

**Problem 1: HCL (7 points)**

```
1. bool and = !(!a || !b);

2. Sol1: bool out = (v0 && v1) || (v1 && v2) || (v2 && v0);
   Sol2: bool out = [v0 && v1 && v2 : 1;
                     v0 && v1      : 1;
                     v1 && v2      : 1;
                     v2 && v0      : 1;
                     1             : 0;
                     ];
```

\*Or other solutions satisfied the truth table.

**Problem 2: Y86 (10 points)**

```
1. [1] irmovl Stack, %esp      [2] 30f154000000
   [3] 2676                    [4] 0x048
   [5] jne Loop                [6] 0x054
   [7] cefa0000                [8].pos 0x160

2. %eax = 0x0000fade (Sum of absolute value of four values)
```

**Problem 3 : Processor (18 points)**

1.

Field	ssjxx
Fetch	$icode:ifun \leftarrow M_1[PC]$ $rA:rB \leftarrow M_1[PC + 1]$ $valC \leftarrow M_4[PC + 2]$ $valP \leftarrow PC + 6$
Decode	$valA \leftarrow R[rA]$
Execute	$Cnd \leftarrow Cond(CC, ifun)$
Memory	--
Write Back	--
PC update	$PC \leftarrow Cnd? valA : valC$

```
2. [1] E_icode == ISSJXX && e_Cnd
   [2] -- [3] Bubble [4] Bubble
```

3. The origin design is better. The new design cannot take advantage of branch prediction. Because address in register needs to be read at DECODE period, it still need stall one cycle even though prediction is correct.

Problem 4: Cache (32 points)

- 1 [1] 11 [2] 3 [3] 2
- 2 64 bytes
- 3 [1] 0 [2] Miss [3] --  
 [4] 6 [5] Miss [6] --  
 [7] 1 [8] Hit [9] 0x22  
 [10] 0 [11] Miss [12] --  
 [13] 0.6
- 4 1)  $3 * 4 * 16 = 192$   
 2)  $(28 + 28 + 1) / 192 = 19/64$   
 3) Both C1 and C2 can reduce the miss rate, the miss rate is 11/64.

Problem 5: Memory Allocation (16 points)

1.

16/1			16/1	8/0	8/0	24/1			
	24/1	16/1			16/1	16/1			16/1
16/1			16/1	24/0					24/0

Internal: 46bytes

2.

16/1			16/1	8/0	8/0	16/1			16/1
8/0	8/0	16/1			16/1	16/1			16/1
16/0			16/0	24/1					24/1

Internal: 46bytes

3. Let's count the read needed for each operation:

First-fit:  $1+6+2+4+2+0=15$

Best-fit:  $6+6+2+5+2+6=27$

We can find that first-fit needs fewer reads to allocate the memory. Therefore, first-fit enjoys better performance.

**Problem 6: Optimization (17 points)**

```
1  int tmp1 = 0, tmp2 = 0, i;
    // reduce loop overhead
    int len = get_length(ra);
    // reduce function call
    state_result *sa = ra->results;
    for (i = 0; i < len - 1; i += 2) {
        // two way loop unrolling + reassociation
        tmp1 = tmp1 + (sa[i].trump + sa[i].clinton);
        // two way loop unrolling + two accumulators
        tmp2 = tmp2 + (sa[i+1].trump + sa[i+1].clinton);
    }
    for (; i < len; i++)
        tmp1 += sa[i].trump + sa[i].clinton;
    // reduce memory access
    *sum = tmp1 + tmp2;

2  int res = states[i].clinton > states[i].trump;
    int winner = (res > 0) ? 0 : 1;
    int high = (res > 0) ? states[i].clinton : states[i].trump;
    int low = (res > 0) ? states[i].trump : states[i].clinton;
    states[i].winner = winner;
    states[i].gap = high - low;

3  // optimized code
    int total_clinton(state_result *r, int len) {
        if (len <= 0) return 0;
        return r[0].clinton + r[1].clinton + total_clinton(r + 2, len - 2);
    }

    // invocation example
    extern int length;
    extern state_result *array;
    int res;
    if (length % 2 == 0)
        res = total_clinton(array, length);
    else
        res = total_clinton(array + 1, length - 1);
```