

Group Signature by David Chaum's 4th Proposal

- KeyGen($1^n, n$); choose a prime p and $g \in \mathbb{Z}_p^*$.

Each member chooses $s_i \in \mathbb{Z}_p$ and computes $y_i \equiv g^{s_i} \pmod{p}$.

output $pk = (p, g, \{y_i\})$ and $sk_i = (s_i)$ for all $1 \leq i \leq n$.

- Sign(m, pk, s_i); compute $\sigma \equiv m^{s_i} \pmod{p}$

output (m, σ)

σ is a valid signature of m iff $\sigma \equiv m^{s_i} \pmod{p}$ and $g^{s_i} \in \{y_i\}$.

To prove this statement without leaking s_i , ZKP is used.

Privacy Problem in BTC

i) anonymity: hiding identities of sender and receiver \rightarrow by ring signature

ii) confidentiality: hiding the amount transferred.

\hookrightarrow by confidential transaction: every transaction amount is hidden using a commitment to the amount.

to prove: i) the sum of inputs is greater than the sum of outputs.

ii) all transactions values are positive.

) range proof.

Current Proposals for CT ZKP

{ [PBF+] : large proof size or required a trusted setup.

{ SNARKs : required a trusted setup.

STARKs : range proof size is large

\Rightarrow [BBB+H] Bulletproof @ S&P'18.