Problems 1, 2, and 3 require you to have read the lecture notes or watched the lecture, but should be doable.

Problem 4 is important to have seen and attempted: you will go through it during the tutorial (but it is worth giving it some thought before). Problem 5 is recommended, and while the analysis to get the final bound on *C* is technical and somewhat annoying, it is a good idea to try to attempt the rest.

Problem 6 is quite technical and long (especially the question marked with a (\*)): it is alright to skip it, or to skip that subquestion if you attempt the problem. Attempt Problem 7 if you have time: it is not necessary, but gives perspective on the use of LSH.

## Warm-up

**Problem 1.** Give a data structure for the Nearest Neighbour problem over a d-dimensional universe using space O(nd), for which Query runs in time O(nd)). (Also, show that it can maintain S dynamically, and implement Insert and Remove methods running in time O(nd).)

**Problem 2.** Give a data structure for the Nearest Neighbour problem over  $\{0,1\}^d$  using space  $O(2^d)$ , for which Query runs in time  $O(2^d)$  (independent of n). (Also, can maintain S dynamically, and implement Insert and Remove methods running in time O(1).)

**Problem 3.** Check your understanding: since we want very efficient lookups and are willing to accept a small probability of failure for QUERY, can we use Bloom filters for the "baby version" of LSH instead of hash tables? What fails?

## Problem solving

**Problem 4.** (\*) Prove a simplified version of Theorem 38 from the lecture notes, showing how to solve the "general" ANN from the "baby version," at the cost of only a logarithmic factor in the ratio

$$\Delta = \frac{\max_{x,x' \in S} \operatorname{dist}(x,x')}{\min_{x,x' \in S} \operatorname{dist}(x,x')}$$

Note that, for the Hamming space  $\{0,1\}^d$ ,  $\Delta = O(d)$ , where d is the dimension.

**Problem 5.** Analyse the LSH family described in the lecture notes for the Euclidean case, where a locally-sensitive hash function  $h_g: \mathbb{R}^d \to \{-1,1\}$  is obtained by drawing a d-dimensional Gaussian random vector  $g \sim \mathcal{N}(0_d, I_d)$  (all coordinates are independent  $\mathcal{N}(0,1)$  normal random variables) and setting

$$h_g: x \in \mathbb{R}^d \to \operatorname{sign}\left(\sum_{i=1}^d g_i x_i\right)$$

We will make the (restrictive) assumption that all data points and query points have unit norm:  $||x||_2 = 1$ . Show that, for every r > 0, C > 1, this defines an (r, C, p, q)-LSH family with p, q such that  $\rho \le 1/C$ . [Note: this is called the SimHash scheme.]

**Problem 6.** ( $\star$ ) For the set  $[d] = \{1, 2, ..., d\}$ , let the universe  $\mathcal{X}$  be the set of all  $2^d$  subsets of [d], along with the *Jaccard distance*:

$$\operatorname{dist}(A,B) = 1 - \frac{|A \cap B|}{|A \cup B|}, \quad A,B \in \mathcal{X}$$

Consider the following hash family  $\mathcal{H}$ : for every permutation  $\pi$ :  $[d] \to [d]$ , define  $h_{\pi} \colon \mathcal{X} \to [d]$  by setting

$$h_{\pi}(A) = \min_{a \in A} \pi(a)$$

and  $\mathcal{H} = \{h_{\pi}\}_{\pi}$ .

- a) ( $\star$ ) Verify that the Jaccard distance is a metric on  $\mathcal{X}$ . What is its range?
- b) What is the size of  $\mathcal{H}$ ?
- c) Show that, for every  $r \in (0,1]$  and C > 1,  $\mathcal{H}$  is an (r,C,p,q)-LSH family for p = 1 r and q = 1 Cr. What is its sensitivity parameter  $\rho$ ?

## Advanced

**Problem 7.** Give a data structure for the Nearest Neighbour problem over the Euclidean space  $(\mathbb{R}^d, \ell_2)$  based on kd-trees. Analyse the space complexity of the data structure and its query time.