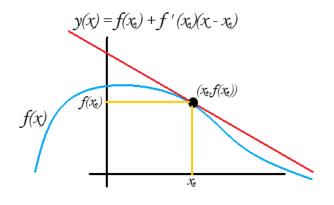


## Cálculo Diferencial Solucionario PC1 25-2

Realizado por: Jesús Paucar





## UNIVERSIDAD NACIONAL DE INGENIERÍA

Facultad de Ingeniería Industrial y de Sistemas **DEPARTAMENTO DE CIENCIAS BÁSICAS** 

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## PRIMERA PRÁCTICA CALIFICADA

Tiempo de duración: 110 minutos

1. Dadas las proposiciones P y Q. Determine si la proposición  $P \wedge Q$  es una tautología,[5,0 puntos] donde

$$P \equiv \left\{ \left[ (p \to q) \land \sim (p \to \sim r) \right] \to \sim (r \to \sim q) \right\} \ \land \ \left\{ \left[ (q \to p) \land (\sim p) \land (\sim s) \right] \to \sim (q \lor s) \right\}$$

$$Q \equiv \left\{ \left[ (p \to q) \vee (q \wedge r) \right] \ \vee \ \sim \left[ (q \to r) \wedge (p \vee (p \wedge m)) \right] \right\} \ \vee \ \sim \left\{ (p \to q) \wedge (r \to \sim p) \right\}$$

- **2.** Sean A, B, C y D subconjuntos no vacíos de U.
  - a) Si  $A \subset B$ , simplifique.

[3,0 puntos]

$$\left\{\left\{A\cap \left[(B\cap D)\cup (D\cap A)\right]\,\cup\, \left[(B\cup A)\cap D'\right]\right\}\cap (B\triangle C)\right\}\,\cup\, \left\{(A\cap C)\cap \left((A\cap C)\triangle (A'\cup C')\right)\right\}$$

b) Simplifique.

[2,0 puntos]

$$\left\{ \left[A' \cup (B \cap C)'\right] \cap \left[A \cap (B' \cup C)\right] \right\} \ \cup \ \left\{ \left[(A \cap B) \cap (A \cup B)\right] \ \cup \ \left[C \cap ((C' \cup B) \cap A)\right] \right\}$$

- 3. Sean A, B, C y D subconjuntos no vacíos de U. Demostrar por elementos.
  - **a**)

[3,0 puntos]

$$\left\{\left[\left(A'\cup B\right)\cap\left(A'\cup C'\right)\right]'\ \cup\ \left(C'\cup B'\right)'\right\}\ \subset\ \left\{\left[\left(B'\cup A\right)\cap\left(A'\cap C'\right)\right]'\ \cup\ \left(B\cup D\right)'\right\}$$

b)

[2,0 puntos]

$$A - B \subset (A - C) \cup (C - B).$$



1) Determinar si  $\mathcal{P} \wedge \mathcal{Q}$  es una tautología.

$$\mathcal{P} \equiv \left\{ \left[ (p \to q) \land \sim (p \to \sim r) \right] \to \sim (r \to \sim q) \right\} \land \left\{ \left[ (q \to p) \land (\sim p) \land (\sim s) \right] \to \sim (q \lor s) \right\}$$

$$\underbrace{\left[ (\sim p \lor q) \land (p \land r) \right] \to (r \land q) \land \left[ (\sim q \lor p) \land (\sim p) \land (\sim s) \right] \to (\sim q \land \sim s)}_{\text{Semi-absorción}}$$

$$\left[ \left( \underbrace{(p \land q) \land r} \right) \right] \to (r \land q) \land \left[ \underbrace{(\sim p \land \sim q) \land (\sim s)}_{\text{Ley Asociativa}} \right] \to (\sim q \land \sim s)$$

$$\underbrace{\left\{ p \land (q \land r) \to (r \land q) \right\}}_{\text{Ley de Exportación}} \land \underbrace{\left\{ \left[ \sim p \land (\sim q \land \sim s) \right] \to (\sim q \land \sim s) \right\}}_{\text{Ley de Exportación}}$$

$$\left\{ p \to \left[ (q \land r) \to (r \land q) \right] \right\} \land \left\{ \left[ \sim p \to \left[ (\sim q \land \sim s) \to (\sim q \land \sim s) \right] \right] \right\}$$

$$\left\{ p \to (V) \right\} \land \left\{ \sim p \to V \right\} \equiv \left\{ V \right\} \land \left\{ V \right\} \equiv V$$

$$\mathcal{Q} \ \equiv \ \big\{ \left[ (p \to q) \lor (q \land r) \right] \ \lor \ \sim \left[ (q \to r) \land \underbrace{\left( p \lor (p \land m) \right)}_{\text{Absorción}} \right] \big\} \ \lor \ \sim \big\{ (p \to q) \land (r \to \sim p) \big\}$$
 
