.

We can do this for $\text{Th}(\text{SKI})$ by adding a reduction context marker as a unary operation, and a structural congruence rule which pushes the marker to the left-hand side of an application; lastly we modify the rewrite rules to be valid only when the marker is present:

Th(SKI+R)		
sort	t	
term constructors	$S, K, I:$	$1 \rightarrow t$
	$R:$	$t \rightarrow t$
	$(- -):$	$t^2 \rightarrow t$
structural congruence	$R(x\ y) = (R\ x\ y)$	
rewrites	$\sigma_r:$	$((RS -) =) \Rightarrow ((R -) (=))$
	$\kappa_r:$	$((RK -) =) \Rightarrow R -$
	$\iota_r:$	$(RI -) \Rightarrow R -$

The *SKI*-calculus is thereby equipped with “lazy evaluation”, an essential paradigm in modern programming. This represents a broad potential application of equipping theories with computational methods, such as evaluation strategies.

Moreover, these equipments can be added or removed as needed: using change-of-theory, we can utilize a “free reduction” **Gph**-functor $f_R: \text{Th}(\text{SKI}) \rightarrow \text{Th}(\text{SKI} + R)$:

objects	t^n	\mapsto	t^n
hom-vertices	S, K, I	\mapsto	S, K, I
	$(- -)$	\mapsto	$R(- -)$
hom-edges	σ, κ, ι	\mapsto	$\sigma_r, \kappa_r, \iota_r$

This essentially interprets ordinary *SKI* as having every subterm be a reduction context. This is a **Gph**-functor because its hom component consists of graph-homomorphisms:

$$f_{n,m}: \text{Th}(\text{SKI})(t^n, t^m) \rightarrow \text{Th}(\text{SKI} + R)(t^n, t^m)$$

which simply send each application to its postcomposition with R , and each rewrite to its “marked” correspondent; and this is all coherent: for example, even though $((S\ x)\ y)\ z \mapsto R(R(R(S\ x)\ y)\ z)$, the extra markers are ignored by σ_r , because they are now just a part of the lefthand terms.

So, by precomposition this induces the change of theory on categories of models:

$$f_R^*: \text{Mod}(\text{Th}(\text{SKI} + R), \mathbf{C}) \rightarrow \text{Mod}(\text{Th}(\text{SKI}), \mathbf{C})$$

for all semantic categories \mathbf{C} , which forgets the reduction contexts.

Similarly, there is a **Gph**-functor $u_R: \text{Th}(\text{SKI} + R) \rightarrow \text{Th}(\text{SKI})$ which forgets reduction contexts, by sending $\sigma_r, \kappa_r, \iota_r \mapsto \sigma, \kappa, \iota$ and $R \mapsto id_t$; this latter is the only way that the marked reductions can be mapped coherently to the unmarked. However, this means that u_R^* does not give the desired change-of-theory of “freely adjoining contexts”, because collapsing R to the identity eliminates the significance of the marker.

This illustrates a key aspect of categorical universal algebra: because change-of-theory is given by precomposition and is thus contravariant, *properties* (equations) and *structure* (operations) can only be removed.

This is a necessary limitation, at least in the present setup, but there are ways of working around it: of course, these abstract theories are not floating in isolation but are implemented in code. One can simply use a “maximal theory” with all pertinent structure, then selectively forget as needed.

7.4. Multisorted: the ρ -calculus. Many algebraic theories involve multiple sorts in an essential way. In concurrency theory, *process calculi* exhibit an ontology which is fundamentally distinct from that of sequential computing — rather than simply expressing a series of terms and rewrites, these calculi represent dynamical systems of communicating processes.

The π -calculus, designed by Milner [22], consists of **names** and **processes**, or *channels* and *agents* which communicate on those channels. Far more than a sequence of instructions on a single machine, computation develops through the interaction of independent participants in a network.

This powerful idea of modern computer science is being utilized by Greg Meredith and Mike Stay to design a deeply cooperative distributive computing system, called RChain. The “R” stands for “reflective higher-order π -calculus”, or **ρ -calculus**. It is like Milner’s original language, with one crucial difference: “reflection” is a formal system’s ability to turn code into data and vice versa. This is a powerful idea which replaces opaque, atomic variables with transparent, anatomical names, or “quoted processes” [20].

Utilizing both reflection and combinators in a theory requires special type discipline; there is a designated auxiliary sort T for combinatory terms which are analogous to “machine code”, as contrasted with the sorts N and P which are to be thought of as the actual language: see [30] §7.3 for the details of the translation. The presentation below is only a fragment; it has yet to be determined how to best represent the full algebraic theory of the ρ -calculus; we expect that a true mathematical characterization of reflection will call for original and enlightening ideas.

Th(RHO)

Sorts	N P	names processes	T	terms
Operations	$0:$ $1 \rightarrow P$	null process	$S:$ $1 \rightarrow T$	combinator
	$\&:$ $P \rightarrow N$	code to data	$K:$ $1 \rightarrow T$	combinator
	$*$: $N \rightarrow P$	data to code	$(- -):$ $T^2 \rightarrow T$	application
	$!:$ $N \times P \rightarrow P$	send		
	$?:$ $N^2 \times P \rightarrow P$	receive		
	$- \mid -:$ $P^2 \rightarrow P$	parallel		
Equations	$(P, \mid, 0)$	commutative monoid		
Rewrites	$\gamma:$ $x?(y).P \mid x!(z) \mid Q \Rightarrow P[z/y] \mid Q$		$\sigma:$ $((S -) =) \equiv \Rightarrow ((- \equiv) (= \equiv))$	
	$\epsilon:$ $*(&(P)) \Rightarrow P$		$\kappa:$ $((K -) =) \Rightarrow -$	

8. CONCLUSION

We have established the basics of how enriched Lawvere theories provide a framework for unifying the syntax and semantics, the structure and behavior of formal languages. Enriching theories in category-like structures reifies operational semantics by incorporating rewrites between terms; and cartesian functors between enriching categories induce change-of-semantics functors between categories of models—this simplified condition is obtained by using only finite cardinal arities.

This base-change, along with change-of-theory and change-of-modelling, can be assimilated into one category using an iterated Grothendieck construction Thy , which consists of all enriched Lawvere theories. Finally, enriched theories can be used not only for computational effects but the actual design of concrete programming languages, through the use of combinators.

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