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Todo mundo adora o Chris

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Contest (1)

.bashrc   13 lines

```
b()(g++ -DLOC -O2 -std=c++20 -Wall -W -Wfatal-errors \
    -Wconversion -Wshadow -Wlogical-op \
    -Wfloat-equal -o $1 $1.cpp )
d()(b $@ -O0 -g -D_GLIBCXX_DEBUG -fsanitize=address,undefined)
run()( $@ && echo start >&2 && time ./$2)
hash-cpp()(sed -n $2','$3' p' $1 | sed '/^#w/d' | cpp -dD -P \
  -fpreprocessed | tr -d '[:space:]' | md5sum |cut -c-6)
# Other compilation flags:
# -Wformat=2 -Wshift-overflow=2 -Wcast-qual
# -Wcast-align -Wduplicated-cond
# -D_GLIBCXX_DEBUG_PEDANTIC -D_FORTIFY_SOURCE=2
# -fno-sanitize-recover -fstack-protector
# Print optimization info: -fopt-info-all
```

Mathematics (2)

2.1 Equations

$$ax^2 + bx + c = 0 \Rightarrow x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

The extremum is given by $x = -b/2a$.

$$\begin{aligned} ax + by &= e \\ cx + dy &= f \end{aligned} \Rightarrow \begin{aligned} x &= \frac{ed - bf}{ad - bc} \\ y &= \frac{af - ec}{ad - bc} \end{aligned}$$

In general, given an equation $Ax = b$, the solution to a variable x_i is given by

$$x_i = \frac{\det A_i'}{\det A}$$

where A_i' is A with the i'th column replaced by b.

2.2 Recurrences

If $a_n = c_1 a_{n-1} + \cdots + c_k a_{n-k}$, and $r_1, \ldots, r_k$ are distinct roots of $x^k - c_1 x^{k-1} - \cdots - c_k$, there are $d_1, \ldots, d_k$ s.t.

$$a_n = d_1 r_1^n + \cdots + d_k r_k^n.$$

Non-distinct roots r become polynomial factors, e.g. $a_n = (d_1 n + d_2) r^n$.

2.3 Trigonometry

$$\sin(v + w) = \sin v \cos w + \cos v \sin w$$
$$\cos(v + w) = \cos v \cos w - \sin v \sin w$$

$$\tan(v + w) = \frac{\tan v + \tan w}{1 - \tan v \tan w}$$

$$\sin v + \sin w = 2 \sin \frac{v + w}{2} \cos \frac{v - w}{2}$$

$$\cos v + \cos w = 2 \cos \frac{v + w}{2} \cos \frac{v - w}{2}$$

$$(V + W) \tan(v - w)/2 = (V - W) \tan(v + w)/2$$

where V, W are lengths of sides opposite angles v, w.

$$a \cos x + b \sin x = r \cos(x - \phi)$$
$$a \sin x + b \cos x = r \sin(x + \phi)$$

where $r = \sqrt{a^2 + b^2}, \phi = \operatorname{atan2}(b, a)$.

2.4 Geometry

2.4.1 Triangles

Side lengths: a, b, c

Semiperimeter: $p = \dfrac{a + b + c}{2}$

Area: $A = \sqrt{p(p - a)(p - b)(p - c)}$

Circumradius: $R = \dfrac{abc}{4A}$

Inradius: $r = \dfrac{A}{p}$

Length of median (divides triangle into two equal-area triangles): $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$

Length of bisector (divides angles in two):

$$s_a = \sqrt{bc\left[1 - \left(\frac{a}{b + c}\right)^2\right]}$$

Law of sines: $\dfrac{\sin \alpha}{a} = \dfrac{\sin \beta}{b} = \dfrac{\sin \gamma}{c} = \dfrac{1}{2R}$

Law of cosines: $a^2 = b^2 + c^2 - 2bc \cos \alpha$

Law of tangents: $\dfrac{a + b}{a - b} = \dfrac{\tan \dfrac{\alpha + \beta}{2}}{\tan \dfrac{\alpha - \beta}{2}}$

Pick's: A polygon on an integer grid strictly containing i lattice points and having b lattice points on the boundary has area $i + \frac{b}{2} - 1$. (Nothing similar in higher dimensions)
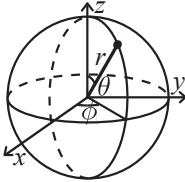
2.4.2 Quadrilaterals

With side lengths a, b, c, d, diagonals e, f, diagonals angle θ, area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2 f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is $180°$, $ef = ac + bd$, and $A = \sqrt{(p - a)(p - b)(p - c)(p - d)}$.

2.4.3 Spherical coordinates

$$\begin{aligned} x &= r \sin \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \sin \theta \sin \phi & \theta &= \operatorname{acos}(z / \sqrt{x^2 + y^2 + z^2}) \\ z &= r \cos \theta & \phi &= \operatorname{atan2}(y, x) \end{aligned}$$

2.5 Derivatives/Integrals

$$\frac{d}{dx} \arcsin x = \frac{1}{\sqrt{1 - x^2}} \qquad \frac{d}{dx} \arccos x = -\frac{1}{\sqrt{1 - x^2}}$$

$$\frac{d}{dx} \tan x = 1 + \tan^2 x \qquad \frac{d}{dx} \arctan x = \frac{1}{1 + x^2}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \qquad \int x \sin ax = \frac{\sin ax - ax \cos ax}{a^2}$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2} \operatorname{erf}(x) \qquad \int x e^{ax} dx = \frac{e^{ax}}{a^2}(ax - 1)$$

Integration by parts:

$$\int_a^b f(x) g(x) dx = [F(x) g(x)]_a^b - \int_a^b F(x) g'(x) dx$$

Green's theorem:
Let C be a positive, smooth, simple curve. D is a region bounded by C.

$$\oint_C (P dx + Q dy) = \int \int_D (\frac{\partial Q}{\partial x} - \frac{\partial P}{\partial y})$$

To calculate area, $\frac{\partial Q}{\partial x} - \frac{\partial P}{\partial y} = 1$, usually, picking $Q = \frac{1}{2}x$ and $P = -\frac{1}{2}y$ suffice.

Then we have

$$\frac{1}{2} \oint_C x dy - \frac{1}{2} \oint_C y dx$$

Line integral:
C given by $x = x(t), y = y(t), t \in [a, b]$, then

$$\oint_C f(x, y) ds = \int_a^b f(x(t), y(t)) ds$$

where, $ds = \sqrt{(\frac{dx}{dt})^2 + (\frac{dy}{dt})^2} dt$ or $\sqrt{(1 + (\frac{dy}{dx})^2} dx$

2.5.1 XOR sum

$$\bigoplus_{x=0}^{n-1} x = \{0, n-1, 1, n\}[n \bmod 4]$$

$$\bigoplus_{x=l}^{r-1} x = \bigoplus_{a=0}^{r-1} a \oplus \bigoplus_{b=0}^{l-1} b$$

2.6 Sums

$$c^a + c^{a+1} + \cdots + c^b = \frac{c^{b+1} - c^a}{c - 1}, c \neq 1$$

$$1 + 2 + 3 + \cdots + n = \frac{n(n+1)}{2}$$

$$1^2 + 2^2 + 3^2 + \cdots + n^2 = \frac{n(2n+1)(n+1)}{6}$$

$$1^3 + 2^3 + 3^3 + \cdots + n^3 = \frac{n^2(n+1)^2}{4}$$

$$1^4 + 2^4 + 3^4 + \cdots + n^4 = \frac{n(n+1)(2n+1)(3n^2 + 3n - 1)}{30}$$

2.7 Series

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \ldots, (-\infty < x < \infty)$$

$$\ln(1 + x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \ldots, (-1 < x \leq 1)$$

$$\sqrt{1 + x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \ldots, (-1 \leq x \leq 1)$$

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \ldots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \ldots, (-\infty < x < \infty)$$

2.8 Probability theory

Let X be a discrete random variable with probability $p_X(x)$ of assuming the value x. It will then have an expected value (mean) $\mu = \mathbb{E}(X) = \sum_x x p_X(x)$ and variance $\sigma^2 = V(X) = \mathbb{E}(X^2) - (\mathbb{E}(X))^2 = \sum_x (x - \mathbb{E}(X))^2 p_X(x)$ where σ is the standard deviation. If X is instead continuous it will have a probability density function $f_X(x)$ and the sums above will instead be integrals with $p_X(x)$ replaced by $f_X(x)$.

Expectation is linear:

$$\mathbb{E}(aX + bY) = a\mathbb{E}(X) + b\mathbb{E}(Y)$$

For independent X and Y,

$$V(aX + bY) = a^2 V(X) + b^2 V(Y).$$

2.8.1 Discrete distributions
Binomial distribution

The number of successes in n independent yes/no experiments, each which yields success with probability p is $\text{Bin}(n, p)$, $n = 1, 2, \ldots, 0 \leq p \leq 1$.

$$p(k) = \binom{n}{k} p^k (1 - p)^{n-k}$$

$$\mu = np, \sigma^2 = np(1 - p)$$

$\text{Bin}(n, p)$ is approximately $\text{Po}(np)$ for small p.
First success distribution

The number of trials needed to get the first success in independent yes/no experiments, each wich yields success with probability p is $\text{Fs}(p)$, $0 \leq p \leq 1$.

$$p(k) = p(1 - p)^{k-1}, k = 1, 2, \ldots$$

$$\mu = \frac{1}{p}, \sigma^2 = \frac{1 - p}{p^2}$$

Poisson distribution

The number of events occurring in a fixed period of time t if these events occur with a known average rate κ and independently of the time since the last event is $\text{Po}(\lambda)$, $\lambda = t\kappa$.

$$p(k) = e^{-\lambda} \frac{\lambda^k}{k!}, k = 0, 1, 2, \ldots$$

$$\mu = \lambda, \sigma^2 = \lambda$$

2.8.2 Continuous distributions
Uniform distribution

If the probability density function is constant between a and b and 0 elsewhere it is $\text{U}(a, b)$, $a < b$.

$$f(x) = \begin{cases} \frac{1}{b-a} & a < x < b \\ 0 & \text{otherwise} \end{cases}$$

$$\mu = \frac{a + b}{2}, \sigma^2 = \frac{(b-a)^2}{12}$$

Exponential distribution

The time between events in a Poisson process is $\text{Exp}(\lambda)$, $\lambda > 0$.

$$f(x) = \begin{cases} \lambda e^{-\lambda x} & x \geq 0 \\ 0 & x < 0 \end{cases}$$

$$\mu = \frac{1}{\lambda}, \sigma^2 = \frac{1}{\lambda^2}$$

Normal distribution

Most real random values with mean μ and variance σ^2 are well described by $\mathcal{N}(\mu, \sigma^2)$, $\sigma > 0$.

$$f(x) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$

If $X_1 \sim \mathcal{N}(\mu_1, \sigma_1^2)$ and $X_2 \sim \mathcal{N}(\mu_2, \sigma_2^2)$ then

$$aX_1 + bX_2 + c \sim \mathcal{N}(\mu_1 + \mu_2 + c, a^2\sigma_1^2 + b^2\sigma_2^2)$$

2.9 Markov chains

A *Markov chain* is a discrete random process with the property that the next state depends only on the current state. Let $X_1, X_2, \ldots$ be a sequence of random variables generated by the Markov process. Then there is a transition matrix $\mathbf{P} = (p_{ij})$, with $p_{ij} = \text{Pr}(X_n = i | X_{n-1} = j)$, and $\mathbf{p}^{(n)} = \mathbf{P}^n \mathbf{p}^{(0)}$ is the probability distribution for X_n (i.e., $p_i^{(n)} = \text{Pr}(X_n = i)$), where $\mathbf{p}^{(0)}$ is the initial distribution.

π is a stationary distribution if $\pi = \pi\mathbf{P}$. If the Markov chain is *irreducible* (it is possible to get to any state from any state), then $\pi_i = \frac{1}{\mathbb{E}(T_i)}$ where $\mathbb{E}(T_i)$ is the expected time between two visits in state i. π_j / π_i is the expected number of visits in state j between two visits in state i.

For a connected, undirected and non-bipartite graph, where the transition probability is uniform among all neighbors, π_i is proportional to node i's degree.

A Markov chain is *ergodic* if the asymptotic distribution is independent of the initial distribution. A finite Markov chain is ergodic iff it is irreducible and *aperiodic* (i.e., the gcd of cycle lengths is 1). $\lim_{k \to \infty} \mathbf{P}^k = \mathbf{1}\pi$.

A Markov chain is an absorbing chain if
1. there is at least one absorbing state and
2. it is possible to go from any state to at least one absorbing state in a finite number of steps.
A Markov chain is an A-chain if the states can be partitioned into two sets **A** and **G**, such that all states in **A** are absorbing ($p_{ii} = 1$), and all states in **G** leads to an absorbing state in **A**. The probability for absorption in state $i \in \mathbf{A}$, when the initial state is j, is $a_{ij} = p_{ij} + \sum_{k \in \mathbf{G}} a_{ik} p_{kj}$. The expected time until absorption, when the initial state is i, is
$t_i = 1 + \sum_{k \in \mathbf{G}} p_{ki} t_k$.

Data Structures (3)

order-statistic-tree.h
Description: A set (not multiset!) with support for finding the n'th element, and finding the index of an element.
Time: $\mathcal{O}(\log N)$
<bits/extc++.h>   acfa21, 17 lines
```cpp
template <typename K, typename V, typename Comp = std::less<K>>
using ordered_map = __gnu_pbds::tree<
  K, V, Comp,
  __gnu_pbds::rb_tree_tag,
  __gnu_pbds::tree_order_statistics_node_update
>;
template <typename K, typename Comp = std::less<K>>
using ordered_set = ordered_map<K, __gnu_pbds::null_type, Comp
  >;
void example() {
  ordered_set<int> t, t2; t.insert(8);
  auto it = t.insert(10).first;
  assert(it == t.lower_bound(9));
  assert(t.order_of_key(10) == 1); // num strictly smaller
  assert(t.order_of_key(11) == 2);
  assert(*t.find_by_order(0) == 8);
  t.join(t2); // assuming T < T2 or T > T2, merge t2 into t
}
```

dsu.h
Description: Disjoint-set data structure.
Time: $\mathcal{O}(\alpha(N))$
                                                        7d5db8, 14 lines
```cpp
struct UF {
  vector<int> e;
  UF(int n) : e(n, -1) {}
  bool same_set(int a, int b) { return find(a) == find(b); }
  int size(int x) { return -e[find(x)]; }
  int find(int x) { return e[x] < 0 ? x : e[x] = find(e[x]); }
  bool unite(int a, int b) {
    a = find(a), b = find(b);
    if (a == b) return 0;
    if (e[a] > e[b]) swap(a, b);
    e[a] += e[b]; e[b] = a;
    return 1;
  }
};
```

bipartite-dsu.h
Description: Disjoint-set data structure.
Time: $\mathcal{O}(\alpha(N))$
                                                        07774d, 31 lines
```cpp
struct DSU {
  vector<int> p, rk, color, bipartite;
```

```cpp
  DSU(int n) : p(n), rk(n), color(n), bipartite(n, 1) {
    iota(p.begin(), p.end(), 0);
  };
  int find(int u) {
    if (u == p[u]) return u;
    int v = find(p[u]);
    color[u] ^= color[p[u]];
    return p[u] = v;
  }
  int find_color(int u) {
    find(u);
    return color[u];
  }
  // check if it doesn't create an odd cycle
  bool can(int u, int v) {
    return find(u) != find(v) || color[u] != color[v];
  }
  void unite(int u, int v) {
    int pu = find(u), pv = find(v);
    if (pu == pv) {
      if (color[u] == color[v]) bipartite[pu] = false;
      return;
    }
    if (rk[pu] < rk[pv]) swap(pu, pv);
    if (color[u] == color[v]) color[pv] ^= 1;
    p[pv] = pu, rk[pu] += (rk[pu] == rk[pv]);
    if (not bipartite[pv]) bipartite[pu] = false;
  }
};
```

dsu-rollback.h
Description: Disjoint-set data structure with undo.
Usage: int t = uf.time(); ...; uf.rollback(t);
Time: $\mathcal{O}(\log(N))$
                                                        7ddf1d, 21 lines
```cpp
struct RollbackUF {
  vector<int> e; vector<pair<int,int>> st;
  RollbackUF(int n) : e(n, -1) {}
  int size(int x) { return -e[find(x)]; }
  int find(int x) { return e[x] < 0 ? x : find(e[x]); }
  int time() { return st.size(); }
  void rollback(int t) { // hash-1
    for (int i = time(); i --> t;)
      e[st[i].first] = st[i].second;
    st.resize(t);
  } // hash-1 = 30bb61
  bool unite(int a, int b) { // hash-2
    a = find(a), b = find(b);
    if (a == b) return false;
    if (e[a] > e[b]) swap(a, b);
    st.push_back({a, e[a]});
    st.push_back({b, e[b]});
    e[a] += e[b]; e[b] = a;
    return true;
  } // hash-2 = a7445e
};
```

monotonic-queue.h
Description: Supports pop and push queue-like, and add function adds a constant to all elements currently in the queue.
Time: $\mathcal{O}(1)$
                                                        8b6ad8, 18 lines
```cpp
template<typename T, typename Comp> struct monotonic_queue {
  int lo, hi; T S;
  deque<pair<T, T>> q;
  monotonic_queue() : lo(0), hi(0), S(0) {}
  void push(T val) {
    while (!q.empty() && Comp()(val, q.back().first + S))
      q.pop_back();
```

```cpp
    q.emplace_back(val - S, hi++);
  }
  void pop() {
    if (!q.empty() && q.front().second == lo++) q.pop_front();
  }
  void add(T val) { S += val; }
  T get_val() const { return q.front().first + S; }
  int size() const { return hi-lo; }
};
template<typename T> using min_monotonic_queue =
    monotonic_queue<T, std::less_equal<T>>;
template<typename T> using max_monotonic_queue =
    monotonic_queue<T, std::greater_equal<T>>;
```

point-context.h
Description: Examples of Segment Tree
                                                        70d417, 32 lines
```cpp
struct seg_node {
  int val, mi, ma;
  seg_node() : mi(INT_MAX), ma(INT_MIN), val(0) {}
  seg_node(int x) : mi(x), ma(x), val(x) {}
  void merge(const seg_node& l, const seg_node& r) {
    val = l.val + r.val;
    mi = min(l.mi, r.mi), ma = max(l.ma, r.ma);
  }
  void update(int x) {
    mi = ma = val = x;
  }
  bool acc_min(int& acc, int x) const {
    if (x >= mi) return true;
    if (acc > mi) acc = mi;
    return false;
  }
  bool acc_max(int& acc, int x) const {
    if (x <= ma) return true;
    if (acc < ma) acc = ma;
    return false;
  }
};
// 1 + min of [a, N) <= x
auto find_min_right = [&](segtree<seg_node>& sg, int a, int x)
    -> int {
  int acc = INT_MAX;
  return sg.find_first(a, N, &seg_node::acc_min, acc, x);
};
// max of [0, a) >= x
auto find_max_left = [&](segtree<seg_node>& sg, int a, int x)
    -> int {
  int acc = INT_MIN;
  return sg.find_last(0, a, &seg_node::acc_max, acc, x);
};
```

rec-lazy-segtree.h
Description: Segment Tree with Lazy update (half-open interval).
Time: $\mathcal{O}(\lg(N) * Q)$
                                                        22037e, 48 lines
```cpp
template<class T> struct segtree_range { // hash-1
  int N; vector<T> ts;
  segtree_range() {}
  segtree_range(int M) : segtree_range(vector<T>(M, T(0))) {}
  template<class Q> segtree_range(const vector<Q>& A) {
    const int N_ = int(A.size());
    N = (1 << __lg(2*N_-1)); ts.resize(2*N);
    for (int i = 0; i < N_; ++i) at(i) = T(A[i]);
    build();
  } // hash-1 = 37f0ae
  T& at(int a) { return ts[a + N]; }
  void build() { for (int a = N; --a; ) merge(a); }
  inline void push(int a) { ts[a].push(ts[2*a], ts[2*a+1]); }
```

```cpp
  inline void merge(int a) { ts[a].merge(ts[2*a], ts[2*a+1]); }
  T query(int v, int l, int r, int a, int b) { // hash-2
    if (l >= b || r <= a) return T();
    if (l >= a && r <= b) return ts[v];
    int m = (l + r)/2; push(v); T t;
    t.merge(query(2*v, l, m, a, b), query(2*v+1, m, r, a, b));
    return t;
  } // hash-2 = eb3198
  T query(int a, int b) { return query(1, 0, N, a, b); }
  template<class F, class... Args> void update(int v, int l,
      int r, int a, int b, F f, Args&&... args) { // hash-3
    if (l >= b || r <= a) return;
    if (l >= a && r <= b && (ts[v].*f)(args...)) return;
    int m = (l + r)/2; push(v);
    update(2*v, l, m, a, b, f, args...);
    update(2*v+1, m, r, a, b, f, args...);
    merge(v);
  } // hash-3 = 02f5df
  template<class F, class... Args>
  void update(int a, int b, F f, Args&&... args) {
    update(1, 0, N, a, b, f, args...);
  }
  template<class F, class... Args> int find_first(int v, int l,
      int r, int a, int b, F f, Args&&... args) { // hash-4
    if (l >= b || r <= a || !(ts[v].*f)(args...)) return -1;
    if (l + 1 == r) return l;
    int m = (l + r)/2; push(v);
    int cur = find_first(2*v, l, m, a, b, f, args...);
    if (cur == -1)
      cur = find_first(2*v+1, m, r, a, b, f, args...);
    return cur;
  } // hash-4 = 68aa30
  template<class F, class... Args>
  int find_first(int a, int b, F f, Args&&... args) {
    return find_first(1, 0, N, a, b, f, args...);
  }
};
```

lazy-context.h

Description: Examples of Segment Tree with Lazy update bd0d51, 173 lines

```cpp
template<typename T = int64_t> struct seg_node {
  T val, lz_add, lz_set;
  int sz; bool to_set;
  seg_node(T n = 0) : val(n), lz_add(0), lz_set(0), sz(1),
      to_set(0) {}
  void push(seg_node& l, seg_node& r) {
    if (to_set) {
      l.assign(lz_set), r.assign(lz_set);
      lz_set = 0; to_set = false;
    }
    if (lz_add != 0) {
      l.add(lz_add), r.add(lz_add), lz_add = 0;
    }
  }
  void merge(const seg_node& l, const seg_node& r) {
    sz = l.sz + r.sz; val = l.val + r.val;
  }
  bool add(T v) {   // update range a[i] <- a[i] + v
    val += v * sz; lz_add += v; return true;
  }
  bool assign(T v) {   // update range a[i] <- v
    val = v * sz; lz_add = 0;
    lz_set = v; to_set = true; return true;
  }
  T get_sum() const { return val; } // sum a[l, r)
};
// update range a[i] <- a[i] + b * (i - s) + c
// assuming b and c are non zero, be careful
```

```cpp
// get sum a[l, r)
template<typename T = int64_t> struct seg_node {
  T sum, lzB, lzC;
  int sz, idx;
  seg_node(int id = 0, T v = 0, int s = 0, T b = 0, T c = 0) :
    sum(v), lzB(b), lzC(c - s * b), idx(id), sz(1) {}
  void push(seg_node& l, seg_node& r) {
    l.add(lzB, lzC), r.add(lzB, lzC);
    lzB = lzC = 0;
  }
  void merge(const seg_node& l, const seg_node& r) {
    idx = min(l.idx, r.idx), sz = l.sz + r.sz;
    sum = l.sum + r.sum;
  }
  T sum_idx(T n) const { return n * (n + 1) / 2; }
  bool add(T b, T c) {
    sum += b * (sum_idx(idx + sz) - sum_idx(idx)) + sz * c;
    lzB += b, lzC += c; return true;
  }
  T get_sum() const { return sum; }
};

// update range a[i] <- b * a[i] + c
// get sum a[l, r)
struct seg_node {
  int sz; i64 sum, lzB, lzC;
  seg_node() : sz(1), sum(0), lzB(1), lzC(0) {}
  seg_node(i64 v) : sz(1), sum(v), lzB(1), lzC(0) {}
  void push(seg_node& l, seg_node& r) {
    l.add(lzB, lzC), r.add(lzB, lzC);
    lzB = 1, lzC = 0;
  }
  void merge(const seg_node& l, const seg_node& r) {
    sz = l.sz + r.sz, sum = l.sum + r.sum;
  }
  bool add(i64 b, i64 c) {
    sum = (b * sum + c * sz), lzB = (lzB * b);
    lzC = (lzC * b + c); return true;
  }
  i64 get_sum() const { return sum; }
};
// update range a[i] <- min(a[i], b);
// update range a[i] <- max(a[i], b);
// get val a[i]
struct seg_node {
  int mn, mx;
  int lz0, lz1;
  seg_node() : mn(INT_MAX), mx(INT_MIN), lz0(INT_MAX), lz1(
      INT_MIN) {}
  void push(seg_node& l, seg_node& r) {
    l.minimize(lz0), l.maximize(lz1);
    r.minimize(lz0), r.maximize(lz1);
    lz0 = INT_MAX, lz1 = INT_MIN;
  }
  void merge(const seg_node& l, const seg_node& r) {
    mn = min(l.mn, r.mn), mx = max(l.mx, r.mx);
  }
  bool minimize(int val) {
    mn = lz0 = min(lz0, val);
    mx = lz1 = min(lz0, lz1); return true;
  }
  bool maximize(int val) {
    mx = lz1 = max(lz1, val);
    mn = lz0 = max(lz0, lz1); return true;
  }
  pair<int, int> get() const { return {mx, mn}; }
};

template<typename T> struct lazy_t {
```

```cpp
  T a, b, c;
  lazy_t() : a(0), b(-INF), c(+INF) {}
  lazy_t(T a, T b, T c) : a(a), b(b), c(c) {}
  void add(T val) {
    a += val, b += val, c += val;
  }
  void upd_min(T val) {
    if (b > val) b = val;
    if (c > val) c = val;
  }
  void upd_max(T val) {
    if (b < val) b = val;
    if (c < val) c = val;
  }
};
template<typename T = int64_t> struct seg_node {
  T mi, mi2, ma, ma2, sum;
  T cnt_mi, cnt_ma, sz;
  lazy_t<T> lz;
  seg_node() : mi(INF), mi2(INF), ma(-INF), ma2(-INF), sum(0),
      cnt_mi(0), cnt_ma(0), sz(0), lz() {}
  seg_node(T n) : mi(n), mi2(INF), ma(n), ma2(-INF), sum(n),
      cnt_mi(1), cnt_ma(1), sz(1), lz() {}
  void push(seg_node& l, seg_node& r) {
    if (!l.can_apply(lz) || !r.can_apply(lz)) return;
    lz = lazy_t<T>();
  }
  bool can_apply(const lazy_t<T>& f) {
    if (!add(f.a) || !upd_max(f.b) || !upd_min(f.c)) return
        false;
    return true;
  }
  void merge(const seg_node& l, const seg_node& r) {
    mi = min(l.mi, r.mi);
    mi2 = min((l.mi == mi) ? l.mi2 : l.mi, (r.mi == mi) ? r.mi2
        : r.mi);
    cnt_mi = ((l.mi == mi) ? l.cnt_mi : 0) + ((r.mi == mi) ? r.
        cnt_mi : 0);
    ma = max(l.ma, r.ma);
    ma2 = max((l.ma == ma) ? l.ma2 : l.ma, (r.ma == ma) ? r.ma2
        : r.ma);
    cnt_ma = ((l.ma == ma) ? l.cnt_ma : 0) + ((r.ma == ma) ? r.
        cnt_ma : 0);
    sum = l.sum + r.sum;
    sz = l.sz + r.sz;
  }
  bool add(T v) { // a_i = a_i + v
    if (v) {
      mi += v;
      if (mi2 < INF) mi2 += v;
      ma += v;
      if (ma2 > -INF) ma2 += v;
      sum += sz * v;
      lz.add(v);
    }
    return true;
  }
  bool upd_max(T v) { // a_i = max(a_i, v)
    if (v > -INF) {
      if (v >= mi2) return false;
      else if (v > mi) {
        if (ma == mi) ma = v;
        if (ma2 == mi) ma2 = v;
        sum += cnt_mi * (v - mi);
        mi = v;
        lz.upd_max(v);
      }
    }
    return true;
```

```cpp
  }
  bool upd_min(T v) { // a_i = min(a_i, v)
    if (v < INF) {
      if (v <= ma2) return false;
      else if (v < ma) {
        if (ma == mi) mi = v;
        if (mi2 == ma) mi2 = v;
        sum -= cnt_ma * (ma - v);
        ma = v;
        lz.upd_min(v);
      }
    }
    return true;
  }
  T get_sum() const { return sum; } // sum a[l, r)
};
```

sparse-segtree.h

Description: Sparse Segment Tree with point update. Doesnt allocate storage for nodes with no data. Use BumpAllocator for better performance!

014f68, 39 lines

```cpp
const int SZ = 1<<19;
template<class T> struct node_t {
  T delta = 0; node_t<T>* c[2];
  node_t() { c[0] = c[1] = nullptr; }
  void upd(int pos, T v, int L = 0, int R = SZ-1) { // add v
    if (L == pos && R == pos) { delta += v; return; }
    int M = (L + R)>>1;
    if (pos <= M) {
      if (!c[0]) c[0] = new node_t();
      c[0]->upd(pos, v, L, M);
    } else {
      if (!c[1]) c[1] = new node_t();
      c[1]->upd(pos, v, M+1, R);
    }
    delta = 0;
    for (int i = 0; i < 2; ++i) if (c[i]) delta += c[i]->delta;
  }
  T query(int lx, int rx, int L = 0, int R = SZ-1) { // query
      sum of segment
    if (rx < L || R < lx) return 0;
    if (lx <= L && R <= rx) return delta;
    int M = (L + R)>>1; T res = 0;
    if (c[0]) res += c[0]->query(lx, rx, L, M);
    if (c[1]) res += c[1]->query(lx, rx, M+1, R);
    return res;
  }
  void upd(int pos, node_t *a, node_t *b, int L = 0, int R = SZ
      -1) {
    if (L != R) {
      int M = (L + R)>>1;
      if (pos <= M) {
        if (!c[0]) c[0] = new node_t();
        c[0]->upd(pos, a ? a->c[0] : nullptr, b ? b->c[0] :
            nullptr, L, M);
      } else {
        if (!c[1]) c[1] = new node_t();
        c[1]->upd(pos,a ? a->c[1] : nullptr, b ? b->c[1] :
            nullptr, M+1, R);
      }
    }
    delta = (a ? a->delta : 0)+(b ? b->delta : 0);
  }
};
```

segtree-2d.h

Description: 2D Segment Tree.
Time: $\mathcal{O}\left(N \log^2 N\right)$ of memory, $\mathcal{O}\left(\log^2 N\right)$ per query

"sparse_seg_tree.h"   09098e, 25 lines

```cpp
template<class T> struct Node {
  node_t<T> seg; Node* c[2];
  Node() { c[0] = c[1] = nullptr; }
  void upd(int x, int y, T v, int L = 0, int R = SZ-1) {//add v
    if (L == x && R == x) { seg.upd(y,v); return; }
    int M = (L+R)>>1;
    if (x <= M) {
      if (!c[0]) c[0] = new Node();
      c[0]->upd(x,y,v,L,M);
    } else {
      if (!c[1]) c[1] = new Node();
      c[1]->upd(x,y,v,M+1,R);
    }
    seg.upd(y,v); // only for addition
            // seg.upd(y,c[0]?&c[0]->seg:nullptr,c[1]?&c[1]->
                seg:nullptr);
  }
  T query(int x1, int x2, int y1, int y2, int L = 0, int R = SZ
      -1) { // query sum of rectangle
    if (x1 <= L && R <= x2) return seg.query(y1,y2);
    if (x2 < L || R < x1) return 0;
    int M = (L+R)>>1; T res = 0;
    if (c[0]) res += c[0]->query(x1, x2, y1, y2, L, M);
    if (c[1]) res += c[1]->query(x1, x2, y1, y2, M+1, R);
    return res;
  }
};
```

persistent-segtree.h

Description: Persistent implementation of a segment tree. This one compute the kth smallest element in a subarray $[a, b]$.

d277eb, 31 lines

```cpp
struct segtree_t {
  struct snapshot {
    int cnt, linkl, linkr;
    snapshot() : cnt(0), linkl(0), linkr(0) {}
    snapshot(int _cnt, int l, int r) : cnt(_cnt), linkl(l),
        linkr(r) {}
  };
  int id;
  vector<snapshot> tree;
  segtree_t() {}
  segtree_t(int n) : id(1), tree(20*n) {}
  int update(int v, int l, int r, int x) {
    if (x < l || x > r) return v;
    if (l == r) {
      tree[id] = snapshot(1, 0, 0);
      return id++;
    }
    int m = (l + r) >> 1;
    int lx = update(tree[v].linkl, l, m, x);
    int rx = update(tree[v].linkr, m+1, r, x);
    tree[id] = snapshot(tree[lx].cnt + tree[rx].cnt, lx, rx);
    return id++;
  }
  int query(int a, int b, int l, int r, int k) { // kth
    if (l == r) return l;
    int m = (l + r) >> 1;
    int cnt = tree[tree[b].linkl].cnt - tree[tree[a].linkl].cnt
        ;
    if (k <= cnt)
      return query(tree[a].linkl, tree[b].linkl, l, m, k);
    return query(tree[a].linkr, tree[b].linkr, m+1, r, k-cnt);
  }
};
```

merge-sort-tree.h

Description: Build segment tree where each node stores a sorted version of the underlying range.
Time: $\mathcal{O}\left(\log^2 N\right)$

9216c7, 36 lines

```cpp
struct merge_sort_tree {
  vector<int> v, ids;
  vector<vector<int>> tree;
  merge_sort_tree(vector<int> &v) : v(v), tree(4*(v.size()+1))
      {
    for(int i = 0; i < v.size(); ++i) ids.push_back(i);
    sort(ids.begin(), ids.end(), [&v](int i, int j) { return v[
        i] < v[j]; });
    build(1, 0, v.size()-1);
  }
  void build(int id, int left, int right) {
    if (left == right) tree[id].push_back(ids[left]);
    else {
      int mid = (left + right)>>1;
      build(id<<1, left, mid);
      build(id<<1|1, mid+1, right);
      tree[id] = vector<int>(right - left + 1);
      merge(tree[id<<1].begin(), tree[id<<1].end(),
        tree[id<<1|1].begin(), tree[id<<1|1].end(),
        tree[id].begin());
    }
  }
  // how many elements in this node have id in the range [a,b]
  int how_many(int id, int a, int b) {
    return (int)(upper_bound(tree[id].begin(), tree[id].end(),
        b)
      - lower_bound(tree[id].begin(), tree[id].end(), a));
  }
  int query(int id, int left, int right, int a, int b, int x) {
    if (left == right) return v[tree[id].back()];
    int mid = (left + right)>>1;
    int lcount = how_many(id<<1, a, b);
    if (lcount >= x) return query(id<<1, left, mid, a, b, x);
    else return query(id<<1|1, mid+1, right, a, b, x - lcount);
  }
  int kth(int a, int b, int k) {
    return query(1, 0, v.size()-1, a, b, k);
  }
};
```

rmq.h

Description: Range Minimum/Maximum Queries on an array. Returns $min(V[a], V[a + 1], ... V[b])$ in constant time. Returns a pair that holds the answer, first element is the value and the second is the index.
Usage: `rmq_t<pair<int, int>> rmq(values);`
```cpp
// values is a vector of pairs {val(i), index(i)}
rmq.query(inclusive, exclusive);
rmq_t<pair<int, int>, greater<pair<int, int>>> rmq(values)
//max query
```
Time: $\mathcal{O}\left(|V| \log |V| + Q\right)$

8c53c5, 19 lines

```cpp
template<typename T, typename Cmp=less<T>>
struct rmq_t : private Cmp {
  int N = 0;
  vector<vector<T>> table;
  const T& min(const T& a, const T& b) const { return Cmp::
      operator()(a, b) ? a : b; }
  rmq_t() {}
  rmq_t(const vector<T>& values) : N(int(values.size())), table
      (__lg(N) + 1) {
    table[0] = values;
    for (int a = 1; a < int(table.size()); ++a) {
      table[a].resize(N - (1 << a) + 1);
      for (int b = 0; b + (1 << a) <= N; ++b)
```

```
      table[a][b] = min(table[a-1][b], table[a-1][b + (1 << (
          a-1))]);
      }
    }
    T query(int a, int b) const {
      int lg = __lg(b - a);
      return min(table[lg][a], table[lg][b - (1 << lg) ]);
    }
};
```

fenwick-tree.h

Description: Computes partial sums a[0] + a[1] + ... + a[pos - 1], and updates single elements a[i], taking the difference between the old and new value.
Time: Both operations are $\mathcal{O}(\log N)$.

2ee6d4, 26 lines

```
template<typename T> struct FT {
  vector<T> s;
  FT(int n) : s(n) {}
  FT(const vector<T>& A) : s(A) {
    const int N = int(s.size());
    for (int a = 0; a < N; ++a)
      if ((a | (a + 1)) < N) s[a | (a + 1)] += s[a];
  }
  void update(int pos, T dif) { // a[pos] += dif
    for (; pos < (int)s.size(); pos |= pos + 1) s[pos] += dif;
  }
  T query(int pos) { // sum of values in [0, pos)
    T res = 0;
    for (; pos > 0; pos &= pos - 1) res += s[pos-1];
    return res;
  }
  // min pos st sum of [0, pos] >= sum. Returns n if no sum
  int lower_bound(T sum) { //is >= sum, or -1 if empty sum is.
    if (sum <= 0) return -1;
    int pos = 0;
    for (int pw = 1 << 25; pw; pw >>= 1)
      if (pos + pw <= (int)s.size() && s[pos + pw-1] < sum)
        pos += pw, sum -= s[pos-1];
    return pos;
  }
};
```

fenwick-tree-2d.h

Description: Computes sums a[i,j] for all i<I, j<J, and increases single elements a[i,j]. Requires that the elements to be updated are known in advance (call fakeUpdate() before init()).
Time: $\mathcal{O}(\log^2 N)$. (Use persistent segment trees for $\mathcal{O}(\log N)$.)

