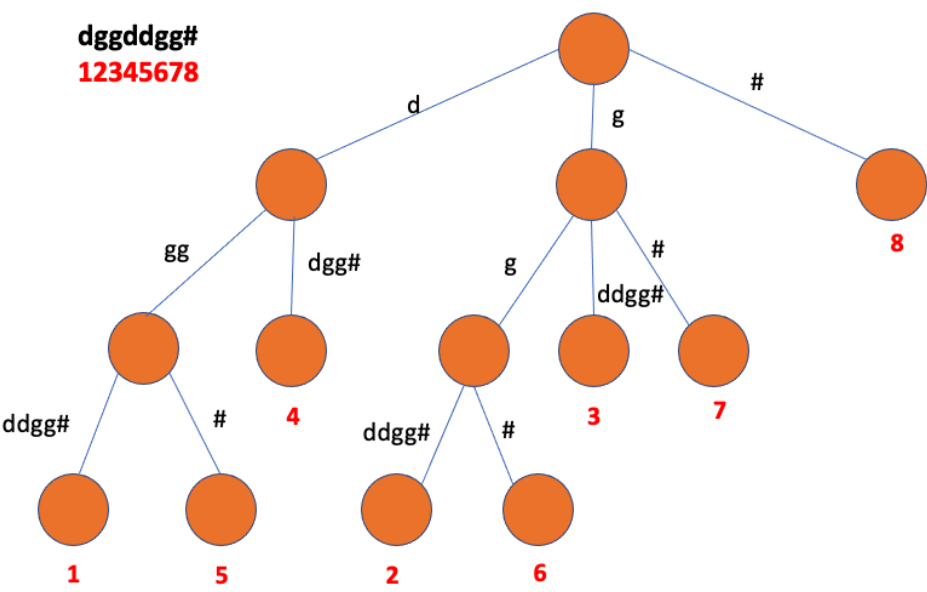


COP 5536 Fall 2020 Assignment 2

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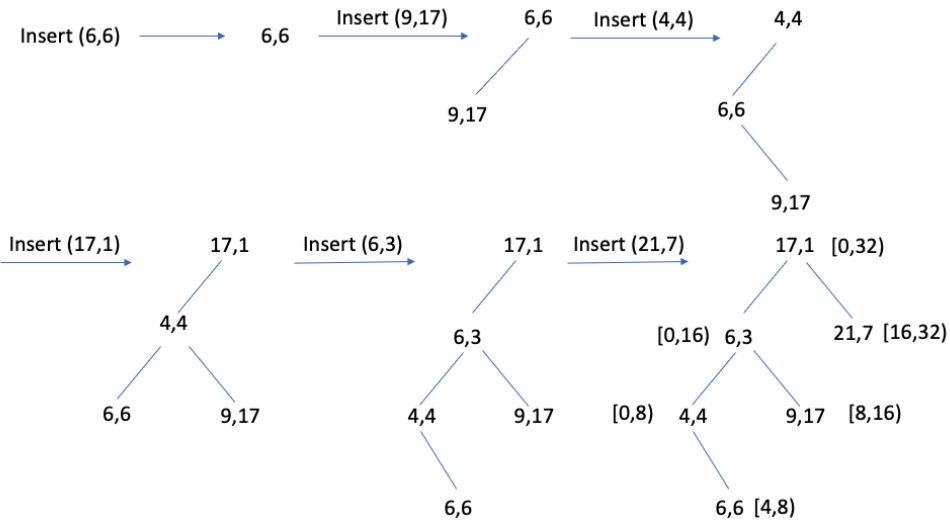
Q1: (a)



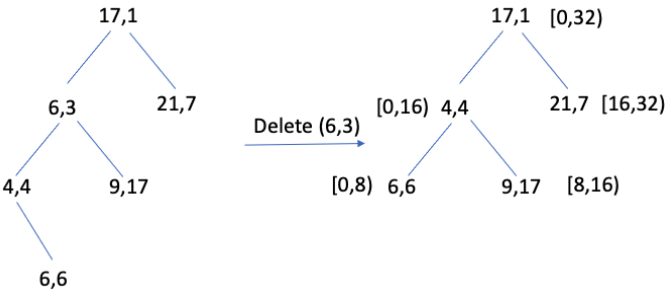
(b) : For two strings **S**; **T**,

1. create a generalized string **U = S\$T#**.
2. construct suffix tree for the generalized string **U** in $O(|S| + |T|)$. Each leaf of the tree represents either a suffix from one of the two strings or a suffix that occurs in both the strings
3. Mark each internal node with a **s(t)** if there is a leaf in the subtree of **v** representing a suffix from **S (T)**. The path-label of any internal node marked both **s** and **t** is a substring common to both S1 and S2, and the longest such sting is the longest common substrting.
4. So we find the node with the greatest string-depth that is marked both **s** and **t**. The node marking and calculations of string-depth can be done by standard linear-time tree traversal methods

Q2: (a)



(b)

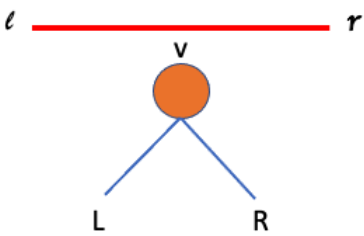


Q3: (a) $n = 100000$ is large, $h = (\ln 2)m/u \sim 0.693 * (5000/1000) = 3.47$, so $h = 3$ or 4 .

(b) $h = 3 \rightarrow p(u) = (1 - 1/n)^u * [1 - (1 - 1/m)^{uh}]^h = (1 - 1/100000)^{1000} * (1 - (1 - 1/5000)^{3000})^3 \sim 0.09095$

$h = 4 \rightarrow p(u) = (1 - 1/n)^u * [1 - (1 - 1/m)^{uh}]^h = (1 - 1/100000)^{1000} * (1 - (1 - 1/5000)^{4000})^4 \sim 0.09106$

Q4:



1. Query interval is $[l, r]$;
2. if $v.key \in [l, r]$, all intervals in v overlap and then search L ; R for additional overlapping intervals.

3. if $v.\text{key} < l$, intervals in v with $r_i \geq l$ overlap and then search R for additional overlapping intervals.
 4. if $v.\text{key} > r$, intervals in v with $l_i \leq r$ overlap and then search L for additional overlapping intervals.
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