



CSCIT 2021 - Lecture 1

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Sydney)

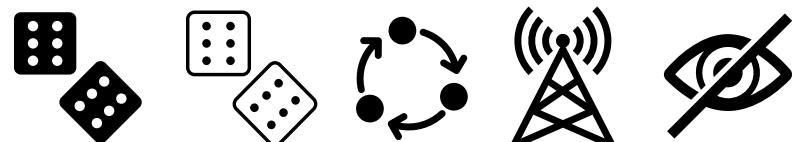
Estimation and hypothesis
testing under information
constraints

Contents of this lecture

1. What are learning and testing?
2. Baseline: the "centralised" setting
3. Beyond the centralised setting: 3 flavours
 - Private-coin protocols
 - Public-coin protocols
 - Interactive protocols
4. What are information constraints?
 - Two guiding examples: communication and privacy

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What are learning and testing?

Standard statistical setting: n iid samples from some
unknown probability distribution p

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→ learn p : output \hat{p} such that

$$\mathbb{E}[\ell(\hat{p}, p)] \leq \epsilon$$

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Goal: estimate something about p

→ learn p : output \hat{p} such that

$E_{\hat{P}}[l(\hat{p}, p)] \leq \epsilon$

depends on the
 n samples

target rate

loss function

What are learning and testing?

Examples of loss functions:

- $\text{KL}(p \parallel q) = - \sum_{\omega} p(\omega) \log \frac{q(\omega)}{p(\omega)}$
- $\ell_2(p, q) = \sum_{\omega} (p(\omega) - q(\omega))^2$
- $\chi^2(p, q) = \sum_{\omega} \frac{(p(\omega) - q(\omega))^2}{q(\omega)}$
- $\text{TV}(p, q) = \sup_S (p(S) - q(S)) = \frac{1}{2} \sum_{\omega} |p(\omega) - q(\omega)|$

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↳ learn a parameter/functional ϑ of p
output $\hat{\theta}$ such that

$$\underset{P}{\mathbb{E}}[l(\hat{\theta}, \theta(p))] \leq \varepsilon$$

estimating What are learning and testing?

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Goal: estimate something about p

for instance,
the mean of p

learn a parameter/functional θ of p
output $\hat{\theta}$ such that

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Goal: estimate something about p

↳ is p what I thought it was?

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Goal: estimate something about p

↳ Hypothesis: $\mathcal{H}_0 = "p=q"$ (null)

$\mathcal{H}_1 = "TV(p,q) > \varepsilon"$ (altern.)

Output $\hat{b} \in \{0,1\}$ st. $\underset{q}{P}\{\hat{b}=1\} + \sup_{p \in \mathcal{H}_1} P\{\hat{b}=0\} \leq \frac{1}{10}$

