

# Geometric realization of diagrammatic sets

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## 1 Directed simplex category

### 1.1 Combinatorics and lattice structure

We call  $\Delta$  the simplex category whose objects are finite linear orders  $[n] = \{0 < \dots < n\}$  and morphisms (weakly) monotonic functions. It is generated by the face maps  $\delta_i : [n-1] \rightarrow [n]$  and the coface maps  $\sigma_j : [n+1] \rightarrow [n]$ .

**Definition 1 (Ordering):** An *ordering*  $P$  on  $[n] \in \Delta$  is a reflexive relation  $P \subseteq [n] \times [n]$ .

**Definition 2 (Directed simplex):** The *directed simplex* category  $d\Delta$  has object dependant pairs  $([n], P \subseteq [n] \times [n])$  with  $P$  an ordering, and morphisms are the one of  $\Delta$  that preserve orderings. More precisely,  $f : ([n], P) \rightarrow ([m], Q)$  is a morphism  $f : [n] \rightarrow [m]$  in  $\Delta$  such that  $f(P) \subseteq Q$ , or equivalently,  $P \subseteq f^{-1}(Q)$ .

We consider  $\Delta$  to be the subcategory of  $d\Delta$  by sending  $[n]$  to  $([n], [n] \times [n])$ . As in  $\Delta$ , we have the factorisation lemma of the maps.

**Lemma 3 (Factorisation):** Any map  $f : ([n], P) \rightarrow ([m], Q)$  factorizes as

$$f = d_{i_1} \dots d_{i_p} s_{j_1} \dots s_{j_q}$$

where the orderings are constructed iteratively, starting from  $P$ , that is:

$$s_{j_k} : ([n], s_{j_{k+1}} \dots s_{j_q}(P)) \rightarrow ([n - k], s_{j_k} s_{j_{k+1}} \dots s_{j_q}(P))$$

and similarly for  $d_{i_k}$ , except for  $d_{i_1} : ([n], d_{i_2} \dots d_{i_p} s_{j_1} \dots s_{j_q}(P)) \rightarrow ([m], Q)$  (or for  $s_{j_1}$  if  $p = 0$ ).

*Proof.* (**Lemma 3**) We use the decomposition from the underlying map  $f$  in  $\Delta$ , and by construction, each  $s_{j_k}, d_{i_k}$  lift to  $d\Delta$ .  $\square$

## 2 Geometric realisation

### 2.1 Directed geometric realisation

In this section, we describe how to realize a combinatorial computad (atoms + inclusions) into a directed topological space. The need to realize only the atoms, rather than all the category of diagrammatic sets comes from the absence of good functorial properties from the map we will describe below, in the presence of surjections between diagrams. Moreover, without the surjections, we consider the category  $d\Delta$  to be the directed semi-simplicial category, meaning that it is generated only by the maps  $\delta_i$  with the equation  $\delta_i \delta_j = \delta_{j+1} \delta_i$  whenever  $i \leq j$ . This is not too much of a restriction as Simon Henry's work on weak model structures shows that one can do homotopy theory just fine in this way. All the results of the previous section also apply.

We want to turn a directed semi-simplicial set into a directed topological space. For that, we realize the underlying semi-simplicial set the usual way, and we add a set of paths according to the ordering.

**Definition 4 (d-spaces and d-maps):** A *d-space*  $(X, dX)$  is a topological space together with a set  $dX$  of dipaths such that:

1. every constant path is a dipath,
2. If  $\varphi : [0, 1] \rightarrow [0, 1]$  is a directed partial reparametrisation (i.e. weakly increasing), then  $f \circ \varphi$  is a dipath for all  $f \in dX$ ,
3. The concatenation of two dipaths is a dipath.

A *d-map* (or dimap)  $f : (X, dX) \rightarrow (Y, dY)$  is a continuous map such that for all  $h \in dX$ ,  $f \circ h \in dY$ . This gathers into the category **dTop**.

If  $x_i, x_j$  are points in  $\mathbb{R}$ , we call  $\gamma_{ij}$  the path defined by  $\gamma_{ij}(t) = tx_i + (1 - t)x_j$ .