"fenwick-tree.h" aebbdc, 25 lines

```
template<typename T> struct FT2 {
  vector<vector<int>> ys; vector<FT<T>> ft;
  FT2(int limx) : ys(limx) {}
  void fakeUpdate(int x, int y) {
    for (; x < (int)ys.size(); x |= x + 1) ys[x].push_back(y);
  }
  void init() {
    for(auto &v : ys){
      sort(v.begin(), v.end());
      v.resize(unique(v.begin(),v.end()) - v.begin());
      ft.emplace_back(v.size());
    }
  }
  int ind(int x, int y) {
    return (int)(lower_bound(ys[x].begin(), ys[x].end(), y) -
        ys[x].begin()); }
  void update(int x, int y, T dif) {
    for (; x < ys.size(); x |= x + 1)
      ft[x].update(ind(x, y), dif);
  }
```

```
  T query(int x, int y) {
    T sum = 0;
    for (; x; x &= x - 1) sum += ft[x-1].query(ind(x-1, y));
    return sum;
  }
};
```

mo.h

Description: Answer interval or tree path queries by finding an approximate TSP through the queries, and moving from one query to the next by adding/removing points at the ends. If values are on tree edges, change step to add/remove the edge (a, c) and remove the initial add call (but keep in).
Time: $\mathcal{O}(N\sqrt{Q})$

5ef29d, 49 lines

```
void add(int ind, int end) { ... } // add a[ind] (end = 0 or 1)
void del(int ind, int end) { ... } // remove a[ind]
int calc() { ... } // compute current answer

vector<int> mo(vector<pair<int, int>> Q) { // hash-1
  int L = 0, R = 0, blk = 350; // ~N/sqrt(Q)
  vector<int> s(int(Q.size())), res = s;
#define K(x) pair<int, int>(x.first/blk, x.second ^ -(x.first/
    blk & 1))
  iota(s.begin(), s.end(), 0);
  sort(s.begin(), s.end(), [&](int s, int t){ return K(Q[s]) <
      K(Q[t]); });
  for (int qi : s) {
    auto q = Q[qi];
    while (L > q.first) add(--L, 0);
    while (R < q.second) add(R++, 1);
    while (L < q.first) del(L++, 0);
    while (R > q.second) del(--R, 1);
    res[qi] = calc();
  }
  return res;
} // hash-1 = d9247c

vector<int> moTree(vector<array<int, 2>> Q, vector<vector<int
    >>& ed, int root=0){ // hash-2
  int N = int(ed.size()), pos[2] = {}, blk = 350; // ~N/sqrt(Q)
  vector<int> s(int(Q.size())), res = s, I(N), L(N), R(N), in(N
      ), par(N);
  add(0, 0), in[0] = 1;
  auto dfs = [&](int x, int p, int dep, auto& f) -> void {
    par[x] = p;
    L[x] = N;
    if (dep) I[x] = N++;
    for (int y : ed[x]) if (y != p) f(y, x, !dep, f);
    if (!dep) I[x] = N++;
    R[x] = N;
  };
  dfs(root, -1, 0, dfs);
#define K(x) pii(I[x[0]] / blk, I[x[1]] ^ -(I[x[0]] / blk & 1))
  iota(s.begin(), s.end(), 0);
  sort(s.begin(), s.end(), [&](int s, int t){ return K(Q[s]) <
      K(Q[t]); });
  for (int qi : s) for (int end = 0; end < 2; ++end) {
    int &a = pos[end], b = Q[qi][end], i = 0;
#define step(c) { if (in[c]) { del(a, end); in[a] = 0; } \
    else { add(c, end); in[c] = 1; } a = c; }
    while (!(L[b] <= L[a] && R[a] <= R[b]))
      I[i++] = b, b = par[b];
    while (a != b) step(par[a]);
    while (i--) step(I[i]);
    if (end) res[qi] = calc();
  }
  return res;
} // hash-2 = bbf891
```

line-container.h

Description: Container where you can add lines of the form kx+m, and query maximum values at points x. Useful for dynamic programming ("convex hull trick").
Time: $\mathcal{O}(\log N)$

cd3f16, 27 lines

```
struct Line {
  mutable ll k, m, p;
  bool operator<(const Line& o) const { return k < o.k; }
  bool operator<(ll x) const { return p < x; }
};
struct LineContainer : multiset<Line, less<>> {
  static const ll inf = LLONG_MAX; //for doubles 1/.0
  ll div(ll a, ll b) { //for doubles a/b
    return a / b - ((a ^ b) < 0 && a % b); }
  bool isect(iterator x, iterator y) { // hash-1
    if (y == end()) { x->p = inf; return false; }
    if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
    else x->p = div(y->m - x->m, x->k - y->k);
    return x->p >= y->p;
  } // hash-1 = ea7809
  void add(ll k, ll m) { // hash-2
    auto z = insert({k, m, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() && isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin() && (--x)->p >= y->p)
      isect(x, erase(y));
  } // hash-2 = 08625f
  ll query(ll x) { // hash-3
    assert(!empty()); auto l = *lower_bound(x);
    return l.k * x + l.m;
  } // hash-3 = d21e2f
};
```

lichao.h

Description: Line Segments Li Chao Tree. Allows line add, segment add and point query.
Time: $\mathcal{O}(\log N)$ except for segment add $\mathcal{O}(\log^2 N)$

42896f, 62 lines

```
template<typename T, T L, T R>
struct lichao_t{
  static const T inf = numeric_limits<T>::max() / 2;
  bool first_best( T a, T b ){ return a < b; }
  T get_best( T a, T b ){   return first_best(a, b) ? a : b;   }
  struct line{
    T m, b;
    T operator()( T x ){ return m*x + b; }
  };
  struct node{
    line li;
    node *left, *right;
    node( line _li = {0, inf}): li(_li), left(nullptr), right(
        nullptr){}
    ~node(){
      delete left;
      delete right;
    }
  };
  node *root;
  lichao_t( line li = {0, inf} ): root ( new node(li) ) {}
  ~lichao_t(){ delete root; }
  T query( T x , node *cur , T l, T r){
    if(cur == nullptr) return inf;
    if(x < l || x > r) return inf;
    T mid = ( l + r ) >> 1;
    T ans =  cur->li(x);
    ans = get_best( ans , query(x, cur->left, l, mid) );
    ans = get_best( ans , query(x, cur->right, mid+1, r) );
    return ans;
  }
```

```cpp
  T query( T x ){ return query( x, root, L, R ); }
  void add( line li, node *&cur, T l, T r){
    if(cur == nullptr){
      cur = new node(li);
      return;
    }
    T mid = ( l + r ) >> 1;
    if( first_best( li(mid), cur->li(mid) ) )
      swap(li, cur->li);
    if( first_best( li(l), cur->li(l) ) )
      add(li, cur->left, l, mid);
    if( first_best( li(r), cur->li(r) ) )
      add(li, cur->right, mid + 1, r);
  }
  void add( T m, T b ){ add( {m, b}, root, L, R ); }
  void addSegment( line li, node *&cur, T l, T r, T lseg, T
      rseg){
    if(r < lseg || l > rseg) return;
    if(cur == nullptr) cur =  new node;
    if(lseg <= l && r <= rseg){
      add(li, cur, l, r);
      return;
    }
    T mid = ( l + r ) >> 1;
    if(l != r){
      addSegment(li, cur->left, l, mid, lseg, rseg);
      addSegment(li, cur->right, mid+1, r, lseg, rseg);
    }
  }
  void addSegment( T m, T b, T left, T right){
    addSegment( {m, b}, root, L, R, left, right);
  }
};
```

lichao-lazy.h

Description: Lazy Li Chao Tree. Allows line add, segment add, segment update and point query.

Time: $\mathcal{O}(\log N)$ except for segment add $\mathcal{O}(\log^2 N)$

```cpp
template<typename T, T L, T R>
struct lichao_lazy{
  static const T inf = numeric_limits<T>::max() / 2;
  bool first_best( T a, T b ){ return a < b; }
  T get_best( T a, T b ){  return first_best(a, b) ? a : b;  }
  struct line{ // hash-1
    T m, b;
    T operator()( T x ){ return m*x + b; }
    void apply(line other){
      m += other.m;
      b += other.b;
    }
  }; // hash-1 = 88f949
  struct node{ // hash-2
    line li, lazy;
    node *left, *right;
    node( line _li = {0, inf}): li(_li), lazy({0,0}), left(
        nullptr), right(nullptr){}
    void apply(line other){
      li.apply(other);
      lazy.apply(other);
    }
    ~node(){
      delete left;
      delete right;
    }
  }; // hash-2 = e6c99b
  node *root; // hash-3
  lichao_lazy( line li = {0, inf} ): root ( new node(li) ) {}
  ~lichao_lazy(){ delete root; }
```

```cpp
  void propagateLazy(node *&cur){
    if(cur == nullptr) return;
    if(cur->left == nullptr) cur->left = new node;
    if(cur->right == nullptr) cur->right = new node;
    cur->left->apply( cur-> lazy);
    cur->right->apply( cur-> lazy);
    cur->lazy = {0, 0};
  } // hash-3 = 4f0a91
  T query( T x , node *cur , T l, T r){ // hash-4
    if(x < l || x > r || l > r) return inf;
    if(cur == nullptr) return inf;
    T mid = ( l + r ) >> 1;
    if(l != r) propagateLazy(cur);
    T ans =  cur->li(x);
    ans = get_best( ans , query(x, cur->left, l, mid) );
    ans = get_best( ans , query(x, cur->right, mid+1, r) );
    return ans;
  } // hash-4 = f56802
  T query( T x ){ return query( x, root, L, R ); }
  void add( line li, node *&cur, T l, T r){ // hash-5
    if(cur == nullptr){
      cur = new node(li);
      return;
    }
    T mid = ( l + r ) >> 1;
    propagateLazy(cur);
    if( first_best( li(mid), cur->li(mid) ) )
      swap(li, cur->li);
    if( first_best( li(l), cur->li(l) ) )
      add(li, cur->left, l, mid);
    if( first_best( li(r), cur->li(r) ) )
      add(li, cur->right, mid + 1, r);
  } // hash-5 = 4191d1
  void add( T m, T b ){ add( {m, b}, root, L, R ); }
  void propagateLine(node *&cur, T l, T r){ // hash-6
    if(cur == nullptr) return;
    T mid = ( l + r ) >> 1;
    add(cur->li, cur->left, l, mid);
    add(cur->li, cur->right, mid+1, r);
    cur->li = {0, inf};
  } // hash-6 = 8d3255
  void addSegment( line li, node *&cur, T l, T r, T lseg, T
      rseg){ // hash-7
    if(r < lseg || l > rseg) return;
    if(cur == nullptr) cur =  new node;
    if(lseg <= l && r <= rseg){
      add(li, cur, l, r);
      return;
    }
    T mid = ( l + r ) >> 1;
    if(l != r){
      propagateLazy(cur);
      addSegment(li, cur->left, l, mid, lseg, rseg);
      addSegment(li, cur->right, mid+1, r, lseg, rseg);
    }
  } // hash-7 = 1a6dd3
  void addSegment( T m, T b, T left, T right){
    addSegment( {m, b}, root, L, R, left, right);
  }
  void updateSegment( line li, node *&cur, T l, T r, T lseg, T
      rseg){ // hash-8
    if(r < lseg || l > rseg) return;
    if(cur == nullptr) cur =  new node;
    if(lseg <= l && r <= rseg){
      cur->apply(li);
      return;
    }
    T mid = ( l + r ) >> 1;
    propagateLazy(cur);
```

```cpp
    propagateLine(cur, l, r);
    updateSegment(li, cur->left, l, mid, lseg, rseg);
    updateSegment(li, cur->right, mid+1, r, lseg, rseg);
  } // hash-8 = cce50c
  void updateSegment( T m, T b, T left, T right){
    updateSegment( {m, b}, root, L, R, left, right);
  }
};
```

lichao-range.h

Description: Lazy Li Chao Tree. Allows line add, segment add, segment update (only linear coeficient) and range query.

Time: $\mathcal{O}(\log N)$ except for segment add $\mathcal{O}(\log^2 N)$

```cpp
template<typename T, T L, T R>
struct lichao_range{
  static const T inf = numeric_limits<T>::max() / 2;
  static bool first_best( T a, T b ){ return a < b; }
  static T get_best( T a, T b ){  return first_best(a, b) ? a :
      b;  }
  struct line{
    T m, b;
    T operator()( T x ){ return m*x + b; }
    void apply(line other){
      m += other.m;
      b += other.b;
    }
  };
  struct node{
    line li, lazy;
    node *left, *right;
    T answer;
    node( line _li = {0, inf}): li(_li), lazy({0,0}), left(
        nullptr), right(nullptr), answer(inf){}
    void apply(T l, T r, line other){
      li.apply(other);
      lazy.apply(other);
      answer = get_best(inf, answer + other.b);
    }
    ~node(){
      delete left;
      delete right;
    }
  };
  node *root;
  lichao_range( line li = {0, inf} ): root ( new node(li) ) {}
  ~lichao_range(){ delete root; }
  void updateAnswer(node *&cur, T l, T r){
    if(cur == nullptr) return;
    cur->answer = inf;
    if(cur->left != nullptr) cur->answer = get_best(cur->answer
        , cur->left->answer);
    if(cur->right != nullptr) cur->answer = get_best(cur->
        answer, cur->right->answer);
    cur->answer = get_best(cur->answer, cur->li(l));
    cur->answer = get_best(cur->answer, cur->li(r));
  }
  void propagateLazy(node *&cur, T l, T r){
    if(cur == nullptr) return;
    if(cur->left == nullptr) cur->left = new node;
    if(cur->right == nullptr) cur->right = new node;
    T mid = ( l + r ) >> 1;
    cur->left->apply( l, mid, cur-> lazy);
    cur->right->apply( mid+1, r, cur-> lazy);
    cur->lazy = {0, 0};
  }
  T query( node *cur , T l, T r, T lseg, T rseg){
    if(r < lseg || l > rseg) return inf;
    if(cur == nullptr) return inf;
```

```cpp
    if(lseg <= l && r <= rseg) return cur->answer;
    T answer = get_best(cur->li(max(l, lseg)), cur->li(min(r,
        rseg)));
    if(l != r) propagateLazy(cur, l, r);
    T mid = ( l + r ) >> 1;
    answer = get_best(answer, query(cur->left, l, mid, lseg,
        rseg));
    answer = get_best(answer, query(cur->right, mid+1, r, lseg,
        rseg));
    updateAnswer(cur, l, r);
    return answer;
  }
  T query( T l, T r){ return query( root, L, R, l, r); }
  void add( line li, node *&cur, T l, T r){
    if(cur == nullptr){
      cur = new node(li);
      return;
    }
    T mid = ( l + r ) >> 1;
    propagateLazy(cur, l, r);
    if( first_best( li(mid), cur->li(mid) ) )
      swap(li, cur->li);
    if( first_best( li(l), cur->li(l) ) )
      add(li, cur->left, l, mid);
    if( first_best( li(r), cur->li(r) ) )
      add(li, cur->right, mid + 1, r);
    updateAnswer(cur, l, r);
  }
  void add( T m, T b ){ add( {m, b}, root, L, R ); }
  void propagateLine(node *&cur, T l, T r){
    if(cur == nullptr) return;
    T mid = ( l + r ) >> 1;
    add(cur->li, cur->left, l, mid);
    add(cur->li, cur->right, mid+1, r);
    cur->li = {0, inf};
  }
  void addSegment( line li, node *&cur, T l, T r, T lseg, T
      rseg){
    if(r < lseg || l > rseg) return;
    if(cur == nullptr) cur = new node;
    if(lseg <= l && r <= rseg){
      add(li, cur, l, r);
      return;
    }
    T mid = ( l + r ) >> 1;
    if(l != r){
      propagateLazy(cur, l, r);
      addSegment(li, cur->left, l, mid, lseg, rseg);
      addSegment(li, cur->right, mid+1, r, lseg, rseg);
    }
    updateAnswer(cur, l, r);
  }
  void addSegment( T m, T b, T left, T right){
    addSegment( {m, b}, root, L, R, left, right);
  }
  void updateSegment( T b, node *&cur, T l, T r, T lseg, T rseg
      ){
    if(r < lseg || l > rseg) return;
    if(cur == nullptr) cur = new node;
    if(lseg <= l && r <= rseg){
      cur->apply(l, r, {0, b});
      return;
    }
    T mid = ( l + r ) >> 1;
    propagateLazy(cur, l, r);
    propagateLine(cur, l, r);
    updateSegment(b, cur->left, l, mid, lseg, rseg);
    updateSegment(b, cur->right, mid+1, r, lseg, rseg);
    updateAnswer(cur, l, r);
  }
}
  void updateSegment( T b, T left, T right){
    updateSegment( b , root, L, R, left, right);
  }
};
```

matrix.h

Description: Basic operations on square matrices.
Usage: Matrix<int> A(N, vector<int>(N));
447637, 28 lines

```cpp
template <typename T> struct Matrix : vector<vector<T>> {
  using vector<vector<T>>::vector;
  using vector<vector<T>>::size;
  int h() const { return int(size()); }
  int w() const { return int((*this)[0].size()); }
  Matrix operator*(const Matrix& r) const {
    assert(w() == r.h()); Matrix res(h(), vector<T>(r.w()));
    for(int i = 0; i < h(); ++i) for(int j = 0; j < r.w(); ++j)
      for (int k = 0; k < w(); ++k)
        res[i][j] += (*this)[i][k] * r[k][j];
    return res;
  }
  friend auto operator*(const Matrix<T>& A,const vector<T>& b){
    int N = int(A.size()), M = int(A[0].size());
    vector<T> y(N);
    for (int i = 0; i < N; ++i)
      for (int j = 0; j < M; ++j) y[i] += A[i][j] * b[j];
    return y;
  }
  Matrix& operator*=(const Matrix& r){return *this= *this * r;}
  Matrix pow(ll n) const {
    assert(h() == w()); assert(n >= 0);
    Matrix x = *this, r(h(), vector<T>(w()));
    for (int i = 0; i < h(); ++i) r[i][i] = T(1);
    while (n) { if (n & 1) r *= x; x *= x; n >>= 1; }
    return r;
  }
};
```

submatrix.h

Description: Calculate submatrix sums quickly, given upper-left and lower-right corners (half-open).
Usage: SubMatrix<int> m(matrix);
m.sum(0, 0, 2, 2); // top left 4 elements
Time: $\mathcal{O}\left(N^2 + Q\right)$
cd3f87, 13 lines

```cpp
template<class T> struct SubMatrix {
  vector<vector<T>> p;
  SubMatrix(vector<vector<T>>& v) {
    int R = v.size(), C = v[0].size();
    p.assign(R+1, vector<T>(C+1));
    for (int r = 0; r < R; ++r)
      for (int c = 0; c < C; ++c)
        p[r+1][c+1] = v[r][c] + p[r][c+1] + p[r+1][c] - p[r][c
            ];
  }
  T sum(int u, int l, int d, int r) {
    return p[d][r] - p[d][l] - p[u][r] + p[u][l];
  }
};
```

wavelet.h

Description: Segment tree on values instead of indices.
Time: $\mathcal{O}\left(\log(n)\right)$
80ec5e, 130 lines

```cpp
struct wavelet_t { // hash-1
  struct BitVector { // space: 32N bits
    vector<int> _rank = {0};
    BitVector(vector<char> v = vector<char>()) {
      _rank.reserve(v.size() + 1);
      for (int d : v) _rank.push_back(_rank.back() + d);
    }
    int rank(bool f, int k) { return f ? _rank[k] : (k - _rank[
        k]); }
    int rank(bool f, int l, int r) { return rank(f, r) - rank(f
        , l); }
  }; // hash-1 = 637778
  /*
    struct BitVector {  // space: 1.5N bits
    vector<ull> v;
    vector<int> _rank;
    BitVector(vector<char> _v = vector<char>()) {
    int n = int(_v.size());
    v = vector<ull>((n + 63) / 64);
    _rank = vector<int>(v.size() + 1);
    for (int i = 0; i < n; i++)
    if (_v[i]) {
    v[i / 64] |= 1ULL << (i % 64);
    _rank[i / 64 + 1]++;
    }
    }
    for (int i = 0; i < int(v.size()); i++) {
    _rank[i+1] += _rank[i];
    }
    }
    int rank(int k) {
    int a = _rank[k / 64];
    if (k % 64) a += __builtin_popcountll(v[k / 64] << (64 - k
        % 64));
    return a;
    }
    int rank(bool f, int k) { return f ? rank(k) : k - rank(k)
        ; }
    int rank(bool f, int l, int r) { return rank(f, r) - rank(
        f, l); }
  };
  */
  int n, lg = 1; // hash-2
  vector<int> mid;
  vector<BitVector> data;
  wavelet_t(vector<int> v = vector<int>()) : n(int(v.size())) {
    int ma = 0;
    for (int x : v) ma = max(ma, x);
    while ((1 << lg) <= ma) lg++;
    mid = vector<int>(lg);
    data = vector<BitVector>(lg);
    for (int lv = lg - 1; lv >= 0; lv--) {
      vector<char> buf;
      vector<vector<int>> nx(2);
      for (int d : v) {
        bool f = (d & (1 << lv)) > 0;
        buf.push_back(f);
        nx[f].push_back(d);
      }
      mid[lv] = int(nx[0].size());
      data[lv] = BitVector(buf);
      v.clear();
      v.insert(v.end(), nx[0].begin(), nx[0].end());
      v.insert(v.end(), nx[1].begin(), nx[1].end());
    }
  } // hash-2 = 0339b1
  pair<int, int> succ(bool f, int a, int b, int lv) {
    int na = data[lv].rank(f, a) + (f ? mid[lv] : 0);
    int nb = data[lv].rank(f, b) + (f ? mid[lv] : 0);
    return {na, nb};
  }
  // count i, s.t. (a <= i < b) && (v[i] < u)
  int rank(int a, int b, int u) { // hash-3
```

```cpp
      if ((1 << lg) <= u) return b - a;
      int ans = 0;
      for (int lv = lg - 1; lv >= 0; lv--) {
        bool f = (u & (1 << lv)) > 0;
        if (f) ans += data[lv].rank(false, a, b);
        tie(a, b) = succ(f, a, b, lv);
      }
      return ans;
    } // hash-3 = b2983b
    // k-th(0-indexed!) number in v[a..b]
    int select(int a, int b, int k) { // hash-4
      int u = 0;
      for (int lv = lg - 1; lv >= 0; lv--) {
        int le = data[lv].rank(false, a, b);
        bool f = (le <= k);
        if (f) {
          u += (1 << lv);
          k -= le;
        }
        tie(a, b) = succ(f, a, b, lv);
      }
      return u;
    } // hash-4 = c98faa
    // k-th(0-indexed!) largest number in v[a..b]
    int large_select(int a, int b, int k) {
      return select(a, b, b - a - k - 1);
    }
    // count i s.t. (a <= i < b) && (x <= v[i] < y)
    int count(int a, int b, int x, int y) {
      return rank(a, b, y) - rank(a, b, x);
    }
    // max v[i] s.t. (a <= i < b) && (v[i] < x)
    int pre_count(int a, int b, int x) {
      int cnt = rank(a, b, x);
      return cnt == 0 ? -1 : select(a, b, cnt - 1);
    }
    // min v[i] s.t. (a <= i < b) && (x <= v[i])
    int nxt_count(int a, int b, int x) {
      int cnt = rank(a, b, x);
      return cnt == b - a ? -1 : select(a, b, cnt);
    }
};

struct CompressWavelet { // hash-5
  wavelet_t wt;
  vector<int> v, vidx;
  int zip(int x) {
    return int(lower_bound(vidx.begin(), vidx.end(), x) - vidx.
        begin());
  }
  CompressWavelet(vector<int> _v = vector<int>()) : v(_v), vidx
      (v) {
    sort(vidx.begin(), vidx.end());
    vidx.erase(unique(vidx.begin(), vidx.end()), vidx.end());
    for (auto& d : v) d = zip(d);
    wt = Wavelet(v);
  }
  int rank(int a, int b, int u) { return wt.rank(a, b, zip(u));
        }
  int select(int a, int b, int k) { return vidx[wt.select(a, b,
       k)]; }
  int largest(int a, int b, int k) { return wt.large_select(a,
      b, k); }
  int count(int a, int b, int mi, int ma) { return wt.count(a,
      b, mi, ma); }
  int find_max(int a, int b, int x) { return wt.pre_count(a, b,
       x); }
  int find_min(int a, int b, int x) { return wt.nxt_count(a, b,
       x); }
```

```cpp
}; // hash-5 = 2447db
```

range-color.h
Description: RangeColor structure, supports point queries and range updates, if C isn't int32_t change $freq$ to map
Time: $\mathcal{O}(\lg(L) * Q)$ 3d860e, 35 lines

```cpp
template<class T, class C> struct RangeColor {
    struct Node{
        T lo, hi; C color;
        bool operator<(const Node &n) const { return hi < n.hi; }
    };
    C minInf; set<Node> st; vector<T> freq;
    RangeColor(T first, T last, C maxColor, C iniColor = C(0)):
        minInf(first - T(1)), freq(maxColor + 1) {
        freq[iniColor] = last - first + T(1);
        st.insert({first, last, iniColor});
    }
    C query(T i) { //get color in position i
        return st.upper_bound({T(0), i - T(1), minInf})->color;
    }
    void upd(T a, T b, C x) { //set x in [a, b]
        auto p = st.upper_bound({T(0), a - T(1), minInf});
        assert(p != st.end());
        T lo = p->lo, hi = p->hi; C old = p->color;
        freq[old] -= (hi - lo + T(1)); p = st.erase(p);
        if (lo < a)
            freq[old] += (a-lo), st.insert({lo, a-T(1), old});
        if (b < hi)
            freq[old] += (hi-b), st.insert({b+T(1), hi, old});
        while ((p != st.end()) && (p->lo <= b)) {
            lo = p->lo, hi = p->hi; old = p->color;
            freq[old] -= (hi - lo + T(1));
            if (b < hi){
                freq[old] += (hi - b); st.erase(p);
                st.insert({b + T(1), hi, old});
                break;
            } else p = st.erase(p);
        }
        freq[x] += (b - a + T(1)); st.insert({a, b, x});
    }
    T countColor(C x){ return freq[x]; }
};
```

implicit-treap.h
Description: A short self-balancing tree. It acts as a sequential container with log-time splits/joins, and is easy to augment with additional data.
Time: $\mathcal{O}(\log N)$ 8c89db, 76 lines

```cpp
struct node {
    int val, p, sz; bool rev;
    array<node*, 2> c{nullptr, nullptr};
    node(int k) : val(k), p(rng()), sz(0), rev(false) {}
    ~node() {
        delete c[0];
        delete c[1];
    }
};
inline int sz(node *t) {
    return (!t ? 0 : t->sz);
}
inline void push(node *t) {
    if (!t) return;
    if (t->rev) {
        swap(t->c[0], t->c[1]);
        if (t->c[0]) t->c[0]->rev ^= t->rev;
        if (t->c[1]) t->c[1]->rev ^= t->rev;
        t->rev = 0;
    }
```

```cpp
}
inline void pull(node *t) {
    if (!t) return;
    push(t); push(t->c[0]); push(t->c[1]);
    t->sz = sz(t->c[0]) + sz(t->c[1]) + 1;
}
inline void split(node *t, node *&a, node *&b, int k) { //k on
    left
    push(t);
    if (!t) a = b = nullptr;
    else if (k <= sz(t->c[0])) {
        split(t->c[0], a, t->c[0], k);
        b = t;
    } else {
        split(t->c[1], t->c[1], b, k-1-sz(t->c[0]));
        a = t;
    }
    pull(t);
}
inline void merge(node *&t, node *a, node *b) {
    push(a); push(b);
    if (!a) t = b;
    else if (!b) t = a;
    else if (a->p <= b->p) {
        merge(a->c[1], a->c[1], b);
        t = a;
    } else {
        merge(b->c[0], a, b->c[0]);
        t = b;
    }
    pull(t);
}
inline void add(node *&t, node *a, int k) {
    push(t);
    if (!t) t = a;
    else if (a->p >= t->p) {
        split(t, a->c[0], a->c[1], k);
        t = a;
    } else if (sz(t->c[0]) >= k) add(t->c[0], a, k);
    else add(t->c[1], a, k-1-sz(t->c[0]));
    pull(t);
}
void del(node *&t, int k) {
    push(t);
    if (!t) return;
    if (sz(t->c[0]) == k) merge(t, t->c[0], t->c[1]);
    else if (sz(t->c[0]) > k) del(t->c[0], k);
    else del(t->c[1], k);
    pull(t);
}
inline void dump_treap(node *t) {
    if (!t) return;
    push(t);
    dump_treap(t->c[0]);
    cerr << t->val << ' ';
    dump_treap(t->c[1]);
}
```

Numerical (4)

polynomial.h
84593c, 17 lines

```cpp
struct Poly {
    vector<double> a;
    double operator()(double x) const {
        double val = 0;
        for(int i = a.size(); i--;) (val *= x) += a[i];
        return val;
    }
```

```
    }
    void diff() {
        for(int i = 1; i < a.size(); ++i) a[i-1] = i*a[i];
        a.pop_back();
    }
    void divroot(double x0) {
        double b = a.back(), c; a.back() = 0;
        for(int i = a.size()-1; i--;) c = a[i],a[i]=a[i+1]*x0+b, b=
            c;
        a.pop_back();
    }
};
```

poly-roots.h
Description: Finds the real roots to a polynomial.
Usage: `poly_roots({{2,-3,1}},-1e9,1e9) // solve x^2-3x+2 = 0`
Time: $\mathcal{O}\left(n^2 \log(1/\epsilon)\right)$
"polynomial.h"   49396a, 20 lines

```
vector<double> poly_roots(Poly p, double xmin, double xmax) {
    if ((p.a).size() == 2) { return {-p.a[0]/p.a[1]}; }
    vector<double> ret;
    Poly der = p; der.diff();
    auto dr = poly_roots(der, xmin, xmax);
    dr.push_back(xmin-1); dr.push_back(xmax+1);
    sort(dr.begin(), dr.end());
    for(int i = 0; i < dr.size()-1; ++i) {
        double l = dr[i], h = dr[i+1]; bool sign = p(l) > 0;
        if (sign^(p(h) > 0)) {
            for(int it = 0; it < 60; ++it) { // while (h - l > 1e-8)
                double m = (l + h) / 2, f = p(m);
                if ((f <= 0) ^ sign) l = m;
                else h = m;
            }
            ret.push_back((l + h) / 2);
        }
    }
    return ret;
}
```

poly-interpolate.h
Description: Given n points (x[i], y[i]), computes an n-1-degree polynomial p that passes through them: $p(x) = a[0]*x^0 + ... + a[n-1]*x^{n-1}$. For numerical precision, pick $x[k] = c*\cos(k/(n-1)*\pi), k = 0\ldots n-1$.
Time: $\mathcal{O}\left(n^2\right)$
                                                              97a266, 12 lines

```
typedef vector<double> vd;
vd interpolate(vd x, vd y, int n) {
    vd res(n), temp(n);
    for(int k = 0; k < n-1; ++k) for(int i = k+1; i < n; ++i)
        y[i] = (y[i] - y[k]) / (x[i] - x[k]);
    double last = 0; temp[0] = 1;
    for(int k = 0; k < n; ++k) for(int i = 0; i < n; ++i) {
        res[i] += y[k] * temp[i]; swap(last, temp[i]);
        temp[i] -= last * x[k];
    }
    return res;
}
```

lagrange.h
Description: Lagrange interpolation over a finite field and some combo stuff
Time: $\mathcal{O}(N)$
"../number-theory/modular-arithmetic.h", "../number-theory/preparator.h"   ad3879, 13 lines

```
template<typename T> struct interpolator_t {
    vector<T> S;
    interpolator_t(int N): S(N) {}
    T interpolate(const vector<T>& y, T x) {
        int N = int(y.size()); int sgn = (N & 1 ? 1 : -1);
```

```
    T res = 0, P = 1; S[N - 1] = 1;
    for (int i = N-1; i > 0; --i) S[i-1] = S[i] * (x-i);
    for (int i = 0; i < N; ++i, sgn *= -1, P *= (x - i + 1)) {
        res += y[i] * sgn * P * S[i] * ifact[i] * ifact[N-1-i];
    }
    return res;
    }
};
```

berlekamp-massey.h
Description: Recovers any n-order linear recurrence relation from the first $2n$ terms of the recurrence. Useful for guessing linear recurrences after brute-forcing the first terms. Should work on any field, but numerical stability for floats is not guaranteed. Output will have size $\leq n$.
Usage: `BerlekampMassey({0, 1, 1, 3, 5, 11}) // {1, 2}`
Time: $\mathcal{O}\left(N^2\right)$
"../number-theory/modular-arithmetic.h"   66d78a, 17 lines

```
template <typename num>
vector<num> BerlekampMassey(const vector<num>& s) {
    int n = int(s.size()), L = 0, m = 0; num b = 1;
    vector<num> C(n), B(n), T; C[0] = B[0] = 1;
    for(int i = 0; i < n; i++) { ++m;
        num d = s[i];
        for (int j = 1; j <= L; j++) d += C[j] * s[i - j];
        if (d == 0) continue;
        T = C; num coef = d / b;
        for (int j = m; j < n; j++) C[j] -= coef * B[j - m];
        if (2 * L > i) continue;
        L = i + 1 - L; B = T; b = d; m = 0;
    }
    C.resize(L + 1); C.erase(C.begin());
    for (auto& x : C) x = -x;
    return C;
}
```

linear-recurrence.h
Description: Bostan-Mori algorithm. Generates the k'th term of an n-order linear recurrence $S[i] = \sum_j S[i - j - 1]tr[j]$, given $S[0\ldots n-1]$ and $tr[0\ldots n-1]$. Faster than matrix multiplication. Useful together with Berlekamp–Massey.
Usage: `linear_rec({0, 1}, {1, 1}, k) // k'th Fibonacci number`
Time: $\mathcal{O}(n \log n \log k)$
"../number-theory/modular-arithmetic.h"   aa7314, 16 lines

```
template<typename T>
T linear_rec(const vector<T>& S, const vector<T>& tr, ll K) {
    const int N = int(tr.size());
    vector<T> qs(N + 1); qs[0] = 1;
    for (int i = 0; i < N; ++i) qs[i + 1] = -tr[i];
    auto fs = fft.convolve(S, qs); fs.resize(N);
    for (; K; K /= 2) {
        auto qneg = qs;
        for (int i = 1; i <= N; i += 2) qneg[i] = -qneg[i];
        fs = fft.convolve(fs, qneg), qs = fft.convolve(qs, qneg);
        for (int i = 0; i < N; ++i)
            fs[i] = fs[2 * i + (K & 1)], qs[i] = qs[2 * i];
        qs[N] = qs[2*N]; fs.resize(N), qs.resize(N+1);
    }
    return fs[0];
}
```

integrate.h
Description: Simple integration of a function over an interval using Simpson's rule. The error should be proportional to h^4, although in practice you will want to verify that the result is stable to desired precision when epsilon changes.
                                                              7bb98e, 7 lines

```
template<class F>
double quad(double a, double b, F& f, const int n = 1000) {
```

```
    double h = (b - a) / 2 / n, v = f(a) + f(b);
    for(int i = 1; i < n*2; ++i)
        v += f(a + i*h) * (i&1 ? 4 : 2);
    return v * h / 3;
}
```

integrate-adaptive.h
Description: Fast integration using an adaptive Simpson's rule.
Usage: `double sphereVolume = quad(-1, 1, [](double x) {`
`return quad(-1, 1, [&](double y) {`
`return quad(-1, 1, [&](double z) {`
`return x*x + y*y + z*z < 1; });});});`
                                                              cfcad2, 13 lines

```
typedef double d;
#define S(a,b) (f(a) + 4*f((a+b) / 2) + f(b)) * (b-a) / 6
template <class F>
d rec(F& f, d a, d b, d eps, d S) {
    d c = (a + b) / 2, S1 = S(a, c), S2 = S(c, b), T = S1 + S2;
    if (abs(T - S) <= 15 * eps || b - a < 1e-10)
        return T + (T - S) / 15;
    return rec(f, a, c, eps / 2, S1) + rec(f, c, b, eps / 2, S2);
}
template<class F>
d quad(d a, d b, F f, d eps = 1e-8) {
    return rec(f, a, b, eps, S(a, b));
}
```

gaussian-elimination.h
Time: $\mathcal{O}(\min(N, M)NM)$
"../data-structures/matrix.h"   a5570d, 61 lines

```
template<typename T> struct gaussian_elimination {
    int N, M; Matrix<T> A, E;
    vector<int> pivot; int rank, nullity, sgn;
    gaussian_elimination(const Matrix<T>& A_) : A(A_) { // hash-1
        N = A.size(), M = A[0].size(), E=Matrix<T>(N, vector<T>(N))
            ;
        for (int i = 0; i < N; ++i) E[i][i] = 1;
        rank = 0, nullity = M, sgn = 0; pivot.assign(M, -1);
        for (int col = 0, row = 0; col < M && row < N; ++col) {
            int sel = -1;
            for (int i = row; i < N; ++i) if (A[i][col] != 0) {
                sel = i; break;
            }
            if (sel == -1) continue;
            if (sel != row) {
                sgn += 1;
                swap(A[sel], A[row]); swap(E[sel], E[row]);
            }
            for (int i = 0; i < N; ++i) {
                if (i == row) continue;
                T c = A[i][col] / A[row][col];
                for (int j = col; j < M; ++j)
                    A[i][j] -= c*A[row][j];
                for (int j = 0; j < N; ++j)
                    E[i][j] -= c*E[row][j];
            }
            pivot[col] = row++; ++rank, --nullity;
        }
    } // hash-1 = d52725
    pair<bool, vector<T>> solve(vector<T> b, bool reduced = false
        ) const { // hash-2
        if (reduced == false) b = E * b;
        vector<T> x(M);
        for (int j = 0; j < M; ++j) {
            if (pivot[j] == -1) continue;
            x[j] = b[pivot[j]] / A[pivot[j]][j];
            b[pivot[j]] = 0;
        }
        for (int i = 0; i < N; ++i)
```

```cpp
    if (b[i] != 0) return {false, x};
    return {true, x};
  } // hash-2 = 479018
  vector<vector<T>> kernel_basis() const { // hash-3
    vector<vector<T>> basis; vector<T> e(M);
    for (int j = 0; j < M; ++j) {
      if (pivot[j] != -1) continue;
      e[j] = 1; auto y = solve(A * e, true).second;
      e[j] = 0, y[j] = -1; basis.push_back(y);
    }
    return basis;
  } // hash-3 = f715d3
  Matrix<T> inverse() const { // hash-4
    assert(N == M); assert(rank == N);
    Matrix<T> res(N, vector<T>(N));
    vector<T> e(N);
    for (int i = 0; i < N; ++i) {
      e[i] = 1; auto x = solve(e).second;
      for (int j = 0; j < N; ++j) res[j][i] = x[j];
      e[i] = 0;
    }
    return res;
  } // hash-4 = 422806
};
```

linear-solver-z2.h

Description: Solves $Ax = b$ over $\mathbb{F}_2$. If there are multiple solutions, one is returned arbitrarily. Returns true, or false if no solutions. Last column of a is b. c is the rank.
Time: $\mathcal{O}\left(n^2 m\right)$