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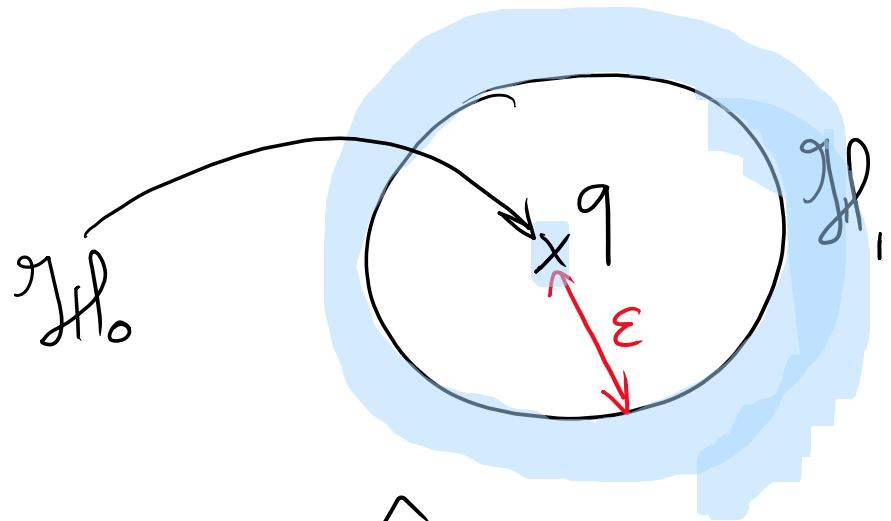
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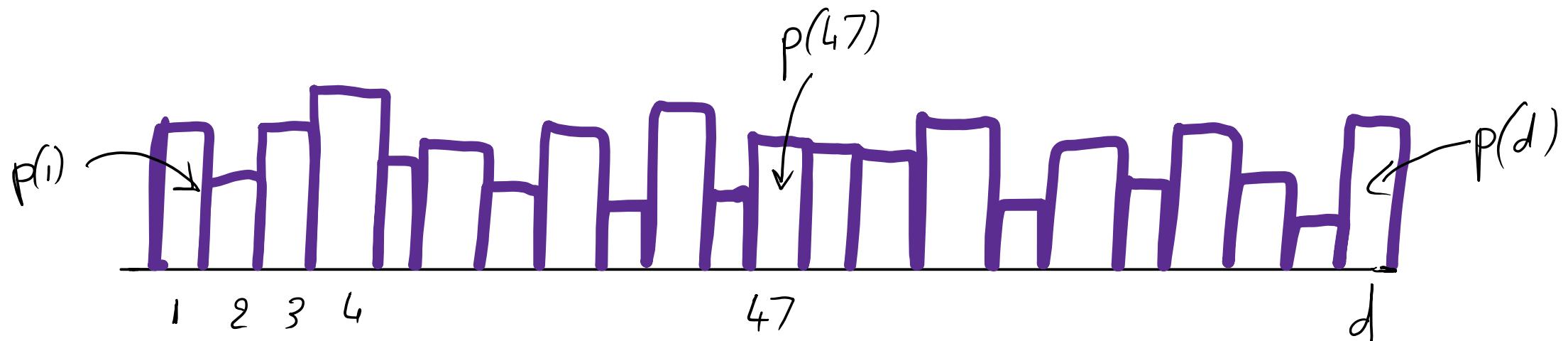
TYPE I

TYPE II \approx arbitrary

Now, what are we learning and testing?

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- ① Discrete distributions over d elements: $[d] := \{1, 2, \dots, d\}$
- Learning under TV loss → Testing if uniform on $[d]$



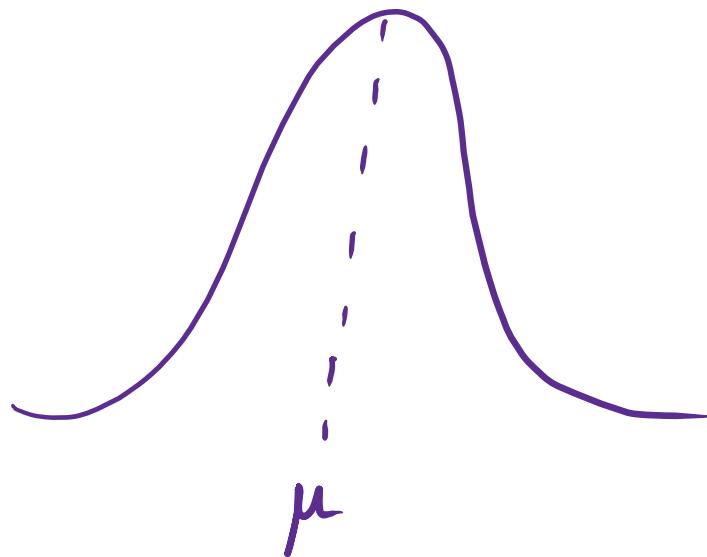
Now, what are we learning and testing?

- ② High-dimensional Gaussians (with identity covariance)
↑ dimension d

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Learning the mean under
 ℓ_2 loss



Testing if the mean
is zero (also ℓ_2)

$$p = \mathcal{N}(\mu, I_d) \\ \mu \in \mathbb{R}^d$$

Baseline: the "centralised" setting

$X_1, X_2, \dots, X_n \sim P$ **fully accessible** to the algorithm.

How **large** must n be to solve the learning or testing question?

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$X_1, X_2, \dots, X_n \sim P$ **fully accessible** to the algorithm.

How **large** must n be to solve the learning or testing question? (as a function of d, ε)

"Minimax sample complexity"

Baseline: the "centralised" setting

Discrete distributions

$d \gg 1$
 $\epsilon \in (0, 1]$

Theorem. Learning an arbitrary p over $[d]$ to TV loss ϵ has sample complexity .