**Definition 5 (Directed geometric realisation):** The *directed geometric realisation* of an element  $([n], P) \in d\Delta$  is the directed topological space  $(\Delta^n, dP)$  where  $\Delta^n$  is the convex hull spanned by  $n + 1$  affine independent points  $x_0, \dots, x_n$  of  $\mathbb{R}^{n+1}$ , and  $dP$  is the closure under concatenation and reparametrisation of the set of paths:

$$\left\{ \sum_{(i,j) \in P} c_{ij} \gamma_{ij} \mid \sum_{(i,j) \in P} c_{ij} = 1 \right\}$$

The directed geometric realisation has underlying topological space the standard geometric realisation. The set  $dP$  forms a set of dipaths. Indeed by definition it is closed under concatenation and reparametrisation. It also contains all the constant path, as if  $x$  is in  $\Delta^n$ , then  $x = \sum_{i=0}^n c_i x_i$  with  $\sum_{i=0}^n c_i = 1$ , thus by setting  $c_{ij} = 0$  if  $i \neq j$  and  $c_{ij} = c_i$  if  $i = j$ , and by noticing that  $P$  is reflexive, we obtain the constant path at  $x$ . To make it a functor  $|-|_d : d\Delta \rightarrow \mathbf{dTop}$ , we only need to check that the underlying continuous maps from the standard simplex lifts to maps of directed topological space. Consider  $f : ([n], P) \rightarrow ([m], Q)$ , and take  $\gamma = \sum_{(i,j) \in P} c_{ij} \gamma_{ij}$  in  $dP$ . We want to show that  $|f|\gamma$  is in  $dQ$ . By definition,  $|f|$  distributes over barycentric combinations, and  $|f|\gamma_{ij} = \gamma_{f(i)f(j)}$ . For  $(i, j) \in Q$ , set

$$c'_{ij} = \sum_{(k,l) \in f^{-1}(i,j)} c_{kl}$$

where  $f^{-1}(i, j)$  is shorthand for  $(f^{-1}(i) \times f^{-1}(j)) \cap P$ , and

$$|f| \circ \gamma = \sum_{(k,l) \in P} c_{kl} \gamma_{f(k)f(l)} = \sum_{(i,j) \in Q} c'_{ij} \gamma_{ij}$$

with  $\sum_{(i,j) \in Q} c'_{ij} = \sum_{(k,l) \in P} c_{kl} = 1$ , as  $f(P) \subseteq Q$ . Thus,  $|f| \circ \gamma$  is an admissible path of  $dQ$ .

We can left-Kan extend this construction and obtain a functor  $|-|_{d\Delta} : \mathbf{dsSet} \rightarrow \mathbf{dTop}$ .

## 2.2 From combinatorial computad to directed simplicial sets

We call  $\odot_i$  the skeleton of the category of combinatorial computad, that is atoms and inclusions. We have the following picture:

$$\begin{array}{ccc}
 & d\Delta & \xrightarrow{|-|_{d\Delta}} \mathbf{dTop} \\
 & \downarrow y & \nearrow |-|_{d\Delta} \\
 \odot_i & \xrightarrow{?} \mathbf{dsSet} & \\
 \downarrow y & \nearrow \text{Lan}_y(?) & \\
 \odot_i \mathbf{Set} & & 
 \end{array}$$

We would like to find the dashed functor, thus giving a geometric realization of a diagrammatic set. For that, recall from [Had20] that we have a functor:

$$k : \odot_i \rightarrow \mathbf{sSet}$$

sending a molecule  $U$  to the simplicial nerve of its underlying poset. If  $U$  was already a simplex, it would correspond to constructing its barycentric subdivision. We would like to extend it to a functor  $k_d : \odot_i \rightarrow \mathbf{dsSet}$ . As  $\Delta \subseteq d\Delta$ , we already have a functor  $\mathbf{sSet} \rightarrow \mathbf{dsSet}$ , however doing so would lose all the relevant information about the orientation of the diagram  $U$  in  $\odot_i$  that we try to convey. Thus, we need to refine the functor  $k$ , by adding manually orientations. To that end, we introduce the intermediate category of preordered posets.

