上 海 交 通 大 学 试 卷(<u>A</u>卷)

(2016 至 2017 学年 第 1 学期)

	课程名称	计算机系统基础(1)	成绩	
Prob	lem 1: HCL (7points)		
1.				
2.				
	lem 2: Y86 (
1. [1		[2]		
	3]	[4]		
	5]	[6]		
[7	7]	[8]		
2.				
		essor (18 points)		
1. [1	1]			
F-	21			
	2]			
[3	3]			
[4	4]			
[5	5]			
[6	5]			

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承诺人:	批阅人(流水阅 卷教师签名处)							
	心 教师 並 石 入)							

2. [1]

[2]

[3]

[4]

3.

Problem 4: Cache (32 points)

1. [1]

[2]

[3]

2.

3. [1]

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[7]

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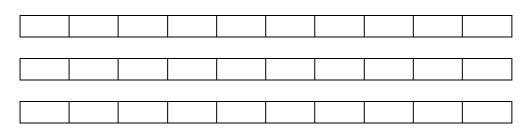
2)

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Problem 5: Memory Allocation (16 points)

1.

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3.

Problem 6: Optimization (17 points)

1.

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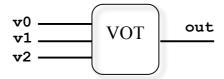
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Problem 1: HCL (7 points)

Please write down the HCL expressions for the following signals (HINT: you can refer to the **Section 4.2.2** in the CSAPP book). (3'+4'=7')

EXAMPLE: Show if the two input signals a and b are equal bool eq = (a&&b) || (!a && !b);

- 1. The HCL expression for a signal or, which equals to the OR of inputs a and b, and you should only use AND (&&) and NOT (!) operators.
- 2. The HCL expression for a three-way voter **VOT**. It means, if and only if there are equal or greater than 2 inputs are true (1), output will be true (1). Each input and output is one-bit wise. (Hints: You can use boolean expressions or case expressions.)



Problem 2: Y86 (10 points)

```
0x000:
                              .pos 0
0x000: 30f440010000
                     | init:
                             irmovl Stack, %esp
0x006: 30f540010000
                                ___[1]
0x00c: 8012000000
                             call
                                    Main
0x011: 00
                             halt
0x012: a05f
                    | Main: pushl %ebp
0x014: 2045
                             rrmovl %esp,
                                            %ebp
0x016: [2]
                             irmovl Array, %ecx
0x01c: 30f204000000
                              irmovl $4,
                                            %edx
0x022: 30f000000000
                              irmovl $0,
                                            %eax
0x028: 506100000000
                             mrmovl (%ecx),%esi
                     | Loop:
0x02e: 30f700000000 |
                              irmovl $0,
                                            %edi
0x034: 6167
                              subl
                                   %esi,
                                            %edi
0x036: ___[3]__
                              cmovg %edi,
                    1
                                           %esi
0 \times 038:6060
                              addl %esi,
                                           %eax
0x03a: 30f304000000 |
                              irmovl $4,
                                            %ebx
[4] : 6031
                     Ι
                              addl %ebx,
                                            %ecx
0x042: 30f3ffffffff |
                              irmovl $-1,
                                            %ebx
0x048: 6032
                              addl
                                    %ebx,
                                            %edx
0x04a: 7428000000
                     1
                                 [5]
0x04f: 2054
                              rrmovl %ebp, %esp
                     | End:
0x051: b05f
                     Τ
                             popl %ebp
0 \times 053:90
                              ret
```

		1		
[6]:		1	.align	4
[6]:	feffffff	Array:	.long	0xfffffffe
0x058:	0a000000	1	.long	0x0000000a
0x05c:	[7]	_ I	.long	0x0000beef
0x060:	fcffffff	1	.long	0xffffffc
		1		
0x140:		1		[8]
0x140:		Stack:		
0x140:		Stack:		

- 1. Please fill in the blanks within above Y86 binary and assembly code. (1'*8=8')
- 2. Please calculate the value of %eax after the program HALT. (2')

Problem 3: Processor (18 points)

Suppose we are using hardware structure of **PIPE-2016** which is modified from **PIPE** (Figure 4.52 in CSAPP book). Now, we want to add a new instruction: simply selective jump, **ssjxx**, to the original Y86 instruction set, using the following encoding:

	Byte	0		1		2	3	4	5
ssjxx		E	Fn	rA	F	Dest			

The **Fn** field is the same as that of jxx. For example, 0xE5 stands for ssjge. If the condition described by **Fn** is satisfied, we will jump to the address stored in rA, we call this case as **TAKEN**. Otherwise, we will jump to Dest, we call this case as **NOT_TAKEN**.