7a24e1, 24 lines

```cpp
typedef bitset<2010> bs;
bool gauss(vector<bs> a, bs& ans, int n) {
  int m = int(a.size()), c = 0;
  bs pos; pos.set();
  for (int j = n-1, i; j >= 0; --j) {
    for (i = c; i < m; ++i)
      if (a[i][j]) break;
    if (i == m) continue;
    swap(a[c], a[i]);
    i = c++; pos[j] = 0;
    for (int k = 0; k < m; ++k)
      if (a[k][j] && k != i) a[k] ^= a[i];
  } ans = pos;
  for(int i = 0; i < m; ++i) {
    int ac = 0;
    for (int j = 0; j < n ; ++j) {
      if (!a[i][j]) continue;
      if (!pos[j]) pos[j] = 1, ans[j] = ac^a[i][n];
      ac ^= ans[j];
    }
    if (ac != a[i][n]) return false;
  }
  return true;
}
```

char-poly.h

Description: Calculates the characteristic polynomial of a matrix.
$\sum_{k=0}^{n} p(k)(-1)^{n-k}$
Time: $\mathcal{O}\left(N^3\right)$ and div-free is $\mathcal{O}\left(N^4\right)$

30bd65, 55 lines

```cpp
// det(x I + a)
template<class T> vector<T> char_poly(const vector<vector<T>>&
    a) {
  const int N = int(a.size()); auto b = a;
  for (int j = 0; j < N - 2; ++j) {
    for (int i = j + 1; i < N; ++i) {
      if (b[i][j]) {
        swap(b[j + 1], b[i]);
```

```cpp
        for (int k = 0; k < N; ++k) swap(b[k][j + 1], b[k][i]);
        break;
      }
    }
    if (b[j + 1][j]) {
      const T r = 1 / b[j + 1][j];
      for (int i = j + 2; i < N; ++i) {
        const T s = r * b[i][j];
        for (int q = j; q < N; ++q) b[i][q] -= s * b[j + 1][q];
        for (int p = 0; p < N; ++p) b[p][j + 1] += s * b[p][i];
      }
    }
  }
  // fss[i] := det(x I_i + b[0..i][0..i])
  vector<vector<T>> fss(N + 1);
  fss[0] = {1};
  for (int i = 0; i < N; ++i) {
    fss[i + 1].assign(i + 2, 0);
    for (int k = 0; k <= i; ++k) fss[i + 1][k + 1] = fss[i][k];
    for (int k = 0; k <= i; ++k) fss[i + 1][k] += b[i][i] * fss
        [i][k];
    T q = 1;
    for (int j = i - 1; j >= 0; --j) {
      q *= -b[j + 1][j];
      const T s = q * b[j][i];
      for (int k = 0; k <= j; ++k) fss[i + 1][k] += s * fss[j][
          k];
    }
  }
  return fss[N];
}
// det(x I + a), division free
template<class T> vector<T> char_poly_div_free(const vector<
    vector<T>>& a) {
  const int N = int(a.size());
  vector<T> ps(N + 1, 0);
  ps[N] = 1;
  for (int h = N - 1; h >= 0; --h) {
    vector<vector<T>> sub(N, vector<T>(h + 1, 0));
    for (int i = N; i >= 1; --i)
      sub[i - 1][h] += ps[i];
    for (int i = N - 1; i >= 1; --i) for (int u = 0; u <= h; ++
        u) {
      for (int v = 0; v < h; ++v)
        sub[i - 1][v] -= sub[i][u] * a[u][v];
    }
    for (int i = N - 1; i >= 1; --i) for (int u = 0; u <= h; ++
        u) {
      ps[i] += sub[i][u] * a[u][h];
    }
  }
  return ps;
}
```

simplex.h

Description: Solves a general linear maximization problem: maximize $c^T x$ subject to $Ax \leq b$, $x \geq 0$.
Time: $\mathcal{O}\left(NM * \#pivots\right)$, where a pivot may be e.g. an edge relaxation. $\mathcal{O}\left(2^n\right)$ in the general case. WARNING- segfaults on empty (size 0) max cx st Ax<=b, x>=0 do 2 phases; 1st check feasibility; 2nd check boundedness and ans

83fec6, 39 lines

```cpp
using dbl = double; using vd = vector<dbl>;
vd simplex(vector<vd> A, vd b, vd c) { const dbl E = 1e-9;
  int n = A.size(), m = A[0].size() + 1, r = n, s = m-1;
  auto D = vector<vd>(n+2, vd(m+1));
  vector<int> ix = vector<int>(n + m);
  for (int i = 0; i < n + m; ++i) ix[i] = i;
  for (int i = 0; i < n; ++i) {
```

```cpp
    for (int j = 0; j < m-1; ++j) D[i][j] = -A[i][j];
    D[i][m - 1] = 1; D[i][m] = b[i];
    if (D[r][m] > D[i][m]) r = i;
  }
  for (int j = 0; j < m-1; ++j) D[n][j] = c[j];
  D[n + 1][m - 1] = -1; int z = 0;
  for (dbl d;;) {
    if (r < n) { swap(ix[s], ix[r + m]);
      D[r][s] = 1.0/D[r][s];
      for (int j = 0; j <= m; ++j) if (j != s)
        D[r][j] *= -D[r][s];
      for (int i = 0; i <= n+1; ++i) if (i != r) {
        for (int j = 0; j <= m; ++j)
          if (j != s) D[i][j] += D[r][j] * D[i][s];
        D[i][s] *= D[r][s];
      }
    }
    r = -1; s = -1;
    for (int j = 0; j < m; ++j) if (s < 0 || ix[s] > ix[j])
      if (D[n+1][j]>E || D[n+1][j]>-E && D[n][j]>E) s = j;
    if (s < 0) break;
    for (int i = 0; i < n; ++i) if (D[i][s] < -E) {
      if (r<0 || (d = D[r][m]/D[r][s]-D[i][m]/D[i][s]) < -E
        || d < E && ix[r+m] > ix[i+m]) r = i;
    }
    if (r < 0) return vd(); // unbounded
  }
  if (D[n+1][m] < -E) return vd(); // infeasible
  vd x(m-1);
  for(int i=m; i < n+m; ++i) if(ix[i]<m-1) x[ix[i]]=D[i-m][m];
  dbl result = D[n][m]; return x; // ans: D[n][m]
}
```

tridiagonal.h

Description: $x = \text{tridiagonal}(d, p, q, b)$ solves the equation system

$$\begin{pmatrix} b_0 \\ b_1 \\ b_2 \\ b_3 \\ \vdots \\ b_{n-1} \end{pmatrix} = \begin{pmatrix} d_0 & p_0 & 0 & 0 & \cdots & 0 \\ q_0 & d_1 & p_1 & 0 & \cdots & 0 \\ 0 & q_1 & d_2 & p_2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & 0 & \cdots & q_{n-3} & d_{n-2} & p_{n-2} \\ 0 & 0 & \cdots & 0 & q_{n-2} & d_{n-1} \end{pmatrix} \begin{pmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_{n-1} \end{pmatrix}.$$

This is useful for solving problems on the type

$$a_i = b_i a_{i-1} + c_i a_{i+1} + d_i, \ 1 \leq i \leq n,$$

where a_0, a_{n+1}, b_i, c_i and d_i are known. a can then be obtained from

$$\{a_i\} = \text{tridiagonal}(\{1, -1, -1, ..., -1, 1\}, \{0, c_1, c_2, \ldots, c_n\},$$
$$\{b_1, b_2, \ldots, b_n, 0\}, \{a_0, d_1, d_2, \ldots, d_n, a_{n+1}\}).$$

Fails if the solution is not unique.
If $|d_i| > |p_i| + |q_{i-1}|$ for all i, or $|d_i| > |p_{i-1}| + |q_i|$, or the matrix is positive definite, the algorithm is numerically stable and neither tr nor the check for diag[i] == 0 is needed.
Time: $\mathcal{O}\left(N\right)$

be9642, 23 lines

```cpp
typedef double T;
vector<T> tridiagonal(vector<T> diag, const vector<T> &super,
    const vector<T> &sub, vector<T> b) {
  int n = b.size(); vector<int> tr(n);
  for(int i = 0; i < n-1; ++i)
    if (abs(diag[i]) < 1e-9 * abs(super[i])) { // diag[i] == 0
      b[i+1] -= b[i] * diag[i+1] / super[i];
      if (i+2 < n) b[i+2] -= b[i] * sub[i+1] / super[i];
      diag[i+1] = sub[i]; tr[++i] = 1;
    } else {
      diag[i+1] -= super[i]*sub[i]/diag[i];
      b[i+1] -= b[i]*sub[i]/diag[i];
    }
```

```cpp
  for (int i = n; i--;)
    if (tr[i]) {
      swap(b[i], b[i-1]); diag[i-1] = diag[i];
      b[i] /= super[i-1];
    } else {
      b[i] /= diag[i];
      if (i) b[i-1] -= b[i]*super[i-1];
    }
  return b;
}
```

polyominoes.h

Description: Generate all fixed polyominoes with at most n squares. poly[x] gives the polyominoes with x squares. Takes less than a sec if $n < 10$, around 2s if $n = 10$ and around 6s if $n = 11$.

<div align="right">580a1b, 34 lines</div>

```cpp
const int LIM = 11;
using pii = pair<int,int>;
int dx[] = {0, 1, 0, -1};
int dy[] = {1, 0, -1, 0};
vector<vector<pii>> poly[LIM + 1];
void generate(int n = LIM){
  poly[1] = { { { 0 , 0 } } };
  for(int i = 2 ; i <= n; ++i) {
    set<vector<pii>> cur_om;
    for(auto &om : poly[i-1]) for(auto &p : om)
      for(int d = 0; d < 4; ++d) {
        int x = p.first + dx[d];
        int y = p.second + dy[d];
        if( ! binary_search(om.begin(), om.end(), pii(x,y)) ) {
          pii m = min(om[0], {x, y});
          pii new_cell(x - m.first, y - m.second);
          vector<pii> norm;
          norm.reserve(i);
          bool new_in = false;
          for(pii &c : om) {
            pii cur(c.first - m.first, c.second - m.second);
            if( ! new_in && cur > new_cell ) {
              new_in = true;
              norm.push_back(new_cell);
            }
            norm.push_back(cur);
          }
          if( ! new_in ) norm.push_back(new_cell);
          cur_om.insert(norm);
        }
      }
    poly[i].assign(cur_om.begin(), cur_om.end());
  }
}
```

4.1 Fourier transforms

fast-fourier-transform.h

Description: For convolution of complex numbers or more than two vectors: FFT, multiply pointwise, FFT inverse back.

Time: $\mathcal{O}(N \log N)$ with $N = |A| + |B|$ ($\sim$1s for $N = 2^{22}$)

<div align="right">366399, 123 lines</div>

```cpp
"../number-theory/modular-arithmetic.h"

inline int nxt_pow2(int s) { return 1 << (s > 1 ? 32 -
    __builtin_clz(s-1) : 0); }
template <typename T> struct root_of_unity {};

template <typename dbl> struct cplx { // hash-1
  dbl x, y; using P = cplx;
  cplx(dbl x_ = 0, dbl y_ = 0) : x(x_), y(y_) { }
  friend P operator+(P a, P b) { return P(a.x+b.x, a.y+b.y); }
  friend P operator-(P a, P b) { return P(a.x-b.x, a.y-b.y); }
  friend P operator*(P a, P b) { return P(a.x*b.x - a.y*b.y, a.
    x*b.y + a.y*b.x); }
```

```cpp
  friend P conj(P a) { return P(a.x, -a.y); }
  friend P inv(P a) { dbl n = (a.x*a.x+a.y*a.y); return P(a.x/n
    ,-a.y/n); }
};
template <typename dbl> struct root_of_unity<cplx<dbl>> {
  static cplx<dbl> f(int k) {
    static const dbl PI = acos(-1); dbl a = 2*PI/k;
    return cplx<dbl>(cos(a),sin(a));
  }
}; // hash-1 = 84c651

//(MOD 3) := (M1:897581057),(M3:985661441),(M5:935329793)
using M0 = modnum<998244353U>; // hash-2
constexpr unsigned primitive_root(unsigned M) {
  if (M == 880803841U) return 26U; // (M2)
  else if (M == 943718401U) return 7U; // (M4)
  else if (M == 918552577U) return 5U; // (M6)
  else return 3U;
}
template<unsigned MOD> struct root_of_unity<modnum<MOD>> {
  static constexpr modnum<MOD> g0 = primitive_root(MOD);
  static modnum<MOD> f(int K) {
    assert((MOD-1)%K == 0); return g0.pow((MOD-1)/K);
  }
}; // hash-2 = de8785
template <typename T> struct FFT {
  vector<T> rt; vector<int> rev;
  FFT() : rt(2, T(1)) {}
  void init(int N) { // hash-3
    N = nxt_pow2(N);
    if (N > int(rt.size())) {
      rev.resize(N); rt.reserve(N);
      for (int a = 0; a < N; ++a)
        rev[a] = (rev[a/2] | ((a&1)*N)) >> 1;
      for (int k = int(rt.size()); k < N; k *= 2) {
        rt.resize(2*k);
        T z = root_of_unity<T>::f(2*k);
        for (int a = k/2; a < k; ++a)
          rt[2*a] = rt[a], rt[2*a+1] = rt[a] * z;
      }
    }
  } // hash-3 = fcd279
  void fft(vector<T>& xs, bool inverse) const { // hash-4
    int N = int(xs.size());
    int s = __builtin_ctz(int(rev.size())/N);
    if (inverse) reverse(xs.begin() + 1, xs.end());
    for (int a = 0; a < N; ++a) {
      if (a < (rev[a] >> s)) swap(xs[a], xs[rev[a] >> s]);
    }
    for (int k = 1; k < N; k *= 2) {
      for (int a = 0; a < N; a += 2*k) {
        int u = a, v = u + k;
        for (int b = 0; b < k; ++b, ++u, ++v) {
          T z = rt[b + k] * xs[v];
          xs[v] = xs[u] - z, xs[u] = xs[u] + z;
        }
      }
    }
    if (inverse)
      for (int a = 0; a < N; ++a) xs[a] = xs[a] * inv(T(N));
  } // hash-4 = e2ea4a
  vector<T> convolve(vector<T> as, vector<T> bs) { // hash-5
    int N = int(as.size()), M = int(bs.size());
    int K = N + M - 1, S = nxt_pow2(K); init(S);
    if (min(N, M) <= 64) {
      vector<T> res(K);
      for (int u = 0; u < N; ++u) for (int v = 0; v < M; ++v)
        res[u + v] = res[u + v] + as[u] * bs[v];
      return res;
```

```cpp
    } else {
      as.resize(S), bs.resize(S);
      fft(as, false); fft(bs, false);
      for (int i = 0; i < S; ++i) as[i] = as[i] * bs[i];
      fft(as, true); as.resize(K); return as;
    }
  } // hash-5 = e4b903
}; FFT<M0> FFT0;
// T = {unsigned, unsigned long long, modnum<M>}
// Remark: need to satisfy |poly| * mod^2 < \prod_{i} M_i
template<class T, unsigned M0, unsigned M1, unsigned M2,
    unsigned M3, unsigned M4>
T garner(modnum<M0> a0, modnum<M1> a1, modnum<M2> a2, modnum<M3
    > a3, modnum<M4> a4) { // hash-6
  static const modnum<M1> INV_M0_M1 = modnum<M1>(M0).inv();
  static const modnum<M2> INV_M0M1_M2 = (modnum<M2>(M0) * M1).
      inv();
  // static const modnum<M3> INV_M0M1M2_M3 = (modnum<M3>(M0) *
      M1 * M2).inv();
  // static const modnum<M4> INV_M0M1M2M3_M4 = (modnum<M4>(M0)
      * M1 * M2 * M3).inv();
  const modnum<M1> b1 = INV_M0_M1 * (a1 - a0.x);
  const modnum<M2> b2 = INV_M0M1_M2 * (a2 - (modnum<M2>(b1.x) *
      M0 + a0.x));
  // const modnum<M3> b3 = INV_M0M1M2_M3 * (a3 - ((modnum<M3>(
      b2.x) * M1 + b1.x) * M0 + a0.x));
  // const modnum<M4> b4 = INV_M0M1M2M3_M4 * (a4 - (((modnum<M4
      >(b3.x) * M2 + b2.x) * M1 + b1.x) * M0 + a0.x));
  return (T(b2.x) * M1 + b1.x) * M0 + a0.x;
  // return (((T(b4.x) * M3 + b3.x) * M2 + b2.x) * M1 + b1.x) *
      M0 + a0.x;
} // hash-6 = b93180
// results must be in [-448002610255888384, 448002611254132736]
vector<long long> convolve(const vector<long long>& as, const
    vector<long long>& bs) { // hash-7
  static constexpr unsigned M0 = M0::M, M1 = M1::M;
  static const modnum<M1> INV_M0_M1 = modnum<M1>(M0).inv();
  if (as.empty() || bs.empty()) return {};
  const int len_as = int(as.size()), len_bs = int(bs.size());
  vector<modnum<M0>> as0(len_as), bs0(len_bs);
  for (int i = 0; i < len_as; ++i) as0[i] = as[i];
  for (int i = 0; i < len_bs; ++i) bs0[i] = bs[i];
  const vector<modnum<M0>> cs0 = FFT0.convolve(as0, bs0);
  vector<modnum<M1>> as1(len_as), bs1(len_bs);
  for (int i = 0; i < len_as; ++i) as1[i] = as[i];
  for (int i = 0; i < len_bs; ++i) bs1[i] = bs[i];
  const vector<modnum<M1>> cs1 = FFT1.convolve(as1, bs1);
  vector<long long> cs(len_as + len_bs - 1);
  for (int i = 0; i < len_as + len_bs - 1; ++i) {
    const modnum<M1> d1 = INV_M0_M1 * (cs1[i] - cs0[i].x);
    cs[i] = (d1.x > M1 - d1.x)
      ? (-1ULL - (static_cast<unsigned long long>(M1 - 1U - d1.
        x) * M0 + (M0 - 1U - cs0[i].x)))
      : (static_cast<unsigned long long>(d1.x) * M0 + cs0[i].x)
        ;
  }
  return cs;
} // hash-7 = bb50e0
```

fast-subset-transform.h

Description: Transform to a basis with fast convolutions of the form $c[z] = \sum_{z=x \oplus y} a[x] \cdot b[y]$, where $\oplus$ is one of AND, OR, XOR. The size of a must be a power of two.

Time: $\mathcal{O}(N \log N)$

<div align="right">5b9574, 16 lines</div>

```cpp
void FST(vector<int> &a, bool inv) { // hash-1
  for (int n = a.size(), step = 1; step < n; step *= 2) {
    for (int i = 0; i < n; i += 2 * step) for(int j = i; j < i+
      step; ++j) {
```

```cpp
      int &u = a[j], &v = a[j + step]; tie(u, v) =
        inv ? pii(v - u, u) : pii(v, u + v);   // AND
        inv ? pii(v, u - v) : pii(u + v, u);   // OR
        pii(u + v, u - v);                     // XOR
      }
    }
    if (inv) for(auto &x : a) x /= a.size(); // XOR only
} // hash-1 = a4980d
vector<int> conv(vector<int> a, vector<int> b) { // hash-2
  FST(a, 0); FST(b, 0);
  for(int i = 0; i < a.size(); ++i) a[i] *= b[i];
  FST(a, 1); return a;
} // hash-2 = 733c60
```

poly-998244353.h

`"finite-field-fft.h", "../number-theory/mod-sqrt.h", "../number-theory/preparator.h"`

`9cb4ca, 239 lines`

```cpp
using num = modnum<998244353U>; FFT<num> fft_data;
template<unsigned M> struct Poly : public vector<modnum<M>> {
    // hash-1
  Poly() {}
  explicit Poly(int n) : vector<modnum<M>>(n) {}
  Poly(const vector<modnum<M>> &vec) : vector<modnum<M>>(vec)
      {}
  Poly(std::initializer_list<modnum<M>> il) : vector<modnum<M
      >>(il) {}
  int size() const { return vector<modnum<M>>::size(); }
  num at(long long k) const { return (0 <= k && k < size()) ?
      (*this)[k] : 0U; }
  int ord() const { for (int i = 0; i < size(); ++i) if (int((*
      this)[i])) return i; return -1; }
  int deg() const { for (int i = size(); --i >= 0;) if (int((*
      this)[i])) return i; return -1; }
  Poly mod(int n) const { return Poly(vector<modnum<M>>(this->
      data(), this->data() + min(n, size()))); }
  friend std::ostream &operator<<(std::ostream &os, const Poly
      &fs) {
    os << "[";
    for (int i = 0; i < fs.size(); ++i) { if (i > 0) os << ", "
        ; os << fs[i]; }
    return os << "]";
  } // hash-1 = 37d6a1
  Poly &operator+=(const Poly &fs) { // hash-2
    if (size() < fs.size()) this->resize(fs.size());
    for (int i = 0; i < fs.size(); ++i) (*this)[i] += fs[i];
    return *this;
  } // hash-2 = d36bec
  Poly &operator-=(const Poly &fs) { // hash-3
    if (size() < fs.size()) this->resize(fs.size());
    for (int i = 0; i < fs.size(); ++i) (*this)[i] -= fs[i];
    return *this;
  } // hash-3 = 1f5853
  Poly &operator*=(const Poly &fs) { // hash-4
    if (this->empty() || fs.empty()) return *this = {};
    *this = fft_data.convolve(*this, fs);
    return *this;
  } // hash-4 = 24a997
  Poly &operator*=(const num &a) { // hash-5
    for (int i = 0; i < size(); ++i) (*this)[i] *= a;
    return *this;
  } // hash-5 = ea9fbe
  Poly &operator/=(const num &a) { // hash-6
    const num b = a.inv();
    for (int i = 0; i < size(); ++i) (*this)[i] *= b;
    return *this;
  } // hash-6 = 71618f
  Poly &operator/=(const Poly &fs) { // hash-7
    auto ps = fs;
    if (size() < ps.size()) return *this = {};
```

```cpp
    int s = int(size()) - int(ps.size()) + 1;
    int nn = 1; for (; nn < s; nn <<= 1) {}
    reverse(this->begin(), this->end());
    reverse(ps.begin(), ps.end());
    this->resize(nn); ps.resize(nn);
    ps = ps.inv();
    *this = *this * ps;
    this->resize(s); reverse(this->begin(), this->end());
    return *this;
  } // hash-7 = 291cdc
  Poly &operator%=(const Poly& fs) { // hash-8
    if (size() >= fs.size()) {
      Poly Q = (*this / fs) * fs;
      this->resize(fs.size() - 1);
      for (int x = 0; x < int(size()); ++x) (*this)[x] -= Q[x];
    }
    while (size() && this->back() == 0) this->pop_back();
    return *this;
  } // hash-8 = d6a389
  Poly inv() const { // hash-9
    if (this->empty()) return {};
    Poly b({(*this)[0].inv()}), fs;
    b.reserve(2 * int(this->size()));
    while (b.size() < this->size()) {
      int len = 2 * int(b.size());
      b.resize(2 * len, 0);
      if (int(fs.size()) < 2 * len) fs.resize(2 * len, 0);
      fill(fs.begin(), fs.begin() + 2 * len, 0);
      copy(this->begin(), this->begin() + min(len, int(this->
          size())), fs.begin());
      fft_data.fft(b, false);
      fft_data.fft(fs, false);
      for (int x = 0; x < 2*len; ++x) b[x] = b[x] * (2 - fs[x]
          * b[x]);
      fft_data.fft(b, true);
      b.resize(len);
    }
    b.resize(this->size()); return b;
  } // hash-9 = 49a504
  Poly differential() const { // hash-10
    if (this->empty()) return {};
    Poly f(max(size() - 1, 1));
    for (int x = 1; x < size(); ++x) f[x - 1] = x * (*this)[x];
    return f;
  } // hash-10 = 0b718c
  Poly integral() const { // hash-11
    if (this->empty()) return {};
    Poly f(size() + 1);
    for (int x = 0; x < size(); ++x) f[x + 1] = invs[x + 1] *
        (*this)[x];
    return f;
  } // hash-11 = 71d33a
  Poly log() const { // hash-12
    if (this->empty()) return {};
    Poly f = (differential() * inv()).integral();
    f.resize(size()); return f;
  } // hash-12 = 6a365e
  Poly exp() const { // hash-13
    Poly f = {1};
    if (this->empty()) return f;
    while (f.size() < size()) {
      int len = min(f.size() * 2, size());
      f.resize(len);
      Poly d(len);
      copy(this->begin(), this->begin() + len, d.begin());
      Poly g = d - f.log();
      g[0] += 1;
      f *= g;
      f.resize(len);
```

```cpp
    }
    return f;
  } // hash-13 = 25174b
  Poly pow(int N) const { // hash-14
    Poly b(size());
    if (N == 0) { b[0] = 1; return b; }
    int p = 0;
    while (p < size() && (*this)[p] == 0) ++p;
    if (1LL * N * p >= size()) return b;
    num mu = ((*this)[p]).pow(N), di = ((*this)[p]).inv();
    Poly c(size() - N*p);
    for (int x = 0; x < int(c.size()); ++x) {
      c[x] = (*this)[x + p] * di;
    }
    c = c.log();
    for (auto& val : c) val *= N;
    c = c.exp();
    for (int x = 0; x < int(c.size()); ++x) {
      b[x + N*p] = c[x] * mu;
    }
    return b;
  } // hash-14 = 48fee9
  Poly sqrt(int N) const { // hash-15
    if (!size()) return {};
    if (deg() == -1) return Poly(N);
    int p = 0;
    while (at(p) == 0 && p < size()) ++p;
    if (p >= N) return {0};
    Poly fs(2*N);
    copy(this->begin() + p, this->end(), fs.begin());
    auto v = mod_sqrt(fs.at(0).x, M);
    if (p & 1 || v.empty()) return {};
    fs.resize(size() - p/2);
    fs *= fs.front().inv();
    fs = v[0] * (fs.log() / 2).exp();
    fs.insert(fs.begin(), p/2, 0);
    return fs;
  } // hash-15 = 262e0f
  Poly operator+() const { return *this; }
  Poly operator-() const {
    Poly fs(size());
    for (int i = 0; i < size(); ++i) fs[i] = -(*this)[i];
    return fs;
  }
  Poly operator+(const Poly &fs) const { return (Poly(*this) +=
      fs); }
  Poly operator-(const Poly &fs) const { return (Poly(*this) -=
      fs); }
  Poly operator*(const Poly &fs) const { return (Poly(*this) *=
      fs); }
  Poly operator%(const Poly &fs) const { return (Poly(*this) %=
      fs); }
  Poly operator/(const Poly &fs) const { return (Poly(*this) /=
      fs); }
  Poly operator*(const num &a) const { return (Poly(*this) *= a
      ); }
  Poly operator/(const num &a) const { return (Poly(*this) /= a
      ); }
  friend Poly operator*(const num &a, const Poly &fs) { return
      fs * a; }
  // multipoint evaluation/interpolation
  friend Poly eval(const Poly& fs, const Poly& qs) { // hash-16
    int N = int(qs.size());
    if (N == 0) return {};
    vector<Poly> up(2 * N);
    for (int x = 0; x < N; ++x)
      up[x + N] = Poly({0-qs[x], 1});
    for (int x = N-1; x >= 1; --x)
      up[x] = up[2 * x] * up[2 * x + 1];
```

```cpp
  vector<Poly> down(2 * N);
  down[1] = fs % up[1];
  for (int x = 2; x < 2*N; ++x)
    down[x] = down[x/2] % up[x];
  Poly y(N);
  for (int x = 0; x < N; ++x)
    y[x] = (down[x + N].empty() ? 0 : down[x + N][0]);
  return y;
} // hash-16 = d2cd78
friend Poly interpolate(const Poly& fs, const Poly& qs) { //
    hash-17
  int N = int(fs.size());
  vector<Poly> up(2 * N);
  for (int x = 0; x < N; ++x)
    up[x + N] = Poly({0-fs[x], 1});
  for (int x = N-1; x >= 1; --x)
    up[x] = up[2 * x] * up[2 * x + 1];
  Poly E = eval(up[1].differential(), fs);
  vector<Poly> down(2 * N);
  for (int x = 0; x < N; ++x)
    down[x + N] = Poly({qs[x] * E[x].inv()});
  for (int x = N-1; x >= 1; --x)
    down[x] = down[2*x] * up[2*x+1] + down[2*x+1] * up[2*x];
  return down[1];
} // hash-17 = 874db6
friend Poly convolve_all(const vector<Poly>& fs, int l, int r
    ) {
  if (r - l == 1) return fs[l];
  else {
    int md = (l + r) / 2;
    return convolve_all(fs, l, md) * convolve_all(fs, md, r);
  }
}
Poly bernoulli(int N) const { // hash-18
  N += 5; Poly fs(N); fs[1] = 1;
  fs = fs.exp();
  copy(fs.begin()+1, fs.end(), fs.begin());
  fs = fs.inv();
  for (int x = 0; x < N; ++x) fs[x] *= fact[x];
  fs.resize(N - 5);
  return fs;
} // hash-18 = 237eef
// x(x - 1)(x - 2)...(x - N + 1)
Poly stirling_first(int N) const {
  if (N == 0) return {1};
  vector<Poly> P(N);
  for (int x = 0; x < N; ++x) P[x] = {-x, 1};
  return convolve_all(P, 0, N);
}
Poly stirling_second(int N) const {
  if (N == 0) return {1};
  Poly P(N), Q(N);
  for (int x = 0; x < N; ++x) {
    P[x] = (x & 1 ? -1 : 1) * ifact[x];
    Q[x] = num(x).pow(N-1) * ifact[x];
  }
  P *= Q; P.resize(N);
  return P;
}
Poly taylor_shift(int N, int K) const {
  Poly P(N), Q = *this; P[0] = 1;
  for (int i = 1; i < N; ++i) P[i] += P[i-1] * K;
  for (int i = 1; i < N; ++i) P[i] *= ifact[i];
  reverse(P.begin(), P.end());
  for (int i = 1; i < N; ++i) Q[i] *= fact[i];
  P *= Q;
  for (int i = 0; i < N; ++i) P[i] = P[N - 1 + i] * ifact[i];
  P.resize(N);
  return P;
```

```cpp
  }
};
```

sum-of-powers.h

Description: Computes monomials and sum of powers product certain polynomials. Check "General purpose numbers" section for more info. (Monomials) $pw(x) = x^d$ for a fixed d. $\sum_{x=0}^{\infty} r^x f(x)$. (degree of $f \le d$). $\sum_{x=0}^{N-1} r^x f(x)$. (degree of $f \le d$).

```cpp
vector<num> get_monomials(int N, long long d) {
  vector<int> pfac(N);
  for (int i = 2; i < N; ++i) pfac[i] = i;
  for (int p = 2; p < N; ++p) if (pfac[p] == p)
    for(int m = 2*p; m < N; m += p) if(pfac[m] > p) pfac[m]=p;
  vector<num> pw(N);
  for (int i = 0; i < N; ++i)
    if (i <= 1 || pfac[i] == i) pw[i] = num(i).pow(d);
    else pw[i] = (pw[pfac[i]] * pw[i / pfac[i]]);
  return pw;
}
num sum_of_power_limit(num r, int d, const vector<num>& fs) {
  interpolator_t<num> M(d + 2); num s = 1; auto gs = fs;
  for (int x = 0; x <= d; ++x, s *= r) gs[x] *= s;
  num ans = 0, cur_sum = 0; s = 1;
  for (int x = 0; x <= d; ++x, s *= -r) {
    cur_sum += choose(d+1, x) * s; ans += cur_sum * gs[d-x];
  } ans *= (1 - r).pow(-(d + 1));
  return ans;
}
num sum_of_power(num r, int d, vector<num>& fs, ll N) {
  if (r == 0) return (0 < N) ? fs[0] : 0;
  interpolator_t<num> M(d + 2);
  vector<num> gs(d + 2); gs[0] = 0; num s = 1;
  for (int x = 0; x <= d; ++x, s *= r)
    gs[x + 1] = gs[x] + s * fs[x];
  if (r == 1) return M.interpolate(gs, N);
  const num c = sum_of_power_limit(r, d, fs);
  const num r_inv = r.inv(); num w = 1;
  for (int x = 0; x <= d + 1; ++x, w *= r_inv)
    gs[x] = w * (gs[x] - c);
  return c + r.pow(N) * M.interpolate(gs, N);
}
```

4.1.1 General linear recurrences

If $a_n = \sum_{k=0}^{n-1} a_k b_{n-k}$, then $A(x) = \frac{a_0}{1-B(x)}$.

4.1.2 Polyominoes

How many free (rotation, reflection), one-sided (rotation) and fixed n-ominoes are there?

n	3	4	5	6	7	8	9	10
free	2	5	12	35	108	369	1.285	4.655
one-sided	2	7	18	60	196	704	2.500	9.189
fixed	6	19	63	216	760	2.725	9.910	36.446

4.1.3 LP formulation

Standard form: maximize $\mathbf{c}^T \mathbf{x}$ subject to $A\mathbf{x} \le \mathbf{b}$ and $\mathbf{x} \ge 0$.

Dual LP: minimize $\mathbf{b}^T \mathbf{y}$ subject to $A^T \mathbf{y} \ge \mathbf{c}$ and $\mathbf{y} \ge 0$.

$\bar{\mathbf{x}}$ and $\bar{\mathbf{y}}$ are optimal if and only if for all $i \in [1, n]$, either $\bar{x}_i = 0$ or $\sum_{j=1}^{m} A_{ji} \bar{y}_j = c_i$ holds and for all $i \in [1, m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^{n} A_{ij} \bar{x}_j = b_i$ holds.

