

7/26 ElGamal Encryption and Its Variants.

ElGamal Encryption

- KeyGen(\mathcal{P}); Choose a large prime p and $g \xleftarrow{\$} \mathbb{Z}_p^*$.

Choose a secret key $s \xleftarrow{\$} \mathbb{Z}_p^*$

Compute $y = g^s \bmod p$ (\because Discrete Logarithm Problem)

output $pk = (p, g, y)$ and $sk = (s)$

- Enc(pk, m); choose a random $r \xleftarrow{\$} \mathbb{Z}_p^*$.

output $(g^r, m \cdot y^r)$

- Dec(sk, c); parse c as (c_1, c_2) .

output $c_2 \cdot (c_1^s)^{-1}$.

$$\left(\begin{array}{l} m \cdot y^r \cdot (g^r)^{-s} = m \cdot (g^s)^r \cdot (g^r)^{-s} \\ \qquad \qquad \qquad = m \end{array} \right)$$

finite field.

- Mult(c_1, c_2); parse c_1 as $(c_{11}, c_{12}) = (g^{r_1}, m_1 \cdot y^{r_1})$

c_2 as $(c_{21}, c_{22}) = (g^{r_2}, m_2 \cdot y^{r_2})$

output $(c_{11} \cdot c_{21}, c_{12} \cdot c_{22}) = (g^{r_1} \cdot g^{r_2}, m_1 \cdot y^{r_1} \cdot m_2 \cdot y^{r_2})$
 $= (g^{r_1+r_2}, m_1 \cdot m_2 \cdot y^{r_1+r_2})$