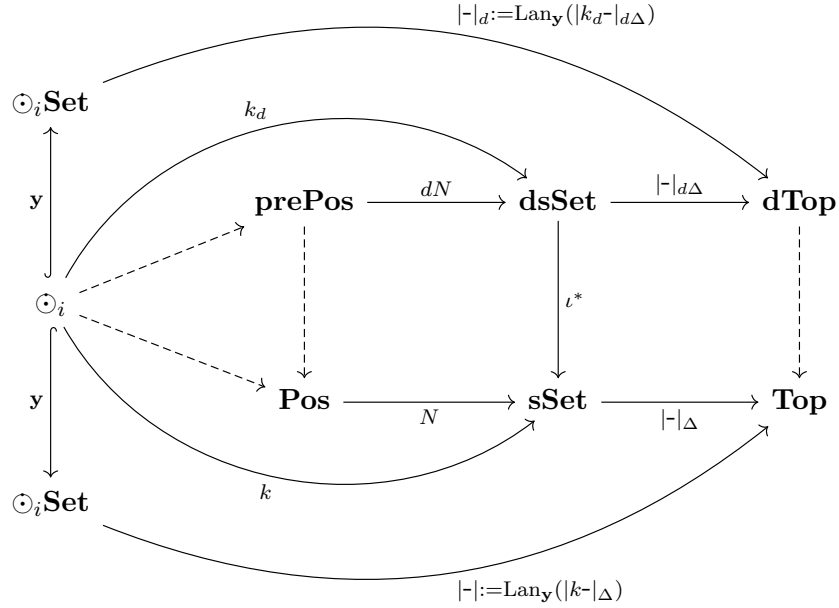
**Definition 6 (Preordered poset):** A *preordered poset* is a triple  $(P, \leq, \preceq)$  where  $(P, \leq)$  is a poset and  $(P, \preceq)$  is a preorder. A map  $f : (P, \leq, \preceq) \rightarrow (Q, \leq, \preceq)$  is a map respecting both the poset and the preorder structure. This form the category **prePos**.

We have a forgetful functor  $\mathbf{prePos} \rightarrow \mathbf{Pos}$  and another  $\odot_i \rightarrow \mathbf{Pos}$ . We now describe how to factorize them via another forgetful functor from  $\mathcal{U} : \odot_i \rightarrow \mathbf{prePos}$ . Let  $V$  be an atom, we define  $\mathcal{U}V := (\mathcal{H}V, \leq, \preceq)$ , where the posetal structure is the one of its Hasse diagram, and the preorder is the transitive and reflexive closure of the relation given by:

$$x \preceq y \iff o(x \leftarrow y) = - \text{ or } o(y \leftarrow x) = +$$

As inclusions preserve orientation, this construction is functorial. We may now construct the directed nerve functor  $dN : \mathbf{prePos} \rightarrow \mathbf{dsSet}$ . Let  $(X, \leq, \preceq)$  in **prePos**. We define  $dN(X)([n], P)$  to be the set of  $(n+1)$  chains  $(x_0 \leq \dots \leq x_n)$  (like in the ordinary nerve) that are moreover such that for all  $(i, j) \in P$ , we have  $x_i \preceq x_j$ . We call  $k_d : \odot_i \rightarrow \mathbf{dsSet}$  the functor  $dN \circ \mathcal{U}$ .

We summarize our constructions with the following diagram where everything commutes. A dashed arrow indicates a forgetful functor.



## 2.3 From diagrammatic sets to directed simplicial sets

In this part, we study a different version of the geometric realisation of diagrammatic sets. We will take the same approach as in the previous section, but we will have to deal with surjective maps between molecules, and thus we will need to take more care when constructing the forgetful functor to **prePos**. We notice that we can leave unchanged the **Definition 5** on geometric realisation, as the proof of functoriality equally applies if we add degeneracies. What breaks is the construction of the preorder  $\preceq$ , that does not behave functorially as the surjections do not preserve orientations in general. Thus, we will adapt it here.