$$\big\{ \left[ (\sim p \lor q) \lor (q \land r) \right] \ \lor \ \sim \big[ (\sim q \lor r) \land (p) \big] \big\} \ \lor \ \sim \big\{ (\sim p \lor q) \land (\sim r \lor \sim p) \big\}$$
 
$$\big\{ \left[ \sim p \lor q \right] \ \lor \ \left[ \sim p \lor (q \land \sim r) \right] \big\} \ \lor \ \big\{ \left[ p \land \sim q \lor r \right) \big\} \big\}$$
 
$$\big\{ \left[ \sim p \lor q \right] \ \lor \ \left[ \sim p \lor (q \land \sim r) \right] \ \lor \ \big\{ p \land (\sim q \lor r) \big\} \big\}$$

Por Asociativa y porque  $[\sim p \lor (q \land \sim r)] \equiv a, [p \land (\sim q \lor r)] \equiv \sim a$  $\Rightarrow$  Quedaría  $\{[\sim p \lor q] \lor (V)\} \equiv V$ 

Hallando  $\mathcal{P} \wedge \mathcal{Q} \equiv (V) \wedge (V) \equiv V$  tautología.

Nota:  $A \subset B \Rightarrow A \cap B' \neq \emptyset$ 

$$\left\{ A \cap \left[ (B \cap D) \cup (D \cap A) \cup ((B \cup A) \cap D') \right] \cap (B \triangle C) \right\} \cup \left\{ (A \cap C) \cap ((A \cap C) \triangle (A' \cup C')) \right\}$$

$$\left\{ A \cap \left[ (B \cup A) \cup (B \cup A) \cap D' \right] \cap (B \triangle C) \right\} \cup \left\{ (A \cap C) \cap (A \cap C) \triangle \varnothing \right\}$$

$$\left\{ A \cap \left[ (B \cup A) \cap (D \cup D') \right] \cap (B \triangle C) \right\} \cup \left\{ A \cap C \right\}$$

$$\left\{ A \cap (B \cup A) \cap (B \triangle C) \right\} \cup \left\{ A \cap C \right\}$$

$$\left\{ (A \cap B) \triangle (A \cap C) \right\} \cup \left\{ A \cap C \right\}$$

$$\left\{ (A \cap B) \triangle (A \cap C) \right\} \cup \left\{ (A \cap C) \right\} \cup \left\{ A \cap C \right\}$$

$$\equiv \left\{ ((A \cap B) \cap (A' \cup C') \cup (A \cap C) \cap (A' \cup B') \right\} \cup \left\{ A \cap C \right\}$$

$$\left\{ (A \cap B) \cap (A' \cup C') \cup (A \cap C) \cap (A' \cup B') \right\} \cup \left\{ A \cap C \right\}$$

$$\equiv \left\{ (A \cap B \cap C') \cup \left\{ (A \cap C) \cap (A \cap B') \right\} \cup \left\{ (A \cap C) \cap (B \cap C') \cup (A \cap B) \cup (A \cap B) \cup (A \cap C) \cup (A \cap B) \cup (A \cap C) \cup (A \cap B) \cup (A \cap C) \cup (A \cap C)$$



## 3) a) Demostración por elementos.

$$\underbrace{\left\{ \left[ (A' \cup B) \cap (A' \cup C') \right]' \ \cup \ (C' \cup B')' \right\}}_{\mathcal{Y}} \subseteq \underbrace{\left\{ \left[ (B' \cup A) \cap (A' \cap C') \right]' \ \cup \ (B \cup D)' \right\}}_{\mathcal{Z}}$$

$$x \in \mathcal{Y} \quad \Rightarrow \quad x \in \mathcal{Z}$$

$$\left\{ \left[ x \in (A' \cup B) \land x \in (A' \cup C') \right]' \lor x \in (C' \cup B')' \right\} \Rightarrow \left\{ \left[ x \in (B' \cup A) \land x \in (A' \cap C') \right]' \lor x \in (B \cup D)' \right\}$$

$$\left\{ \left[ (x \in A' \lor x \in B) \land (x \in A' \lor x \in C') \right]' \lor (x \in C' \lor x \in B')' \right\} \Rightarrow$$

$$\left\{ \left[ (x \in B' \lor x \in A) \land (x \in A' \land x \in C') \right]' \lor (x \in B \lor x \in D)' \right\}$$

$$\left\{\left[x\in A'\vee(x\in B\wedge x\in C')\right]'\vee(x\in C\wedge x\in B)\right\} \Rightarrow \left\{\left[x\in A'\wedge x\in B'\wedge x\in C'\right]'\vee(x\in B'\wedge x\in D')\right\}$$

