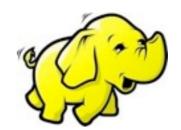
Why Prismatic Goes Faster With Clojure



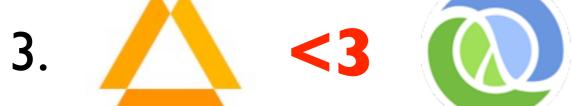
One Slide Summary

Fine-grained Composable > **A**bstractions

Monolithic Frameworks



lets you make FCA







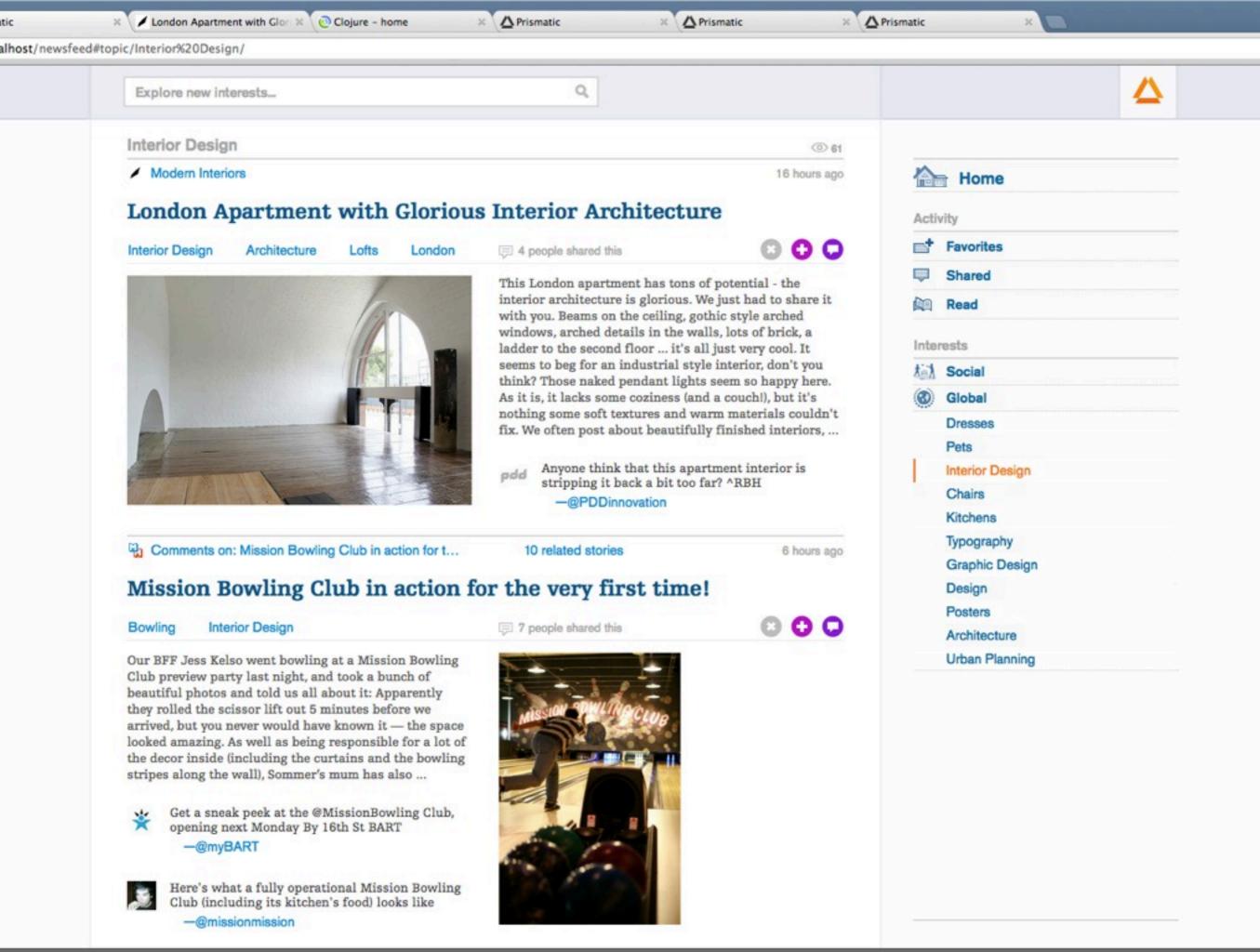
About Prismatic

We learn about your interests

Personalized feeds based on interests

Explore new interests

Live Demo At End



Our Backend Team

Three CS PhDs in Al

Me

Aria

Jenny Jason









Zero **Brogrammers**



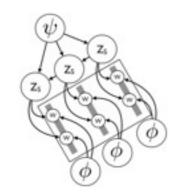
What We Build



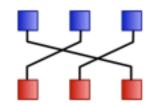
Web crawlers



Social Graph Analysis



Topic Models



Relevance Ranking

Newsfeeds

Real-time indexing of social, entity elements

Online clustering of related stories

Real-time personalized reranking of feeds

Must serve requests in under about 200ms

Our Design Approach



We tend to roll our own.



Libraries >> Frameworks





99.9% Clojure, 0.1% Java

Flop Library

- We do lots of double[] processing
- For efficiency, often in-place mutation
- Native Clojure makes this a PITA

```
;; Add (5.0+j) to j-th element of array
(dotimes [j (alength arr)]
  (aset xs j (+ 5.0 j (aget xs j))))
```

Flop Library

- Even type-hinting can yield inefficient code
- In Flop,

```
// Add (5.0+j) to j-th element of array (afill! [[j v] arr] (+ v 5.0 j))
```

- Succinct and efficient!
- Can't yield code with reflection

Flop Library

- Rare use of macros in our code
- doarr: doseq for double[]



$$w \cdot x = \sum_{i=1}^{n} w_i x_i$$

Inner loop in machine learning

prediction

 $\arg\max_{\ell} w \cdot x_{\ell}$

training

 $P(x; w) \propto \exp\{w \cdot x\}$



$$w \cdot x = \sum_{i=1}^{n} w_i x_i$$



$$w \cdot x = \sum_{i=1}^{n} w_i x_i$$

```
double dotProd(double[] ws, double[] xs) {
  double sum = 0.0;