1. In case of minimization, let $c_i' = -c_i$
2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
3. $\sum_{1 \le i \le n} A_{ji} x_i = b_j$
   - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$
   - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
4. If x_i has no lower bound, replace x_i with $x_i - x_i'$

4.1.4 Generating functions

A list of generating functions for useful sequences:

$(1, 1, 1, 1, 1, 1, \ldots)$	$\frac{1}{1-z}$
$(1, -1, 1, -1, 1, -1, \ldots)$	$\frac{1}{1+z}$
$(1, 0, 1, 0, 1, 0, \ldots)$	$\frac{1}{1-z^2}$
$(1, 0, \ldots, 0, 1, 0, 1, 0, \ldots, 0, 1, 0, \ldots)$	$\frac{1}{1-z^2}$
$(1, 2, 3, 4, 5, 6, \ldots)$	$\frac{1}{(1-z)^2}$
$(1, \binom{m+1}{m}, \binom{m+2}{m}, \binom{m+3}{m}, \ldots)$	$\frac{1}{(1-z)^{m+1}}$
$(1, c, \binom{c+1}{2}, \binom{c+2}{3}, \ldots)$	$\frac{1}{(1-z)^c}$
$(1, c, c^2, c^3, \ldots)$	$\frac{1}{1-cz}$
$(0, 1, \frac{1}{2}, \frac{1}{3}, \frac{1}{4}, \ldots)$	$\ln \frac{1}{1-z}$

A neat manipulation trick is:

$$\frac{1}{1-z} G(z) = \sum_n \sum_{k \le n} g_k z^n$$

Number theory (5)

5.1 Modular arithmetic

modular-arithmetic.h

Description: Operators for modular arithmetic.

```cpp
template<unsigned M_> struct modnum {
  static constexpr unsigned M = M_; using num = modnum;
  using ll = int64_t; using ull = uint64_t; unsigned x;
  num& norm(unsigned a){x = a<M ? a : a-M;return *this;}
  constexpr modnum(ll a = 0U) : x(unsigned((a %= ll(M)) < 0 ? a
      + ll(M) : a)) {}
  explicit operator int() const { return x; }
  num& operator+=(const num& a){ return norm(x+a.x); }
  num& operator-=(const num& a){ return norm(x-a.x+M); }
  num& operator*=(const num& a){ x = unsigned(ull(x)*a.x%M);
      return *this; }
  num& operator/=(const num& a){ return (*this *= a.inv());}
  num operator+(const num& a) const {return (num(*this) += a);}
  num operator-(const num& a) const {return (num(*this) -= a);}
  num operator*(const num& a) const {return (num(*this) *= a);}
  num operator/(const num& a) const {return (num(*this) /= a);}
  template<typename T> friend num operator+(T a, const num& b){
      return (num(a) += b); }
  template<typename T> friend num operator-(T a, const num& b){
      return (num(a) -= b); }
  template<typename T> friend num operator*(T a, const num& b){
      return (num(a) *= b); }
  template<typename T> friend num operator/(T a, const num& b){
      return (num(a) /= b); }
  num operator+() const { return *this; }
  num operator-() const { return num() - *this; }
  num pow(ll e) const {
```

```cpp
    if (e < 0) { return inv().pow(-e); } num b = x, xe = 1U;
    for (; e; e >>= 1) { if (e & 1) xe *= b; b *= b; }
    return xe;
  }
  num inv() const { return minv(x, M); }
  friend num inv(const num& a) { return a.inv(); }
  explicit operator bool() const { return x; }
  friend bool operator==(const num& a, const num& b){return a.x
        == b.x;}
  friend bool operator!=(const num& a, const num& b){return a.x
        != b.x;}
};
```

pairnum-template.h
Description: Support pairs operations using modnum template. Pretty good for string hashing.
229a89, 42 lines

```cpp
template <typename T, typename U> struct pairnum {
  T t; U u;
  pairnum() : t(0), u(0) {}
  pairnum(long long v) : t(v), u(v) {}
  pairnum(const T& t_, const U& u_) : t(t_), u(u_) {}
  friend std::ostream& operator << (std::ostream& out, const
        pairnum& n) { return out << '(' << n.t << ',' << ' ' <<
        n.u << ')'; }
  friend std::istream& operator >> (std::istream& in, pairnum&
        n) { long long v; in >> v; n = pairnum(v); return in; }
  friend bool operator == (const pairnum& a, const pairnum& b)
        { return a.t == b.t && a.u == b.u; }
  friend bool operator != (const pairnum& a, const pairnum& b)
        { return a.t != b.t || a.u != b.u; }
  pairnum inv() const {
    return pairnum(t.inv(), u.inv());
  }
  pairnum neg() const {
    return pairnum(t.neg(), u.neg());
  }
  pairnum operator- () const {
    return pairnum(-t, -u);
  }
  pairnum operator+ () const {
    return pairnum(+t, +u);
  }
  pairnum& operator += (const pairnum& o) {
    t += o.t; u += o.u;
    return *this;
  }
  pairnum& operator -= (const pairnum& o) {
    t -= o.t; u -= o.u;
    return *this;
  }
  pairnum& operator *= (const pairnum& o) {
    t *= o.t; u *= o.u;
    return *this;
  }
  pairnum& operator /= (const pairnum& o) {
    t /= o.t; u /= o.u;
    return *this;
  }
  friend pairnum operator + (const pairnum& a, const pairnum& b
        ) { return pairnum(a) += b; }
  friend pairnum operator - (const pairnum& a, const pairnum& b
        ) { return pairnum(a) -= b; }
  friend pairnum operator * (const pairnum& a, const pairnum& b
        ) { return pairnum(a) *= b; }
  friend pairnum operator / (const pairnum& a, const pairnum& b
        ) { return pairnum(a) /= b; }
};
```

preparator.h
Description: Precompute factorials and inverses
"modular-arithmetic.h" 4bfc8c, 11 lines

```cpp
constexpr int V = 1 << 20;
num invs[V], fact[V], ifact[V];
void prepare() {
  invs[1] = 1;
  for (int i = 2; i < V; ++i) invs[i] = -((num::M / i) * invs[
        num::M % i]);
  fact[0] = ifact[0] = 1;
  for (int i = 1; i < V; ++i) {
    fact[i] = fact[i - 1] * i;
    ifact[i] = ifact[i - 1] * invs[i];
  }
}
```

mod-inv.h
Description: Find x such that $ax \equiv 1(\bmod\ m)$. The inverse only exist if a and m are coprimes.
3b5ae9, 4 lines

```cpp
template<typename T> T minv(T a, T m) {
  a %= m; assert(a);
  return a == 1 ? 1 : T(m - i64(minv(m, a)) * m / a);
}
```

mod-sum.h
Description: Sums of mod'ed arithmetic progressions.
modsum(to, c, k, m) $= \sum_{i=0}^{to-1} (ki+c)\%m$. divsum is similar but for floored division.
Time: $\log(m)$, with a large constant.
decfb8, 16 lines

```cpp
typedef unsigned long long ull;
ull sumsq(ull to) { return to / 2 * ((to-1) | 1); }
ull divsum(ull to, ull c, ull k, ull m) {
  ull res = k / m * sumsq(to) + c / m * to;
  k %= m; c %= m;
  if (k) {
    ull to2 = (to * k + c) / m;
    res += to * to2;
    res -= divsum(to2, m-1 - c, m, k) + to2;
  }
  return res;
}
lint modsum(ull to, lint c, lint k, lint m) {
  c = ((c % m) + m) % m; k = ((k % m) + m) % m;
  return to * c + k * sumsq(to) - m * divsum(to, c, k, m);
}
```

mod-mul.h
Description: Calculate $a \cdot b \bmod c$ (or $a^b \bmod c$) for $0 \le a, b \le c \le 7.2 \cdot 10^{18}$.
Time: $\mathcal{O}(1)$ for modmul, $\mathcal{O}(\log b)$ for modpow
59afa8, 11 lines

```cpp
typedef unsigned long long ull;
ull modmul(ull a, ull b, ull M) {
  lint ret = a * b - M * ull(1.L / M * a * b);
  return ret + M * (ret < 0) - M * (ret >= (lint)M);
}
ull modpow(ull b, ull e, ull mod) {
  ull ans = 1;
  for (; e; b = modmul(b, b, mod), e /= 2)
    if (e & 1) ans = modmul(ans, b, mod);
  return ans;
}
```

mod-sqrt.h
Description: Tonelli-Shanks algorithm for modular square roots. Finds x s.t. $x^2 = a \pmod p$ ($-x$ gives the other solution).
Time: $\mathcal{O}(\log^2 p)$ worst case, $\mathcal{O}(\log p)$ for most p
abffbe, 33 lines

```cpp
int jacobi(ll a, ll m) { // Jacobi symbol (a/m)
  int s = 1;
  if (a < 0) a = a % m + m;
  for (; m > 1; ) {
    a %= m; if (a == 0) return 0;
    const int r = __builtin_ctzll(a);
    if ((r & 1) && ((m + 2) & 4)) s = -s;
    a >>= r; if (a & m & 2) s = -s;
    swap(a, m);
  } return s;
}
vector<ll> mod_sqrt(ll a, ll p) {
  if (p == 2) return {a & 1};
  const int j = jacobi(a, p);
  if (j == 0) return {0};
  if (j == -1) return {};
  ll b, d;
  while (true) {
    b = xrand() % p; d = (b * b - a) % p;
    if (d < 0) d += p;
    if (jacobi(d, p) == -1) break;
  }
  ll f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (ll e = (p + 1) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (g0 * f0 + d * ((g1 * f1) % p)) % p;
      g1 = (g0 * f1 + g1 * f0) % p; g0 = tmp;
    }
    tmp = (f0 * f0 + d * ((f1 * f1) % p)) % p;
    f1 = (2 * f0 * f1) % p; f0 = tmp;
  }
  return (g0<p-g0) ? vector<ll>{g0,p-g0} : vector<ll>{p-g0,g0};
}
```

mul-order.h
Description: Find the smallest integer k such that $a^k \pmod m = 1$. $0 < k < m$.
Time: $\mathcal{O}(\log(N))$
"prime-factors.h", "mod-pow.h" 3d20e1, 12 lines

```cpp
template<typename T> T mul_order(T a, T m) {
  if (__gcd(a, m) != 1) return 0;
  auto N = phi(m);
  auto primes = prime_factorize(N);
  T res = 1;
  for (auto &[p, e] : primes) {
    while (N % p == 0 && modpow(a, N/p, m) == 1) {
      N /= p;
    }
  }
  return N;
}
```

mod-range.h
Description: min $x \ge 0$ s.t. $l \le ((ax)\ \bmod m) \le r$, $m > 0$, $a \ge 0$.
eb665e, 10 lines

```cpp
template<typename T> T mod_range(T m, T a, T l, T r) {
  l = max(1, T(0)); r = min(r, m - 1);
  if (l > r) return -1;
  a %= m;
  if (a == 0) return (l > 0) ? -1 : 0;
  const T k = (l + a - 1) / a;
  if (a * k <= r) return k;
  const T y = mod_range(a, m, a * k - r, a * k - l);
  return (y == -1) ? -1 : ((m * y + r) / a);
}
```

5.2   Primality

sieve.h

Description: Prime sieve for generating all primes up to a certain limit. $pfac[i]$ is the lowest prime factor of i. Also useful if you need to compute any multiplicative function.

Time: $\mathcal{O}(N)$

<div align="right">a0b9c3, 17 lines</div>

```cpp
vector<int> run_sieve(int N) {
  vector<int> pfac(N+1), primes, mu(N+1,-1), phi(N+1);
  primes.reserve(N+1); mu[1] = phi[1] = 1;
  for (int i = 2; i <= N; ++i) {
    if (!pfac[i])
      pfac[i] = i, phi[i] = i-1, primes.push_back(i);
    for (int p : primes) {
      if (p > N/i) break;
      pfac[p * i] = p; mu[p * i] *= mu[i];
      phi[p * i] = phi[i] * phi[p];
      if (i % p == 0) {
        mu[p * i] = 0; phi[p * i] = phi[i] * p;
        break;
      }
    }
  } return primes;
}
```

segmented-sieve.h

Description: Prime sieve for generating all primes smaller than S.

Time: S=1e9 $\approx$ 1.5s

<div align="right">68455e, 20 lines</div>

```cpp
const int S = 1e6;
bitset<S> isPrime;
vector<int> eratosthenes() {
  const int S = round(sqrt(S)), R = S/2;
  vector<int> pr = {2}, sieve(S+1); pr.reserve(int(S/log(S)
      *1.1));
  vector<pair<int,int>> cp;
  for (int i = 3; i <= S; i += 2) if (!sieve[i]) {
    cp.push_back({i, i*i/2});
    for (int j = i*i; j <= S; j += 2*i) sieve[j] = 1;
  }
  for (int L = 1; L <= R; L += S) {
    array<bool, S> block{};
    for (auto &[p, idx] : cp)
      for (int i=idx; i < S+L; idx = (i+=p)) block[i-L] = 1;
    for (int i = 0; i < min(S, R - L); ++i)
      if (!block[i]) pr.push_back((L + i)*2 + 1);
  }
  for (int i : pr) isPrime[i] = 1;
  return pr;
}
```

miller-rabin.h

Description: Deterministic Miller-Rabin primality test. Guaranteed to work for numbers up to 2^{64}; for larger numbers, extend A randomly.

Time: 7 times the complexity of $a^b \bmod c$.

<div align="right">"mod-mul.h"   bbee97, 12 lines</div>

```cpp
bool isPrime(ull n) {
  if (n < 2 || n % 6 % 4 != 1) return (n | 1) == 3;
  vector<ull> A = {2, 325, 9375, 28178, 450775, 9780504,
      1795265022};
  ull s = __builtin_ctzll(n-1), d = n >> s;
  for(ull a : A) {    // ^ count trailing zeroes
    ull p = modpow(a % n, d, n), i = s;
    while (p != 1 && p != n - 1 && a % n && i--)
      p = modmul(p, p, n);
    if (p != n-1 && i != s) return 0;
  }
  return 1;
}
```

pollard-rho.h

Description: Pollard-rho randomized factorization algorithm. Returns prime factors of a number, in arbitrary order (e.g. 2299 -> {11, 19, 11}).

Time: $\mathcal{O}\left(n^{1/4}\right)$, less for numbers with small factors.

<div align="right">"mod-mul.h", "extended-euclid.h", "miller-rabin.h"   6bf31f, 17 lines</div>

```cpp
ull pollard(ull n) {
  auto f = [n](ull x, ull k) { return modmul(x, x, n) + k; };
  ull x = 0, y = 0, t = 30, prd = 2, i = 1, q;
  while (t++ % 40 || gcd(prd, n) == 1) {
    if (x == y) x = ++i, y = f(x, i);
    if ((q = modmul(prd, max(x,y) - min(x,y), n))) prd = q;
    x = f(x, i), y = f(f(y, i), i);
  }
  return gcd(prd, n);
}
vector<ull> factor(ull n) {
  if (n == 1) return {};
  if (isPrime(n)) return {n};
  ull x = pollard(n); auto l = factor(x), r = factor(n/x);
  l.insert(l.end(), r.begin(), r.end());
  return l;
}
```

5.3   Divisibility

extended-euclid.h

Description: Finds two integers x and y, such that $ax + by = \gcd(a, b)$. If a and b are coprime, then x is the inverse of $a \pmod{b}$.

<div align="right">8b62c4, 6 lines</div>

```cpp
template<typename T>
T egcd(T a, T b, T &x, T &y) {
  if (!a) { x = 0, y = 1; return b; }
  T g = egcd(b % a, a, y, x);
  x -= y * (b/a); return g;
}
```

division-lemma.h

Description: This lemma let us exploit the fact tha the sequence (harmonic on integer division) has at most $2\sqrt{N}$ distinct elements, so we can iterate through every possible value of $\lfloor \frac{N}{i} \rfloor$, using the fact that the greatest integer j satisfying $\lfloor \frac{N}{i} \rfloor = \lfloor \frac{N}{j} \rfloor$ is $\lfloor \frac{N}{\lfloor \frac{N}{i} \rfloor} \rfloor$. This one computes the $\sum_{i=1}^{N} \lfloor \frac{N}{i} \rfloor i$.

Time: $\mathcal{O}\left(\sqrt{N}\right)$

<div align="right">b2c1ab, 15 lines</div>

```cpp
int res = 0;
for (int a = 1, b; a <= N; a = b + 1) { // floor
  b = N / (N / a);
  // quotient (N/a) and there are (b - a + 1) elements
  int l = b - a + 1, r = a + b; // l * r / 2 = sum(i, j)
  if (l & 1) r /= 2;
  else l /= 2;
  res += l * r * (N / a);
}
// [1, N), need to deal with case where a = N separately
for (int a = 1, b; a < N; a = b + 1) { // ceil
  const int k = (N - 1) / a + 1; // quotient k
  b = (N - 1) / (k - 1);
  int cnt = b - a + 1; // occur cnt times on interval [a, b]
}
```

prime-factors.h

Description: Find all prime factors of n.

Time: $\mathcal{O}(log(n))$

<div align="right">"sieve.h"   7a803a, 25 lines</div>

```cpp
template<typename T>
vector<pair<T, int>> prime_factorize(T n) {
```

```cpp
}
  vector<pair<T, int>> factors;
  while (n != 1) {
    T p = pfac[n];
    int exp = 0;
    do {
      n /= p;
      ++exp;
    } while(n % p == 0);
    factors.push_back({p, exp});
  }
  for (T p : primes) {
    if (p * p > n) break;
    if (p * p == 0) {
      factors.push_back({p, 0});
      do {
        n /= p;
        ++factors.back().second;
      } while(n % p == 0);
    }
  }
  if (n > 1) factors.push_back({n, 1});
  return factors;
}
```

divisors.h

Description: Generate all factors of n given it's prime factorization.

Time: $\mathcal{O}\left(\frac{\sqrt{N}}{\log N}\right)$

<div align="right">"prime-factors.h"   5de75c, 14 lines</div>

```cpp
template<typename T> vector<T> get_divisors(T N) {
  auto factors = prime_factorize(N);
  vector<T> ans; ans.reserve(int(sqrtl(N) + 1));
  auto dfs = [&](auto&& self, auto& ans, T val, int d) -> void{
    auto& [P, E] = factors[d];
    if (d == int(factors.size())) ans.push_back(val);
    else {
      T X = 1;
      for (int pw = 0; pw <= E; ++pw, X *= P)
        self(self, ans, val * X, d + 1);
    }
  }; dfs(dfs, ans, 1, 0);
  return ans;
}
```

num-div.h

Description: Count the number of divisors of n. Requires having run Sieve up to at least sqrt(n).

Time: $\mathcal{O}(log(N))$

<div align="right">"sieve.h"   be1146, 15 lines</div>

```cpp
template<typename T> T numDiv(T n) {
  T how_many = 1, prime_factors = 0;
  while (n != 1) {
    T p = lp[n];
    int exp = 0;
    do {
      n /= p;
      ++exp;
      ++prime_factors; //count prime factors!
    } while(n % p == 0);
    how_many *= 1ll * (exp + 1);
  }
  if (n != 1) ++prime_factors;
  return how_many;
}
```

sum-div.h

Description: Sum of all divisors of n.

Time: $\mathcal{O}(log(N))$

"sieve.h", "mod-pow.h"   0448f4, 13 lines
```cpp
template<typename T> T divSum(T n) {
  T sum = 1;
  while (n > 1) {
    int exp = 0;
    T p = pfac[n];
    do {
      n /= p;
      ++exp;
    } while(n % p == 0);
    sum *= (modpow(p, exp + 1, mod) - 1)/(p - 1);
  }
  return sum;
}
```

phi-function.h

da7671, 6 lines
```cpp
const int n = int(1e5)*5; vector<int> phi(n);
void calculatePhi() {
  for(int i = 0; i < n; ++i) phi[i] = i&1 ? i : i/2;
  for(int i = 3; i < n; i += 2) if (phi[i] == i)
    for(int j = i; j < n; j += i) phi[j] -= phi[j]/i;
}
```

discrete-log.h

Description: Returns the smallest $x >= 0$ s.t. $a^x = b \pmod{m}$, or -1 if no such x exists. modLog(a,1,m) can be used calculate the order of a. Assumes that $0^0 = 1$.
Time: $\mathcal{O}\left(\sqrt{m}\right)$

"extended-euclid.h"   62fc5e, 15 lines
```cpp
template<typename T> T modLog(T a, T b, T m) {
  T k = 1, it = 0, g;
  while ((g = gcd(a, m)) != 1) {
    if (b == k) return it;
    if (b % g) return -1;
    b /= g; m /= g; ++it; k = k * a / g % m;
  }
  T n = sqrtl(m) + 1, f = 1, j = 1;
  unordered_map<T, T> A;
  while (j <= n)
    f = f * a % m, A[f * b % m] = j++;
  for(int i = 1; i <= n; ++i) if (A.count(k = k * f % m))
    return n * i - A[k] + it;
  return -1;
}
```

primitive-roots.h

Description: a is a primitive root mod n if for every number x coprime to n there is an integer z s.t. $x \equiv a^z \pmod{n}$. The number of primitive roots mod n, if there are any, is equal to $phi(phi(N))$. If m isnt prime, replace $m - 1$ by $phi(m)$.
Time: $\mathcal{O}(\log(N))$

<sieve.h>, <prime-factors.h>, <mod-pow.h>   05729f, 6 lines
```cpp
template<typename T> bool is_primitive(T a, T m) {
  vector<pair<T, T>> D = prime_factorize(m-1);
  for (auto p : D)
    if (modpow(a, (m-1)/p.first, m) == 1) return false;
  return true;
}
```

prime-counting.h

Description: Count the number of primes up to N. Also useful for sum of primes.
Time: $\mathcal{O}(N^{3/4}/\log N)$

c26239, 20 lines
```cpp
struct primes_t {
  vector<ll> dp, w;
```

```cpp
  ll pi(ll N) {
    const int sqrtN = int(sqrt(N));
    for (ll a = 1, b; a <= N; a = b+1)
      b = N / (N / a), w.push_back(N/a);
    auto get = [&](ll x) {
      if (x <= sqrtN) return int(x-1);
      return int(w.size() - N/x);
    };
    reverse(w.begin(), w.end()); dp.reserve(w.size());
    for (auto& x : w) dp.push_back(x-1);
    for (ll i = 2; i*i <= N; ++i) {
      if (dp[i-1] == dp[i-2]) continue;
      for (int j = int(w.size())-1; w[j] >= i*i; --j)
        dp[j] -= dp[get(w[j]/i)] - dp[i-2];
    }
    return dp.back();
  }
};
```

5.4 Chinese remainder theorem

chinese-remainder.h

Description: Chinese Remainder Theorem. crt(a, m, b, n) computes x such that $x \equiv a \pmod{m}$, $x \equiv b \pmod{n}$. If $|a| < m$ and $|b| < n$, x will obey $0 \le x < \operatorname{lcm}(m, n)$. Assumes $mn < 2^{62}$.
Time: $\mathcal{O}(\log(LCM(m, n)))$

"extended-euclid.h"   ece59a, 7 lines
```cpp
pair<ll, ll> crt(ll a, ll m, ll b, ll n) {
  if (n > m) swap(a, b), swap(m, n);
  ll x, y, g = egcd(m, n, x, y);
  if ((a - b) % g != 0) return {0, -1};
  x = (b - a) % n * x % n / g * m + a;
  return {x + (x < 0 ? m*n/g : 0), m*n/g};
}
```

5.5 Fractions

fractions.h

Description: Template that helps deal with frtions.

596163, 28 lines
```cpp
template<typename num> struct fraction_t {
  num p, q; using fr = fraction_t;
  fraction_t() : p(0), q(1) { }
  fraction_t(num _n, num _d = 1): p(_n), q(_d){
    num g = gcd(p, q); p /= g, q /= g;
    if (q < 0) p *= -1, q *= -1; assert(q != 0);
  }
  friend bool operator<(const fr& l, const fr& r){
    return l.p*r.q < r.p*l.q;}
  friend bool operator==(const fr& l, const fr& r){return l.p
    == r.p && l.q == r.q;}
  friend bool operator!=(const fr& l, const fr& r){return !(l
    == r);}
  friend fr operator+(const fr& l, const fr& r){
    num g = gcd(l.q, r.q);
    return fr(r.q / g * l.p + l.q / g * r.p, l.q / g * r.q);
  }
  friend fr operator-(const fr& l, const fr& r) {
    num g = gcd(l.q, r.q);
    return fr( r.q / g * l.p - l.q / g * r.p, l.q / g * r.q);
  }
  friend fr operator*(const fr& l, const fr& r){
    return fr(l.p*r.p, l.q*r.q);}
  friend fr operator/(const fr& l, const fr& r){
    return l*fr(r.q,r.p);}
  friend fr& operator+=(fr& l, const fr& r){return l=l+r;}
  friend fr& operator-=(fr& l, const fr& r){return l=l-r;}
  template<class T> friend fr& operator*=(fr& l, const T& r){
    return l=l*r;}
```

```cpp
  template<class T> friend fr& operator/=(fr& l, const T& r){
    return l=l/r;}
};
```

continued-fractions.h

Description: Given N and a real number $x \ge 0$, finds the closest rational approximation p/q with $p, q \le N$. It will obey $|p/q - x| \le 1/qN$.
For consecutive convergents, $p_{k+1}q_k - q_{k+1}p_k = (-1)^k$. ($p_k/q_k$ alternates between $> x$ and $< x$.) If x is rational, y eventually becomes ∞; if x is the root of a degree 2 polynomial the a's eventually become cyclic.
Time: $\mathcal{O}(\log N)$

3dfe17, 18 lines
```cpp
typedef double dbl; // for N ~ 1e7; long double for N ~ 1e9
pair<ll, ll> approximate(dbl x, ll N) { // hash-1
  ll LP = 0, LQ = 1, P = 1, Q = 0, inf = ll(1e18); dbl y = x;
  for (;;) {
    ll lim = min(P ? (N-LP) / P : inf, Q ? (N-LQ) / Q : inf),
       a = (ll)floor(y), b = min(a, lim),
       NP = b*P + LP, NQ = b*Q + LQ;
    if (a > b) {
      // If b > a/2, we have a semi-convergent that gives us a
      // better approximation; if b = a/2, we *may* have one.
      // Return {P, Q} here for a more canonical approximation.
      return (abs(x - (dbl)NP / (dbl)NQ) < abs(x - (dbl)P / (
        dbl)Q)) ?
      make_pair(NP, NQ) : make_pair(P, Q);
    }
    if (abs(y = 1/(y - (dbl)a)) > 3*N) return {NP, NQ};
    LP = P; P = NP; LQ = Q; Q = NQ;
  }
} // hash-1 = 67b717
```

frac-binary-search.h

Description: Given f and N, finds the smallest fraction $p/q \in [0, 1]$ such that $f(p/q)$ is true, and $p, q \le N$. You may want to throw an exception from f if it finds an exact solution, in which case N can be removed.
Usage: fracBS([](Frac f) { return f.p>=3*f.q; }, 10); // {1,3}
Time: $\mathcal{O}(\log(N))$

74ffd8, 20 lines
```cpp
struct Frac { ll p, q; };
template<class F> Frac fracBS(F f, ll N) {
  bool dir = 1, A = 1, B = 1;
  Frac left{0, 1}, right{1, 1};//right{1,0} to search (0,N]
  assert(!f(left)); assert(f(right));
  while (A || B) {
    ll adv = 0, step = 1; // move right if dir, else left
    for (int si = 0; step; (step *= 2) >>= si) {
      adv += step;
      Frac mid{left.p * adv + right.p, left.q * adv + right.q};
      if (abs(mid.p) > N || mid.q > N || dir == !f(mid)) {
        adv -= step; si = 2;
      }
    }
    right.p += left.p * adv; right.q += left.q * adv;
    dir = !dir; swap(left, right);
    A = B; B = !!adv;
  }
  return dir ? right : left;
}
```

5.5.1 Bézout's identity

For $a \ne 0$, $b \ne 0$, then $d = gcd(a, b)$ is the smallest positive integer for which there are integer solutions to

$$ax + by = d$$

If (x, y) is one solution, then all solutions are given by

$$\left(x + \frac{kb}{\gcd(a,b)}, y - \frac{ka}{\gcd(a,b)}\right), \quad k \in \mathbb{Z}$$

5.5.2 Pythagorean Triples

The Pythagorean triples are uniquely generated by

$$a = k \cdot (m^2 - n^2), \ \ b = k \cdot (2mn), \ \ c = k \cdot (m^2 + n^2),$$

with $m > n > 0$, $k > 0$, $m \perp n$, and either m or n even.

5.5.3 Chicken McNugget theorem

Let x and y be two coprime integers, the greater integer that can't be written in the form of $ax + by$ is $\frac{(x-1)(y-1)}{2}$

5.6 Primes

$p = 962592769$ is such that $2^{21} \mid p - 1$, which may be useful. For hashing use 970592641 (31-bit number), 31443539979727 (45-bit), 3006703054056749 (52-bit). There are 78498 primes less than 1 000 000.

Primitive roots exist modulo any prime power p^a, except for $p = 2, a > 2$, and there are $\phi(\phi(p^a))$ many. For $p = 2, a > 2$, the group $\mathbb{Z}_{2^a}^\times$ is instead isomorphic to $\mathbb{Z}_2 \times \mathbb{Z}_{2^{a-2}}$.

5.6.1 Prime counting function ($\pi(x)$)

The prime counting function is asymptotic to $\frac{x}{\log x}$, by the prime number theorem.

x	10	10^2	10^3	10^4	10^5	10^6	10^7	10^8
$\pi(x)$	4	25	168	1.229	9.592	78.498	664.579	5.761.455

5.6.2 Sum of primes

For any multiplicative f:

$$S(n, p) = S(n, p - 1) - f(p) \cdot (S(n/p, p - 1) - S(p - 1, p - 1))$$

5.6.3 Moebius Function

$$\mu(n) = \begin{cases} 0 & n \text{ is not square free} \\ 1 & n \text{ has even number of prime factors} \\ -1 & n \text{ has odd number of prime factors} \end{cases}$$

Moebius Inversion:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu(d) g(n/d)$$

Other useful formulas/forms:

$$\sum_{d|n} \phi(d) = n$$

$$\sum_{\substack{i < n \\ \gcd(i,n)=1}} i = n \frac{\phi(n)}{2}$$

$$\sum_{a=1}^n \sum_{b=1}^n [\gcd(a,b) = 1] = \sum_{d=1}^n \mu(d) \lfloor \frac{n}{d} \rfloor^2$$

$$\sum_{a=1}^n \sum_{b=1}^n \gcd(a,b) = \sum_{d=1}^n d \sum_{d|x} \lfloor \frac{n}{x} \rfloor^2 \mu(\frac{x}{d})$$

$$\sum_{a=1}^n \sum_{b=a}^n \gcd(a,b) = \sum_{d=1}^n \sum_{d|x} \phi(\frac{x}{d}) d$$

$$\sum_{a=1}^n \sum_{b=1}^n \text{lcm}(a,b) = \sum_{d=1}^n \mu(d) d \sum_{d|x}^n x \left(\frac{\lfloor \frac{n}{x} \rfloor + 1}{2}\right)^2$$

$$\sum_{a=1}^n \sum_{b=a+1}^n \text{lcm}(a,b) = \sum_{d=1}^n \sum_{d|x \wedge x > d} \phi(\frac{x}{d}) \frac{x^2}{2d}$$

$$\sum_{a \in S} \sum_{b \in S} \gcd(a,b) = \sum_{d=1}^n (\sum_{x|d} \frac{d}{x} \mu(x))(\sum_{d|v} \text{freq}[v])^2$$

$$\sum_{a \in S} \sum_{b \in S} \text{lcm}(a,b) = \sum_{d=1}^n \frac{1}{d} (\sum_{x|d} x\mu(x))(\sum_{d|v} v \cdot \text{freq}[v])^2$$

$$\sum_{d|n} \mu(d) = [n = 1] \text{ (very useful)}$$

$$g(n) = \sum_{n|d} f(d) \Leftrightarrow f(n) = \sum_{n|d} \mu(d/n) g(d)$$

$$g(n) = \sum_{1 \le m \le n} f(\lfloor \frac{n}{m} \rfloor) \Leftrightarrow f(n) = \sum_{1 \le m \le n} \mu(m) g(\lfloor \frac{n}{m} \rfloor)$$

5.6.4 Dirichlet Convolution

Given a function $f(x)$, let

$$(f * g)(x) = \sum_{d|x} g(d) f(x/d)$$

If the partial sums $s_{f*g}(n), s_g(n)$ can be computed in $O(1)$ and $s_f(1 \ldots n^{2/3})$ can be computed in $O\left(n^{2/3}\right)$ then all $s_f\left(\frac{n}{d}\right)$ can as well. Use

$$s_{f*g}(n) = \sum_{d=1}^n g(d) s_f(n/d).$$

$$\implies s_f(n) = \frac{s_{f*g}(n) - \sum_{d=2}^n g(d) s_f(n/d)}{g(1)}$$

1. If $f(x) = \mu(x)$ then $g(x) = 1, (f * g)(x) = (x == 1)$, and $s_f(n) = 1 - \sum_{i=2}^n s_f(n/i)$

2. If $f(x) = \phi(x)$ then $g(x) = 1, (f * g)(x) = x$, and $s_f(n) = \frac{n(n+1)}{2} - \sum_{i=2}^n s_f(n/i)$

dirichlet-convolution.h
Description: Dirichlet convolution. Change f, gs and fgs accordingly. This example calculates $\phi(N)$.
Time: $\mathcal{O}\left(N^{\frac{2}{3}}\right)$
eac754, 21 lines

```
template<typename T, typename V> struct mertens {
  V N; T inv; // ~ N^{2/3}
  vector<V> fs; vector<T> psum;
  unordered_map<V, T> mapa;
  V f(V x) { return fs[x]; }
  T gs(V x) { return x; }
  T fgs(V x) { return T(x) * (x + 1) / 2; }
  mertens(V M, const vector<V>& F) : N(M+1), fs(F), psum(M+1){
    inv = gs(1);
    for (V a = 0; a + 1 < N; ++a)
      psum[a + 1] = f(a + 1) + psum[a];
  }
```

```
  T query(V x) {
    if (x < N) return psum[x];
    if (mapa.find(x) != mapa.end()) return mapa[x];
    T ans = fgs(x);
    for (V a = 2, b; a <= x; a = b + 1)
      b = x/(x/a), ans -= (gs(b)-gs(a-1)) * query(x/a);
    return mapa[x] = (ans / inv);
  }
};
```

5.6.5 Estimates

$$\sum_{d|n} d = O(n \log \log n)$$

The number of divisors of n is at most around 100 for $n < 5e4$, 500 for $n < 1e7$, 2000 for $n < 1e10$, 200 000 for $n < 1e19$.

Combinatorial (6)

6.1 Permutations

6.1.1 Factorial

n	1	2	3	4	5	6	7	8	9	10
$n!$	1	2	6	24	120	720	5040	40320	362880	3628800

n	11	12	13	14	15	16	17
$n!$	4.0e7	4.8e8	6.2e9	8.7e10	1.3e12	2.1e13	3.6e14

n	20	25	30	40	50	100	150	171
$n!$	2e18	2e25	3e32	8e47	3e64	9e157	6e262	>DBL_MAX

int-perm.h
Description: Permutation -> integer conversion. (Not order preserving.)
Time: $\mathcal{O}(n)$
06f786, 6 lines

```
int permToInt(vector<int>& v) {
  int use = 0, i = 0, r = 0;
  for (auto &x : v) r=r* ++i +__builtin_popcount(use&-(1<<x)),
    use |= 1 << x;              // (note: minus, not ~!)
  return r;
}
```

6.1.2 Binomials

- Sum of every element in the n-th row of pascal triangle is 2^n.

- The product of the elements in each row is $\frac{(n+1)^n}{n!}$

- $\sum_{k=0}^n \binom{n}{k}^2 = \binom{2n}{n}$

- In a row p where p is a prime number, all the terms in that row except the 1s are multiples of p

- To count odd terms in row n, convert n to binary. Let x be the number of 1s in the binary representation. Then the number of odd terms will be 2^x

- Every entry in row $2^n - 1$ is odd

$$\binom{n}{k} = \frac{n}{k}\binom{n-1}{k-1}$$
$$\sum_{k=0}^{n}\binom{n}{k}k = n2^{n-1}$$
$$\sum_{m=0}^{n}\binom{m}{j}\binom{n-m}{k-j} = \binom{n+1}{k+1}$$
$$\binom{n}{k}\binom{n-k}{v} = \binom{n}{v}\binom{n-v}{k}$$
$$\sum_{j=0}^{m}\binom{m}{j}^2 = \binom{2m}{m}$$

$$(x+y)^n = \sum_{k=0}^{n}\binom{n}{k}x^{n-k}y^k$$
$$\sum_{j=0}^{k}\binom{m}{j}\binom{n-m}{k-j} = \binom{n}{k}$$
$$\sum_{r=0} m\binom{n+r}{r} = \binom{n+m+1}{m}$$
$$\sum_{m=k}^{n}\binom{m}{k} = \binom{n+1}{k+1}$$
$$2\sum_{i=L}^{R}\binom{n}{i} - \binom{n}{L} - \binom{n}{R} = \sum_{i=L+1}^{R}\binom{n+1}{i}$$

rolling-binomial.h
Description: $\binom{n}{k} \pmod{m}$ in time proportional to the difference between (n,k) and the previous (n,k).

`"../number-theory/preparator.h"`   de359e, 12 lines

```cpp
struct Bin {
  int N = 0,  K = 0; ll r = 1;
  void m(int a, int b) { r = r * a % mod * invs[b] % mod; }
  ll choose(int n, int k) {
    if (k > n || k < 0) return 0;
    while(N < n) ++N, m(N, N - K);
    while(K < k) ++K, m(N - K + 1, K);
    while(K > k) m(K, N - K + 1), --K;
    while(N > n) m(N - K, N), --N;
    return r;
  }
};
```

lucas.h
Description: Lucas' thm: Let n, m be non-negative integers and p a prime. Write $n = n_k p^k + \ldots + n_1 p + n_0$ and $m = m_k p^k + \ldots + m_1 p + m_0$. Then $\binom{n}{m} \equiv \prod_{i=0}^{k}\binom{n_i}{m_i} \pmod{p}$. fact and ifact must hold pre-computed factorials / inverse factorials, e.g. from ModInv.h.
Time: $\mathcal{O}\left(\log_p m\right)$

multinomial.h
Description: Computes $\binom{k_1 + \cdots + k_n}{k_1, k_2, \ldots, k_n} = \frac{(\sum k_i)!}{k_1! k_2! \ldots k_n!}$.   864cdb, 7 lines

```cpp
lint multinomial(vector<int>& v) {
  lint c = 1, m = v.empty() ? 1 : v[0];
  for (int i = 1 < v.size(); ++i)
    for (int j = 0; j < v[i]; ++j)
      c = c * ++m / (j+1);
  return c;
}
```

6.1.3 Cycles
Let the number of n-permutations whose cycle lengths all belong to the set S be denoted by $g_S(n)$

$$\sum_{n=0}^{\infty} g_S(n)\frac{x^n}{n!} = \exp\left(\sum_{n \in S}\frac{x^n}{n}\right)$$

6.1.4 Involutions
An involution is a permutation with maximum cycle length 2, and it is its own inverse.

$$a(n) = a(n-1) + (n-1)a(n-2)$$

$$a(0) = a(1) = 1$$

1, 1, 2, 4, 10, 26, 76, 232, 764, 2620, 9496, 35696, 140152

6.1.5 Inclusion-Exclusion Principle
Let $A_1, A_2, \ldots, A_n$ be finite sets. Then $A_1 \cup A_2 \cup \ldots \cup A_n$ is

$$\left|\bigcup_{i=1}^{n} A_i\right| = \sum_{\substack{I \subseteq \{1,2,\ldots,n\} \\ I \neq \emptyset}} (-1)^{|I|+1}\left|\bigcap_{i \in I} A_i\right|$$

6.1.6 The twelvefold way (from Stanley)
How many functions $f: N \to X$ are there?

N	X	Any f	Injective	Surjective
dist.	dist.	x^n	$\frac{x!}{(x-n)!}$	$x!\left\{{n \atop x}\right\}$
indist.	dist.	$\binom{x+n-1}{n}$	$\binom{x}{n}$	$\binom{n-1}{n-x}$
dist.	indist.	$\left\{{n \atop 1}\right\} + \ldots + \left\{{n \atop x}\right\}$	$[n \leq x]$	$\left\{{n \atop k}\right\}$
indist.	indist.	$p_1(n) + \ldots p_x(n)$	$[n \leq x]$	$p_x(n)$

Where $\binom{a}{b} = \frac{1}{b!}(a)_b$, $p_x(n)$ is the number of ways to partition the integer n using x summand and $\left\{{n \atop x}\right\}$ is the number of ways to partition a set of n elements into x subsets (aka Stirling number of the second kind).

6.1.7 Burnside
Given a group G of symmetries and a set X, the number of elements of X *up to symmetry* equals

$$\frac{1}{|G|}\sum_{g \in G}|X^g|,$$

where X^g are the elements fixed by g ($g.x = x$).

If $f(n)$ counts "configurations" (of some sort) of length n, we can ignore rotational symmetry using $G = \mathbb{Z}_n$ to get

$$g(n) = \frac{1}{n}\sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n}\sum_{k|n} f(k)\phi(n/k).$$

6.1.8 Derangements
Permutations of a set such that none of the elements appear in their original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor\frac{n!}{e}\right\rfloor$$

6.2 Partitions and subsets
6.2.1 Partition function
Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z}\backslash\{0\}} (-1)^{k+1} p(n - k(3k-1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

n	0	1	2	3	4	5	6	7	8	9	20	50	100
$p(n)$	1	1	2	3	5	7	11	15	22	30	627	~2e5	~2e8

6.3 General purpose numbers
6.3.1 Bernoulli numbers
EGF of Bernoulli numbers is $B(t) = \frac{t}{e^t - 1}$ (FFT-able).
$B[0, \ldots] = [1, -\frac{1}{2}, \frac{1}{6}, 0, -\frac{1}{30}, 0, \frac{1}{42}, \ldots]$

Sums of powers:

$$\sum_{i=1}^{n} n^m = \frac{1}{m+1}\sum_{k=0}^{m}\binom{m+1}{k}B_k(n+1)^{m+1-k}$$

Euler-Maclaurin formula for infinite sums:

$$\sum_{i=m}^{\infty} f(i) = \int_{m}^{\infty} f(x)dx - \sum_{k=1}^{\infty}\frac{B_k}{k!}f^{(k-1)}(m)$$

$$\approx \int_{m}^{\infty} f(x)dx + \frac{f(m)}{2} - \frac{f'(m)}{12} + \frac{f'''(m)}{720} + O(f^{(5)}(m))$$

6.3.2 Stirling numbers of the first kind
Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k), c(0,0) = 1$$

$$\sum_{k=0}^{n} c(n,k)x^k = x(x+1)\ldots(x+n-1)$$

$c(8,k) = 8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1$
$c(n,2) = 0, 0, 1, 3, 11, 50, 274, 1764, 13068, 109584, \ldots$

6.3.3 Eulerian numbers
Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, $k+1$ j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k}(-1)^j\binom{n+1}{j}(k+1-j)^n$$

6.3.4 Stirling numbers of the second kind
Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!}\sum_{j=0}^{k}(-1)^{k-j}\binom{k}{j}j^n$$

6.3.5 Bell numbers
Total number of partitions of n distinct elements. $B(n) = 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \ldots$.

$$\mathcal{B}_{n+1} = \sum_{k=0}^{n}\binom{n}{k}\mathcal{B}_k$$

Also possible to calculate using Stirling numbers of the second kind,

$$B_n = \sum_{k=0}^{n} S(n,k)$$

If p is prime:

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

6.3.6 Labeled unrooted trees

on n vertices: n^{n-2}
on k existing trees of size n_i: $n_1 n_2 \cdots n_k n^{k-2}$
with degrees d_i: $(n-2)!/((d_1-1)! \cdots (d_n-1)!)$ # forests with exactly k rooted trees:

$$\binom{n}{k} k \cdot n^{n-k-1}$$

.

