Given a symmetric, Cartesian, monoidal category  $(C, \circ, id, \times, 1, \pi_1, \pi_2, \top, \sigma)$  with a strong monad  $(P, \mu, \eta, \tau)$  the natural transformations have the following types:

$$id: a \rightarrow a$$

$$\pi_1: a_1 \times a_2 \rightarrow a_1$$

$$\pi_2: a_1 \times a_2 \rightarrow a_2$$

$$T: a \rightarrow 1$$

$$\sigma: a_1 \times a_2 \rightarrow a_2 \times a_1$$

$$\mu: P(P(a)) \rightarrow P(a)$$

$$\eta: a \rightarrow P(a)$$

$$\tau: a_1 \times P(a_2) \rightarrow P(a_1 \times a_2)$$

Define *comprehensions*, morphisms  $[\cdots]: a \to P(b)$ , as

• Given  $t: a \to b$ ,

$$[t] := \eta \circ t : a \to P(b)$$

• Given  $t: b \to c$  and  $u: a \to P(b)$ ,

$$[t|b\leftarrow u]\coloneqq P\left(t
ight)\circ u:a
ightarrow P\left(c
ight)$$

• Given  $t: a_n \to b$  and  $u_i: a_{i-1} \to P(a_i)$  for  $i = 1, \ldots, n$ ,

$$[t|a_1 \leftarrow u_1, \dots, a_n \leftarrow u_n] \\ :=$$

$$\mu \circ [[t|a_2 \leftarrow u_2, \dots, a_n \leftarrow u_n] | a_1 \leftarrow u_1] : a_0 \rightarrow P(b)$$

• Given  $u: a \to P(b)$ 

$$[b|b \leftarrow u] \coloneqq u : a \rightarrow P(b)$$

• Given  $u: a \to P(b_1 \times b_2)$ ,

$$[b_1|b_1 \times b_2 \leftarrow u] := P(\pi_1) \circ u : a \to P(b_1)$$
$$[b_2|b_1 \times b_2 \leftarrow u] := P(\pi_2) \circ u : a \to P(b_1)$$
$$[b_2 \times b_1|b_1 \times b_2 \leftarrow u] := P(\sigma) \circ u : a \to P(b_2 \times b_1)$$

• Given  $t: c \to d$  and  $u: a \to P(b)$ ,

$$[t \times b | b \leftarrow u] \coloneqq \tau \circ (t \times u) : c \times a \to P(d \times b)$$
$$[b \times t | b \leftarrow u] \coloneqq P(\sigma) \circ \tau \circ (t \times u) \circ \sigma : a \times c \to P(b \times d)$$

Conjecture. Extending to co-Cartesian categories will extend the comprehension calculus by pattern matching on sum types, much like  $\pi_1, \pi_2$  perform pattern matching on product types.