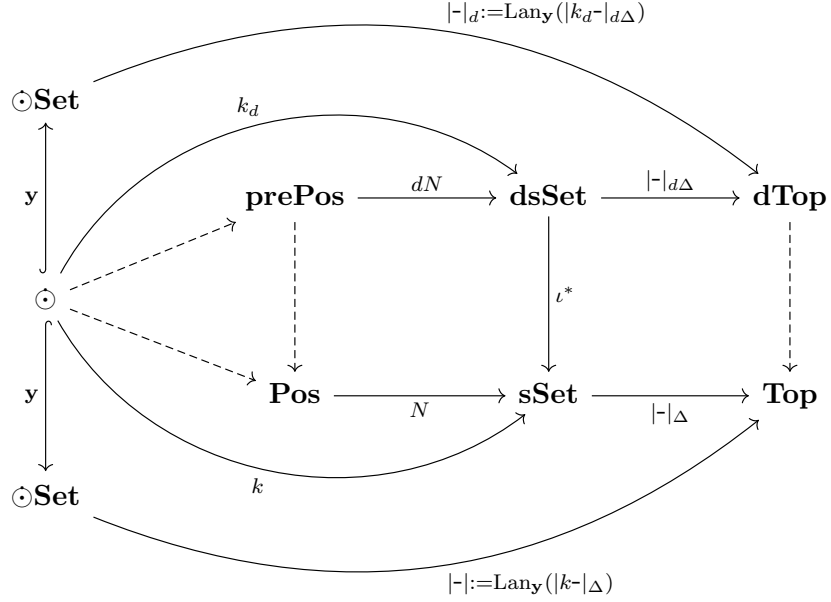
We aim to construct a functor  $\mathcal{U} : \odot \rightarrow \mathbf{prePos}$ . For that we consider  $V \in \odot$ , and as usual, we associate it its underlying poset  $(\mathcal{H}V, \leq)$ . Now, the preorder is the reflexive and transitive closure of the relation given by:

$$x \preceq y \iff x \in \partial_0^- y \text{ or } y \in \partial_0^+ x$$

That is whenever  $x \preceq y$  (and  $x \neq y$ ), then  $x$  or  $y$  is a point (or both by transitivity). Now, a map  $f : P \rightarrow Q$  of  $\odot$  satisfies in particular

$$\partial_0^\alpha f(x) = f(\partial_0^\alpha x)$$

for all  $x \in P$ , thus if  $x \in \partial_0^- y$ , then  $f(x) \in \partial_0^- f(y)$ , making  $\mathcal{U}$  a functor  $\odot \rightarrow \mathbf{prePos}$ . We now have the same picture as previously



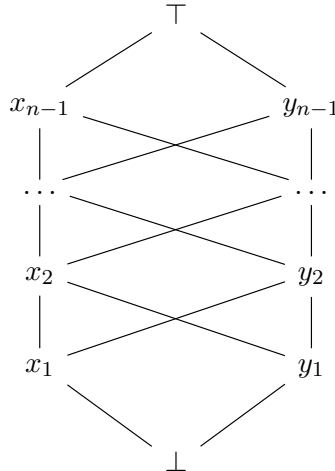
with the whole category of diagrammatic sets.

### 3 Examples

We may now work with the latter geometric realisation.

#### 3.1 The cone

Consider the poset  $C_n$ :



and preorder  $\perp \preceq x_i, y_i$  for all  $1 \leq x_i, y_i \leq n-1$ . Notice that this is the pre-posetal set  $\mathcal{U}(O^n)$  where we remove the output 0-boundary. Abusing notations, we also call  $C_n$  the geometrical cone of  $\mathbb{R}^{n+1}$  whose base is a  $(n-2)$ -square, whose apex  $\perp$  is at the origin, and whose height is parallel to an axis that we call  $z$ . If a point  $x$  is in the cone, we call  $\theta$  the angle between the  $z$  axis and the line  $\perp x$ . This way, each path determines another continuous path  $t \mapsto (z(t), \theta(t))$ . We endow  $C_n$  with a directed topological structure by taking the dipaths  $dC_n$  to be all the piecewise-linear paths whose  $(z, \theta)$  value is weakly increasing.