6.3.7 Catalan numbers

$$C_n = \frac{1}{n+1}\binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \; C_{n+1} = \frac{2(2n+1)}{n+2}C_n, \; C_{n+1} = \sum C_i C_{n-i}$$

$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \ldots$

- sub-diagonal monotone paths in a $n \times n$ grid.
- strings with n pairs of parenthesis, correctly nested.
- binary trees with with $n+1$ leaves (0 or 2 children) or $2n+1$ elements.
- ordered trees with $n+1$ vertices.
- # ways a convex polygon with $n+2$ sides can be cut into triangles by connecting vertices with straight lines.
- permutations of $[n]$ with no 3-term increasing subsequence.

6.3.8 Super Catalan numbers

The number of monotonic lattice paths of a nxn grid that do not touch the diagonal.

$$S(n) = \frac{3(2n-3)S(n-1) - (n-3)S(n-2)}{n}$$
$$S(1) = S(2) = 1$$

1, 1, 3, 11, 45, 197, 903, 4279, 20793, 103049, 518859

6.3.9 Motzkin numbers

Number of ways of drawing any number of nonintersecting chords among n points on a circle. Number of lattice paths from (0, 0) to $(n, 0)$ never going below the x-axis, using only steps NE, E, SE.

$$M(n) = \frac{3(n-1)M(n-2) + (2n+1)M(n-1)}{n+2}$$

$$M(0) = M(1) = 1$$

1, 1, 2, 4, 9, 21, 51, 127, 323, 835, 2188, 5798, 15511, 41835, 113634

6.3.10 Narayana numbers

Number of lattice paths from (0,0) to $(2n,0)$ never going below the x-axis, using only steps NE and SE, and with k peaks.

$$N(n,k) = \frac{1}{n}\binom{n}{k}\binom{n}{k-1}$$

$$N(n,1) = N(n,n) = 1$$

$$\sum_{k=1}^{n} N(n,k) = C_n$$

1, 1, 1, 1, 3, 1, 1, 6, 6, 1, 1, 10, 20, 10, 1, 1, 15, 50

6.3.11 Schroder numbers

Number of lattice paths from (0, 0) to (n, n) using only steps N,NE,E, never going above the diagonal. Number of lattice paths from (0, 0) to $(2n, 0)$ using only steps NE, SE and double east EE, never going below the x-axis. Twice the Super Catalan number, except for the first term.
1, 2, 6, 22, 90, 394, 1806, 8558, 41586, 206098

6.3.12 Triangles

Given rods of length 1, ..., n,

$$T(n) = \frac{1}{24} \left\{ \begin{array}{ll} n(n-2)(2n-5) & n \text{ even} \\ (n-1)(n-3)(2n-1) & n \text{ odd} \end{array} \right\}$$

is the number of distinct triangles (positive are) that can be constructed, i.e., the # of 3-subsets of $[n]$ s.t. $x \le y \le z$ and $z \ne x + y$.

6.4 Fibonacci

$Fib(x+y) = Fib(x+1)Fib(y) + Fib(x)Fib(y-1)$

$Fib(n+1)Fib(n-1) - Fib(n)^2 = (-1)^n$

$Fib(2n-1) = Fib(n)^2 - Fib(n-1)^2$

$\sum_{i=0}^{n} Fib(i) = Fib(n+2) - 1$

$\sum_{i=0}^{n} Fib(i)^2 = Fib(n)Fib(n+1)$

$\sum_{i=0}^{n} Fib(i)^3 = \frac{Fib(n)Fib(n+1)^2 - (-1)^n Fib(n-1)+1}{2}$

6.5 Linear Recurrences

(i) $F_n = F_{n-1} + F_{n-2}$

$$\begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}\begin{bmatrix} F_0 \\ F_1 \end{bmatrix} = \begin{bmatrix} F_1 \\ F_2 \end{bmatrix}$$

(ii) $F_i = \sum_{j=1}^{K} C_j F_{i-j}$

$$\begin{bmatrix} 0 & 1 & 0 & 0 & . & 0 \\ 0 & 0 & 1 & 0 & . & 0 \\ 0 & 0 & 0 & 1 & . & 0 \\ . & . & . & . & . & . \\ C_K & C_{K-1} & C_{K-2} & C_{K-3} & . & C_1 \end{bmatrix}\begin{bmatrix} F_0 \\ F_1 \\ F_2 \\ . \\ F_{K-1} \end{bmatrix} = \begin{bmatrix} F_1 \\ F_2 \\ F_3 \\ . \\ F_K \end{bmatrix}$$

(iii) $F_i = \sum_{j=1}^{K} C_j F_{i-j} + D$

$$\begin{bmatrix} 0 & 1 & 0 & 0 & . & 0 & 0 \\ 0 & 0 & 1 & 0 & . & 0 & 0 \\ 0 & 0 & 0 & 1 & . & 0 & 0 \\ . & . & . & . & . & . & . \\ C_K & C_{K-1} & C_{K-2} & C_{K-3} & . & C_1 & 1 \\ 0 & 0 & 0 & 0 & . & 0 & 1 \end{bmatrix}\begin{bmatrix} F_0 \\ F_1 \\ F_2 \\ . \\ F_{K-1} \\ D \end{bmatrix} = \begin{bmatrix} F_1 \\ F_2 \\ F_3 \\ . \\ F_K \\ D \end{bmatrix}$$

6.6 Game Theory

A game can be reduced to Nim if it is a finite impartial game. Nim and its variants include:

6.6.1 Nim

Let $X = \bigoplus_{i=1}^{n} x_i$, then $(x_i)_{i=1}^{n}$ is a winning position iff $X \ne 0$. Find a move by picking k such that $x_k > x_k \oplus X$.

6.6.2 Misère Nim

Regular Nim, except that the last player to move *loses*. Play regular Nim until there is only one pile of size larger than 1, reduce it to 0 or 1 such that there is an odd number of piles. The second player wins $(a_1, \ldots, a_n)$ if 1) there is a pile $a_i > 1$ and $\oplus_{i=1}^{n} a_i = 0$ or 2) all $a_i \le 1$ and $\oplus_{i=1}^{n} a_i = 1$.

6.6.3 Staircase Nim

Stones are moved down a staircase and only removed from the last pile. $(x_i)_{i=1}^{n}$ is an L-position if $(x_{2i-1})_{i=1}^{n/2}$ is (i.e. only look at odd-numbered piles).

6.6.4 Moore's Nim$_k$

The player may remove from at most k piles (Nim = Nim$_1$). Expand the piles in base 2, do a carry-less addition in base $k+1$ (i.e. the number of ones in each column should be divisible by $k+1$).

6.6.5 Dim$^+$

The number of removed stones must be a divisor of the pile size. The Sprague-Grundy function is $k+1$ where 2^k is the largest power of 2 dividing the pile size.

6.6.6 Aliquot Game

Same as above, except the divisor should be proper (hence 1 is also a terminal state, but watch out for size 0 piles). Now the Sprague-Grundy function is just k.

6.6.7 Nim (at most half)

Write $n+1 = 2^m y$ with m maximal, then the Sprague-Grundy function of n is $(y-1)/2$.

6.6.8 Lasker's Nim

Players may alternatively split a pile into two new non-empty piles. $g(4k+1) = 4k+1$, $g(4k+2) = 4k+2$, $g(4k+3) = 4k+4$, $g(4k+4) = 4k+3$ ($k \ge 0$).

6.6.9 Hackenbush on Trees

A tree with stalks $(x_i)_{i=1}^n$ may be replaced with a single stalk with length $\bigoplus_{i=1}^n x_i$.

nim-product.cpp

Description: Product of nimbers is associative, commutative, and distributive over addition (xor). Forms finite field of size 2^{2^k}. Application: Given 1D coin turning games G_1, G_2 $G_1 \times G_2$ is the 2D coin turning game defined as follows. If turning coins at $x_1, x_2, \ldots, x_m$ is legal in G_1 and $y_1, y_2, \ldots, y_n$ is legal in G_2, then turning coins at all positions (x_i, y_j) is legal assuming that the coin at (x_m, y_n) goes from heads to tails. Then the grundy function $g(x, y)$ of $G_1 \times G_2$ is $g_1(x) \times g_2(y)$.
Time: 64^2 xors per multiplication, memorize to speed up.   *f55947, 24 lines*

```
ull nim_prod[64][64];
ull nim_prod2(int i, int j) {
    if (nim_prod[i][j]) return nim_prod[i][j];
    if ((i & j) == 0) return nim_prod[i][j] = 1ull << (i|j);
    int a = (i&j) & -(i&j);
    return nim_prod[i][j] = nim_prod2(i ^ a, j) ^ nim_prod2((i ^
        a) | (a-1), (j ^ a) | (i & (a-1)));
}
void all_nim_prod() {
    for (int i = 0; i < 64; i++)
        for (int j = 0; j < 64; j++)
            if ((i & j) == 0) nim_prod[i][j] = 1ull << (i|j);
            else {
                int a = (i&j) & -(i&j);
                nim_prod[i][j] = nim_prod[i ^ a][j] ^ nim_prod[(i ^ a)
                    | (a-1)][(j ^ a) | (i & (a-1))];
            }
}
ull get_nim_prod(ull x, ull y) {
    ull res = 0;
    for (int i = 0; i < 64 && (x >> i); ++i)
        if ((x >> i) & 1)
            for (int j = 0; j < 64 && (y >> j); ++j)
                if ((y >> j) & 1) res ^= nim_prod2(i, j);
    return res;
}
```

partitions.h

   3af1e7, 14 lines

```
vector<int64_t> prep(int N) {
    vector<int64_t> dp(N); dp[0] = 1;
    for (int n = 1; n < N; ++n) {
        int64_t sum = 0;
        for (int k = 0, l = 1, m = n - 1; ;) {
            sum += dp[m]; if ((m -= (k += 1)) < 0) break;
            sum += dp[m]; if ((m -= (l += 2)) < 0) break;
            sum -= dp[m]; if ((m -= (k += 1)) < 0) break;
            sum -= dp[m]; if ((m -= (l += 2)) < 0) break;
        }
        if ((sum %= M) < 0) sum += M;
        dp[n] = sum;
    } return dp;
}
```

Graph (7)

7.1 Fundamentals

bellman-ford.h

Description: Calculates shortest paths from s in a graph that might have negative edge weights. Unreachable nodes get dist = inf; nodes reachable through negative-weight cycles get dist = -inf. Assumes $V^2 \max |w_i| < \sim 2^{63}$.
Time: $\mathcal{O}(VE)$ *d37b03, 37 lines*

```
const lint inf = LLONG_MAX;
struct edge_t { int a, b, w, s() { return a < b ? a : -a; }};
struct node_t { lint dist = inf; int prev = -1; };

void bellmanFord(vector<node_t>& nodes, vector<edge_t>& eds,
        int s) {
    nodes[s].dist = 0;
    sort(eds.begin(), eds.end(), [](edge_t a, edge_t b) { return
        a.s() < b.s(); });
    int lim = nodes.size() / 2 + 2; // /3+100 with shuffled
        vertices
    for(int i = 0; i < lim; ++i) for(auto &ed : eds) {
        node_t cur = nodes[ed.a], &dest = nodes[ed.b];
        if (abs(cur.dist) == inf) continue;
        lint d = cur.dist + ed.w;
        if (d < dest.dist) {
            dest.prev = ed.a;
            dest.dist = (i < lim-1 ? d : -inf);
        }
    }
    for(int i = 0; i < lim; ++i) for(auto &e : eds)
        if (nodes[e.a].dist == -inf) nodes[e.b].dist = -inf;
}
vector<int> negCyc(int n, vector<edge_t>& edges) {
    vector<int64_t> d(n); vector<int> p(n);
    int v = -1;
    for (int i = 0; i < n; ++i) {
        v = -1;
        for (edge_t &u : edges)
            if (d[u.b] > d[u.a] + u.w) {
                d[u.b] = d[u.a] + u.w;
                p[u.b] = u.a, v = u.b;
            }
        if (v == -1) return {};
    }
    for (int i = 0; i < n; ++i) v = p[v]; // enter cycle
    vector<int> cycle = {v};
    while (p[cycle.back()] != v) cycle.push_back(p[cycle.back()]);
    return {cycle.rbegin(), cycle.rend()};
}
```

floyd-warshall.h

Description: Calculates all-pairs shortest path in a directed graph that might have negative edge distances. Input is an distance matrix m, where $m[i][j] = $ inf if i and j are not adjacent. As output, $m[i][j]$ is set to the shortest distance between i and j, inf if no path, or -inf if the path goes through a negative-weight cycle.
Time: $\mathcal{O}(N^3)$ *578e31, 16 lines*

```
const lint inf = 1LL << 62;
void floydWarshall(vector<vector<lint>>& m) {
    int n = m.size();
    for (int i = 0; i < n; ++i) m[i][i] = min(m[i][i], {});
    for (int k = 0; k < n; ++k)
        for (int i = 0; i < n; ++i)
            for (int j = 0; j < n; ++j)
                if (m[i][k] != inf && m[k][j] != inf) {
                    auto newDist = max(m[i][k] + m[k][j], -inf);
                    m[i][j] = min(m[i][j], newDist);
                }
    for (int k = 0; k < n; ++k) if (m[k][k] < 0)
        for (int i = 0; i < n; ++i)
            for (int j = 0; j < n; ++j)
                if (m[i][k] != inf && m[k][j] != inf) m[i][j] = -inf;
}
```

dijkstra.h

Description: Faster implementation of Dijkstra's algorithm. Makes very easy to handle SSSP on state graphs.
Time: $\mathcal{O}(N \log N)$ *67beaf, 31 lines*

```
#include<bits/extc++.h> // keep-include!!

template <class D> struct MinDist {
    vector<D> dist; vector<int> from;
};
template <class D, class E> // Weight type and Edge info
MinDist<D> Dijkstra(const vector<vector<E>>& g, int s, D inf =
        numeric_limits<D>::max()) {
    int N = int(g.size());
    vector<D> dist = vector<D>(N, inf);
    vector<int> par = vector<int>(N);
    struct state_t {
        D key;
        int to;
        bool operator<(state_t r) const { return key > r.key; }
    };
    __gnu_pbds::priority_queue<state_t> q;
    q.push(state_t{0, s});
    dist[s] = D(0);
    while (!q.empty()) {
        state_t p = q.top(); q.pop();
        if (dist[p.to] < p.key) continue;
        for (E nxt : g[p.to]) {
            if (p.key + nxt.second < dist[nxt.first]) {
                dist[nxt.first] = p.key + nxt.second;
                par[nxt.first] = p.to;
                q.push(state_t{dist[nxt.first], nxt.first});
            }
        }
    }
    return MinDist<D>{dist, par};
}
```

euler-walk.h

Description: Eulerian undirected/directed path/cycle algorithm. Input should be a vector of (dest, global edge index), where for undirected graphs, forward/backward edges have the same index. Returns a list of nodes in the Eulerian path/cycle with src at both start and end, or empty list if no cycle/path exists. To get edge indices back, add .second to s and ret.
Time: $\mathcal{O}(V + E)$ *643df6, 14 lines*

```
vector<int> eulerWalk(vector<vector<pii>>& gr, int nedges, int
        src=0) {
    int n = gr.size();
    vector<int> D(n), its(n), eu(nedges), ret, s = {src};
    D[src]++; // to allow Euler paths, not just cycles
    while (!s.empty()) {
        int x = s.back(), y, e, &it = its[x], end = int(gr[x].size
            ());
        if (it == end){ ret.push_back(x); s.pop_back(); continue; }
        tie(y, e) = gr[x][it++];
        if (!eu[e])
            D[x]--, D[y]++, eu[e] = 1, s.push_back(y);
    }
    for(auto &x : D) if (x < 0 || int(ret.size()) != nedges+1)
        return {};
    return {ret.rbegin(), ret.rend()};
}
```

7.2   Network flow

push-relabel.h
Description: Push-relabel using the highest label selection rule and the gap heuristic. Quite fast in practice. To obtain the actual flow, look at positive values only. id can be used to restore each edge and its amount of flow used.
Time: $\mathcal{O}\left(V^2\sqrt{E}\right)$ Better for dense graphs - Slower than Dinic (in practice)

387a0c, 42 lines

```cpp
template<typename flow_t = int> struct PushRelabel {
  struct edge_t { int dest, back; flow_t f, c; };
  vector<vector<edge_t>> g;
  vector<flow_t> ec;
  vector<edge_t*> cur;
  vector<vector<int>> hs; vector<int> H;
  PushRelabel(int n) : g(n), ec(n), cur(n), hs(2*n), H(n) {}
  void addEdge(int s, int t, flow_t cap, flow_t rcap = 0) { //
      d58501
    if (s == t) return;
    g[s].push_back({t, (int)g[t].size(), 0, cap});
    g[t].push_back({s, (int)g[s].size()-1, 0, rcap});
  }
  void addFlow(edge_t& e, flow_t f) { // 2f7969
    edge_t &back = g[e.dest][e.back];
    if (!ec[e.dest] && f) hs[H[e.dest]].push_back(e.dest);
    e.f += f; e.c -= f; ec[e.dest] += f;
    back.f -= f; back.c += f; ec[back.dest] -= f;
  }
  flow_t maxflow(int s, int t) { // 21100c
    int v = int(g.size()); H[s] = v; ec[t] = 1;
    vector<int> co(2*v); co[0] = v-1;
    for(int i = 0; i < v; ++i) cur[i] = g[i].data();
    for(auto e : g[s]) addFlow(e, e.c);
    for (int hi = 0;;) {
      while (hs[hi].empty()) if (!hi--) return -ec[s];
      int u = hs[hi].back(); hs[hi].pop_back();
      while (ec[u] > 0)   // discharge u
        if (cur[u] == g[u].data() + g[u].size()) {
          H[u] = 1e9;
          for(auto &e : g[u]) if (e.c && H[u] > H[e.dest]+1)
            H[u] = H[e.dest]+1, cur[u] = &e;
          if (++co[H[u]], !--co[hi] && hi < v)
            for(int i = 0; i < v; ++i) if (hi < H[i] && H[i] <
                v)
              --co[H[i]], H[i] = v + 1;
          hi = H[u];
        } else if (cur[u]->c && H[u] == H[cur[u]->dest]+1)
          addFlow(*cur[u], min(ec[u], cur[u]->c));
        else ++cur[u];
    }
  }
  bool leftOfMinCut(int a) { return H[a] >= int(g.size()); }
};
```

dinitz.h
Description: Flow algorithm with complexity $O(VE\log U)$ where $U = \max|cap|$. $O(\min(E^{1/2}, V^{2/3})E)$ if $U = 1$; $O(\sqrt{V}E)$ for bipartite matching. To obtain each partition A and B of the cut look at lvl, for $v \subset A$, $lvl[v] > 0$, for $u \subset B$, $lvl[u] = 0$.

de8d6b, 44 lines

```cpp
template<typename T = int> struct Dinitz {
  struct edge_t { int to, rev; T c, f; };
  vector<vector<edge_t>> adj;
  vector<int> lvl, ptr, q;
  Dinitz(int n) : lvl(n), ptr(n), q(n), adj(n) {}
  inline void addEdge(int a, int b, T c, T rcap = 0) {
    adj[a].push_back({b, (int)adj[b].size(), c, 0});
    adj[b].push_back({a, (int)adj[a].size() - 1, rcap, 0});
  }
  T dfs(int v, int t, T f) { // hash-1
```

```cpp
    if (v == t || !f) return f;
    for (int &i = ptr[v]; i < int(adj[v].size()); ++i) {
      edge_t &e = adj[v][i];
      if (lvl[e.to] == lvl[v] + 1)
        if (T p = dfs(e.to, t, min(f, e.c - e.f))) {
          e.f += p, adj[e.to][e.rev].f -= p;
          return p;
        }
    }
    return 0;
  } // hash-1 = 8ffe6b
  T maxflow(int s, int t) { // hash-2
    T flow = 0; q[0] = s;
    for (int L = 0; L < 31; ++L) do { // consider L = 30
      lvl = ptr = vector<int>(q.size());
      int qi = 0, qe = lvl[s] = 1;
      while (qi < qe && !lvl[t]) {
        int v = q[qi++];
        for (edge_t &e : adj[v])
          if (!lvl[e.to] && (e.c - e.f) >> (30 - L))
            q[qe++] = e.to, lvl[e.to] = lvl[v] + 1;
      }
      while (T p =dfs(s, t, numeric_limits<T>::max()/4))flow+=p;
    } while (lvl[t]);
    return flow;
  } // hash-2 = db2141
  bool leftOfMinCut(int v) { return bool(lvl[v] != 0); }
  auto minCut(int s, int t) { // hash-3
    T cost = maxflow(s,t);
    vector<edge_t> cut;
    for (int i = 0; i < int(adj.size()); i++) for(edge_t &e :
        adj[i])
      if (lvl[i] && !lvl[e.to]) cut.push_back(e);
    return make_pair(cost, cut);
  } // hash-3 = 1843d5
};
```

min-cost-max-flow.h
Description: Min-cost max-flow. Assumes there is no negative cycle.
Time: $\mathcal{O}\left(F(V + E)\log V\right)$, being F the amount of flow.

36ddeb, 57 lines

```cpp
template<class flow_t, class cost_t> struct min_cost {
  static constexpr flow_t FLOW_EPS = flow_t(1e-10);
  static constexpr flow_t FLOW_INF = numeric_limits<flow_t>::
      max();
  static constexpr cost_t COST_EPS = cost_t(1e-10);
  static constexpr cost_t COST_INF = numeric_limits<cost_t>::
      max();
  int n, m{}; vector<int> ptr, nxt, zu;
  vector<flow_t> capa; vector<cost_t> cost;
  min_cost(int N) : n(N),ptr(n,-1),dist(n),vis(n),pari(n) {}
  void add_edge(int u, int v, flow_t w, cost_t c) {
    nxt.push_back(ptr[u]); zu.push_back(v); capa.push_back(w);
        cost.push_back(c); ptr[u] = m++;
    nxt.push_back(ptr[v]); zu.push_back(u); capa.push_back(0);
        cost.push_back(-c); ptr[v] = m++;
  }
  vector<cost_t> pot, dist; vector<bool> vis; vector<int> pari;
  vector<flow_t> flows; vector<cost_t> slopes;
  // You can pass t = -1 to find a shortest
  void shortest(int s, int t) {//path to each vertex. // hash-1
    using E = pair<cost_t, int>;
    priority_queue<E, vector<E>, greater<E>> que;
    for(int u = 0; u < n; ++u){dist[u]=COST_INF; vis[u]=false;}
    for (que.emplace(dist[s] = 0, s); !que.empty(); ) {
      const cost_t c = que.top().first;
      const int u = que.top().second; que.pop();
      if (vis[u]) continue;
      vis[u] = true; if (u == t) return;
```

```cpp
      for (int i = ptr[u]; ~i; i = nxt[i]) if (capa[i] >
          FLOW_EPS) {
        const int v = zu[i];
        const cost_t cc = c + cost[i] + pot[u] - pot[v];
        if(dist[v] > cc){que.emplace(dist[v]=cc,v);pari[v]=i;}
      }
    }
  } // hash-1 = 89f16a
  auto run(int s, int t, flow_t limFlow = FLOW_INF) { // hash-2
    pot.assign(n, 0); flows = {0}; slopes.clear();
    while (true) {
      bool upd = false;
      for (int i = 0; i < m; ++i) if (capa[i] > FLOW_EPS) {
        const int u = zu[i ^ 1], v = zu[i];
        const cost_t cc = pot[u] + cost[i];
        if(pot[v] > cc + COST_EPS) { pot[v] = cc; upd = true; }
      } if (!upd) break;
    }
    flow_t flow = 0; cost_t tot_cost = 0;
    while (flow < limFlow) {
      shortest(s, t); flow_t f = limFlow - flow;
      if (!vis[t]) break;
      for(int u = 0; u < n; ++u)pot[u] += min(dist[u],dist[t]);
      for (int v = t; v != s; ) { const int i = pari[v];
        if (f > capa[i]) { f = capa[i]; } v = zu[i^1];
      }
      for (int v = t; v != s; ) { const int i = pari[v];
        capa[i] -= f; capa[i^1] += f; v = zu[i^1];
      }
      flow += f; tot_cost += f * (pot[t] - pot[s]);
      flows.push_back(flow); slopes.push_back(pot[t] - pot[s]);
    } return make_pair(flow, tot_cost);
  } // hash-2 = 285527
};
```

7.3   Matching

hopcroft-karp.h
Description: Fast bipartite matching algorithm. Graph g should be a list of neighbors of the left partition, and $btoa$ should be a vector full of -1's of the same size as the right partition. Returns the size of the matching. $btoa[i]$ will be the match for vertex i on the right side, or -1 if it's not matched.
Usage: vector<int> btoa(m, -1); hopcroftKarp(g, btoa);
Time: $\mathcal{O}\left(\sqrt{V}E\right)$

d9a55d, 35 lines

```cpp
using vi = vector<int>;
bool dfs(int a, int L, const vector<vi> &g, vi &btoa, vi &A, vi
    &B) {
  if (A[a] != L) return 0;
  A[a] = -1;
  for(auto &b : g[a]) if (B[b] == L + 1) {
    B[b] = 0;
    if (btoa[b] == -1 || dfs(btoa[b], L+1, g, btoa, A, B))
      return btoa[b] = a, 1;
  }
  return 0;
}
int hopcroftKarp(const vector<vi> &g, vi &btoa) {
  int res = 0;
  vector<int> A(g.size()), B(int(btoa.size())), cur, next;
  for (;;) {
    fill(A.begin(), A.end(), 0), fill(B.begin(), B.end(), 0);
    cur.clear();
    for(auto a : btoa) if (a != -1) A[a] = -1;
    for (int a = 0; a < g.size(); ++a) if (A[a] == 0) cur.
        push_back(a);
    for (int lay = 1;; ++lay) {
      bool islast = 0; next.clear();
      for(auto &a : cur) for(auto &b : g[a]) {
```

```
        if (btoa[b] == -1) B[b] = lay, islast = 1;
        else if (btoa[b] != a && !B[b])
          B[b] = lay, next.push_back(btoa[b]);
      }
      if (islast) break;
      if (next.empty()) return res;
      for(auto &a : next) A[a] = lay;
      cur.swap(next);
    }
    for(int a = 0; a < int(g.size()); ++a)
      res += dfs(a, 0, g, btoa, A, B);
  }
}
```

bipartite-matching.h
Description: Fast Kuhn! Simple maximum cardinality bipartite matching algorithm. Better than hopcroftKarp in practice. Worst case is $O(VE)$ on an hairy tree. Shuffling the edges and vertices ordering should break some worst-case inputs.
Time: $\Omega(VE)$

                                               1b4d72, 31 lines

```cpp
struct bm_t {
  int N, M, T;
  vector<vector<int>> adj;
  vector<int> match, seen;
  bm_t(int a, int b) : N(a), M(a+b), T(0), adj(M),
  match(M, -1), seen(M, -1) {}
  void add_edge(int a, int b) { adj[a].push_back(b + N); }
  bool dfs(int cur) {
    if (seen[cur] == T) return false;
    seen[cur] = T;
    for (int nxt : adj[cur]) if (match[nxt] == -1) {
      match[nxt] = cur, match[cur] = nxt;
      return true;
    }
    for (int nxt : adj[cur]) if (dfs(match[nxt])) {
      match[nxt] = cur, match[cur] = nxt;
      return true;
    }
    return false;
  }
  int solve() {
    int res = 0;
    for (int cur = 1; cur; ) {
      cur = 0; ++T;
      for (int i = 0; i < N; ++i) if (match[i] == -1)
        cur += dfs(i);
      res += cur;
    }
    return res;
  }
};
```

weighted-matching.h
Description: Given a weighted bipartite graph, matches every node on the left with a node on the right such that no nodes are in two matchings and the sum of the edge weights is minimal. Takes cost[N][M], where cost[i][j] = cost for L[i] to be matched with R[j] and returns (min cost, match), where L[i] is matched with R[match[i]]. Negate costs for max cost.
Time: $\mathcal{O}\left(N^2 M\right)$

                                          f0ea90, 28 lines

```cpp
pair<int, vector<int>> hungarian(const vector<vector<int>> &a){
  if (a.empty()) return {0, {}};
  int n = a.size() + 1, m = a[0].size() + 1;
  vector<int> u(n), v(m), p(m), ans(n - 1);
  for(int i = 1; i < n; ++i) {
    p[0] = i; int j0 = 0; // add "dummy" worker 0
    vector<int> dist(m, INT_MAX), pre(m, -1);
    vector<bool> done(m + 1);
    do {
      done[j0] = true;
      int i0 = p[j0], j1, delta = INT_MAX;
      for(int j = 1; j < m; ++j) if (!done[j]) {
        auto cur = a[i0-1][j-1] - u[i0] - v[j];
        if (cur < dist[j]) dist[j] = cur, pre[j] = j0;
        if (dist[j] < delta) delta = dist[j], j1 = j;
      }
      for(int j = 0; j < m; ++j)
        if (done[j]) u[p[j]] += delta, v[j] -= delta;
        else dist[j] -= delta;
      j0 = j1;
    } while (p[j0]);
    while (j0) { // update alternating path
      int j1 = pre[j0]; p[j0] = p[j1], j0 = j1;
    }
  }
  for(int j = 1; j < m; ++j) if (p[j]) ans[p[j]-1] = j-1;
  return {-v[0], ans}; // min cost
}
```

general-matching.h
Description: Maximum Matching for general graphs (undirected and non bipartite) using Edmond's Blossom Algorithm.
Time: $\mathcal{O}\left(EV^2\right)$

                                          e5db8e, 47 lines

```cpp
struct blossom_t {
  int N, M; vector<vector<int>> adj;
  vector<int> match, ts, ps; vector<array<int, 2>> fs;
  blossom_t(auto& G) : N(int(G.size())), M(0), adj(G), match(N
  ,-1), ts(N,-1), ps(N,-1), fs(N,{-1,-1}) {}
  int root(int a) {
    return (ts[a] != M || !~ps[a]) ? a : (ps[a] = root(ps[a]));
  }
  void rematch(int a, int b) {
    const int w = match[a]; match[a] = b; auto [x, y] = fs[a];
    if (~w && match[w] == a) {
      if (~y) rematch(x, y), rematch(y, x);
      else match[w] = x, rematch(x, w);
    }
  }
  bool augment(int src) {
    vector<int> bfs = {src}; bfs.reserve(N);
    ts[src] = M; ps[src] = -1; fs[src] = {-1,-1};
    for (int z = 0; z < int(bfs.size()); ++z) {
      int cur = bfs[z];
      for (int nxt : adj[cur]) if (nxt != src) {
        if (match[nxt] == -1) {
          match[nxt] = cur; rematch(cur, nxt); return true;
        }
        if (ts[nxt] == M) {
          int a = root(cur), b = root(nxt), m = src;
          if (a == b) continue;
          while (a != src || b != src) {
            if (b != src) swap(a, b);
            if (fs[a][0]==cur&&fs[a][1]==nxt) { m = a; break; }
            fs[a] = {cur, nxt}; a = root(fs[match[a]][0]);
          }
          for (const int r : {root(cur), root(nxt)})
            for (int v = r; v != m; v = root(fs[match[v]][0]))
              ts[v] = M, ps[v] = m, bfs.push_back(v);
        } else if (ts[match[nxt]] != M) {
          fs[nxt] = {-1, -1}; ts[match[nxt]] = M;
          ps[match[nxt]] = nxt; fs[match[nxt]] = {cur, -1};
          bfs.push_back(match[nxt]);
        }
      }
    }
    return false;
  }
}
```

max-independent-set.h
Description: To obtain a maximum independent set of a graph, find a max clique of the complement. If the graph is bipartite, see MinimumVertex-Cover.

min-vertex-cover.h
Description: Finds a minimum vertex cover in a bipartite graph. The size is the same as the size of a maximum matching, and the complement is a maximum independent set.
"bipartite-matching.h"   31f695, 20 lines

```cpp
vector<int> cover(bm_t& B, int N, int M) {
  int ma = B.solve();
  vector<bool> lfound(N, true), seen(N+M);
  for (int i = N; i < N+M; ++i) if (B.match[i] != -1)
    lfound[B.match[i]] = false;
  vector<int> q, cover;
  for (int i = 0; i < N; ++i) if (lfound[i]) q.push_back(i);
  while (!q.empty()) {
    int v = q.back(); q.pop_back();
    lfound[v] = true;
    for(int e : B.adj[v]) if (!seen[e] && B.match[e] != -1) {
      seen[e] = true;
      q.push_back(B.match[e]);
    }
  }
  for (int i = 0; i < N; ++i) if (!lfound[i]) cover.push_back(i
    );
  for (int i = N; i < N+M; ++i) if (seen[i]) cover.push_back(i)
    ;
  assert(cover.size() == ma);
  return cover;
}
```

min-edge-cover.h
Description: Finds a minimum edge cover in a bipartite graph. The size is the same as the number of vertices minus the size of a maximum matching. The mark vector represents who the vertices of set B has an edge to.
Usage: vector<int> mark(n+m, -1);
auto cover = minEdgeCover(g, mark, n, m);
"bipartite-matching.h"   ad86e9, 13 lines

```cpp
vector<pair<int,int>> minEdgeCover(bm_t& g, vector<int>& mark,
    int N, int M) {
  int ma = g.solve();
  vector<pair<int,int>> cover;
  for (int i = 0; i < N; ++i) {
    if (g.match[i] >= 0) cover.push_back({i, g.match[i]-N});
    else if (int(g.adj[i].size()))
      cover.push_back({i, g.adj[i][0] - N});
  }
  for (int i = N; i < N + M; ++i)
    if (g.match[i] == -1 && mark[i] >= 0)
      cover.push_back({mark[i], i - N});
  return cover;
}
```

min-path-cover.h
Description: Finds a minimum vertex-disjoint path cover in a dag. The size is the same as the number of vertices minus the size of a maximum matching.
"bipartite-matching.h"   212c5c, 15 lines

```cpp
vector<vector<int>> minPathCover(bm_t& g, int N) {
```

```cpp
int run() {
  for(int v = 0; v < N; ++v) if(!~match[v]) M += augment(v);
  return M;
}
};
```

```cpp
    int how_many = int(g.adj.size()) - g.solve();
    vector<vector<int>> paths;
    for (int i = 0; i < N; ++i)
      if (g.match[i + N] == -1) {
        vector<int> path = {i};
        int cur = i;
        while (g.match[cur] >= 0) {
          cur = g.match[cur] - N;
          path.push_back(cur);
        }
        paths.push_back(path);
      }
    return paths;
}
```