Theorem. Testing if an arbitrary p over $[d]$ is u or has $\text{TV}(p, u) > \epsilon$ has sample complexity .
uniform over $[d]$

Baseline: the "centralised" setting

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Theorem. Learning an arbitrary p over $[d]$ to TV loss ϵ has sample complexity $\Theta\left(\frac{d}{\epsilon^2}\right)$.

Theorem. Testing if an arbitrary p over $[d]$ is u or has $\text{TV}(p, u) > \epsilon$ has sample complexity _____.

Baseline: the "centralised" setting

Discrete distributions

$$d \gg 1$$
$$\varepsilon \in (0, 1]$$

Theorem. Learning an arbitrary p over $[d]$ to TV loss ε has sample complexity $\Theta\left(\frac{d}{\varepsilon^2}\right)$.

Theorem. Testing if an arbitrary p over $[d]$ is u or has $\text{TV}(p, u) > \varepsilon$ has sample complexity $\Theta\left(\frac{\sqrt{d}}{\varepsilon^2}\right)$.

Baseline: the "centralised" setting

Discrete distributions

$d \gg 1$
 $\varepsilon \in (0, 1]$

Proof.

Baseline: the "centralised" setting

$$d \gg 1$$

$$\varepsilon \in (0, 1]$$

Identity-covariance Gaussians

Theorem. Learning the mean of an unknown $\mathcal{N}(\mu, \text{Id})$ to ℓ_2 loss ε^2 has sample complexity .

Theorem. Testing if an unknown $\mathcal{N}(\mu, \text{Id})$ has $\mu = \text{O}_d$ vs. $\|\mu\|_2 > \varepsilon$ has sample complexity .

Baseline: the "centralised" setting

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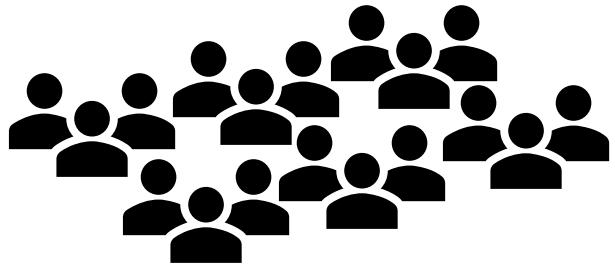
Baseline: the "centralised" setting

$$d \gg 1$$
$$\varepsilon \in (0, 1]$$

Identity - covariance Gaussians

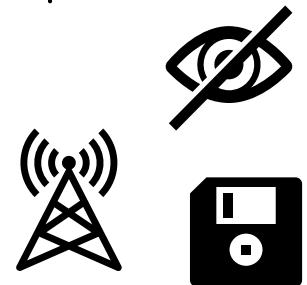
Proof.