**Proposition 7 (Realisation of the cone):** We have

$$|C_n|_d \simeq (C_n, dC_n)$$

*Proof.* (**Proposition 7**) First, it is the case that this is the realisation as (non-directed) topological spaces, indeed  $C_n$  will be the glueing along the correct faces of  $2^{n-1}$   $n$ -simplices  $\perp u_1 \dots u_{n-1} \top$  where  $u_i \in \{x_i, y_i\}$ , each one determined by a path from  $\perp$  to  $\top$  in the Hasse diagram of the cone. Each of this  $n$  simplex share the edge  $\top \perp$  and it suffice to check that in each of these simplex, the admissible path are piecewise linear, weakly increasing along the  $\top \perp$  edge.

Let  $(D, dD) := |\perp u_1 \dots u_{n-1} \top|_d$  be such a directed simplex. To determine its admissible paths, it suffices to determine the maximal orderings  $P$  such that for all  $(i, j) \in P$ ,  $D_i \preceq D_j$ . Indeed, if there is a smaller ordering  $Q$ , then it will induce an inclusion of directed topological spaces via  $\text{id} : ([n], Q) \rightarrow ([n], P)$ . By construction of the preorder of the cone, we have

$$P = \text{Diag} \cup \{0i \mid 1 \leq i \leq n\}$$

Thus,  $dD$  is the closure under concatenation and reparametrisation of

$$\left\{ \sum_{i=1}^n c_i \gamma_{0i} \mid \sum_{i=1}^n c_i = 1 \right\}$$

Note that by reparametrisation, one need not to include the constant paths in the sum. These paths are piecewise linear, and indeed each  $\gamma_{0i}$  is weakly increasing along  $(z, \theta)$ .

Conversely, all linear path of increasing value along  $\perp \top$  are of the form (by reparametrisation)  $t \mapsto tx + (1-t)y$  with  $x \leq_{(z, \theta)} y$ , and such a path can be written (how?) as a sum  $\sum_{i=0}^n c_i \gamma_{0i}$ .  $\square$

### 3.2 Globe, interval, sphere

Let  $C'_n$  be the cone  $C_n$  but with preorder  $x_i, y_i \preceq \perp$ , i.e.  $C'_n = (C_n, \leq, \preceq^{\text{op}})$ . If  $C$  is a preposetal set, then let  $C^*$  be the same preposetal set without its bottom element (when it exists). We have in **prePos**:

$$(C'_n)^* \simeq C_n^*$$

The globe  $O^n$  is then the pushout of inclusions in **prePos**:

$$\begin{array}{ccc} C_n'^* \simeq C_n^* & \hookrightarrow & C_n' \\ \downarrow & & \downarrow \\ C_n & \longrightarrow & O^n \end{array}$$

We call  $\perp \preceq x_i, y_i, \top \preceq \perp'$  the elements of  $O^n$ . Consider  $(I^n, dI^n)$  the directed cube  $[0, 1]^n$  whose dipaths are the piecewise linear weakly increasing ones. In  $C'_n$ , the admissible dipaths are the ones whose  $(z, \theta)$  value is weakly decreasing.

**Lemma 8 (Admissible dipaths):** A path  $\gamma : [0, 1] \rightarrow I^n$  is in  $dI^n$  if and only if the map  $t \mapsto (z(t), \theta(t))$  is weakly increasing while  $z(t) \leq \frac{1}{2}\sqrt{n}$  and the map  $t \mapsto (z(t), -\theta(t))$  is weakly increasing while  $z(t) \geq \frac{1}{2}\sqrt{n}$ .

*Proof.* (**Lemma 8**) Cut the  $n$ -cube  $[0, 1]^n$  along the bisecting hyperplane of the diagonal  $0^n 1^n$  and notice that the bottom part is a directed copy of  $C^n$  and the top part is a directed copy of (a symmetric version of)  $C'_n$ .  $\square$

**Corollary 9 (Realisation of the globes):** We have

$$|O^n|_d \simeq (I^n, dI^n)$$

**Corollary 10 (Realisation of the interval):** The diagrammatic interval  $O^1$  is realized as the directed interval  $[0, 1]$  with weakly increasing paths.