$$\left\{ \left[ x \in A \land (x \in B' \lor x \in C) \right] \lor (x \in C \land x \in B) \right\} \Rightarrow \left\{ (x \in A \lor x \in B \lor x \in C) \lor (x \in B' \land x \in D') \right\}$$

$$\left\{ \left( x \in A \ \lor \ \left( x \in C \land x \in B \right) \right) \ \land \ \left( \left( x \in B' \lor \underbrace{x \in C} \right) \lor \left( x \in C \land x \in B \right) \right) \right\}$$

$$\Rightarrow \left\{ x \in A \ \lor \ x \in C \ \lor \ \left( x \in B \ \lor \ x \in D' \right) \right\}$$

$$\left\{\left.(x\in A\vee x\in C)\wedge(x\in A\vee x\in B)\wedge(x\in B'\vee x\in C)\right.\right\} \Rightarrow \left.\left\{\left.x\in A\vee x\in B\vee x\in C\vee x\in B'\vee x\in D'\right.\right\}\right.$$

$$\left\{\left.(x\in A\vee x\in C)\wedge(x\in A\vee x\in B)\wedge(\sim(x\in B)\vee x\in C)\right.\right\} \ \Rightarrow \ \left(x\in A\vee x\in B\vee x\in C\vee\sim(x\in D)\right)$$

 $\mbox{\it Utilizando \'Algebra de Boole.} \quad x \in A \equiv p, \ x \in B \equiv q, \ x \in C \equiv r, \ x \in D \equiv s.$ 

$$\Big\{ \left( p \vee r \right) \ \land \ \left( p \vee q \right) \ \land \ \left( \sim q \vee r \right) \Big\} \ \Rightarrow \ \Big\{ \left. p \vee q \vee r \vee \sim s \right. \Big\}$$

Por Asociativa y ordenando de forma conveniente

$$\left(\underbrace{(\sim\!\!\!p\wedge\sim\!\!\!q)\ \lor\ (p\vee q)}_{V\ (\text{independiente al valor})}\right)\ \lor\ \left((\sim\!\!\!p\wedge\sim\!\!r)\ \lor\ (q\wedge\sim\!\!r)\right)\ \lor\ (r\vee\sim\!\!s)\ \equiv\ V\ (\text{tautología}).$$

Por lo tanto, si  $\forall x \in \mathcal{Y} \Rightarrow x \in \mathcal{Z} \equiv V$ , concluimos

$$\mathcal{Y}\subseteq\mathcal{Z}$$
 .

**3) b)** 
$$A - B \subset (A - C) \cup (C - B)$$

$$x \in (A - B) \Rightarrow x \in (A - C) \lor x \in (C - B)$$

$$x \in (A \cap B') \implies x \in (A \cap C') \lor x \in (C \cap B')$$

$$x \in A \land x \in B' \Rightarrow (x \in A \land x \in C') \lor (x \in C \land x \in B')$$

$$x \in A \land x \in B' \Rightarrow \underbrace{((x \in A \land x \in C') \lor (x \in C))}_{\text{Semi absorción}} \land ((x \in A \land x \in C') \lor (x \in B'))$$

$$x \in A \land x \in B' \Rightarrow (x \in A \lor x \in C) \land (x \in A \lor x \in B') \land (x \in C' \lor x \in B')$$

$$x \in A \, \wedge \, \sim (x \in B) \, \Rightarrow \, (x \in A \, \vee \, x \in C) \, \wedge \, (x \in A \, \vee \, \sim (x \in B)) \, \wedge \, (\sim (x \in C) \, \vee \, \sim (x \in B))$$

*Utilizando Álgebra de Boole.*  $x \in A \equiv p, x \in B \equiv q, x \in C \equiv r.$ 

Tenemos 
$$(p \land \sim q) \Rightarrow ((p \lor r) \land (p \lor \sim q) \land (\sim r \lor \sim q))$$

$$(\sim p \lor q) \lor ((p \lor r) \land (p \lor \sim q) \land (\sim q \lor \sim r))$$

$$\sim p \ \lor \ \left(\left(q\lor (p\lor r)\right) \ \land \ \underbrace{\left(q\lor (p\lor \sim q)\right)}_{V} \ \land \ \underbrace{\left(q\lor (\sim q\lor \sim r)\right)}_{V}\right)$$

$$\equiv \, \sim p \ \lor \big( q \lor (p \lor r) \big) \qquad \text{Aplicando Asociativo}.$$

$$\equiv (\sim p \vee p) \vee q \vee r$$

$$\equiv V \lor (q \lor r)$$

 $\equiv V$  o tautología.

$$\Rightarrow \forall x \in (A-B): x \in (A-C) \lor x \in (C-B).$$
 Por lo tanto,

$$A - B \subset (A - C) \cup (C - B)$$
 Demostrado.