7.4 DFS algorithms

dfs-tree.h
Description: Builds dfs tree. Find cut vertices and bridges.
Usage: Call solve right after build the graph
<div align="right">72963d, 52 lines</div>

```cpp
struct tree_t {
  int n, timer;
  vector<vector<int>> edges;
  vector<pair<int,int>> bridges;
  vector<int> parent, mindepth, depth, st, child;
  vector<bool> cut;
  tree_t(int N) : n(N), timer(0), edges(n), parent(n,-1),
  mindepth(n,-1), depth(n,-1), st(n,-1) {}
  void addEdge(int a, int b) {
    edges[a].push_back(b); edges[b].push_back(a);
  }
  void dfs(int v) {
    st[v] = timer;
    mindepth[v] = depth[v];
    for (int u : edges[v]) {
      if (u == parent[v]) continue;
      if (st[u] == timer) {
        mindepth[v] = min(mindepth[v], depth[u]);
        continue;
      }
      depth[u] = 1 + depth[v];
      parent[u] = v;
      dfs(u);
      mindepth[v] = min(mindepth[v], mindepth[u]);
    }
  }
  vector<pair<int,int>> find_bridges() {
    for (int i = 0; i < n; ++i)
      if (parent[i] != -1 && mindepth[i] == depth[i])
        bridges.push_back({parent[i], i});
    return bridges;
  }
  vector<bool> find_cut() {
    cut.resize(n), child.resize(n);
    for (int i = 0; i < n; ++i)
      if (parent[i] != -1 && mindepth[i] >= depth[parent[i]])
        cut[parent[i]] = 1;
    for (int i = 0; i < n; ++i)
      if (parent[i] != -1) child[parent[i]]++;
    for (int i = 0; i < n; ++i)
      if (parent[i] == -1 && child[i] < 2) cut[i] = 0;
    return cut;
  }
  void solve() {
    for (int i = 0; i < n; ++i)
      if (depth[i] == -1) {
        depth[i] = 0; parent[i] = -1;
        ++timer;
        dfs(i);
```

```cpp
      }
    }
  }
};
```

centroid-decomposition.h
Description: Divide and Conquer on Trees.
<div align="right">dd21a1, 65 lines</div>

```cpp
template<typename T> struct centroid_t {
  int N;
  vector<vector<int>> adj;
  vector<vector<int>> dist; // dist to all ancestors
  vector<bool> blocked; // processed centroid
  vector<int> sz, depth, parent; // centroid parent
  centroid_t(int _n) : N(_n), adj(_n), dist(32 - __builtin_clz(
    _n), vector<int>(_n)),
  blocked(_n), sz(_n), depth(_n), parent(_n) {}
  void add_edge(int a, int b) {
    adj[a].push_back(b); adj[b].push_back(a);
  }
  void dfs_sz(int cur, int prv) {
    sz[cur] = 1;
    for (int nxt : adj[cur]) {
      if (nxt == prv || blocked[nxt]) continue;
      dfs_sz(nxt, cur); sz[cur] += sz[nxt];
    }
  }
  int find(int cur, int prv, int tsz) {
    for (int nxt : adj[cur])
      if (!blocked[nxt] && nxt != prv && 2*sz[nxt] > tsz)
        return find(nxt, cur, tsz);
    return cur;
  }
  void dfs_dist(int cur, int prv, int layer, int d) {
    dist[layer][cur] = d;
    for (int nxt : adj[cur]) {
      if (blocked[nxt] || nxt == prv) continue;
      dfs_dist(nxt, cur, layer, d + 1);
    }
  }
  void get_path(int cur, int prv, int d,vector<int>& cur_path){
    cur_path.push_back(d);
    for (int nxt : adj[cur]) {
      if (nxt == prv || blocked[nxt]) continue;
      get_path(nxt, cur, d + 1, cur_path);
    }
  }
  // solve for each subtree (cnt := # of paths of length K
  // that goes through vertex cur)
  T solve_subtree(int cur, int prv, int K) {
    vector<T> dp(sz[prv] + 1); dp[0] = 1;
    T cnt = 0;
    for (int nxt : adj[cur]) {
      if (blocked[nxt]) continue;
      vector<int> path; get_path(nxt, cur, 1, path);
      for (int d : path) {
        if (d > K || K - d > sz[prv]) continue;
        cnt += dp[K - d];
      }
      for (int d : path) dp[d] += 1;
    } return cnt;
  }
  T decompose(int cur, int K, int layer=0, int prv_root = -1){
    dfs_sz(cur, -1);
    int root = find(cur, cur, sz[cur]);
    blocked[root] = true; depth[root] = layer;
    parent[root] = prv_root; dfs_dist(root, root, layer, 0);
    T res = solve_subtree(root, cur, K);
    for (int nxt : adj[root]) {
      if (blocked[nxt]) continue;
```

```cpp
      res += decompose(nxt, K, layer + 1, root);
    } return res;
  }
};
```

tarjan.h
Description: Finds all strongly connected components in a directed graph.
Usage: scc_t s(g); s.solve([&](const vector<int>& cc) {...});
visits all components in reverse topological order.
Time: $\mathcal{O}(E + V)$
<div align="right">50f8c4, 29 lines</div>

```cpp
struct scc_t {
  int n, t, scc_num;
  vector<vector<int>> adj;
  vector<int> low, id, stk, in_stk, cc_id;
  scc_t(const vector<vector<int>>& g) : n(int(g.size())), t
    (0), scc_num(0),
  adj(g), low(n,-1), id(n,-1), in_stk(n, false), cc_id(n) {}
  template<class F> void dfs(int cur, F f) {
    id[cur] = low[cur] = t++;
    stk.push_back(cur); in_stk[cur] = true;
    for (int nxt : adj[cur])
      if (id[nxt] == -1)
        dfs(nxt, f), low[cur] = min(low[cur], low[nxt])
          ;
      else if (in_stk[nxt])
        low[cur] = min(low[cur], id[nxt]);
    if (low[cur] == id[cur]) {
      vector<int> cc; cc.reserve(stk.size());
      while (true) {
        int v = stk.back(); stk.pop_back();
        in_stk[v] = false;
        cc.push_back(v); cc_id[v] = scc_num;
        if (v == cur) break;
      } f(cc); scc_num++;
    }
  }
  template<class F> void solve(F f) {
    stk.reserve(n);
    for (int r = 0; r < n; ++r) if (id[r] == -1) dfs(r, f);
  }
};
```

kosaraju.h
Description: Kosaraju's Algorithm, DFS twice to generate strongly connected components in topological order. a, b in same component if both $a \to b$ and $b \to a$ exist.
Time: $\mathcal{O}(V + E)$
<div align="right">25be07, 35 lines</div>

```cpp
struct Kosaraju_t {
  int n;
  vector<vector<int>> edges, redges;
  vector<bool> seen;
  vector<int> cnt_of, cnts;
  Kosaraju_t(const int &N) : n(N), edges(N), redges(N), seen(N)
    , cnt_of(N, -1) {}
  void addEdge(int a, int b) {
    edges[a].push_back(b);
    redges[b].push_back(a);
  }
  void dfs(int v) {
    seen[v] = true;
    for (int u : edges[v]) {
      if (seen[u]) continue;
      dfs(u);
    }
    toposort.push_back(v);
  }
  void dfs_fix(int v, int w) {
```

```cpp
        cnt_of[v] = x;
        for (int u : redges[v]) {
            if (cnt_of[u] == -1) dfs_fix(u, w);
        }
    }
    void solve() {
        for (int i = 0; i < n; ++i)
            if (seen[i] == false) dfs(i);
        reverse(toposort.begin(), toposort.end());
        for (int u : toposort) {
            if (cnt_of[u] != -1) continue;
            dfs_fix(u, u);
            cnts.push_back(u);
        }
    }
};
```

bcc.h

Description: Finds all biconnected components in an undirected graph. In a biconnected component there are at least two distinct paths between any two nodes or the component is a bridge. Note that a node can be in several components. *blockcut* constructs the block cut tree of given graph. The first nodes represents the blocks, the others represents the articulation points.
Usage: int e_id = 0; vector<pair<int, int>> g(N);
for (auto [a,b] : edges) {
g[a].emplace_back(b, e_id);
g[b].emplace_back(a, e_id++); }
bcc_t b(g); b.solve([&](const vector<int>& edges_id) {...});
Time: $\mathcal{O}(E + V)$

f05480, 50 lines

```cpp
struct bcc_t{
    int n, t;
    vector<vector<pii>> adj;
    vector<int> low, id, stk, is_art;
    bcc_t(const vector<vector<pii>> &g) : n(int(g.size())),
    t(0), adj(g), low(n,-1), id(n,-1), is_art(n) {}
    template<class F> void dfs(int cur, int e_par, F f){
        id[cur] = low[cur] = t++;
        stk.push_back(e_par); int c = 0;
        for (auto [nxt, e_id] : adj[cur]) {
            if (id[nxt] == -1) {
                dfs(nxt, e_id, f);
                low[cur] = min(low[cur], low[nxt]); c++;
                if (low[nxt] < id[cur]) continue;
                is_art[cur] = true;
                auto top =find(stk.rbegin(), stk.rend(), e_id);
                vector<int> cc(stk.rbegin(), next(top));
                f(cc); stk.resize(stk.size()-cc.size());
            }
            else if (e_id != e_par) {
                low[cur] = min(low[cur], id[nxt]);
                if (id[nxt] < id[cur]) stk.push_back(e_id);
            }
        } if(e_par == -1) is_art[cur] = (c > 1) ? true : false;
    }
    template<class F> void solve(F f) {
        stk.reserve(n);
        for(int r = 0; r < n; ++r) if(id[r] == -1) dfs(r,-1,f);
    }
    auto blockcut(const vector<vector<pii>> &edges){
        vector<vector<int>> cc; vector<int> cc_id(n,-1);
        solve( [&](const vector<int> &c) {
            set<int> vc;
            for(int e : c){
                auto [a, b] = edges[e];
                cc_id[a] = cc_id[b] = int(cc.size());
                vc.insert(a); vc.insert(b);
            } cc.emplace_back(vc.begin(), vc.end());
        } );
```

```cpp
        for(int a = 0; a < n; a++) if(is_art[a])
            cc_id[a] = int(cc.size()), cc.push_back({a});
        int bcc_num = int(cc.size());
        vector<vector<int>> tree(bcc_num);
        for(int c = 0; c < bcc_num && 1<int(cc[c].size()); ++c)
            for(int a : cc[c]) if(is_art[a]) {
                tree[c].push_back(cc_id[a]);
                tree[cc_id[a]].push_back(c);
            } return make_tuple(cc_id, cc, tree);
    }
};
```

2sat.h

Description: Calculates a valid assignment to boolean variables a, b, c,... to a 2-SAT problem, so that an expression of the type $(a\|\|b)\&\&(!a\|\|c)\&\&(d\|\|\|!b)\&\&...$ becomes true, or reports that it is unsatisfiable. Negated variables are represented by bit-inversions ($\sim$x).
Usage: TwoSat ts(number of boolean variables);
ts.either(0, ~3); // Var 0 is true or var 3 is false
ts.set_value(2); // Var 2 is true
ts.at_most_one({0,~1,2}); // <= 1 of vars 0, ~1 and 2 are true
ts.solve(); // Returns true iff it is solvable
ts.values[0..N-1] holds the assigned values to the vars
Time: $\mathcal{O}(N + E)$, where N is the number of boolean variables, and E is the number of clauses.

"tarjan.h" 74e46b, 36 lines

```cpp
struct TwoSat {
    int N;
    vector<vector<int>> gr;
    vector<int> values; // 0 = false, 1 = true
    TwoSat(int n = 0) : N(n), gr(2*n) {}
    int add_var() { // (optional)
        gr.emplace_back(); gr.emplace_back();
        return N++;
    }
    void either(int f, int j) {
        f = max(2*f, -1-2*f); j = max(2*j, -1-2*j);
        gr[f].push_back(j^1); gr[j].push_back(f^1);
    }
    void implies(int f, int j) { either(~f, j); }
    void set_value(int x) { either(x, x); }
    void at_most_one(const vector<int>& li) { // (optional)
        if (int(li.size()) <= 1) return;
        int cur = ~li[0];
        for (int i = 2; i < int(li.size()); ++i) {
            int next = add_var();
            either(cur, ~li[i]); either(cur, next);
            either(~li[i], next); cur = ~next;
        } either(cur, ~li[1]);
    }
    bool solve() {
        scc s(gr);
        s.solve([](const vector<int> &v){ return; } );
        values.assign(N, -1);
        for (int i = 0; i < N; ++i)
            if (s.cc_id[2*i] == s.cc_id[2*i+1]) return 0;
        for (int i = 0; i < N; ++i)
            if (s.cc_id[2*i] < s.cc_id[2*i+1]) values[i] =
                false;
            else values[i] = true;
        return 1;
    }
};
```

7.5   Heuristics

maximal-cliques.h

Description: Runs a callback for all maximal cliques in a graph (given as a symmetric bitset matrix; self-edges not allowed). Possible optimization: on the top-most recursion level, ignore 'cands', and go through nodes in order of increasing degree, where degrees go down as nodes are removed.
Time: $\mathcal{O}\left(3^{n/3}\right)$, much faster for sparse graphs

57e107, 11 lines

```cpp
typedef bitset<128> B;
template<class F>
void cliques(vector<B> &eds, F f, B P = ~B(), B X={}, B R={}) {
    if (!P.any()) { if (!X.any()) f(R); return; }
    auto q = (P | X)._Find_first();
    auto cands = P & ~eds[q];
    for(int i = 0; i < eds.size(); ++i) if (cands[i]) {
        R[i] = 1; cliques(eds, f, P & eds[i], X & eds[i], R);
        R[i] = P[i] = 0; X[i] = 1;
    }
}
```

maximum-clique.h

Description: Finds a maximum clique of a graph (given as symmetric bitset matrix; self-edges not allowed). Can be used to find a maximum independent set by finding a clique of the complement graph.
Time: Runs in about 1s for n=155 and worst case random graphs (p=.90). Runs faster for sparse graphs.

261d2e, 49 lines

```cpp
using vb = vector<bitset<40>>;
struct Maxclique {
    double limit = 0.025, pk = 0;
    struct Vertex { int i, d = 0; };
    using vv = vector<Vertex>;
    vb e;
    vv V;
    vector<vector<int>> C;
    vector<int> qmax, q, S, old;
    void init(vv& r) {
        for(auto& v : r) v.d = 0;
        for(auto& v : r) for(auto& j : r) v.d += e[v.i][j.i];
        sort(r.begin(), r.end(), [](auto a, auto b) { return a.d >
            b.d; });
        int mxD = r[0].d;
        for(int i = 0; i < int(r.size()); ++i) r[i].d = min(i, mxD)
            + 1;
    }
    void expand(vv& R, int lev = 1) {
        S[lev] += S[lev - 1] - old[lev];
        old[lev] = S[lev - 1];
        while (int(R.size())) {
            if (int(q.size()) + R.back().d <= int(qmax.size()))
                return;
            q.push_back(R.back().i);
            vv T;
            for(auto& v : R) if (e[R.back().i][v.i]) T.push_back({v.i
                });
            if (int(T.size())) {
                if (S[lev]++ / ++pk < limit) init(T);
                int j = 0, mxk = 1, mnk = max(int(qmax.size()) - int(q.
                    size()) + 1, 1);
                C[1].clear(), C[2].clear();
                for(auto& v : T) {
                    int k = 1;
                    auto f = [&](int i) { return e[v.i][i]; };
                    while (any_of(C[k].begin(), C[k].end(), f)) k++;
                    if (k > mxk) mxk = k, C[mxk + 1].clear();
                    if (k < mnk) T[j++].i = v.i;
                    C[k].push_back(v.i);
                }
                if (j > 0) T[j - 1].d = 0;
                for(int k = mnk; k <= mxk; ++k) for(int i : C[k])
                    T[j].i = i, T[j++].d = k;
```

```cpp
        expand(T, lev + 1);
      } else if (int(q.size()) > int(qmax.size())) qmax = q;
      q.pop_back(), R.pop_back();
    }
  }
  vector<int> maxClique() { init(V), expand(V); return qmax; }
  Maxclique(vb conn) : e(conn), C(int(e.size())+1), S(int(C.
    size())), old(S) {
    for(int i = 0; i < int(e.size()); ++i) V.push_back({i});
  }
};
```

chromatic-number.h
Description: Compute the chromatic number of a graph. Minimum number of colors needed to paint the graph in a way s.t. if two vertices share an edge, they must have distinct colors.

Time: $\mathcal{O}\left(N2^N\right)$

```cpp
template<class T> int min_colors(int N, const T& gr) {
  vector<int> adj(N);
  for (int a = 0; a < N; ++a)
    for (int b = a + 1; b < N; ++b) {
      if (!gr[a][b]) continue;
      adj[a] |= (1 << b); adj[b] |= (1 << a);
    }
  static vector<unsigned> dp(1 << N), buf(1 << N), w(1 << N);
  for (int mask = 0; mask < (1 << N); ++mask) {
    bool ok = true;
    for (int i = 0; i < N; ++i) if (mask & 1 << i)
      if (adj[i] & mask) ok = false;
    if (ok) dp[mask]++;
    buf[mask] = 1;
    w[mask] = __builtin_popcount(mask) % 2 == N % 2 ? 1 : -1;
  }
  for (int i = 0; i < N; ++i)
    for (int mask = 0; mask < (1 << N); ++mask)
      if (!(mask & 1 << i)) dp[mask^(1 << i)] += dp[mask];
  for (int colors = 1; colors <= N; ++colors) {
    unsigned S = 0;
    for (int mask = 0; mask < (1 << N); ++mask)
      S += (buf[mask] *= dp[mask]) * w[mask];
    if (S) return colors;
  } assert(false);
}
```

cycle-counting.cpp
Description: Counts 3 and 4 cycles

Time: $\mathcal{O}\left(E\sqrt{E}\right)$

```cpp
using vi = vector<int>;
int count_cycles(const vector<vi>& adj, const vi& deg) {
  const int N = int(adj.size());
  vi idx(N), loc(N); iota(idx.begin(), idx.end(), 0);
  sort(idx.begin(), idx.end(), [&](const int& a, const int& b)
    { return deg[a] < deg[b]; });
  for (int i = 0; i < N; ++i) loc[idx[i]] = i;
  vector<vi> gr(N);
  for (int a = 0; a < N; ++a) for (int b : adj[a])
    if (loc[a] < loc[b]) gr[a].push_back(b);
  int cycle3 = 0, cycle4 = 0;
  {
    vector<bool> seen(N, false);
    for (int a = 0; a < N; ++a) {
      for (int b : gr[a]) seen[b] = true;
      for (int b : gr[a]) for (int c : gr[b])
        if (seen[c]) cycle3 += 1;
      for (int b : gr[a]) seen[b] = false;
    }
  }
```

```cpp
  }
  {
    vi cnt(N);
    for (int a = 0; a < N; ++a) {
      for (int b : adj[a]) for (int c : gr[b])
        if (loc[a] < loc[c]) {
          cycle4 += cnt[c];
          cnt[c]++;
        }
      for (int b : adj[a]) for (int c : gr[b]) cnt[c] = 0;
    }
  } return cycle3;
}
```

edge-coloring.h
Description: Given a simple, undirected graph with max degree D, computes a $(D + 1)$-coloring of the edges such that no neighboring edges share a color. (D-coloring is NP-hard, but can be done for bipartite graphs by repeated matchings of max-degree nodes.)

Time: $\mathcal{O}(NM)$

```cpp
vector<int> misra_gries(int N, vector<pair<int, int>> eds) {
  const int M = int(eds.size());
  vector<int> cc(N + 1), ret(M), fan(N), free(N), loc;
  for (auto e : eds) ++cc[e.first], ++cc[e.second];
  int u, v, ncols = *max_element(cc.begin(), cc.end()) + 1;
  vector<vector<int>> adj(N, vi(ncols, -1));
  for (auto e : eds) {
    tie(u, v) = e; fan[0] = v; loc.assign(ncols, 0);
    int at = u, end = u, d, c = free[u], ind = 0, i = 0;
    while (d = free[v], !loc[d] && (v = adj[u][d]) != -1)
      loc[d] = ++ind, cc[ind] = d, fan[ind] = v;
    cc[loc[d]] = c;
    for (int cd = d; at != -1; cd ^= c ^ d, at = adj[at][cd])
      swap(adj[at][cd], adj[end = at][cd ^ c ^ d]);
    while (adj[fan[i]][d] != -1) {
      int left = fan[i], right = fan[++i], e = cc[i];
      adj[u][e] = left; adj[left][e] = u;
      adj[right][e] = -1; free[right] = e;
    }
    adj[u][d] = fan[i]; adj[fan[i]][d] = u;
    for (int y : {fan[0], u, end})
      for (int& z = free[y] = 0; adj[y][z] != -1; z++);
  }
  for (int i = 0; i < M; ++i)
    for (tie(u, v) = eds[i]; adj[u][ret[i]] != v;) ++ret[i];
  return ret;
}
```

7.6 Trees

lca-binary-lifting.h
Description: Solve lowest common ancestor queries using binary jumps. Can also find the distance between two nodes.

Time: $\mathcal{O}(N \log N + Q \log N)$

```cpp
struct lca_t {
  int logn{0}, preorderpos{0};
  vector<int> invpreorder, height;
  vector<vector<int>> jump, edges;
  lca_t(int n, vector<vector<int>>& adj) :
    edges(adj), height(n), invpreorder(n) {
    while((1<<(logn+1)) <= n) ++logn;
    jump.assign(n+1, vector<int>(logn+1, 0));
    dfs(0, -1, 0);
  }
  void dfs(int v, int p, int h) {
    invpreorder[v] = preorderpos++;
    height[v] = h;
    jump[v][0] = p < 0 ? v : p;
```

```cpp
    for (int l = 1; l <= logn; ++l)
      jump[v][l] = jump[jump[v][l-1]][l-1];
    for (int u : edges[v]) {
      if (u == p) continue;
      dfs(u, v, h + 1);
    }
  }
  int climb(int v, int dist) {
    for (int l = 0; l <= logn; ++l)
      if (dist&(1<<l)) v = jump[v][l];
    return v;
  }
  int query(int a, int b) {
    if (height[a] < height[b]) swap(a, b);
    a = climb(a, height[a] - height[b]);
    if (a == b) return a;
    for (int l = logn; l >= 0; --l)
      if (jump[a][l] != jump[b][l])
        a = jump[a][l], b = jump[b][l];
    return jump[a][0];
  }
  int dist(int a, int b) {
    return height[a] + height[b] - 2 * height[query(a,b)];
  }
  bool is_parent(int p, int v) {
    if (height[p] > height[v]) return false;
    return p == climb(v, height[v] - height[p]);
  }
  bool on_path(int x, int a, int b) {
    int v = query(a, b);
    return is_parent(v, x) && (is_parent(x, a) || is_parent(x,
      b));
  }
  int get_kth_on_path(int a, int b, int k) {
    int v = query(a, b);
    int x = height[a] - height[v], y = height[b] - height[v];
    if (k < x) return climb(a, k);
    else return climb(b, x + y - k);
  }
};
```

lca-euler-tour.h
Description: Data structure for computing lowest common ancestors and build Euler Tour in a tree. Edges should be an adjacency list of the tree, either directed or undirected.

Time: $\mathcal{O}(N \log N + Q + Q \log)$

```cpp
struct small_lca_t {
  int T = 0;
  vector<int> time, path, walk, depth;
  rmq_t<int> rmq;
  small_lca_t(vector<vector<int>> &edges) : time(int(edges.size
      ())),
    depth(time), rmq((dfs(edges,0,-1), walk)) {}
  void dfs(vector<vector<int>> &edges, int v, int p) {
    time[v] = T++;
    for(int u : edges[v]) {
      if (u == p) continue;
      depth[u] = depth[v] + 1;
      path.push_back(v), walk.push_back(time[v]);
      dfs(edges, u, v);
    }
  }
  int lca(int a, int b) {
    if (a == b) return a;
    tie(a, b) = minmax(time[a], time[b]);
    return path[rmq.query(a, b)];
  }
};
```

```cpp
struct lca_t {
  int N;
  vector<vector<int>> adj;
  vector<int> parent, depth, sz;
  vector<int> euler_tour, timer;
  vector<int> tour_in, tour_out, postorder;
  vector<int> idx, rev_idx;
  vector<int> heavy_root;
  rmq_t<pair<int,int>> rmq;
  int next_idx = 0, rev_next_idx = 0;
  bool built = false;
  lca_t() : N(0) {}
  lca_t(vector<vector<int>>& _adj, int root = -1, bool
      build_rmq = true) :
    N(int(_adj.size())), adj(_adj), parent(N, -1), depth(N), sz
        (N), timer(N),
    tour_in(N), tour_out(N), postorder(N), idx(N), heavy_root(N
        ),
    built(false) {

    if (0 <= root && root < N) pre_dfs(root, -1);

    euler_tour.reserve(2 * N);

    for (int i = 0; i < N; ++i)
      if (parent[i] == -1) {
        if (i != root) pre_dfs(i, -1);
        dfs(i, false);
        euler_tour.push_back(-1);
      }

    rev_idx = idx;
    reverse(rev_idx.begin(), rev_idx.end());
    assert(int(euler_tour.size()) == 2 * N);
    vector<pair<int, int>> euler_tour_depths;
    euler_tour_depths.reserve(euler_tour.size());

    int id = 0;
    for (int cur : euler_tour)
      euler_tour_depths.push_back({cur == -1 ? cur : depth[
          cur], id++});
    }

    if (build_rmq) rmq = rmq_t<pair<int, int>>(
        euler_tour_depths);
    built = true;
  }

  void pre_dfs(int cur, int par) {
    parent[cur] = par;
    depth[cur] = (par == -1 ? 0 : 1 + depth[par]);
    adj[cur].erase(remove(adj[cur].begin(), adj[cur].end(), par
        ), adj[cur].end());
    sz[cur] = 1;
    for (int nxt : adj[cur]) {
      pre_dfs(nxt, cur);
      sz[cur] += sz[nxt];
    }
    if (!adj[cur].empty()) {
      auto w = max_element(adj[cur].begin(), adj[cur].end(),
          [&](int a, int b) { return sz[a] < sz[b]; });
      swap(*adj[cur].begin(), *w);
    }
  }

  void dfs(int cur, bool heavy) {
    heavy_root[cur] = heavy ? heavy_root[parent[cur]] : cur;
    timer[cur] = int(euler_tour.size());
```

```cpp
    euler_tour.push_back(cur);
    idx[next_idx] = cur;
    tour_in[cur] = next_idx++;
    bool heavy_child = true;
    for (int next : adj[cur]) {
      dfs(next, heavy_child);
      euler_tour.push_back(cur);
      heavy_child = false;
    }
    tour_out[cur] = next_idx;
    postorder[cur] = rev_next_idx++;
  }

  pair<int, array<int, 2>> get_diameter() const {
    assert(built);
    pair<int, int> u_max = {-1, -1};
    pair<int, int> ux_max = {-1, -1};
    pair<int, array<int, 2>> uxv_max = {-1, {-1, -1}};
    for (int cur : euler_tour) {
      if (cur == -1) break;
      u_max = max(u_max, {depth[cur], cur});
      ux_max = max(ux_max, {u_max.first - 2 * depth[cur], u_max
          .second});
      uxv_max = max(uxv_max, {ux_max.first + depth[cur], {
          ux_max.second, cur}});
    }
    return uxv_max;
  }

  int query(int a, int b) const {
    if (a == b) return a;
    a = timer[a], b = timer[b];
    if (a > b) swap(a, b);
    return euler_tour[rmq.query(a, b).second];
  }

  bool is_ancestor(int a, int b) const {
    return tour_in[a] <= tour_in[b] && tour_in[b] < tour_out[a
        ];
  }

  bool on_path(int x, int a, int b) const {
    return (is_ancestor(x, a) || is_ancestor(x, b)) &&
        is_ancestor(query(a, b), x);
  }

  int dist(int a, int b) const {
    return depth[a] + depth[b] - 2 * depth[query(a, b)];
  }

  int child_ancestor(int a, int b) const {
    assert(a != b); assert(is_ancestor(a, b));
    // Note: this depends on rmq_t breaking ties by latest
    //     index.
    int child = euler_tour[rmq.query(timer[a], timer[b]).second
        + 1];
    assert(parent[child] == a);
    assert(is_ancestor(child, b));
    return child;
  }

  int get_kth_ancestor(int a, int k) const {
    while (a >= 0) {
      int root = heavy_root[a];
      if (depth[root] <= depth[a] - k) return idx[tour_in[a] -
          k];
      k -= depth[a] - depth[root] + 1;
      a = parent[root];
    }
```

```cpp
    return a;
  }

  int get_kth_node_on_path(int a, int b, int k) const {
    int lca = query(a, b);
    int x = depth[a] - depth[lca], y = depth[b] - depth[lca];
    assert(0 <= k && k <= x + y);
    if (k < x) return get_kth_ancestor(a, k);
    else return get_kth_ancestor(b, x + y - k);
  }

  int get_common_node(int a, int b, int c) const {
    // Return the deepest node among lca(a, b), lca(b, c), and
    //     lca(c, a).
    int x = query(a, b), y = query(b, c), z = query(c, a);
    x = depth[y] > depth[x] ? y : x;
    x = depth[z] > depth[x] ? z : x;
    return x;
  }
};
```

heavylight.h

Description: Compress Tree: Given a subset S of nodes, computes the compress tree and returns a list of edges representing a tree rooted at 0.
Time: $\mathcal{O}\left((\log N)^2\right)$

d3dcdc, 77 lines

```cpp
template<bool use_edges> struct hld_t {
  int N, T{};
  vector<vector<int>> adj;
  vector<int> sz, depth, chain, par, in, out, preorder;
  hld_t() {}
  hld_t(const vector<vector<int>>& G, int r = 0) : N(int(G.size
      ())),
    adj(G), sz(N), depth(N), chain(N), par(N), in(N), out(N),
    preorder(N) { dfs_sz(r); chain[r] = r; dfs_hld(r); }
  void dfs_sz(int cur) {
    sz[cur] = 1;
    for (auto& nxt : adj[cur]) {
      par[nxt] = cur; depth[nxt] = 1 + depth[cur];
      adj[nxt].erase(find(adj[nxt].begin(), adj[nxt].end(), cur
          ));
      dfs_sz(nxt); sz[cur] += sz[nxt];
      if (sz[nxt] > sz[adj[cur][0]]) swap(nxt, adj[cur][0]);
    }
  }
  void dfs_hld(int cur) {
    in[cur] = T++; preorder[in[cur]] = cur;
    for (auto& nxt : adj[cur]) {
      chain[nxt] = (nxt == adj[cur][0] ? chain[cur] : nxt);
      dfs_hld(nxt);
    } out[cur] = T;
  }
  int lca(int a, int b) {
    while (chain[a] != chain[b]) {
      if (in[a] < in[b]) swap(a, b);
      a = par[chain[a]];
    } return (in[a] < in[b] ? a : b);
  }
  bool is_ancestor(int a, int b) { return in[a] <= in[b] && in[
      b] < out[a]; }
  int climb(int a, int k) {
    if (depth[a] < k) return -1;
    int d = depth[a] - k;
    while (depth[chain[a]] > d) a = par[chain[a]];
    return preorder[in[a] - depth[a] + d];
  }
  int kth_on_path(int a, int b, int K) {
    int m = lca(a, b);
    int x = depth[a] - depth[m], y = depth[b] - depth[m];
```

```cpp
      if (K > x + y) return -1;
      return (x > K ? climb(a, K) : climb(b, x + y - K));
    }
    // bool is true if path should be reversed (only for
        noncommutative operations)
    const auto& get_path(int a, int b) const {
      static vector<tuple<bool, int, int>> L, R;
      L.clear(); R.clear();
      while (chain[a] != chain[b])
        if (depth[chain[a]] > depth[chain[b]]) {
          L.push_back({true, in[chain[a]], in[a] + 1});
          a = par[chain[a]];
        } else {
          R.push_back({false, in[chain[b]], in[b] + 1});
          b = par[chain[b]];
        }
      if (depth[a] > depth[b])
        L.push_back({true, in[b] + use_edges, in[a] + 1});
      else R.push_back({false, in[a] + use_edges, in[b] + 1});
      L.insert(L.end(), R.rbegin(), R.rend());
      return L;
    }
    auto get_subtree(int a) const {
      return make_pair(in[a] + use_edges, in[a] + sz[a]);
    }
    auto compress(vector<int> s){
      auto cmp = [&](int a, int b){ return in[a] < in[b]; };
      sort(s.begin(), s.end(), cmp); int m = int(s.size())-1;
      for (int i = 0; i < m; ++i)
        s.push_back(lca(s[i], s[i+1]));
      sort(s.begin(), s.end(), cmp);
      s.erase(unique(s.begin(), s.end()), s.end());
      vector<pii> ret = { {0, s[0]} };
      for (int i = 0; i + 1 < int(s.size()); ++i)
        ret.emplace_back(lca(s[i], s[i+1]), s[i+1]);
      return ret;
    }
};
```

tree-isomorphism.h
Time: $\mathcal{O}(N \log(N))$

```cpp
struct tree_t {
  vector<int> cen, sz;
  vector<vector<int>> adj;
  tree_t(vector<vector<int>>& g):cen(2), sz(g.size()), adj(g){}
  int dfs_sz(int v, int p) {
    sz[v] = 1;
    for (int u : adj[v]) if (u != p)
      sz[v] += dfs_sz(u, v);
    return sz[v];
  }
  int dfs(int tsz, int v, int p) {
    for (int u : adj[v]) if (u != p) {
      if (2*sz[u] <= tsz) continue;
      return dfs(tsz, u, v);
    } return cen[0] = v;
  }
  void find_cenroid(int v) {
    int tsz = dfs_sz(v, -1);
    cen[1] = dfs(tsz, v, -1);
    for (int u : adj[cen[0]]) if (2*sz[u] == tsz)
      cen[1] = u;
  }
  int hash_it(int v, int p = -1) {
    static map<vector<int>, int> val;
    vector<int> offset;
    for (int u : adj[v]) if (u != p)
      offset.push_back(hash_it(u, v));
```

```cpp
    sort(offset.begin(), offset.end());
    if (!val.count(offset))  val[offset] = int(val.size());
    return val[offset];
  }
  ll get_hash(int v = 0) {
    find_cenroid(v);
    ll x = hash_it(cen[0]), y=hash_it(cen[1]);
    if (x > y) swap(x, y);
    return (x << 30) + y;
  }
};
```

7.6.1 Sqrt Decomposition
HLD generally suffices. If not, here are some common
strategies:

1. Rebuild the tree after every $\sqrt{N}$ queries.
2. Consider vertices with $>$ or $< \sqrt{N}$ degree separately.
3. For subtree updates, note that there are $O(\sqrt{N})$ distinct sizes among child subtrees of any vertex.

Block Tree: Use a DFS to split edges into contiguous groups of size $\sqrt{N}$ to $2\sqrt{N}$. **Mo's Algorithm for Tree Paths:** Maintain an array of vertices where each one appears twice, once when a DFS enters the vertex (st) and one when the DFS exists (en). For a tree path $u \leftrightarrow v$ such that st[u]<st[v],

1. If u is an ancestor of v, query [st[u],st[v]].
2. Otherwise, query [en[u],st[v]] and consider $lca(u,v)$ separately.

7.7 Other
manhattan-mst.h
Description: Given N points, returns up to 4*N edges, which are guaranteed to contain a minimum spanning tree for the graph with edge weights $w(p,q) = |p.x - q.x| + |p.y - q.y|$. Edges are in the form (distance, src, dst). Use a standard MST algorithm on the result to find the final MST.
Time: $\mathcal{O}(NlogN)$

```cpp
typedef Point<int> P;
pair<vector<array<int, 3>>, int> manhattanMST(vector<P> ps) {
  vector<int> id(ps.size()); iota(id.begin(), id.end(), 0);
  vector<array<int, 3>> edges;
  for(int k = 0; k < 4; ++k) {
    sort(id.begin(), id.end(), [&](int i, int j) {
      return (ps[i]-ps[j]).x < (ps[j]-ps[i]).y;});
    map<int, int> sweep;
    for(auto& i : id) {
      for (auto it = sweep.lower_bound(-ps[i].y);
          it != sweep.end(); sweep.erase(it++)) {
        int j = it->second; P d = ps[i] - ps[j];
        if (d.y > d.x) break;
        edges.push_back({d.y + d.x, i, j});
      } sweep[-ps[i].y] = i;
    }
    if (k & 1) for(auto& p : ps) p.x = -p.x;
    else for(auto& p : ps) swap(p.x, p.y);
  }
  sort(edges.begin(), edges.end());
  UF uf(ps.size()); int cost = 0;
  for (auto e: edges) if (uf.unite(e[1], e[2])) cost += e[0];
  return {edges, cost};
}
```

directed-mst.h
Description: Edmonds' algorithm for finding the weight of the minimum spanning tree/arborescence of a directed graph, given a root node. If no MST exists, returns -1.
Time: $\mathcal{O}(E \log V)$

```cpp
struct Edge { int a, b; ll w; };
struct Node {
  Edge key;
  Node *l, *r; ll delta;
  void prop() {
    key.w += delta;
    if (l) l->delta += delta;
    if (r) r->delta += delta;
    delta = 0;
  }
  Edge top() { prop(); return key; }
};
Node *merge(Node *a, Node *b) {
  if (!a || !b) return a ?: b;
  a->prop(), b->prop();
  if (a->key.w > b->key.w) swap(a, b);
  swap(a->l, (a->r = merge(b, a->r)));
  return a;
}
void pop(Node*& a) { a->prop(); a = merge(a->l, a->r); }
auto dmst(int n, int r, vector<Edge>& g) { // hash-1
  RollbackUF uf(n);
  vector<Node*> heap(n);
  for (Edge e : g) heap[e.b] = merge(heap[e.b], new Node{e});
  ll res = 0;
  vector<int> seen(n, -1), path(n), par(n);
  seen[r] = r;
  vector<Edge> Q(n), in(n, {-1,-1}), comp;
  deque<tuple<int, int, vector<Edge>>> cycs;
  for(int s = 0; s < n; ++s) {
    int u = s, qi = 0, w;
    while (seen[u] < 0) {
      if (!heap[u]) return make_pair(-1L, vector<int>());
      Edge e = heap[u]->top();
      heap[u]->delta -= e.w, pop(heap[u]);
      Q[qi] = e, path[qi++] = u, seen[u] = s;
      res += e.w, u = uf.find(e.a);
      if (seen[u] == s) {
        Node* cyc = 0;
        int end = qi, time = uf.time();
        do cyc = merge(cyc, heap[w = path[--qi]]);
        while (uf.unite(u, w));
        u = uf.find(u), heap[u] = cyc, seen[u] = -1;
        cycs.push_front({u, time, {&Q[qi], &Q[end]}});
      }
    }
    for(int i = 0; i < qi; ++i) in[uf.find(Q[i].b)] = Q[i];
  }
  for (auto& [u,t,comp] : cycs) { // restore sol (optional)
    uf.rollback(t); Edge inEdge = in[u];
    for (auto& e : comp) in[uf.find(e.b)] = e;
    in[uf.find(inEdge.b)] = inEdge;
  }
  for(int i = 0; i < n; ++i) par[i] = in[i].a;
  return make_pair(res, par);
} // hash-1 = 222439
```

7.8 Theorems
7.8.1 Euler's theorem
Let V, A and F be the number of vertices, edges and faces of connected planar graph, $V - A + F = 2$.

7.8.2   Konig's theorem

In any bipartite graph the following holds: Maximum matching and minimum vertex cover have the same cardinality. Minimum edge cover, minimum path cover and maximum independent set have the same cardinality $N - |M|$.

7.8.3   Eulerian Cycles

The number of Eulerian cycles in a *directed* graph G is: $t_w(G) \prod_{v \in G} (\deg v - 1)!$, where $t_w(G)$ is the number of arborescences ("directed spanning" tree) rooted at w (Check Number of Spanning Trees)

7.8.4   Construct minimum vertex cover from maximum matching on bipartite graph

1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.
2. DFS from unmatched vertices in X.
3. $x \in X$ is chosen iff x is unvisited.
4. $y \in Y$ is chosen iff y is visited.

7.8.5   Minimum weight edge cover

1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight $w(u, v)$.
2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
3. Find the minimum weight perfect matching on G'.

7.8.6   Maximum/Minimum flow with lower bound / Circulation problem

1. Construct super source S and sink T.
2. For each edge (x, y, l, u), connect $x \to y$ with capacity $u - l$.
3. For each vertex v, denote by $in(v)$ the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
4. If $in(v) > 0$, connect $S \to v$ with capacity $in(v)$, otherwise, connect $v \to T$ with capacity $-in(v)$.

   - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
   - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.

5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.

7.8.7   Minimum cost cyclic flow

1. Consruct super source S and sink T
2. For each edge (x, y, c), connect $x \to y$ with $(cost, cap) = (c, 1)$ if $c > 0$, otherwise connect $y \to x$ with $(cost, cap) = (-c, 1)$
3. For each edge with $c < 0$, sum these cost as K, then increase $d(y)$ by 1, decrease $d(x)$ by 1

4. For each vertex v with $d(v) > 0$, connect $S \to v$ with $(cost, cap) = (0, d(v))$
5. For each vertex v with $d(v) < 0$, connect $v \to T$ with $(cost, cap) = (0, -d(v))$
6. Flow from S to T, the answer is the cost of the flow $C + K$

7.8.8   Maximum density induced subgraph

1. Binary search on answer, suppose we're checking answer T
2. Construct a max flow model, let K be the sum of all weights
3. Connect source $s \to v$, $v \in G$ with capacity K
4. For each edge (u, v, w) in G, connect $u \to v$ and $v \to u$ with capacity w
5. For $v \in G$, connect it with sink $v \to t$ with capacity $K + 2T - (\sum_{e \in E(v)} w(e)) - 2w(v)$
6. T is a valid answer if the maximum flow $f < K|V|$

7.8.9   Project selection problem

1. If $p_v > 0$, create edge (s, v) with capacity p_v; otherwise, create edge (v, t) with capacity $-p_v$.
2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.
3. The mincut is equivalent to the maximum profit of a subset of projects.