Beyond the centralised setting



Distributed
data

Information or
computational constraints



Limited types of measurements

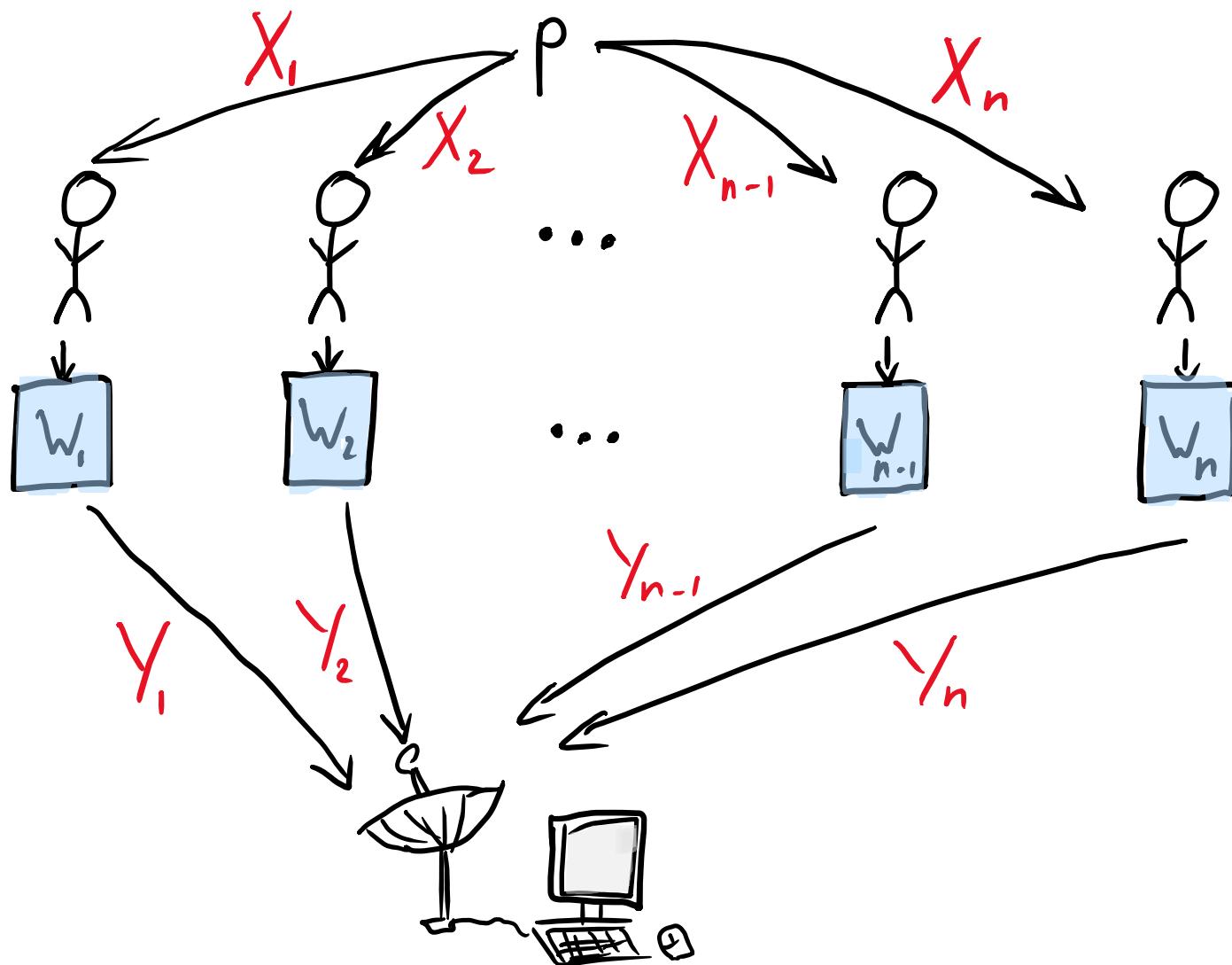
Beyond the centralised setting

- n users, each holding one sample from (same) p
- One center, which has no sample but needs to solve the learning / testing task
- Each user can only send a "limited" type of message

Beyond the centralised setting

- n users, each holding one sample from (same) p
 - One center, which has no sample but needs to solve the learning / testing task
 - Each user can only send a "limited" type of message
- "Local" constraint*

Beyond the centralised setting



Channels
 $W_1, \dots, W_n \in \mathcal{W}$

Beyond the centralised setting

Channel: $W: X \rightarrow Y$ randomised

\uparrow \uparrow
input output
space space
(samples) (messages)