**Corollary 11 (Realisation of the spheres):** The sphere  $\partial O^n$  is realized as  $\partial \vec{I}^n = [0, 1]^n \setminus (0, 1)^n$  whose directed paths are the one inherited from  $\vec{I}^n$

*Proof.* (**Corollary 11**) The same applies, except this time we only glue the boundaries of the  $n$  simplices (whose interior is then empty).  $\square$

### 3.3 Gray product

**Lemma 12:** The forgetful functor  $\odot \rightarrow \mathbf{prePos}$  sends Gray products to cartesian products.

*Proof.* (**Lemma 12**) The forgetful functor sends Gray products to cartesian products in **Pos**. Moreover, by definition

$$x \otimes y \preceq x' \otimes y'$$

if and only if  $x \otimes y \in \partial_0^- x' \otimes y'$  or  $x' \otimes y' \in \partial_0^+ x \otimes y$ , but the 0-boundaries commute with Gray products, that is  $\partial_0^\alpha(U \otimes V) = \partial_0^\alpha U \otimes \partial_0^\alpha V$ , thus  $x \otimes y \preceq x' \otimes y'$  if and only if  $x \preceq x'$  and  $y \preceq y'$ , giving indeed the cartesian product in **prePos**.  $\square$

**Proposition 13 (Realisation of Gray products):** Let  $U, V \in \odot$ , then

$$|U \otimes V|_d \simeq |U|_d \times |V|_d$$

*Proof.* (**Proposition 13**) First,  $k_d(U \otimes V) = U \times V$  in **dsSet**, because the Gray product is sent to the cartesian products by **Lemma 12**, and the directed nerve functor is right adjoint, thus preserves products.

That means it now suffices to show that the realisation of directed simplicial sets preserves finite products. It is the case for the simplicial sets, thus it suffices to show that the admissible paths of the realisation of the product is the product of the admissible paths of the realisation. This is the

case for purely formal reasons. Indeed, as the standard geometric realization preserves products, a point  $x \otimes y$  is realized as the couple  $(|x|, |y|)$ . Thus, the path  $\gamma_{ii',jj'} : (x_i, x_{i'}) \rightarrow (x_j, x_{j'})$  is the path  $(\gamma_{ij}, \gamma_{i'j'})$ . So, for two orderings  $P, Q$ , we can write

$$\begin{aligned} \sum_{(ij,i'j') \in P \times Q} c_{ij,i'j'} \gamma_{ii',jj'} &= \left( \sum_{(ij,i'j') \in P \times Q} c_{ij,i'j'} \gamma_{ij}, \sum_{(ij,i'j') \in P \times Q} c_{ij,i'j'} \gamma_{i'j'} \right) \\ &= \left( \sum_{ij \in P} \left( \sum_{i'j' \in Q} c_{ij,i'j'} \right) \gamma_{ij}, \sum_{i'j' \in Q} \left( \sum_{ij \in P} c_{ij,i'j'} \right) \gamma_{i'j'} \right) \end{aligned}$$

and both paths are admissible in their respective spaces.  $\square$

**Corollary 14 (Realisation of the cubes):** The realisation of the  $n$ -cube is  $[0, 1]^n$  with admissible paths the piecewise linear increasing ones.

*Proof.* (**Corollary 14**) The  $n$ th cube is defined to be  $O^1 \otimes \dots \otimes O^1$   $n$  times, thus by **Proposition 13**, its realisation is the  $\tilde{I}^n$ , which is the cube  $[0, 1]^n$  whose paths are piecewise linear increasing increasing.  $\square$

Notice that we have in **dTop**,  $|O^n|_d \simeq |\square^n|_d$ .

### 3.4 Dihomotopies and reversible paths

A diagram of shape  $f : O^1 \rightarrow X$  in a diagrammatic set exhibits an admissible dipath  $|f| \in dX$ , as it is realized as a map  $|f| : \tilde{I} \rightarrow |X|_d$ , and we have the isomorphism (of topological spaces)  $(|X|_d)^{\tilde{I}} \simeq dX$ . If  $x : 1 \rightarrow X$ , we allow ourselves to name  $x \in |X|_d$  the point determined by the map  $|x|_d : 1 \rightarrow |X|_d$ , thus we have that a diagram of shape  $x \xrightarrow{f} y$  in  $X$ , it is realized as the concatenation of the paths  $|f| = \gamma_{xf} \cdot \gamma_{fy}$ .