7.8.10   0/1 quadratic programming

$$\sum_x c_x x + \sum_y c_y \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x\bar{y} + x'\bar{y}')$$

can be minimized by the mincut of the following graph:

1. Create edge (x, t) with capacity c_x and create edge (s, y) with capacity c_y.
2. Create edge (x, y) with capacity c_{xy}.
3. Create edge (x, y) and edge (x', y') with capacity $c_{xyx'y'}$.

7.8.11   Number of Spanning Trees

Define Laplacian Matrix as $L = D - A$, D being a Diagonal Matrix with $D_{i,i} = deg(i)$ and A an Adjacency Matrix. Create an $N \times N$ Laplacian matrix `mat`, and for each edge $a \to b \in G$, do `mat[a][b]--`, `mat[b][b]++` (and `mat[b][a]--`, `mat[a][a]++` if G is undirected). Remove the ith row and column and take the determinant; this yields the number of directed spanning trees rooted at i (if G is undirected, remove any row/column).

7.8.12   Tutte Matrix

- A graph has a perfect matching iff the *Tutte* matrix has a non-zero determinant.

- The rank of the *Tutte* matrix is equal to twice the size of the maximum matching. The maximum cost matching can be found by polynomial interpolation.

Geometry (8)

8.1   Geometric primitives

Point.h
Description: Class to handle points in the plane. T can be e.g. double or long long. (Avoid int.)
f90ade, 27 lines

```cpp
template <class T> int sgn(T x) { return (x > 0) - (x < 0); }
template<class T>
struct Point {
  typedef Point P;
  T x, y;
  explicit Point(T x=0, T y=0) : x(x), y(y) {}
  bool operator<(P p) const { return tie(x,y) < tie(p.x,p.y); }
  bool operator==(P p) const { return tie(x,y)==tie(p.x,p.y); }
  P operator+(P p) const { return P(x+p.x, y+p.y); }
  P operator-(P p) const { return P(x-p.x, y-p.y); }
  P operator*(T d) const { return P(x*d, y*d); }
  P operator/(T d) const { return P(x/d, y/d); }
  T dot(P p) const { return x*p.x + y*p.y; }
  T cross(P p) const { return x*p.y - y*p.x; }
  T cross(P a, P b) const { return (a-*this).cross(b-*this); }
  T dist2() const { return x*x + y*y; }
  double dist() const { return sqrt((double)dist2()); }
  // angle to x-axis in interval [-pi, pi]
  double angle() const { return atan2(y, x); }
  P unit() const { return *this/dist(); } // makes dist()=1
  P perp() const { return P(-y, x); } // rotates +90 degrees
  P normal() const { return perp().unit(); }
  // returns point rotated 'a' radians ccw around the origin
  P rotate(double a) const {
    return P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a)); }
};
using P = Point<double>;
```

LineDistance.h
Description:
Returns the signed distance between point p and the line containing points a and b. Positive value on left side and negative on right as seen from a towards b. a==b gives nan. P is supposed to be Point<T> or Point3D<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long. Using Point3D will always give a non-negative distance.
"Point.h"
f6bf6b, 4 lines
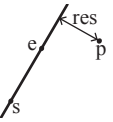
```cpp
template<class P>
double lineDist(const P& a, const P& b, const P& p) {
  return (double)(b-a).cross(p-a)/(b-a).dist();
}
```

SegmentDistance.h
Description:
Returns the shortest distance between point p and the line segment from point s to e.
Usage: Point<double> a, b(2,2), p(1,1);
bool onSegment = segDist(a,b,p) < 1e-10;
"Point.h"
ae751a, 5 lines
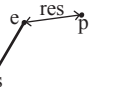
```cpp
double segDist(P& s, P& e, P& p) {
  if (s==e) return (p-s).dist();
  auto d = (e-s).dist2(), t = min(d,max(.0,(p-s).dot(e-s)));
  return ((p-s)*d-(e-s)*t).dist()/d;
}
```

SegmentIntersection.h

Description:
If a unique intersection point between the line segments going from s1 to e1 and from s2 to e2 exists then it is returned. If no intersection point exists an empty vector is returned. If infinitely many exist a vector with 2 elements is returned, containing the endpoints of the common line segment. The wrong position will be returned if P is Point<ll> and the intersection point does not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or long long.

Usage: `vector<P> inter = segInter(s1,e1,s2,e2);`
```
if (sz(inter)==1)
cout << "segments intersect at " << inter[0] << endl;
```
"Point.h", "OnSegment.h"   f6be16, 13 lines
```
template<class P> vector<P> segInter(P a, P b, P c, P d) {
    auto oa = c.cross(d, a), ob = c.cross(d, b),
         oc = a.cross(b, c), od = a.cross(b, d);
    // Checks if intersection is single non-endpoint point.
    if (sgn(oa) * sgn(ob) < 0 && sgn(oc) * sgn(od) < 0)
        return {(a * ob - b * oa) / (ob - oa)};
    set<P> s;
    if (onSegment(c, d, a)) s.insert(a);
    if (onSegment(c, d, b)) s.insert(b);
    if (onSegment(a, b, c)) s.insert(c);
    if (onSegment(a, b, d)) s.insert(d);
    return {s.begin(), s.end()};
}
```

SegmentIntersectionQ.h
Description: Like segmentIntersection, but only returns true/false. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or long long.
"Point.h"   1ff4ba, 16 lines
```
template<class P>
bool segmentIntersectionQ(P s1, P e1, P s2, P e2) {
    if (e1 == s1) {
        if (e2 == s2) return e1 == e2;
        swap(s1,s2); swap(e1,e2);
    }
    P v1 = e1-s1, v2 = e2-s2, d = s2-s1;
    auto a = v1.cross(v2), a1 = d.cross(v1), a2 = d.cross(v2);
    if (a == 0) { // parallel
        auto b1 = s1.dot(v1), c1 = e1.dot(v1),
             b2 = s2.dot(v1), c2 = e2.dot(v1);
        return !a1 && max(b1,min(b2,c2)) <= min(c1,max(b2,c2));
    }
    if (a < 0) { a = -a; a1 = -a1; a2 = -a2; }
    return (0 <= a1 && a1 <= a && 0 <= a2 && a2 <= a);
}
```

LineIntersection.h
Description:
If a unique intersection point of the lines going through s1,e1 and s2,e2 exists {1, point} is returned. If no intersection point exists {0, (0,0)} is returned and if infinitely many exists {-1, (0,0)} is returned. The wrong position will be returned if P is Point<ll> and the intersection point does not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or ll.
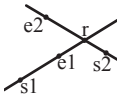
Usage: `auto res = lineInter(s1,e1,s2,e2);`
```
if (res.first == 1)
cout << "intersection point at " << res.second << endl;
```
"Point.h"   a01f81, 8 lines
```
template<class P>
pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
    auto d = (e1 - s1).cross(e2 - s2);
    if (d == 0) // if parallel
        return {-(s1.cross(e1, s2) == 0), P(0, 0)};
```

```
    auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
    return {1, (s1 * p + e1 * q) / d};
}
```

LineProjectionReflection.h
Description: Projects point p onto line ab. Set refl=true to get reflection of point p across line ab insted. The wrong point will be returned if P is an integer point and the desired point doesn't have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow.
"Point.h"   b5562d, 5 lines
```
template<class P>
P lineProj(P a, P b, P p, bool refl=false) {
    P v = b - a;
    return p - v.perp()*(1+refl)*v.cross(p-a)/v.dist2();
}
```

SideOf.h
Description: Returns where *p* is as seen from *s* towards *e*. 1/0/-1 ⇔ left/on line/right. If the optional argument *eps* is given 0 is returned if *p* is within distance *eps* from the line. P is supposed to be Point<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long.
Usage: `bool left = sideOf(p1,p2,q)==1;`
"Point.h"   3af81c, 9 lines
```
template<class P>
int sideOf(P s, P e, P p) { return sgn(s.cross(e, p)); }

template<class P>
int sideOf(const P& s, const P& e, const P& p, double eps) {
    auto a = (e-s).cross(p-s);
    double l = (e-s).dist()*eps;
    return (a > l) - (a < -l);
}
```

OnSegment.h
Description: Returns true iff p lies on the line segment from s to e. Use (segDist(s,e,p)<=epsilon) instead when using Point<double>.
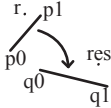"Point.h"   c597e8, 3 lines
```
template<class P> bool onSegment(P s, P e, P p) {
    return p.cross(s, e) == 0 && (s - p).dot(e - p) <= 0;
}
```

LinearTransformation.h
Description:
Apply the linear transformation (translation, rotation and scaling) which takes line p0-p1 to line q0-q1 to point r.

"Point.h"   03a306, 6 lines
```
typedef Point<double> P;
P linearTransformation(const P& p0, const P& p1,
    const P& q0, const P& q1, const P& r) {
    P dp = p1-p0, dq = q1-q0, num(dp.cross(dq), dp.dot(dq));
    return q0 + P((r-p0).cross(num), (r-p0).dot(num))/dp.dist2();
}
```

Angle.h
Description: A class for ordering angles (as represented by int points and a number of rotations around the origin). Useful for rotational sweeping. Sometimes also represents points or vectors.
Usage: `vector<Angle> v = {w[0], w[0].t360() ...};` // sorted
```
int j = 0; rep(i,0,n) { while (v[j] < v[i].t180()) ++j; }
// sweeps j such that (j-i) represents the number of positively
// oriented triangles with vertices at 0 and i
```
   0f0602, 34 lines
```
struct Angle {
    int x, y;
    int t;
```

```
    Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
    Angle operator-(Angle b) const { return {x-b.x, y-b.y, t}; }
    int half() const {
        assert(x || y);
        return y < 0 || (y == 0 && x < 0);
    }
    Angle t90() const { return {-y, x, t + (half() && x >= 0)}; }
    Angle t180() const { return {-x, -y, t + half()}; }
    Angle t360() const { return {x, y, t + 1}; }
};
bool operator<(Angle a, Angle b) {
    // add a.dist2() and b.dist2() to also compare distances
    return make_tuple(a.t, a.half(), a.y * (ll)b.x) <
        make_tuple(b.t, b.half(), b.x * (ll)b.y);
}
// Given two points, this calculates the smallest angle between
// them, i.e., the angle that covers the defined line segment.
pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
    if (b < a) swap(a, b);
    return (b < a.t180() ?
        make_pair(a, b) : make_pair(b, a.t360()));
}
Angle operator+(Angle a, Angle b) { // point a + vector b
    Angle r(a.x + b.x, a.y + b.y, a.t);
    if (a.t180() < r) r.t--;
    return r.t180() < a ? r.t360() : r;
}
Angle angleDiff(Angle a, Angle b) { // angle b - angle a
    int tu = b.t - a.t; a.t = b.t;
    return {a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x, tu - (b < a)};
}
```

AngleCmp.h
Description: Useful utilities for dealing with angles of rays from origin. OK for integers, only uses cross product. Doesn't support (0,0).
   6edd25, 22 lines
```
template <class P>
bool sameDir(P s, P t) {
    return s.cross(t) == 0 && s.dot(t) > 0;
}
// checks 180 <= s..t < 360?
template <class P>
bool isReflex(P s, P t) {
    auto c = s.cross(t);
    return c ? (c < 0) : (s.dot(t) < 0);
}
// operator < (s,t) for angles in [base,base+2pi)
template <class P>
bool angleCmp(P base, P s, P t) {
    int r = isReflex(base, s) - isReflex(base, t);
    return r ? (r < 0) : (0 < s.cross(t));
}
// is x in [s,t] taken ccw? 1/0/-1 for in/border/out
template <class P>
int angleBetween(P s, P t, P x) {
    if (sameDir(x, s) || sameDir(x, t)) return 0;
    return angleCmp(s, x, t) ? 1 : -1;
}
```

8.2 Circles

CircleIntersection.h
Description: Computes the pair of points at which two circles intersect. Returns false in case of no intersection.
"Point.h"   c64785, 10 lines
```
bool circleInter(P a,P b,double r1,double r2,pair<P, P>* out) {
    if (a == b) { assert(r1 != r2); return false; }
    P vec = b - a;
    double d2 = vec.dist2(), sum = r1+r2, dif = r1-r2,
```

```
      p = (d2 + r1*r1 - r2*r2)/(d2*2), h2 = r1*r1 - p*p*d2;
  if (sum*sum < d2 || dif*dif > d2) return false;
  P mid = a + vec*p, per = vec.perp() * sqrt(fmax(0, h2) / d2);
  *out = {mid + per, mid - per};
  return true;
}
```

CircleTangents.h
Description: Finds the external tangents of two circles, or internal if r2 is
negated. Can return 0, 1, or 2 tangents – 0 if one circle contains the other (or
overlaps it, in the internal case, or if the circles are the same); 1 if the circles
are tangent to each other (in which case .first = .second and the tangent line
is perpendicular to the line between the centers). .first and .second give the
tangency points at circle 1 and 2 respectively. To find the tangents of a circle
with a point set r2 to 0.

"Point.h"   b0153d, 13 lines
```cpp
template<class P>
vector<pair<P, P>> tangents(P c1, double r1, P c2, double r2) {
  P d = c2 - c1;
  double dr = r1 - r2, d2 = d.dist2(), h2 = d2 - dr * dr;
  if (d2 == 0 || h2 < 0)  return {};
  vector<pair<P, P>> out;
  for (double sign : {-1, 1}) {
    P v = (d * dr + d.perp() * sqrt(h2) * sign) / d2;
    out.push_back({c1 + v * r1, c2 + v * r2});
  }
  if (h2 == 0) out.pop_back();
  return out;
}
```

Circumcircle.h
Description:
The circumcirle of a triangle is the circle intersecting all
three vertices. ccRadius returns the radius of the circle
going through points A, B and C and ccCenter returns
the center of the same circle.
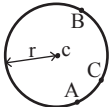
"Point.h"   c98102, 8 lines
```cpp
double ccRadius(const P& A, const P& B, const P& C) {
  return (B-A).dist()*(C-B).dist()*(A-C).dist()/
    abs((B-A).cross(C-A))/2;
}
P ccCenter(const P& A, const P& B, const P& C) {
  P b = C-A, c = B-A;
  return A + (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;
}
```

MinimumEnclosingCircle.h
Description: Computes the minimum circle that encloses a set of points.
Time: expected $\mathcal{O}(n)$

"circumcircle.h"   8ab87f, 19 lines
```cpp
pair<P, double> mec(vector<P> ps) {
  shuffle(ps.begin(),ps.end(), mt19937(time(0)));
  P o = ps[0];
  double r = 0, EPS = 1 + 1e-8;
  for(int i = 0; i < ps.size(); ++i)
    if ((o - ps[i]).dist() > r * EPS) {
      o = ps[i], r = 0;
      for(int j = 0; j < i; ++j) if ((o - ps[j]).dist() > r *
          EPS) {
        o = (ps[i] + ps[j]) / 2;
        r = (o - ps[i]).dist();
        for(int k = 0; k < j; ++k)
          if ((o - ps[k]).dist() > r * EPS) {
            o = ccCenter(ps[i], ps[j], ps[k]);
            r = (o - ps[i]).dist();
          }
      }
    }
```

CircleUnion.h
Description: Computes the circles union total area
      fd65da, 86 lines
```cpp
struct CircleUnion {
    static const int maxn = 1e5 + 5;
    const double PI = acos((double)-1.0);
    int n;
    double x[maxn], y[maxn], r[maxn];
    int covered[maxn];
    vector<pair<double, double>> seg, cover;
    double arc, pol;
    inline int sign(double x) {return x < -EPS ? -1 : x > EPS;}
    inline int sign(double x, double y) {return sign(x - y);}
    inline double sqr(const double x) {return x * x;}
    inline double dist(double x1, double y1, double x2, double
        y2) {return sqrt(sqr(x1 - x2) + sqr(y1 - y2));}
    inline double angle(double A, double B, double C) {
        double val = (sqr(A) + sqr(B) - sqr(C)) / (2 * A * B);
        if (val < -1) val = -1;
        if (val > +1) val = +1;
        return acos(val);
    }
    CircleUnion() {
        n = 0;
        seg.clear(), cover.clear();
        arc = pol = 0;
    }
    void init() {
        n = 0;
        seg.clear(), cover.clear();
        arc = pol = 0;
    }
    void add(double xx, double yy, double rr) {
        x[n] = xx, y[n] = yy, r[n] = rr, covered[n] = 0, n++;
    }
    void getarea(int i, double lef, double rig) {
        arc += 0.5 * r[i] * r[i] * (rig - lef - sin(rig - lef))
            ;
        double x1 = x[i] + r[i] * cos(lef), y1 = y[i] + r[i] *
            sin(lef);
        double x2 = x[i] + r[i] * cos(rig), y2 = y[i] + r[i] *
            sin(rig);
        pol += x1 * y2 - x2 * y1;
    }
    double calc() {
        for (int i = 0; i < n; i++) {
            for (int j = 0; j < i; j++)
                if (!sign(x[i] - x[j]) && !sign(y[i] - y[j]) &&
                    !sign(r[i] - r[j])) {
                    r[i] = 0.0;
                    break;
                }
        }
        for (int i = 0; i < n; i++)
            for (int j = 0; j < n; j++)
                if (i != j && sign(r[j] - r[i]) >= 0 && sign(
                    dist(x[i], y[i], x[j], y[j]) - (r[j] - r[i
                    ])) <= 0) {
                    covered[i] = 1;
                    break;
                }
        for (int i = 0; i < n; i++) {
            if (sign(r[i]) && !covered[i]) {
                seg.clear();
                for (int j = 0; j < n; j++) {
                    if (i != j) {
```

```cpp
                    double d = dist(x[i], y[i], x[j], y[j])
                        ;
                    if (sign(d - (r[j] + r[i])) >= 0 ||
                        sign(d - abs(r[j] - r[i])) <= 0)
                        continue;
                    double alpha = atan2(y[j] - y[i], x[j]
                        - x[i]);
                    double beta = angle(r[i], d, r[j]);
                    pair<double, double> tmp(alpha - beta,
                        alpha + beta);
                    if (sign(tmp.first) <= 0 && sign(tmp.
                        second) <= 0)
                        seg.push_back(pair<double, double
                            >(2 * PI + tmp.first, 2 * PI +
                            tmp.second));
                    else if (sign(tmp.first) < 0) {
                        seg.push_back(pair<double, double
                            >(2 * PI + tmp.first, 2 * PI))
                            ;
                        seg.push_back(pair<double, double
                            >(0, tmp.second));
                    }
                    else seg.push_back(tmp);
                }
                sort(seg.begin(), seg.end());
                double rig = 0;
                for (vector<pair<double, double> >::iterator
                    iter = seg.begin(); iter != seg.end();
                    iter++) {
                    if (sign(rig - iter->first) >= 0)
                        rig = max(rig, iter->second);
                    else {
                        getarea(i, rig, iter->first);
                        rig = iter->second;
                    }
                }
                if (!sign(rig)) arc += r[i] * r[i] * PI;
                else getarea(i, rig, 2 * PI);
            }
        }
        return pol / 2.0 + arc;
    }
} ccu;
```

CircleLine.h
Description: Finds the intersection between a circle and a line. Re-
turns a vector of either 0, 1, or 2 intersection points. P is intended to be
Point<double>

"Point.h", "lineDistance.h", "LineProjectionReflection.h"   debf86, 8 lines
```cpp
template<class P>
vector<P> circleLine(P c, double r, P a, P b) {
  double h2 = r*r - a.cross(b,c)*a.cross(b,c)/(b-a).dist2();
  if (h2 < 0) return {};
  P p = lineProj(a, b, c), h = (b-a).unit() * sqrt(h2);
  if (h2 == 0) return {p};
  return {p - h, p + h};
}
```

CircleCircleArea.h
Description: Calculates the area of the intersection of 2 circles
      8bf2b6, 12 lines
```cpp
template<class P>
double circleCircleArea(P c, double cr, P d, double dr) {
  if (cr < dr) swap(c, d), swap(cr, dr);
  auto A = [&](double r, double h) {
    return r*r*acos(h/r)-h*sqrt(r*r-h*h);
  };
  auto l = (c - d).dist(), a = (l*l + cr*cr - dr*dr)/(2*l);
  if (l - cr - dr >= 0) return 0; // far away
```

```
    if (l - cr + dr <= 0) return M_PI*dr*dr;
    if (l - cr >= 0) return A(cr, a) + A(dr, l-a);
    else return A(cr, a) + M_PI*dr*dr - A(dr, a-l);
}
```

CirclePolygonIntersection.h
Description: Returns the area of the intersection of a circle with a ccw polygon.
Time: $\mathcal{O}(n)$
"Point.h"   cf9deb, 18 lines

```
#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(P c, double r, vector<P> ps) {
    auto tri = [&](P p, P q) {
        auto r2 = r * r / 2;
        P d = q - p;
        auto a = d.dot(p)/d.dist2(), b = (p.dist2()-r*r)/d.dist2();
        auto det = a * a - b;
        if (det <= 0) return arg(p, q) * r2;
        auto s = max(0., -a-sqrt(det)), t = min(1., -a+sqrt(det));
        if (t < 0 || 1 <= s) return arg(p, q) * r2;
        P u = p + d * s, v = p + d * t;
        return arg(p,u) * r2 + u.cross(v)/2 + arg(v,q) * r2;
    };
    auto sum = 0.0;
    for (int i = 0; i < ps.size(); ++i)
        sum += tri(ps[i] - c, ps[(i + 1) % ps.size()] - c);
    return sum;
}
```

8.3   Polygons

InsidePolygon.h
Description: Returns true if p lies within the polygon. If strict is true, it returns false for points on the boundary. The algorithm uses products in intermediate steps so watch out for overflow.
Usage: vector<P> v = {P{4,4}, P{1,2}, P{2,1}};
bool in = inPolygon(v, P{3, 3}, false);
Time: $\mathcal{O}(n)$
"Point.h", "OnSegment.h", "SegmentDistance.h"   f9442d, 12 lines

```
template<class P>
bool inPolygon(vector<P> &p, P a, bool strict = true) {
    int cnt = 0, n = p.size();
    for(int i = 0; i < n; ++i) {
        P q = p[(i + 1) % n];
        if (onSegment(p[i], q, a)) return !strict; // change to
            // -1 if u need to detect points in the boundary
        //or: if (segDist(p[i], q, a) <= eps) return !strict;
        cnt ^= ((a.y<p[i].y) - (a.y<q.y)) * a.cross(p[i], q) > 0;
    }
    return cnt;
}
```

PolygonArea.h
Description: Returns the area of a polygon. Clockwise enumeration gives negative area. Watch out for overflow if using int as T!
"Point.h"   3794ee, 17 lines

```
template<class T>
T polygonArea(vector<Point<T>> &v) {
    T a = v.back().cross(v[0]);
    for(int i = 0; i < v.size()-1; ++i)
        a += v[i].cross(v[i+1]);
    return abs(a)/2.0;
}

Point<T> polygonCentroid(vector<Point<T>> &v) { // not tested
    Point<T> cent(0,0); T area = 0;
    for(int i = 0; i < v.size(); ++i) {
        int j = (i+1) % (v.size()); T a = cross(v[i], v[j]);
```

```
        cent += a * (v[i] + v[j]);
        area += a;
    }
    return cent/area/(T)3;
}
```

PolygonCenter.h
Description: Returns the center of mass for a polygon.
Time: $\mathcal{O}(n)$
"Point.h"   26a00f, 8 lines

```
P polygonCenter(const vector<P>& v) {
    P res(0, 0); double A = 0;
    for (int i = 0, j = v.size() - 1; i < v.size(); j = ++i) {
        res = res + (v[i] + v[j]) * v[j].cross(v[i]);
        A += v[j].cross(v[i]);
    }
    return res / A / 3;
}
```

PolygonCut.h
Description:
Returns a vector with the vertices of a polygon with everything to the left of the line going from s to e cut away.
Usage: vector<P> p = ...;
p = polygonCut(p, P(0,0), P(1,0));
"Point.h", "lineIntersection.h"   7df36f, 11 lines

```
vector<P> polygonCut(const vector<P>& poly, P s, P e) {
    vector<P> res;
    for(int i = 0; i < poly.size(); ++i) {
        P cur = poly[i], prev = i ? poly[i-1] : poly.back();
        bool side = s.cross(e, cur) < 0;
        if (side != (s.cross(e, prev) < 0))
            res.push_back(lineInter(s, e, cur, prev).second);
        if (side) res.push_back(cur);
    }
    return res;
}
```

ConvexHull.h
Description:
Returns a vector of indices of the convex hull in counter-clockwise order. Points on the edge of the hull between two other points are not considered part of the hull.
Time: $\mathcal{O}(n \log n)$
"Point.h"   3612d7, 12 lines

```
vector<P> convexHull(vector<P> pts) {
    if (pts.size() <= 1) return pts;
    sort(pts.begin(), pts.end());
    vector<P> h(pts.size()+1);
    int s = 0, t = 0;
    for (int it = 2; it--; s = --t, reverse(pts.begin(), pts.end()))
        for (P p : pts) {
            while (t >= s + 2 && h[t-2].cross(h[t-1], p) <= 0) t--;
            h[t++] = p;
        }
    return {h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])};
}
```

HullDiameter.h
Description: Returns the two points with max distance on a convex hull (ccw, no duplicate/colinear points).
                                                                   0e0c1f, 11 lines

```
array<P, 2> hullDiameter(vector<P> S) {
    int n = S.size(), j = n < 2 ? 0 : 1;
    pair<lint, array<P, 2>> res({0, {S[0], S[0]}});
    for(int i = 0; i < j; ++i)
        for (;; j = (j + 1) % n) {
```

```
            res = max(res, {(S[i] - S[j]).dist2(), {S[i], S[j]}});
            if ((S[(j + 1) % n] - S[j]).cross(S[i + 1] - S[i]) >= 0)
                break;
        }
    return res.second;
}
```

PointInsideHull.h
Description: Determine whether a point t lies inside a convex hull (CCW order, with no colinear points). Returns true if point lies within the hull. If strict is true, points on the boundary aren't included.
Time: $\mathcal{O}(\log N)$
"Point.h", "sideOf.h", "OnSegment.h"   7b8514, 12 lines

```
bool inHull(const vector<P> &l, P p, bool strict = true) {
    int a = 1, b = l.size() - 1, r = !strict;
    if (l.size() < 3) return r && onSegment(l[0], l.back(), p);
    if (sideOf(l[0], l[a], l[b]) > 0) swap(a, b);
    if (sideOf(l[0], l[a], p) >= r || sideOf(l[0], l[b], p)<= -r)
        return false;
    while (abs(a - b) > 1) {
        int c = (a + b) / 2;
        (sideOf(l[0], l[c], p) > 0 ? b : a) = c;
    }
    return sgn(l[a].cross(l[b], p)) < r;
}
```

minkowski-sum.h
Description: Minkowski sum of two convex polygons given in ccw order.
Time: $\mathcal{O}(N + M)$
                                                                   71a25a, 20 lines

```
vector<P> minkowski_sum(vector<P> A, vector<P> B) {
    if (int(A.size()) > int(B.size())) swap(A, B);
    if (A.empty()) return {};
    if (int(A.size()) == 1) {
        for (auto& b : B) b = b + A.front();
        return B;
    }
    rotate(A.begin(), min_element(A.begin(), A.end()), A.end());
    rotate(B.begin(), min_element(B.begin(), B.end()), B.end());
    A.push_back(A[0]); A.push_back(A[1]);
    B.push_back(B[0]); B.push_back(B[1]);
    const int N = int(A.size()), M = int(B.size());
    vector<P> ans; ans.reserve(N+M);
    for (int i = 0, j = 0; i+2 < N || j+2 < M;) {
        ans.push_back(A[i] + B[j]);
        auto sgn = (A[i+1] - A[i]).cross(B[j+1] - B[j]);
        i += (sgn >= 0); j += (sgn <= 0);
    }
    return ans;
}
```

PolyUnion.h
Description: Calculates the area of the union of n polygons (not necessarily convex). The points within each polygon must be given in CCW order. Guaranteed to be precise for integer coordinates up to 3e7. If epsilons are needed, add them in sideOf as well as the definition of sgn.
Time: $\mathcal{O}(N^2)$, where N is the total number of points
"Point.h", "sideOf.h"   a45bd4, 33 lines

```
double rat(P a, P b) { return sgn(b.x) ? a.x/b.x : a.y/b.y; }
double polyUnion(vector<vector<P>>& poly) {
    double ret = 0;
    for(int i = 0; i < poly.size(); ++i)
        for(int v = 0; v < poly[i].size(); ++v) {
            P A = poly[i][v], B = poly[i][(v + 1) % poly[i].size()];
            vector<pair<double, int>> segs = {{0, 0}, {1, 0}};
            for(int j = 0; j < poly.size(); ++j) if (i != j) {
                for(int u = 0; u < poly[j].size(); ++u) {
```

```
        P C = poly[j][u], D = poly[j][(u + 1) % poly[j].size
            ()];
        int sc = sideOf(A, B, C), sd = sideOf(A, B, D);
        if (sc != sd) {
          double sa = C.cross(D, A), sb = C.cross(D, B);
          if (min(sc, sd) < 0)
            segs.emplace_back(sa / (sa - sb), sgn(sc - sd));
        } else if (!sc && !sd && j<i && sgn((B-A).dot(D-C))
            >0){
          segs.emplace_back(rat(C - A, B - A), 1);
          segs.emplace_back(rat(D - A, B - A), -1);
        }
      }
    }
    sort(segs.begin(), segs.end());
    for(auto& s : segs) s.first = min(max(s.first, 0.0), 1.0);
    double sum = 0;
    int cnt = segs[0].second;
    for(int j = 1; j < segs.size(); ++j) {
      if (!cnt) sum += segs[j].first - segs[j - 1].first;
      cnt += segs[j].second;
    }
    ret += A.cross(B) * sum;
  }
  return ret / 2;
}
```

LineHullIntersection.h

Description: Line-convex polygon intersection. The polygon must be ccw and have no colinear points. lineHull(line, poly) returns a pair describing the intersection of a line with the polygon: ● $(-1, -1)$ if no collision, ● $(i, -1)$ if touching the corner i, ● (i, i) if along side $(i, i+1)$, ● (i, j) if crossing sides $(i, i+1)$ and $(j, j+1)$. In the last case, if a corner i is crossed, this is treated as happening on side $(i, i+1)$. The points are returned in the same order as the line hits the polygon. extrVertex returns the point of a hull with the max projection onto a line.

Time: $\mathcal{O}(N + Q\log n)$

"Point.h"   65ebb6, 39 lines

```
typedef array<P, 2> Line;
#define cmp(i,j) sgn(dir.perp().cross(poly[(i)%n]-poly[(j)%n]))
#define extr(i) cmp(i + 1, i) >= 0 && cmp(i, i - 1 + n) < 0
int extrVertex(vector<P>& poly, P dir) { // hash-1
  int n = poly.size(), left = 0, right = n;
  if (extr(0)) return 0;
  while (left + 1 < right) {
    int m = (left + right) / 2;
    if (extr(m)) return m;
    int ls = cmp(left + 1, left), ms = cmp(m + 1, m);
    (ls < ms || (ls == ms && ls == cmp(left, m)) ? right : left
        ) = m;
  }
  return left;
} // hash-1 = 99da02

#define cmpL(i) sgn(line[0].cross(poly[i], line[1]))
array<int, 2> lineHull(Line line, vector<P>& poly) { // hash-2
  int endA = extrVertex(poly, (line[0] - line[1]).perp());
  int endB = extrVertex(poly, (line[1] - line[0]).perp());
  if (cmpL(endA) < 0 || cmpL(endB) > 0)
    return {-1, -1};
  array<int, 2> res;
  for(int i = 0; i < 2; ++i) {
    int left = endB, right = endA, n = poly.size();
    while ((left + 1) % n != right) {
      int m = ((left + right + (left < right ? 0 : n)) / 2) % n
          ;
      (cmpL(m) == cmpL(endB) ? left : right) = m;
    }
    res[i] = (left + !cmpL(right)) % n;
```

```
    swap(endA, endB);
  }
  if (res[0] == res[1]) return {res[0], -1};
  if (!cmpL(res[0]) && !cmpL(res[1]))
    switch ((res[0] - res[1] + sz(poly) + 1) % poly.size()) {
      case 0: return {res[0], res[0]};
      case 2: return {res[1], res[1]};
    }
  return res;
} // hash-2 = ba025e
```

HalfPlane.h

Description: Halfplane intersection area

"Point.h", "lineIntersection.h"   c0a94b, 70 lines

```
#define eps 1e-8
typedef Point<double> P;

struct Line {
  P P1, P2;
  // Right hand side of the ray P1 -> P2
  explicit Line(P a = P(), P b = P()) : P1(a), P2(b) {};
  P intpo(Line y) {
    pair<int, P> r = lineInter(P1, P2, y.P1, y.P2);
    assert(r.first == 1);
    return r.second;
  }
  P dir() { return P2 - P1; }
  bool contains(P x) {
    return (P2 - P1).cross(x - P1) < eps;
  }
  bool out(P x) { return !contains(x); }
};

template<class T>
bool mycmp(Point<T> a, Point<T> b) {
  // return atan2(a.y, a.x) < atan2(b.y, b.x);
  if (a.x * b.x < 0) return a.x < 0;
  if (abs(a.x) < eps) {
    if (abs(b.x) < eps) return a.y > 0 && b.y < 0;
    if (b.x < 0) return a.y > 0;
    if (b.x > 0) return true;
  }
  if (abs(b.x) < eps) {
    if (a.x < 0) return b.y < 0;
    if (a.x > 0) return false;
  }
  return a.cross(b) > 0;
}

bool cmp(Line a, Line b) { return mycmp(a.dir(), b.dir()); }

double Intersection_Area(vector <Line> b) {
  sort(b.begin(), b.end(), cmp);
  int n = b.size();
  int q = 1, h = 0, i;
  vector<Line> c(b.size() + 10);
  for (i = 0; i < n; i++) {
    while (q < h && b[i].out(c[h].intpo(c[h - 1]))) h--;
    while (q < h && b[i].out(c[q].intpo(c[q + 1]))) q++;
    c[++h] = b[i];
    if (q < h && abs(c[h].dir().cross(c[h - 1].dir())) < eps) {
      if (c[h].dir().dot(c[h - 1].dir()) > 0) {
        h--;
        if (b[i].out(c[h].P1))  c[h] = b[i];
      }else {
        // The area is either 0 or infinite.
        // If you have a bounding box, then the area is
        //     definitely 0.
        return 0;
```

```
      }
    }
  }
  while (q < h-1 && c[q].out(c[h].intpo(c[h - 1]))) h--;
  while (q < h-1 && c[h].out(c[q].intpo(c[q + 1]))) q++;
  // Intersection is empty. This is sometimes different from
  //     the case when
  // the intersection area is 0.
  if (h - q <= 1) return 0;
  c[h + 1] = c[q];
  vector<P> s;
  for (i = q; i <= h; i++) s.push_back(c[i].intpo(c[i + 1]));
  s.push_back(s[0]);
  double ans = 0;
  for (i = 0; i < (int) s.size()-1; i++) ans += s[i].cross(s[i
      + 1]);
  return ans / 2;
}
```

8.4 Misc. Point Set Problems

ClosestPair.h

Description: Finds the closest pair of points.

Time: $\mathcal{O}(n\log n)$

"Point.h"   32b14f, 16 lines

```
pair<P, P> closest(vector<P> v) {
  assert(v.size() > 1);
  set<P> S;
  sort(v.begin(), v.end(), [](P a, P b) { return a.y < b.y; });
  pair<int64_t, pair<P, P>> ret{LLONG_MAX, {P(), P()}};
  int j = 0;
  for(P &p : v) {
    P d{1 + (int64_t)sqrt(ret.first), 0};
    while (v[j].y <= p.y - d.x) S.erase(v[j++]);
    auto lo = S.lower_bound(p - d), hi = S.upper_bound(p + d);
    for (; lo != hi; ++lo)
      ret = min(ret, {(*lo - p).dist2(), {*lo, p}});
    S.insert(p);
  }
  return ret.second;
}
```

KdTree.h

Description: KD-tree (2d, can be extended to 3d)

"Point.h"   915562, 63 lines

```
typedef long long T;
typedef Point<T> P;
const T INF = numeric_limits<T>::max();

bool on_x(const P& a, const P& b) { return a.x < b.x; }
bool on_y(const P& a, const P& b) { return a.y < b.y; }

struct Node {
  P pt; // if this is a leaf, the single point in it
  T x0 = INF, x1 = -INF, y0 = INF, y1 = -INF; // bounds
  Node *first = 0, *second = 0;

  T distance(const P& p) { // min squared distance to a point
    T x = (p.x < x0 ? x0 : p.x > x1 ? x1 : p.x);
    T y = (p.y < y0 ? y0 : p.y > y1 ? y1 : p.y);
    return (P(x,y) - p).dist2();
  }

  Node(vector<P>&& vp) : pt(vp[0]) {
    for (P p : vp) {
      x0 = min(x0, p.x); x1 = max(x1, p.x);
      y0 = min(y0, p.y); y1 = max(y1, p.y);
    }
    if (vp.size() > 1) {
```

```cpp
        // split on x if the box is wider than high (not best
        //      heuristic...)
        sort(vp.begin(),vp.end(), x1 - x0 >= y1 - y0 ? on_x :
            on_y);
        // divide by taking half the array for each child (not
        // best performance with many duplicates in the middle)
        int half = vp.size()/2;
        first = new Node({vp.begin(), vp.begin() + half});
        second = new Node({vp.begin() + half, vp.end()});
      }
    }
};

struct KDTree {
  Node* root;
  KDTree(const vector<P>& vp) : root(new Node({vp.begin(), vp.
      end()})) {}

  pair<T, P> search(Node *node, const P& p) {
    if (!node->first) {
      // uncomment if we should not find the point itself:
      // if (p == node->pt) return {INF, P()};
      return make_pair((p - node->pt).dist2(), node->pt);
    }

    Node *f = node->first, *s = node->second;
    T bfirst = f->distance(p), bsec = s->distance(p);
    if (bfirst > bsec) swap(bsec, bfirst), swap(f, s);

    // search closest side first, other side if needed
    auto best = search(f, p);
    if (bsec < best.first)
      best = min(best, search(s, p));
    return best;
  }

  // find nearest point to a point, and its squared distance
  // (requires an arbitrary operator< for Point)
  pair<T, P> nearest(const P& p) {
    return search(root, p);
  }
};
```

DelaunayTriangulation.h
Description: Computes the Delaunay triangulation of a set of points. Each circumcircle contains none of the input points. If any three points are colinear or any four are on the same circle, behavior is undefined.
Time: $\mathcal{O}\left(n^2\right)$
"Point.h", "3dHull.h"   f6175a, 10 lines
```cpp
template<class P, class F>
void delaunay(vector<P>& ps, F trifun) {
  if (ps.size() == 3) { int d = (ps[0].cross(ps[1], ps[2]) < 0)
      ;
    trifun(0,1+d,2-d); }
  vector<P3> p3;
  for(auto &p : ps) p3.emplace_back(p.x, p.y, p.dist2());
  if (ps.size() > 3) for(auto &t: hull3d(p3)) if ((p3[t.b]-p3[t
      .a]).
      cross(p3[t.c]-p3[t.a]).dot(P3(0,0,1)) < 0)
    trifun(t.a, t.c, t.b);
}
```

FastDelaunay.h
Description: Fast Delaunay triangulation. There must be no duplicate points. If all points are on a line, no triangles will be returned. Should work for doubles as well, though there may be precision issues in 'circ'. Returns triangles in order {t[0][0], t[0][1], t[0][2], t[1][0], ...}, all counter-clockwise.