Notation: $W(y|sc) = P\{W(sc) = y\}$

Beyond the centralised setting

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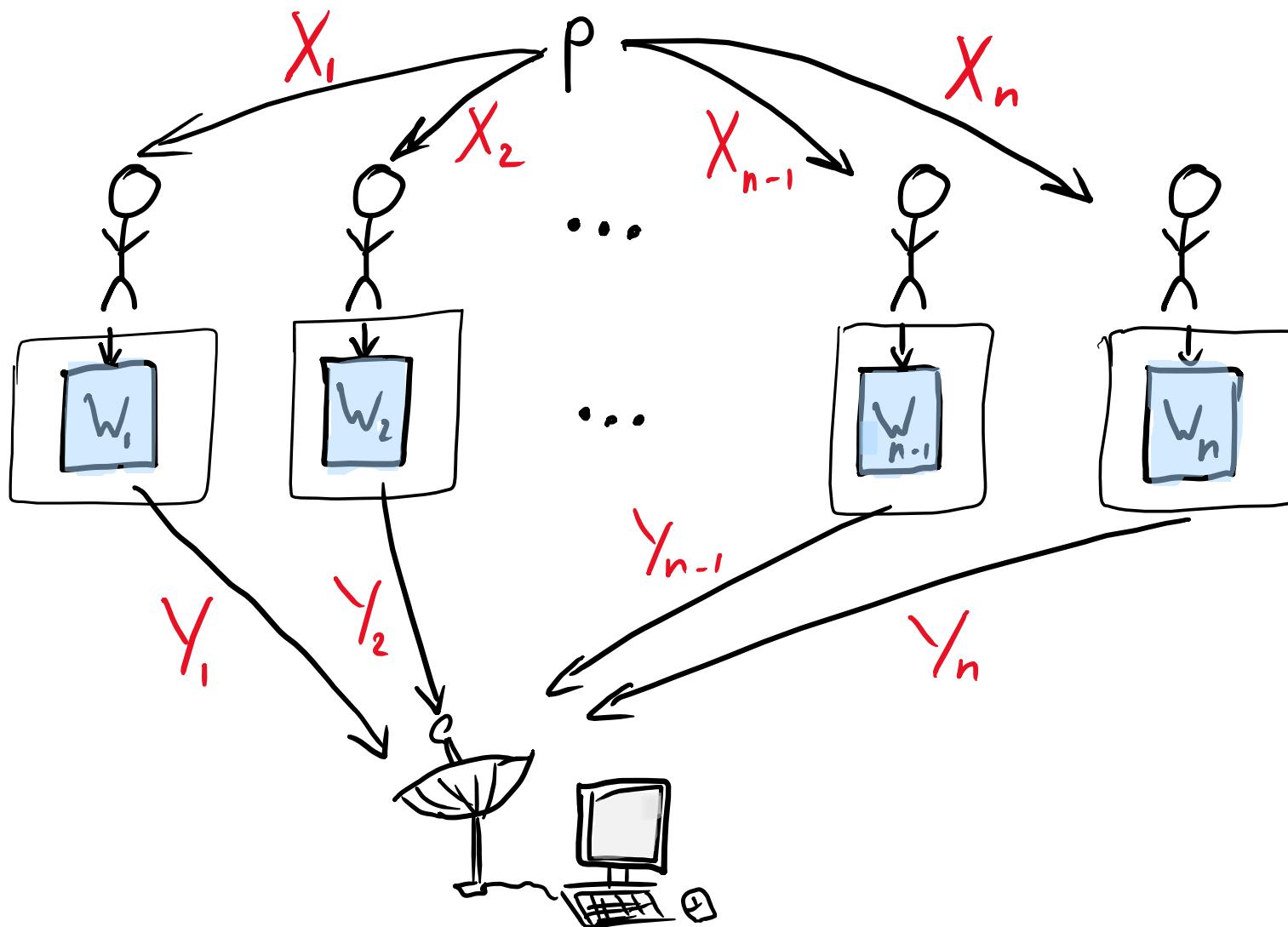
Notation: $W(y|sc) = P\{W(sc) = y\}$

$\mathcal{W} \subseteq \{X \rightarrow Y\}$ set of allowed channels

What happens if \mathcal{W} contains
the identity mapping?