**Definition 15 (Dihomotopies):** Let  $(X, dX)$  be a directed topological space, and let  $f, g : x \rightarrow y$  be two admissible dipaths of  $dX$ . A *dihomotopy*  $\varphi : f \rightarrow g$  is a continuous map  $\varphi : [0, 1] \rightarrow dX$  such that  $\varphi(0) = f$ ,  $\varphi(1) = g$  and for all  $t \in [0, 1]$ , we have  $\varphi(t) \in dX$ . If there is a dihomotopy  $\varphi : f \rightarrow g$ , we write  $f \sim_d g$ . This is an equivalence relation.

A dihomotopy is simply a (particular) homotopy between in the space  $dX$ , thus is a morphism  $\varphi : \tilde{I} \times I \rightarrow X$  in **dTop**.

**Lemma 16 (Diagrams in the realisation):** Let  $X$  be a diagrammatic set.

1. A diagram of shape  $x \xrightarrow{f} y \xrightarrow{g} z$  in  $X$  is realized as the concatenation of the paths in  $dX$  determined by the subdiagrams  $x \xrightarrow{f} y$  and  $y \xrightarrow{g} z$ .
2. A diagram of shape

$$\begin{array}{ccc} & g & \\ x & \xrightarrow{\quad} & y \\ & f & \end{array} \quad \begin{array}{c} \uparrow \alpha \\ \uparrow \end{array}$$

in  $X$  proves that the paths determined by the subdiagrams  $f : O^1 \rightarrow X$  and  $g : O^1 \rightarrow X$  are dihomotopic.



*Proof.* (**Lemma 16**) The first point is clear as  $|O^1 \#_0 O^1|_d = (|O^1|_d \amalg |O^1|_d) / (\partial^+ O^1 = \partial^- O^1)$ , thus the diagram of shape  $f; g : O^1 \#_0 O^1 \rightarrow X$  is indeed the concatenation of two paths and the two inclusions

$$\begin{array}{ccccc} O^1 & \hookrightarrow & O^1 \#_0 O^1 & \hookleftarrow & O^1 \\ & \searrow f & \downarrow f;g & \swarrow g & \\ & & X & & \end{array}$$

ensure the equality of the desired paths.

For the second point, we notice the path  $|f|$  is the concatenation of paths  $\gamma_{xf} \cdot \gamma_{fy}$ , and similarly  $|g| = \gamma_{xg} \cdot \gamma_{gy}$ . It suffices to show that the linear interpolation of those two paths via  $\varphi(t) = t|f| + (1-t)|g|$  is a dihomotopy. To show that  $\varphi(t)$  is a dipath for all  $t$ , we can notice that we can write it as the concatenation:

$$\varphi(t) = (t\gamma_{xf} + (1-t)\gamma_{xg}) \cdot (t\gamma_{fy} + (1-t)\gamma_{gy})$$

and for all  $t$ , both  $t\gamma_{xf} + (1-t)\gamma_{xg}$  and  $t\gamma_{fy} + (1-t)\gamma_{gy}$  constitute admissible dipaths by construction.  $\square$

We could similarly define the notion of  $\vec{I}$ -dihomotopy as maps  $\vec{I} \times \vec{I} \rightarrow X$ , but the realisation we are working with will not detect the difference between

$$\begin{array}{ccc} & g & \\ x & \begin{array}{c} \curvearrowright \\ \uparrow \alpha \\ \curvearrowleft \end{array} & y \\ & f & \end{array}$$

and

$$\begin{array}{ccc} & g & \\ x & \begin{array}{c} \curvearrowright \\ \alpha \downarrow \\ \curvearrowleft \end{array} & y \\ & f & \end{array}$$

and would not allow us to interpolate the paths  $|f|$  and  $|g|$ . Indeed, it would require that the paths  $t \mapsto t|f|(u) + (1-t)|g|(u)$  are directed for all  $0 \leq u \leq 1$ , but then taking for instance  $u = \frac{1}{2}$ , we would end up with the path  $\gamma_{f\alpha} \cdot \gamma_{\alpha g}$  that does not belong to the realisation, as we do not have, for instance,  $f \preceq \alpha$  or  $\alpha \preceq f$ .