  for (int i=0; i < xs.length; ++i) {
    sum += ws[i] * xs[i];
  }
  return sum;
}</pre>
```



$$w \cdot x = \sum_{i=1}^{n} w_i x_i$$

```
(defn dot-product [ws xs]
  (flop/asum [w ws x xs] (* w x)))
```



Expected

Log Probs

Inner loop in topic modeling

$$\psi_i = \mathbb{E}_{\theta}(\lg \theta_i | \alpha)$$
$$\theta \sim \text{Dirichlet}(\alpha)$$

$$\psi_i = \gamma(\alpha_i) - \gamma\left(\sum_{i=1}^n \alpha_i\right)$$

 $\begin{array}{ccc} \gamma(x) & \underline{\text{Digamma Function}} \\ & \text{expensive + gnarly} \\ & \text{Taylor approximation} \end{array}$



Expected

Log Probs

```
\psi_i = \mathbb{E}_{\theta}(\lg \theta_i | \alpha)\theta \sim \text{Dirichlet}(\alpha)
```

```
(defn exp-log-probs [alphas]
  (let [log-z (digamma (asum alphas))]
    (flop/amap [a alphas]
          (- (digamma a) log-z))))
```



- Comparable performance to tuned Java
 - State-of-the-art numerical optimization in
 < 180 lines
 - LDA-style topic modeling with variational inference < 180 lines



- Storage and aggregation abstractions
- Key-value protocol over Memory, File system, S3, BDB, Mongo, SQL
- implementations use specific features of underlying



- Key-value protocol: bucket/get, bucket/put
- the big deal: bucket/update
- can reify IMergeBucket: bucket/merge
- IWriteBucket has bucket/sync



- Automatic hosting for any store. f.ex. HTTP handlers for GET, PUT, MERGE ops for store & bucket.
- Easily test services by swapping persistent stores with memory stores
- Abstract over buffer & flush policies



Store Library

```
;; MERGE 1: index bigrams
(def bigrams
  (bucket/new
   {:type :mem
    :merge (partial merge-with +)}))
;; For each word, count following words
(doseq [[before after]
         (partition-all 2 words)]
   (bucket/merge bigrams
     before {after 1}))
```