Beyond the centralised setting

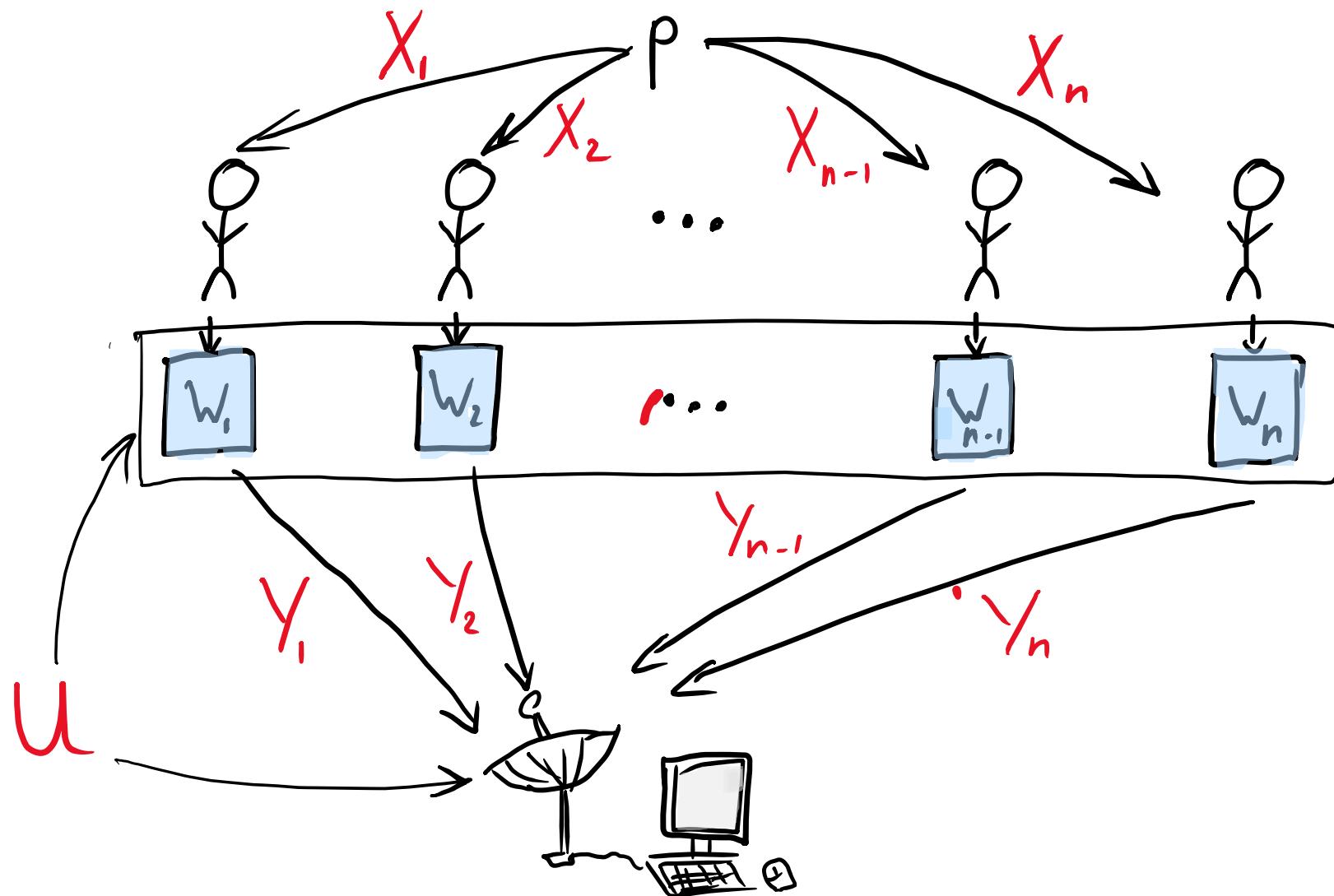
Private-coin



Channels
 $W_1, \dots, W_n \in \mathcal{W}$
independently
randomised

Beyond the centralised setting

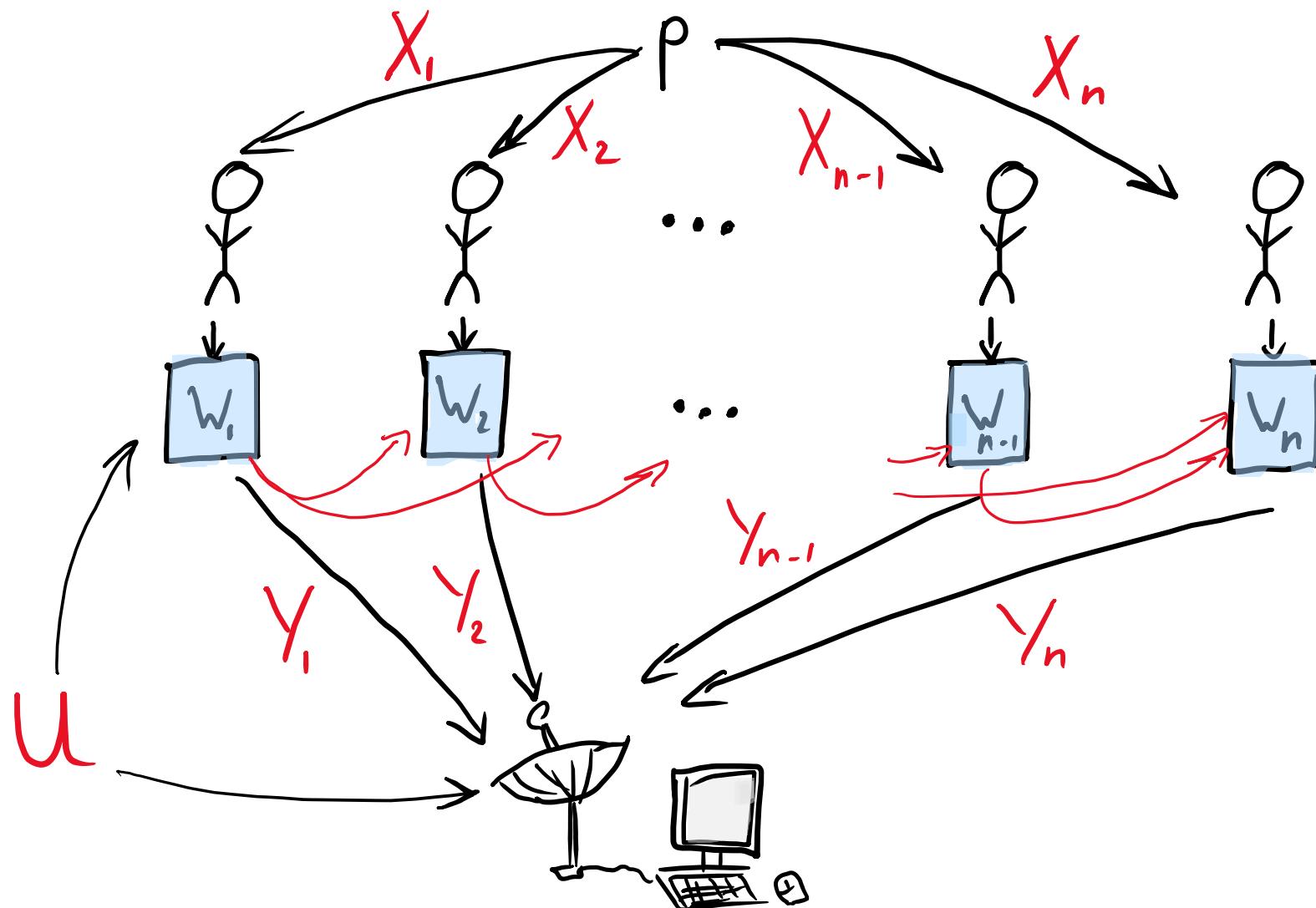
Public-coin



Channels
 $W_1, \dots, W_n \in \mathcal{W}$
- jointly
randomised

Beyond the centralised setting

Interactive



Channels
 $W_1, \dots, W_n \in \mathcal{W}$

$$W_t = W^{Y_1, \dots, Y_{t-1}}$$

depends on previous
messages
(+ public randomness)

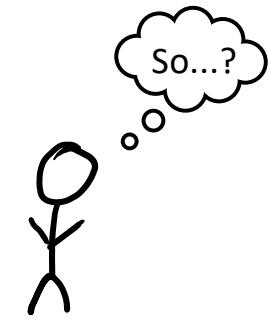
Beyond the centralised setting

Implementation
and deployment

Sample
complexity

Private-coin \leq Public-coin \leq Interactive

Private-coin \geq Public-coin \geq Interactive



Two guiding examples of channel families

Communication



Each user can only send ℓ bits

$$\mathcal{W}_\ell = \{ w: X \rightarrow \{0,1\}^\ell \}$$

Local Privacy



Each user requires ρ -differential privacy

$$\forall w \in \mathcal{W}_\ell$$

$$\forall y, x, x' \quad w(y|x) \leq e^\rho w(y|x')$$

Two guiding examples of channel families

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$$\forall w \in \mathcal{W}_\ell$$

$$\forall y, x, x' \quad w(y|x) \leq e^\rho w(y|x') \approx (1+\rho) w(y|x')$$

(think of $\rho \in (0,1]$)

Two guiding examples of channel families

Communication



Each user can only send ℓ bits

$$\mathcal{W}_\ell = \{ w: X \rightarrow \{0,1\}^\ell \}$$

Can't send too much

Local Privacy



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$$\forall w \in \mathcal{W}_\epsilon$$

$$\forall y, x, x' \quad w(y|x) \leq e^\epsilon w(y|x')$$

Can't reveal too much

Recap: this lecture

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Next lecture:

Learning and testing **discrete distributions** under
those information constraints