However, this difference would be detected by the first realisation, and it would be possible to construct  $\vec{I}$ -dihomotopies  $f \rightarrow g$  in the first diagram and  $g \rightarrow f$  in the second.

**Definition 17 (Contractible):** A path  $f : x \rightarrow x$  is *contractible* when there exists a dihomotopy  $\varphi : f \rightarrow x$ , where  $x$  is the constant path at  $x$ .

**Proposition 18 (Contractible paths):** Let  $X$  be a diagrammatic set, then a diagram of shape

$$\begin{array}{ccc} x & \xrightarrow{\varepsilon x} & x \\ & \uparrow & \\ & y & \\ & \downarrow & \\ x & \xrightarrow{f^{-1}} & x \end{array}$$

in  $X$  proves that the path  $|f| \cdot |f^{-1}| : x \rightarrow x$  is contractible, and in **Top**, we have  $|f^{-1}| \sim |f|^{-1}$

*Proof.* (**Proposition 18**) A diagram of shape  $\varepsilon_x = !; x : O^1 \rightrightarrows 1 \rightarrow X$  exhibit a path of  $dX$ , and the commutation  $|\varepsilon_x| = !; |x|$  indicates it is constant. By **Lemma 16**, we conclude that the path  $|f| \cdot |f^{-1}|$  is dihomotopic to the constant path at  $x$ . Moreover:

$$\begin{aligned} & |f| \cdot |f^{-1}| \sim x \\ \Rightarrow & |f|^{-1} \cdot |f| \cdot |f^{-1}| \sim |f|^{-1} \cdot x \\ \Rightarrow & |f^{-1}| \sim |f|^{-1}. \end{aligned}$$

□

**Definition 19 (Weak dihomotopy equivalence):** Two direct spaces  $(X, dX)$  and  $(Y, dY)$  are *weakly dihomotopical equivalent* if their paths spaces  $dX$  and  $dY$  are homotopical equivalent.

**Proposition 20:** Two directed spaces  $(X, dX)$  and  $(Y, dY)$  are weakly dihomotopical equivalent if and only if there exists two dimaps  $f : (X, dX) \rightarrow (Y, dY)$  and  $g : (Y, dY) \rightarrow (X, dX)$  such that  $f \circ g \sim \text{id}_Y$  and  $g \circ f \sim \text{id}_X$

*Proof.* (**Proposition 20**) Suppose  $\Gamma : dX \rightarrow dY$  and  $\Gamma' : dY \rightarrow dX$  form a homotopy equivalence. In particular, as constant paths are sent to constant paths,  $\Gamma$  and  $\Gamma'$  restricts to continuous maps  $\Gamma_0 : X \rightarrow Y$  and  $\Gamma'_0 : Y \rightarrow X$ . Take a path  $\gamma \in dX$ , then  $t \mapsto \Gamma_0(\gamma(t)) = \Gamma \circ \gamma \in dY$ , thus  $\Gamma_0$  is a dimap. Similarly,  $\Gamma'$  is also a dimap. We have  $\Gamma'\Gamma \sim \text{id}_{dX}$ , and this homotopy restricts to  $\Gamma'_0\Gamma_0 \sim \text{id}_Y$  by applying it to the constant paths. Similarly  $\Gamma_0\Gamma'_0 \sim \text{id}_X$ .

Conversely, suppose we have  $f : (X, dX) \rightarrow (Y, dY)$  and  $g : (Y, dY) \rightarrow (X, dX)$  such that  $f \circ g \sim \text{id}_Y$  and  $g \circ f \sim \text{id}_X$ . We let  $\Gamma = f \circ -$  and  $\Gamma' = g \circ -$  and conclude by  $\Gamma'\Gamma = g \circ f \circ - \sim \text{id}_X \circ - = \text{id}_{dX}$ , similarly for  $\Gamma\Gamma'$ . □

## References

- [Had20] Amar Hadzihasanovic. Diagrammatic sets and rewriting in weak higher categories, July 2020. arXiv:2007.14505 [math].