1. Please fill in the **generic** function of each stage for **ssjxx** on PIPE-2016 like **Figure 4.21**. (1'*6=6')

Field	ssjxx
Fetch	[1]
Decode	[2]
Execute	[3]
Memory	[4]
Write Back	[5]
PC update	[6]

2. Suppose ssjxx reuse the same forwarding circuit for ret and jxx and NOT_TAKEN prediction strategies. There will be some hazards due to this new instruction. Please list new detection conditions like Figure 4.64 and new control action like Figure 4.66. (3'+2'*3)

Condition	Trigger								
	[1]								
agi mianmadiation	Pipeline register								
ssj misprediction	F	D	E	М	W				
	[2]	[3]	[4]						

3. If we change format of **ssjxx**, which reads both two addresses from registers rather than one from a register and another from immediate number. Which design do you think is better? Why? (Hints: think about branch prediction in both designs.) (3')

Problem 4: Cache (32 points)

Jack has a **32-bit** machine with a **2-way** set associative cache. There are **8 sets**. Each block is **4 bytes**. The following table shows the content of the data cache at time T. **Bytex** is the byte value stored at offset **x**.

Set	Tag	Valid	Byte0	Byte1	Byte2	Byte3	Tag	Valid	Byte0	Byte1	Byte2	Byte3
0	0x5df	1	0x11	0x22	0xfe	0x43	0x5d2	1	0xca	0xdb	0xed	0x00
1	0x7cf	1	0xab	0xcd	0xef	0xff	0x34e	1	0xdf	0x11	0x22	0x33
2	0x233	0	0x23	0x32	0x23	0x33	0x34e	1	0xfd	0x44	0x55	0x66
3	0x435	1	0xde	0xad	0xbe	0xef	0x34e	1	0xdf	0x11	0x22	0x33
4		0						0				
5	0x701	1	0xff	0xff	0xcc	0xcc	0x435	1	0xad	0x18	0x24	0x19
6	0x881	1	0xde	0xed	0xbe	0xef	0x781	0	0x23	0x32	0xff	0xdd
7		0						0				

1. How would a **32-bit** physical memory address be split into tag/set-index /block-offset fields in this machine? (2'*3=6')

- 2. What is the size of this cache in bytes? (2')
- 3. Assume the cache line replacement policy is **LRU**. A short program will read memory in the following sequences starting from time T. Each access will read **one byte**. Please fill the following blanks and compute the miss rate. If there is a cache miss, enter '--' for 'Byte Returned'. (1'*12 + 2'*1 = 14')

Order	Address	Set	Hit/Miss	Byte Returned
1	0xbbe0	0	Hit	0x11
2	0x66a3	[1]	[2]	[3]
3	0xf039	[4]	[5]	[6]
4	0x69c6	[7]	[8]	[9]
5	0xbb41	[10]	[11]	[12]

Miss rate: ____[13]____

Jack buys a **NEW** machine with a **64-bit** physical memory address and tests the following program. The machine has a **4-way** set associative cache. There are **4 sets**. Each block is **16 bytes**. The cache line replacement policy is also **LRU**. The size of int value is **4 bytes**. NOTE that i, j and result are stored in registers. The cache is **empty** before each execution. Please only consider **data cache** access.

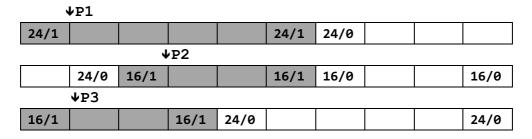
```
int x[4][16];
int y[4];
int z[4][16];
int fun(void) {
  int i, j, result = 0;
  for (j = 0; j < 16; j++)
     for (i = 0; i < 4; i++)
     result += x[i][j] * z[i][j] + y[i];
  return result;
}</pre>
```

Assume the address of x[0][0] is 0x000. The address of y[0] is 0x100. The address of z[0][0] is 0x110. Answer the following questions:

- 1. What is total number of memory accesses? (2')
- 2. Please calculate the miss rate. (2')
- 3. Jack executes **fun** (**void**) on another two different data caches:
 - 1) C1 is a 8-way set associative cache with 4 sets and 16 bytes cache line size;
 - 2) **C2** is a 4-way set associative cache with 8 sets and 16 bytes cache line size; Can **C1** reduce the miss rate? (1') If yes, please calculate the miss rate. If no, please explain the reason. (2')
 - Can **C2** reduce the miss rate? (1') If yes, please calculate the miss rate. If no, please explain the reason. (2')

Problem 5: Memory Allocation (16 points)

Now we organize the heap as a sequence of **contiguous** allocated and free blocks, as shown below. **Allocated** blocks are shaded, and **free** blocks are blank (each block represents 1 word = 4 bytes). **Headers** and **footers** are labeled with the number of bytes and allocated bit. The allocator maintains **double-word** alignment. You are given the execution sequence of memory allocation operations (malloc() or **free()**) from 1 to 6.



- 1.P4 = malloc(3)
- 2.P5 = malloc(14)
- 3.free(P3)
- 4.P6 = malloc(6)
- 5.free(P1)
- 6.P7 = malloc(8)

Please answer the questions below. Assume that **immediate coalescing** strategy and **splitting free blocks** are employed. (*NOTE*: **DON'T** need consider P1, P2 and P3 when calculating **internal** fragments)

- 1. Assume **best-fit** algorithm is used to find free blocks. Please draw the **final** status of memory and mark with block size in headers and footers after the operation sequence is executed (4'). Please also calculate the total bytes of the **internal fragments** (2').
- Assume first-fit algorithm is used to find free blocks. Please draw the final status of memory and mark with block size in headers and footers after the operation sequence is executed (4'). Please also calculate the total bytes of the internal fragments (2').
- 3. According to the final status of memory, we find that best-fit algorithm makes better utilization of memory than first-fit algorithm. However, first-fit algorithm is better in efficiency. Please explain why first-fit enjoys **better performance** based on above execution sequences. (You must use above execution sequences as example.) (4')

Problem 6: Optimization (17 points)

In US presidential election, each state votes separately. **state_result** records how many people vote for Trump and Clinton in a state. According to a state_result array **ra**, function **stat** computes the total votes among all states, and the winner and the gap between two candidates for each state.

```
typedef struct {
2.
    int trump;
    int clinton;
    int winner; // 0 for Trump, 1 for Clinton
4.
    int gap;
5.
   } state result;
6.
7.
8. typedef struct {
9.
    int length;
10. state result *data;
11. } rst;
12.
13. int get_length(rst *ra) { return ra->length; }
14. int get t(rst *ra, int i) {return (ra->data)[i].trump;}
15. int get c(rst *ra, int i) {return (ra->data)[i].clinton;}
16.
17. void stat(rst *ra, int *total) {
18.
     for (int i = 0; i < get_length(ra); i++)</pre>
19.
       *total = *total + get t(ra, i) + get c(ra, i);
20.
21.
     state result *states = ra->data;
22. for (int i = 0; i < get length(ra); i++)</pre>
23.
      if (states[i].trump > states[i].clinton) {
24.
         states[i].winner = 0;
25.
         states[i].gap = states[i].trump - states[i].clinton;
26.
       } else {
27.
        // assume Clinton wins if she gets equal or more
28.
        states[i].winner = 1;
29.
        states[i].gap = states[i].clinton - states[i].trump;
30.
       }
31. }
32.
33. int total trump(state result *r, int len) {
      if (len <= 0) return 0;
35.
      return r->trump + total trump(r + 1, len - 1);
36. }
```

Note: your optimizations cannot change the functionality of code above.

- 1. Please rewrite the loop in line 18-19 with a combination of at least 5 different optimizations you learned in class. Comment briefly on the optimization. (2'*5=10')
- 2. Please rewrite the loop body **in line 23-30** to reduce the branch prediction miss rate, with the fact that Trump won in about half of all US states. You cannot change the code other than line 23-30. (4')
- 3. For an array of length L, the recursion function **total_trump** will recur L times. Please rewrite the function body **in line 34-35** to reduce the depth to about L/2, with an optimization similar to loop unrolling. But you cannot use loop in your solution. NOTE that you can show how to invoke your optimized function if it helps simplify your solution. (3')