Store Example

```
(defn map-reduce [map-fn reduce-fn n xs]
 (let [bspec {:type :mem :merge reduce-fn}
       bs (repeatedly n #(bucket/new bspec))
       work (fn [b x]
              (doseq \lceil \lceil k \lor \rceil (map-fn x)
                (bucket/merge b k v)))
       workers (map #(partial work %) bs)]
     ;; workers process xs in par, blocking
     (do-work workers xs)
     ;; merge all bucket users
     (reduce bucket/merge-all bs)))
```



```
;; MERGE 2: map reduce
(defn map-reduce [map-fn reduce-fn num-workers
input]
(let [pool (workers num-workers))
      agg-bucket #(bucket/new
        {:type :mem :merge reduce-fn})
       res (agg-bucket)
      in-queue (queue/new {:type :mem})
      sentinel (java.util.UUID/randomUUID)]
  (future (do (doseq [x input]
     (queue/offer in-queue x))
         (queue/offer in-queue sentinel)))
```



```
terminal-latch (CountDownLatch. 1)
mapper-latch (CountDownLatch. num-workers)
terminator (fn [x]
             (if (= x sentinel)
               (.countDown terminal-latch)
               (map-fn x))
defaults {:f terminator
          :in #(queue/poll in-queue)}
buckets (repeatedly num-workers agg-bucket)
```



Store Example

```
(doseq [b buckets]
 (exec/submit-to pool
  (let [b (agg-bucket)]
    #(if (= 0 (.getCount terminal-latch))
        (do (try
      (bucket/merge-to! b res)
      (finally (.countDown mapper-latch)))
    :done)
       (assoc defaults :out
         (fn [kvs]
    (doseq [[k v] kvs]
      (bucket/merge b k v)))))))
```



```
;;block on mapper encountering the sentinel
value
(.await terminal-latch)
;;other mappers could still be processing
tasks, ensure they finish.
(.await mapper-latch)
;;ensure all reducers are merged
(doseq [b buckets]
  (bucket/merge-to! b res))
(exec/shutdown-now pool)
res))
```



- wrapper policies
- caching & checkpointing
- buffering & flushing
- checkpoint & drain seqs: coming in store + graph example



Stream graph computation model

Separate specification from execution plan

Optimized for system throughput

```
;; Count entities in documents
(->> (graph)
     (gmap :doc-fetch (juxt :id get-text))
     (gmapcat :ent-tag
      (fn [[id text]]
       (map (fn [ent] [id ent])
        (nlp/extract-entities text))
    ;; Branch output to both nodes
    (>> (gmap :bmerge
         (fn [[id ent]]
           (bucket/merge ent-counts ent 1))
        (gmap:pub
          (publisher :topic "entities"))))
```

Graph Flexibility

- Graph input and outputs play nicely with Store and PubSub libraries
- Execution policies
 - 'compile' to a single fn
 - each node it's own machine/thread-pool
- Real win: monitoring and visibility

Graph Monitoring

 Each node monitors performance: cpu, exceptions, throughput, etc.

node	times	throughput	% cpu	% loss
:doc-fetch	450	1.5	0.10	0.02
:ent-tag	450	5.0	0.88	0.00
:bmerge	5,400	70.0	0.01	0.0
:pub	5,400	1,500	0.01	0.01

Graph + Store

Use graph to compute and monitor

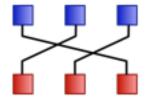
Store as terminal aggregation node

Quickly craft systems for problems

Graph + Store Example

- Online learning over streaming user events
- Collect statistics over time, periodically flush statistics to update existing ranking parameters
- Updating parameters is expensive so trigger batch updates after collecting 'enough' new user events

```
(def params ...)
(def suff-stats (bucket/new ...)
(->> (graph)
     (gmapcat :feature-extract
       (partial event-feats params))
     (gmap : feature-accum
       (fn [[event feat val]]
        (bucket/merge suff-stats
           event {feat val})))
(cron-job
  #(update-params! params
    (bucket/flush suff-stats))
  [60 :minutes])
```



Demo