Time: $\mathcal{O}\left(n \log n\right)$
"Point.h"   a1f392, 90 lines
```cpp
typedef Point<ll> P;
typedef struct Quad* Q;
typedef __int128_t lll; // (can be ll if coords are < 2e4)
P arb(LLONG_MAX,LLONG_MAX); // not equal to any other point

struct Quad { // hash-1
  bool mark; Q o, rot; P p;
  P F() { return r()->p; }
  Q r() { return rot->rot; }
  Q prev() { return rot->o->rot; }
  Q next() { return rot->r()->o->rot; }
}; // hash-1 = ae7c00
// hash-2
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
  lll p2 = p.dist2(), A = a.dist2()-p2,
    B = b.dist2()-p2, C = c.dist2()-p2;
  return p.cross(a,b)*C + p.cross(b,c)*A + p.cross(c,a)*B > 0;
} // hash-2 = 6aff7b
Q makeEdge(P orig, P dest) { // hash-3
  Q q0 = new Quad{0,0,0,orig}, q1 = new Quad{0,0,0,arb},
    q2 = new Quad{0,0,0,dest}, q3 = new Quad{0,0,0,arb};
  q0->o = q0; q2->o = q2; // 0-0, 2-2
  q1->o = q3; q3->o = q1; // 1-3, 3-1
  q0->rot = q1; q1->rot = q2;
  q2->rot = q3; q3->rot = q0;
  return q0;
} // hash-3 = 81016d
void splice(Q a, Q b) { // hash-4
  swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
}
Q connect(Q a, Q b) {
  Q q = makeEdge(a->F(), b->p);
  splice(q, a->next());
  splice(q->r(), b);
  return q;
} // hash-4 = 7e71f7

pair<Q,Q> rec(const vector<P>& s) { // hash-5
  if (sz(s) <= 3) {
    Q a = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.back());
    if (sz(s) == 2) return { a, a->r() };
    splice(a->r(), b);
    auto side = s[0].cross(s[1], s[2]);
    Q c = side ? connect(b, a) : 0;
    return {side < 0 ? c->r() : a, side < 0 ? c : b->r() };
  }

#define H(e) e->F(), e->p
#define valid(e) (e->F().cross(H(base)) > 0)
  Q A, B, ra, rb;
  int half = (sz(s) + 1) / 2;
  tie(ra, A) = rec({s.begin(), s.begin() + half});
  tie(B, rb) = rec({s.begin() + half, s.end()});
  while ((B->p.cross(H(A)) < 0 && (A = A->next())) ||
    (A->p.cross(H(B)) > 0 && (B = B->r()->o)));
  Q base = connect(B->r(), A);
  if (A->p == ra->p) ra = base->r();
  if (B->p == rb->p) rb = base;

#define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
    while (circ(e->dir->F(), H(base), e->F())) { \
      Q t = e->dir; \
      splice(e, e->prev()); \
      splice(e->r(), e->r()->prev()); \
      e = t; \
    }
  for (;;) {
```

```cpp
    DEL(LC, base->r(), o);  DEL(RC, base, prev());
    if (!valid(LC) && !valid(RC)) break;
    if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
      base = connect(RC, base->r());
    else
      base = connect(base->r(), LC->r());
  }
  return { ra, rb };
} // hash-5 = 2488a6

vector<P> triangulate(vector<P> pts) { // hash-6
  sort(pts.begin(), pts.end());  assert(unique(pts.begin(), pts
      .end()) == pts.end());
  if (pts.size() < 2) return {};
  Q e = rec(pts).first;
  vector<Q> q = {e};
  int qi = 0;
  while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
#define ADD { Q c = e; do { c->mark = 1; pts.push_back(c->p); \
  q.push_back(c->r()); c = c->next(); } while (c != e); }
  ADD; pts.clear();
  while (qi < sz(q)) if (!(e = q[qi++])->mark) ADD;
  return pts;
} // hash-6 = 1ebc14
```

RectangleUnionArea.h
Description: Sweep line algorithm that calculates area of union of rectangles in the form $[x1, x2) \times [y1, y2)$
Usage: vector<pair<int,int>>, pair<int,int>> rectangles;
rectangles.push_back({{x1, x2}, {y1, y2}});
lint result = area(rectangles);   529ff1, 51 lines
```cpp
struct seg_node{
  int val, cnt, lz;
  seg_node(int n = INT_MAX, int c = 0): val(n), cnt(c), lz(0)
      {}
  void push(seg_node& l, seg_node& r){
    if(lz){
      l.add(lz);
      r.add(lz);
      lz = 0;
    }
  }
  void merge(const seg_node& l, const seg_node& r){
    if(l.val < r.val) val = l.val, cnt = l.cnt;
    else if(l.val > r.val) val = r.val, cnt = r.cnt;
    else val = l.val, cnt = l.cnt + r.cnt;
  }
  void add(int n){
    val += n;
    lz += n;
  }
  int get_sum(){ return (val ? 0 : cnt); }
};
// x1 y1 x2 y2
lint solve(const vector<array<int, 4>>&v){
  vector<int>ys;
  for(auto& [a, b, c, d] : v){
    ys.push_back(b);
    ys.push_back(d);
  }
  sort(ys.begin(), ys.end());
  ys.erase(unique(ys.begin(), ys.end()), ys.end());
  vector<array<int, 4>>e;
  for(auto [a, b, c, d] : v){
    b = int(lower_bound(ys.begin(), ys.end(), b) - ys.begin());
    d = int(lower_bound(ys.begin(), ys.end(), d) - ys.begin());
    e.push_back({a, b, d, 1});
    e.push_back({c, b, d, -1});
```

```
    }
    sort(e.begin(), e.end());
    int m = (int)ys.size();
    segtree_range<seg_node>seg(m-1);
    for(int i=0;i<m-1;i++) seg.at(i) = seg_node(0, ys[i+1] - ys[i
        ]);
    seg.build();
    int last = INT_MIN, total = ys[m-1] - ys[0];
    lint ans = 0;
    for(auto [x, y1, y2, c] : e){
        ans += (lint)(total - seg.query(0, m-1).get_sum()) * (x -
            last);
        last = x;
        seg.update(y1, y2, &seg_node::add, c);
    }
    return ans;
}
```

8.5   3D

PolyhedronVolume.h
Description: Magic formula for the volume of a polyhedron. Faces should point outwards.

832599, 6 lines

```
template<class V, class L>
double signed_poly_volume(const V &p, const L &trilist) {
    double v = 0;
    for(auto &i : trilist) v += p[i.a].cross(p[i.b]).dot(p[i.c]);
    return v / 6;
}
```

Point3D.h
Description: Class to handle points in 3D space. T can be e.g. double or long long.

8058ae, 33 lines

```
template<class T> struct Point3D { // hash-1
    typedef Point3D P;
    typedef const P& R;
    T x, y, z;
    explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y), z(z) {}
    bool operator<(R p) const {
        return tie(x, y, z) < tie(p.x, p.y, p.z); }
    bool operator==(R p) const {
        return tie(x, y, z) == tie(p.x, p.y, p.z); }
    P operator+(R p) const { return P(x+p.x, y+p.y, z+p.z); }
    P operator-(R p) const { return P(x-p.x, y-p.y, z-p.z); }
    P operator*(T d) const { return P(x*d, y*d, z*d); }
    P operator/(T d) const { return P(x/d, y/d, z/d); }
    T dot(R p) const { return x*p.x + y*p.y + z*p.z; }
    P cross(R p) const {
        return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x);
    } // hash-1 = f914db
// hash-2
    T dist2() const { return x*x + y*y + z*z; }
    double dist() const { return sqrt((double)dist2()); }
    //Azimuthal angle (longitude) to x-axis in interval [-pi, pi]
    double phi() const { return atan2(y, x); }
    //Zenith angle (latitude) to the z-axis in interval [0, pi]
    double theta() const { return atan2(sqrt(x*x+y*y),z); }
    P unit() const { return *this/(T)dist(); } //makes dist()=1
    //returns unit vector normal to *this and p
    P normal(P p) const { return cross(p).unit(); }
    //returns point rotated 'angle' radians ccw around axis
    P rotate(double angle, P axis) const {
        double s = sin(angle), c = cos(angle); P u = axis.unit();
        return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
    }
}; // hash-2 = c9d029
```

3dHull.h
Description: Computes all faces of the 3-dimension hull of a point set. *No four points must be coplanar*, or else random results will be returned. All faces will point outwards.
Time: $\mathcal{O}\left(n^2\right)$

"Point3D.h"   3ed613, 49 lines

```
typedef Point3D<double> P3;

struct PR { // hash-1
    void ins(int x) { (a == -1 ? a : b) = x; }
    void rem(int x) { (a == x ? a : b) = -1; }
    int cnt() { return (a != -1) + (b != -1); }
    int a, b;
}; // hash-1 = cf7c9e

struct F { P3 q; int a, b, c; };

vector<F> hull3d(const vector<P3>& A) { // hash-2
    assert(A.size() >= 4);
    vector<vector<PR>> E(A.size(), vector<PR>(A.size(), {-1, -1})
        );
#define E(x,y) E[f.x][f.y]
    vector<F> FS;
    auto mf = [&](int i, int j, int k, int l) {
        P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
        if (q.dot(A[l]) > q.dot(A[i]))
            q = q * -1;
        F f{q, i, j, k};
        E(a,b).ins(k); E(a,c).ins(j); E(b,c).ins(i);
        FS.push_back(f);
    };
    for(int i=0;i<4;i++) for(int j=i+1;j<4;j++) for(k=j+1;k<4;k
        ++)
        mf(i, j, k, 6 - i - j - k);
// hash-2 = 80e001
    for(int i=4; i<A.size();i++) { // hash-3
        for(int j=0;j<FS.size();++j) {
            F f = FS[j];
            if(f.q.dot(A[i]) > f.q.dot(A[f.a])) {
                E(a,b).rem(f.c);
                E(a,c).rem(f.b);
                E(b,c).rem(f.a);
                swap(FS[j--], FS.back());
                FS.pop_back();
            }
        }
        int nw = FS.size();
        for(int j=0;j<nw;j++) {
            F f = FS[j];
#define C(a, b, c) if (E(a,b).cnt() != 2) mf(f.a, f.b, i, f.c);
            C(a, b, c); C(a, c, b); C(b, c, a);
        }
    }
    for(auto &it: FS) if ((A[it.b] - A[it.a]).cross(
        A[it.c] - A[it.a]).dot(it.q) <= 0) swap(it.c, it.b);
    return FS;
}; // hash-3 = 52653c
```

SphericalDistance.h
Description: Returns the shortest distance on the sphere with radius radius between the points with azimuthal angles (longitude) f1 (ϕ_1) and f2 (ϕ_2) from x axis and zenith angles (latitude) t1 (θ_1) and t2 (θ_2) from z axis. All angles measured in radians. The algorithm starts by converting the spherical coordinates to cartesian coordinates so if that is what you have you can use only the two last rows. dx*radius is then the difference between the two points in the x direction and d*radius is the total distance between the points.

611f07, 8 lines

```
double sphericalDistance(double f1, double t1,
```

```
    double f2, double t2, double radius) {
    double dx = sin(t2)*cos(f2) - sin(t1)*cos(f1);
    double dy = sin(t2)*sin(f2) - sin(t1)*sin(f1);
    double dz = cos(t2) - cos(t1);
    double d = sqrt(dx*dx + dy*dy + dz*dz);
    return radius*2*asin(d/2);
}
```

Strings (9)

kmp.h
Description: failure[x] computes the length of the longest prefix of s that ends at x, other than s[0...x] itself (abacaba -> -1,0,0,1,0,1,2,3). Can be used to find all occurrences of a pattern in a text.
Time: $\mathcal{O}\left(n\right)$

4eb73b, 9 lines

```
vector<int> prefix_function(const string& S) {
    vector<int> fail = {-1}; fail.reserve(S.size());
    for (int i = 0; i < int(S.size()); ++i) {
        int j = fail.back();
        while (j != -1 && S[i] != S[j]) j = fail[j];
        fail.push_back(j+1);
    }
    return fail;
}
```

duval.h
Description: A string is called simple (or a Lyndon word), if it is strictly smaller than any of its own nontrivial suffixes.
Time: $\mathcal{O}\left(N\right)$

19ec0f, 24 lines

```
template<typename T>
pair<int, vector<string>> duval(int n, const T &s) {
    // s += s  // if you need to know the min cyclic string
    vector<string> factors;
    int i = 0, ans = 0;
    while (i < n) { // until n/2 to find min cyclic string
        ans = i; int j = i + 1, k = i;
        while (j < n + n && !(s[j % n] < s[k % n])) {
            if (s[k % n] < s[j % n]) k = i;
            else k++;
            j++;
        }
        while (i <= k) {
            factors.push_back(s.substr(i, j-k));
            i += j - k;
        }
    }
    return {ans, factors};
    // returns 0-indexed position of the least cyclic shift
    // min cyclic string will be s.substr(ans, n/2)
}
template<typename T>pair<int,vector<string>> duval(const T &s){
    return duval((int)s.size(), s);
}
```

z-algorithm.h
Description: z[x] computes the length of the longest common prefix of s[i:] and s, except z[0] = 0. (abacaba -> 0010301)
Time: $\mathcal{O}\left(n\right)$

7c8c64, 13 lines

```
vector<int> Z(const string& S) {
    vector<int> z(S.size()); int l = -1, r = -1;
    for(int i = 1; i < int(S.size()); ++i) {
        z[i] = i >= r ? 0 : min(r - i, z[i - l]);
        while (i + z[i] < int(S.size()) && S[i + z[i]] == S[z[i]])
            z[i]++;
        if (i + z[i] > r) l = i, r = i + z[i];
```

```cpp
  } return z;
}
vector<int> get_prefix(string a, string b) {
  string str = a + '@' + b; vector<int> k = z(str);
  return vector<int>(k.begin() + int(a.size())+1, k.end());
}
```

manacher.h

Description: For each position in a string, computes p[0][i] = half length of longest even palindrome around pos i, p[1][i] = longest odd (half rounded down).

Time: $\mathcal{O}(N)$

```cpp
array<vector<int>, 2> manacher(const string &s) {
  int n = s.size();
  array<vector<int>, 2> p = {vector<int>(n+1), vector<int>(n)};
  for(int z = 0; z < 2; ++z) for(int i=0,l=0,r=0; i < n; i++) {
    int t = r-i+!z;
    if (i<r) p[z][i] = min(t, p[z][l+t]);
    int L = i-p[z][i], R = i+p[z][i]-!z;
    while (L>=1 && R+1<n && s[L-1] == s[R+1])
      p[z][i]++, L--, R++;
    if (R > r) l = L, r = R;
  } return p;
}
```

min-rotation.h

Description: Finds the lexicographically smallest rotation of a string.

Usage: rotate(v.begin(), v.begin()+min_rotation(v), v.end());

Time: $\mathcal{O}(N)$

```cpp
int min_rotation(string s) {
  int a=0, N=s.size(); s += s;
  for(int b = 0; b < N; ++b) for(int i =0; i < N; ++i) {
    if (a+i == b || s[a+i] < s[b+i]) {b += max(0, i-1); break;}
    if (s[a+i] > s[b+i]) { a = b; break; }
  } return a;
}
```

xor-trie.h

Description: Query get the maximum possible xor between an integer X and every possible subarray. Just insert zero and for each prefix xor, insert it in the trie and query for max xor. The answer is the maximum possible value for each prefix query.

```cpp
template<int K = 31> struct trie_t {
  vector<array<int, 2>> trie;
  trie_t() : trie(1, {-1, -1}) {}
  void add(int val) {
    int cur = 0;
    for (int a = K; a >= 0; --a) {
      int b = (val >> a) & 1;
      if (trie[cur][b] == -1) {
        trie[cur][b] = size(trie);
        trie.push_back({-1, -1});
      }
      cur = trie[cur][b];
    }
  }
  int max_xor(int val) {
    int cur = 0, mask = 0;
    for (int a = K; a >= 0; --a) {
      int b = (val >> a) & 1;
      if (trie[cur][!b] == -1) {
        cur = trie[cur][b];
      } else {
        mask |= (1 << a);
        cur = trie[cur][!b];
      }
```

```cpp
    }
    return mask;
  }
};
```

hashing.h

```cpp
const int maxn = 400001;
const int mod = 1004669333, base = 33, inv_base = 121778101;
vector<int> base_pow(maxn + 1), inv_base_pow(maxn + 1);
void prep() {
  base_pow[0] = 1;
  for (int i = 1; i <= maxn; ++i)
    base_pow[i] = (lint)base_pow[i - 1] * base % mod;
  inv_base_pow[0] = 1;
  for (int i = 1; i <= maxn; ++i)
    inv_base_pow[i] = (lint)inv_base_pow[i - 1] * inv_base %
        mod;
}
struct hashes_t {
  string s;
  int n;
  vector<int> acc_hash, acc_inv_hash;
  hashes_t(const string &_s): s(_s), n(s.size()), acc_hash(n +
      1, 0) ,acc_inv_hash(n + 1, 0) { // 127dd9
    for (int i = 0; i < n; ++i) {
      acc_hash[i + 1] =
          (acc_hash[i] + (lint)base_pow[i] * (s[i] - 'a' + 1)) %
              mod;
      acc_inv_hash[i + 1] =
          (acc_inv_hash[i] + (lint)inv_base_pow[i] * (s[i] - 'a'
              + 1)) % mod;
    }
  }
  int get_hash(int a, int b) {
    assert(a <= b);
    int hash = acc_hash[b + 1] - acc_hash[a];
    if (hash < 0) hash += mod;
    hash = (lint)hash * inv_base_pow[a] % mod;
    return hash;
  }
  int get_inv_hash(int a, int b) {
    assert(a <= b);
    int hash = acc_inv_hash[b + 1] - acc_inv_hash[a];
    if (hash < 0) hash += mod;
    hash = (lint)hash * base_pow[b] % mod;
    return hash;
  }
};
```

modnum-double-hashing.h

Description: Simple, short and efficient hashing using pairs to reduce load factor.

```cpp
using num = modnum<int(1e9)+7>;
using hsh = pairnum<num, num>;
const hsh BASE(163, 311);
// uniform_int_distribution<int> MULT_DIST(0.1*MOD,0.9*MOD);
// constexpr hsh BASE(MULT_DIST(rng), MULT_DIST(rng));
struct hash_t { // c9d6c0
  int n;
  string str;
  vector<hsh> hash, basePow;
  hash_t(const string& s) : n(s.size()), str(s), hash(n+1),
      basePow(n+1) { // dd1f3f
    basePow[0] = 1;
    for (int i = 1; i <= n; ++i) basePow[i] = basePow[i-1] *
        BASE;
```

```cpp
    for (int i = 0; i < n; ++i)
      hash[i+1] = hash[i] * BASE + hsh(s[i]);
  }
  hsh get_hash(int left, int right) { // 302ee0
    assert(left <= right);
    return hash[right] - hash[left] * basePow[right - left];
  }
  int lcp(hash_t &other) { // 5eb9e2
    int left = 0, right = min(str.size(), other.str.size());
    while (left < right) {
      int mid = (left + right + 1)/2;
      if (hash[mid] == other.hash[mid]) left = mid;
      else right = mid-1;
    }
    return left;
  }
};
vector<int> rabinkarp(string t, string p) { // c11cfc
  vector<int> matches;
  hsh h(0, 0);
  for (int i = 0; i < p.size(); ++i)
    h = BASE * h + hsh(p[i]);
  hash_t result(t);
  for (int i = 0; i + p.size() <= t.size(); ++i)
    if (result.get_hash(i, i + p.size()) == h)
      matches.push_back(i);
  return matches;
}
```

aho-corasick.h

```cpp
const int sigma = 26;
array<int, sigma> init;
for (int i = 0; i < sigma; i++) init[i] = -1;
vector<array<int, sigma>> trie(1, init);
vector<int> out(1, -1), parent(n, -1), ids(n);
for (int i = 0; i < n; i++) {
  int cur = 0;
  for (char ch : s[i]) {
    int c = ch - 'a';
    if (trie[cur][c] == -1) {
      trie[cur][c] = (int)trie.size();
      trie.push_back(init); out.push_back(-1);
    }
    cur = trie[cur][c];
  }
  if (out[cur] == -1) out[cur] = i;
  ids[i] = out[cur];
}
vector<int> bfs,f(trie.size()); bfs.reserve(trie.size());
for (int c = 0; c < sigma; c++)
  if (trie[0][c] == -1) trie[0][c] = 0;
  else bfs.push_back(trie[0][c]);
for (int z = 0; z < (int)bfs.size() ; z++) {
  int cur = bfs[z];
  for (int c = 0; c < sigma; c++) {
    if (trie[cur][c] == -1)
      trie[cur][c] = trie[f[cur]][c];
    else {
      int nxt = trie[cur][c];
      int fail = trie[f[cur]][c];
      if (out[nxt] == -1) out[nxt] = out[fail];
      else parent[out[nxt]] = out[fail];
      f[nxt] = fail; bfs.push_back(nxt);
    }
  }
}
```

suffix-array.h

Description: Builds suffix array for a string, first element is the size of the string. The `lcp` function calculates longest common prefixes for neighbouring strings in suffix array. The returned vector is of size $n + 1$.

Time: $\mathcal{O}(N \log N)$ where N is the length of the string for creation of the SA. $\mathcal{O}(N)$ for longest common prefixes.

<../data-structures/rmq.h>, <../various/random-numbers.h>   d31829, 66 lines

```cpp
struct suffix_array_t {
  int N, H; vector<int> sa, invsa, lcp;
  rmq_t<pair<int, int>> rmq;
  bool cmp(int a, int b) { return invsa[a+H] < invsa[b+H]; }
  void ternary_sort(int a, int b) {
    if (a == b) return;
    int md = sa[a+rng() % (b-a)], lo = a, hi = b;
    for (int i = a; i < b; ++i) if (cmp(sa[i], md))
      swap(sa[i], sa[lo++]);
    for (int i = b-1; i >= lo; --i) if (cmp(md, sa[i]))
      swap(sa[i], sa[--hi]);
    ternary_sort(a, lo);
    for (int i = lo; i < hi; ++i) invsa[sa[i]] = hi-1;
    if (hi-lo == 1) sa[lo] = -1;
    ternary_sort(hi, b);
  }
  suffix_array_t() {}
  template<typename I>
  suffix_array_t(I begin, I end): N(int(end-begin)+1), sa(N) {
    vector<int> v(begin, end); v.push_back(INT_MIN);
    invsa = v; iota(sa.begin(), sa.end(), 0);
    H = 0; ternary_sort(0, N);
    for(H = 1; H <= N; H *= 2) for(int j=0, i=j; i!=N; i=j)
        if (sa[i] < 0) {
          while (j < N && sa[j] < 0) j += -sa[j];
          sa[i] = -(j - i);
        } else {j = invsa[sa[i]] + 1; ternary_sort(i, j);}
    for (int i = 0; i < N; ++i) sa[invsa[i]] = i;
    lcp.resize(N-1); int K = 0;
    for (int i = 0; i < N-1; ++i) {
      if(invsa[i] > 0) while(v[i+K] == v[sa[invsa[i]-1]+K])++K;
      lcp[invsa[i]-1] = K; K = max(K - 1, 0);
    }
    vector<pair<int, int>> lcp_index(N-1);
    for (int i = 0; i < N-1; ++i) lcp_index[i] = {lcp[i], 1+i};
    rmq = rmq_t<pair<int, int>>(std::move(lcp_index));
  }
  auto rmq_query(int a, int b) const {return rmq.query(a,b);}
  auto get_split(int a, int b) const {return rmq.query(a,b-1);}
  int get_lcp(int a, int b) const {
    if (a == b) return N - a;
    a = invsa[a], b = invsa[b];
    if (a > b) swap(a, b);
    return rmq_query(a, b).first;
  }
};
vector<vector<int>> ch(2*N+1); int V = 0;
vector<array<int, 2>> sa_range(2*N+1);
vector<int> leaves(N+1), par(2*N+1), depth(2*N+1);
auto dfs = [&](auto&& self, int lo, int hi, int prv)-> void{
  int cur = V++; par[cur] = prv;
  if (prv != -1) ch[prv].push_back(cur);
  sa_range[cur] = {lo, hi};
  if (hi - lo == 1) {
    leaves[us.sa[lo]] = cur;
    depth[cur] = N-us.sa[lo] + 1;
  } else {
    int d = us.get_split(lo, hi).first;
    depth[cur] = d; int mi = lo;
    while (hi - mi >= 2) {
      auto [nd, nmi] = us.get_split(mi, hi);
      if (nd != d) break;
```

```cpp
      self(self, mi, nmi, cur); mi = nmi;
    } self(self, mi, hi, cur);
  }
}; dfs(dfs, 0, N+1, -1);
```

suffix-tree.h

Description: Builds suffix-tree informations based by emulating it over the suffix-array and lcp, root of the tree represents the special character (size of string for suffix-array), can therefore be ignored when calculating stuff.

Time: $\mathcal{O}(N \log N)$

<../data-structures/rmq.h>, "suffix-array.h"   31aacf, 46 lines

```cpp
struct suffix_tree_t {
  int N, V;
  vector<vector<int>> ch;
  vector<array<int, 2>> sa_range;
  vector<int> leaves, par, depth;
  vector<int> suff_link;
  vector<bool> is_unique_link, has_unique_child;
  suffix_array_t us;
  suffix_tree_t() {}
  suffix_tree_t(string S) : N(int(S.size())), V(0), ch(2*N+1),
  sa_range(2*N+1), leaves(N+1), par(2*N+1), depth(2*N+1),
  us(S.begin(), S.end()) { dfs(0, N+1, -1); }
  void dfs(int a, int b, int prv) {
    int cur = V++;
    par[cur] = prv;
    if (prv != -1) ch[prv].push_back(cur);
    sa_range[cur] = {a, b};
    if (b - a == 1) {
      leaves[us.sa[a]] = cur;
      depth[cur] = N - us.sa[a];
    } else {
      int d = us.get_split(a, b).first;
      depth[cur] = d;
      int mi = a;
      while (b - mi >= 2) {
        auto [nd, nmi] = us.get_split(mi, b);
        if (nd != d) break;
        dfs(mi, nmi, cur);
        mi = nmi;
      }
      dfs(mi, b, cur);
    }
  }
  void build_links() {
    suff_link.resize(V, -1), is_unique_link.resize(V),
        has_unique_child.resize(V);
    for (int i = 0; i < N; ++i) {
      for (int cur = leaves[i], link = leaves[i+1];; cur = par[
          cur]) {
        if (cur == 0 || suff_link[cur] != -1) break;
        suff_link[cur] = link;
        is_unique_link[cur] = (sa_range[cur][1] - sa_range[cur
            ][0]) == (sa_range[link][1] - sa_range[link][0]);
        if (is_unique_link[cur]) has_unique_child[link] = true;
        while (~link && depth[link] + 1 > depth[par[cur]]) link
            = par[link];
      }
    }
  }
};
```

suffix-automaton.h

Description: Suffix automaton

defb60, 33 lines

```cpp
template<int offset = 'a'> struct array_state {
  array<int, 26> as;
  array_state() { fill(begin(as), end(as), ~0); }
  int& operator[](char c) { return as[c - offset]; }
```

```cpp
  int count(char c) { return (~as[c - offset] ? 1 : 0); }
};
template<typename C, typename state = map<C, int>> struct
    suffix_automaton {
  struct node_t {
    int len, link; int64_t cnt; state next;
  };
  int N, cur; vector<node_t> nodes;
  suffix_automaton() : N(1), cur(0), nodes{node_t{0, -1, 0,
      {}}} {}
  node_t& operator[](int v) { return nodes[v]; };
  void append(C c) {
    int v = cur; cur = N++;
    nodes.push_back(node_t{nodes[v].len + 1, 0, 1, {}});
    for (; ~v && !nodes[v].next.count(c); v = nodes[v].link)
      nodes[v].next[c] = cur;
    if (~v) {
      const int u = nodes[v].next[c];
      if (nodes[v].len + 1 == nodes[u].len) {
        nodes[cur].link = u;
      } else {
        const int clone = N++;
        nodes.push_back(nodes[u]);
        nodes[clone].len = nodes[v].len + 1;
        nodes[u].link = nodes[cur].link = clone;
        for (; ~v && nodes[v].next[c] == u; v = nodes[v].link)
          nodes[v].next[c] = clone;
      }
    }
  }
};
```

9.1 Suffix Automaton

9.1.1 Number of different substrings

Is the number of paths in the automaton starting at the root.

$$d(v) = 1 + \sum_{v \to w} d(w)$$

9.1.2 Total lenght of different substrings

Is the sum of children answers and paths starting at each children.

$$ans(v) = \sum_{v \to w} d(w) + ans(w)$$

9.1.3 Lexicographically K-th substring

Is the K-th lexicographically path, so you can search using the number of paths from each state

9.1.4 Smallest cyclic shift

Construct for string $S + S$. Greedily search the minimal character.

9.1.5 Number of occurrences

For each state not created by cloning, initialize $cnt(v) = 1$. Then, just do a dfs to calculate $cnt(v)$

$$cnt(link(v)) {+} = cnt(v)$$

9.1.6 First occurence position

When we create a new state cur do $first(pos) = len(cur) - 1$. When we clone q as $clone$ do $first(clone) = first(q)$. Answer is $first(v) - size(P) + 1$, where v is the state of string P

9.1.7 All occurence positions

From $first(v)$ do a dfs using suffix link, from $link(u)$ go to u.

Various (10)

10.1 Intervals

interval-container.h
Description: Add and remove intervals from a set of disjoint intervals. Will merge the added interval with any overlapping intervals in the set when adding. Intervals are [inclusive, exclusive).
Time: $\mathcal{O}(\log N)$

edce47, 20 lines

```cpp
set<pii>::iterator addInterval(set<pii> &is, int L, int R) {
  if (L == R) return is.end();
  auto it = is.lower_bound({L, R}), before = it;
  while (it != is.end() && it->first <= R) {
    R = max(R, it->second);
    before = it = is.erase(it);
  }
  if (it != is.begin() && (--it)->second >= L) {
    L = min(L, it->first); R = max(R, it->second);
    is.erase(it);
  } return is.insert(before, {L,R});
}
void removeInterval(set<pii> &is, int L, int R) {
  if (L == R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it->second;
  if (it->first == L) is.erase(it);
  else (int&)it->second = L;
  if (R != r2) is.emplace(R, r2);
}
```

interval-cover.h
Description: Compute indices of smallest set of intervals covering another interval. Intervals should be [inclusive, exclusive). To support [inclusive, inclusive], change (A) to add || R.empty(). Returns empty set on failure (or if G is empty).
Time: $\mathcal{O}(N \log N)$

133eb4, 17 lines

```cpp
template<class T>
vector<int> cover(pair<T, T> G, vector<pair<T, T>> I) {
  vector<int> S(I.size()), R;
  iota(S.begin(), S.end(), 0);
  sort(S.begin(), S.end(), [&](int a, int b) {
    return I[a] < I[b]; });
  T cur = G.first; int at = 0;
  while (cur < G.second) { // (A)
    pair<T, int> mx = {cur, -1};
    while (at < I.size() && I[S[at]].first <= cur) {
      mx = max(mx, {I[S[at]].second, S[at]});
```

```cpp
    at++;
    }
    if (mx.second == -1) return {};
    cur = mx.first; R.push_back(mx.second);
  } return R;
}
```

constant-intervals.h
Description: Split a monotone function on [from, to) into a minimal set of half-open intervals on which it has the same value. Runs a callback g for each such interval.
Usage: constantIntervals(0, sz(v), [&](int x){
return v[x];}, [&](int lo, int hi, T val){...});
Time: $\mathcal{O}\left(k \log \frac{n}{k}\right)$

753a4c, 17 lines

```cpp
template<class F, class G, class T>
void rec(int from, int to, F& f, G& g, int& i, T& p, T q) {
  if (p == q) return;
  if (from == to) {
    g(i, to, p); i = to; p = q;
  } else {
    int mid = (from + to) >> 1;
    rec(from, mid, f, g, i, p, f(mid));
    rec(mid+1, to, f, g, i, p, q);
  }
}
template<class F, class G>
void constantIntervals(int from, int to, F f, G g) {
  if (to <= from) return;
  int i = from; auto p = f(i), q = f(to-1);
  rec(from, to-1, f, g, i, p, q); g(i, to, q);
}
```

10.2 Misc. algorithms

floor.h

e7f4c9, 7 lines

```cpp
template<typename T> T mfloor(T a, T b) {
  return a / b - (((a ^ b) < 0 && a % b != 0) ? 1 : 0);
}
```

```cpp
template<typename T> T mceil(T a, T b) {
  return a / b + (((a ^ b) > 0 && a % b != 0) ? 1 : 0);
}
```

basis-manager.h
Description: A list of basis values sorted in decreasing order, where each value has a unique highest bit.

0d24d9, 34 lines

```cpp
template<typename T, int BITS = 60> struct xor_basis {
  int N = 0; array<T, BITS> basis;
  T min_value(T start) const {
    if (N == BITS) return 0;
    for (int i = 0; i < N; ++i)
      start = min(start, start ^ basis[i]);
    return start;
  }
  T max_value(T start = 0) const {
    if (N == BITS) return ((T) 1 << BITS) - 1;
    for (int i = 0; i < N; ++i)
      start = max(start, start ^ basis[i]);
    return start;
  }
  bool add(T x) {
    x = min_value(x);
    if (x == 0) return false;
    basis[N++] = x;
    for (int k = N - 1; k > 0 && basis[k] > basis[k - 1]; k--)
      swap(basis[k], basis[k - 1]);
    return true;
  }
```

```cpp
  }
  void merge(const xor_basis<T>& other) {
    for (int i = 0; i < other.n && N < BITS; i++)
      add(other.basis[i]);
  }
  void merge(const xor_basis<T>& a, const xor_basis<T>& b) {
    if (a.N > b.N) {
      *this = a; merge(b);
    } else {
      *this = b; merge(a);
    }
  }
};
```

ternary-search.h
Description: Find the smallest i in $[a, b]$ that maximizes $f(i)$, assuming that $f(a) < \ldots < f(i) \geq \cdots \geq f(b)$. To reverse which of the sides allows non-strict inequalities, change the $<$ marked with (A) to $<=$, and reverse the loop at (B). To minimize f, change it to $>$, also at (B). If you are dealing with real numbers, you'll need to pick $m_1 = (2a+b)/3.0$ and $m_2 = (a+2b)/3.0$. Consider setting a constant number of iterations for the search, usually $[200, 300]$ iterations are sufficient for problems with error limit as 10^{-6}.
Usage: int ind = ternSearch(0,n-1,[&](int i){return a[i];});
Time: $\mathcal{O}(\log(b - a))$

35ef73, 11 lines

```cpp
template<class F> int ternSearch(int a, int b, F f) {
  assert(a <= b);
  while (b - a >= 5) {
    int mid = (a + b) / 2;
    if (f(mid) < f(mid+1)) a = mid; // (A)
    else b = mid+1;
  }
  for(int i = a+1; i <= b; ++i)
    if (f(a) < f(i)) a = i; // (B)
  return a;
}
```

merge-sort.h
Time: $\mathcal{O}(N \log(N))$

fac159, 25 lines

```cpp
vector<int> merge(vector<int> &values, int l, int r) {
  static vector<int> result(values.size());
  int i = l, j = l + (r - l)/2;
  int mid = j, k = i, inversions = 0;
  while (i < mid && j < r) {
    if (values[i] < values[j]) result[k++] = values[i++];
    else {
      result[k++] = values[j++];
      inversions += (mid - i);
    }
  }
  while (i < mid) result[k++] = values[i++];
  while (j < r) result[k++] = values[j++];
  for (k = l; k < r; ++k) values[k] = result[k];
  return result;
}
```

```cpp
vector<int> msort(vector<int> &values, int l, int r) {
  if (r - l > 1) {
    int mid = l + (r - l)/2;
    msort(values, l, mid); msort(values, mid, r);
    return merge(values, l, r);
  }
  return {};
}
```

radix-sort.h
Description: Radix Sort Algorithm.
Time: $\mathcal{O}(NK)$ where K is the number of bits in the largest element of the array to be sorted.

889884, 54 lines

```cpp
struct identity {
  template<typename T>
    T operator()(const T &x) const {
      return x;
    }
};
template<typename T, typename T_extract_key = identity>
void radix_sort(vector<T> &data, int bits_per_pass = 10, const
    T_extract_key &extract_key = identity()) {
  if (data.size() < 256) {
    sort(data.begin(), data.end(), [&](const T &a, const T &b)
        {
      return extract_key(a) < extract_key(b);
    });
    return;
  }
  using T_key = decltype(extract_key(data.front()));
  T_key minimum = numeric_limits<T_key>::max();
  for (T &x : data) minimum = min(minimum, extract_key(x));
  int max_bits = 0;
  for (T &x : data) {
    T_key key = extract_key(x);
    max_bits = max(max_bits, key == minimum ? 0 : 64 -
        __builtin_clzll(key - minimum));
  }
  int passes = max((max_bits + bits_per_pass / 2) /
      bits_per_pass, 1);
  if (32 - __builtin_clz(data.size()) <= 1.5 * passes) {
    sort(data.begin(), data.end(), [&](const T &a, const T &b)
        {
      return extract_key(a) < extract_key(b);
    });
    return;
  }
  vector<T> buffer(data.size());
  vector<int> counts;
  int bits_so_far = 0;
  for (int p = 0; p < passes; p++) {
    int bits = (max_bits + p) / passes;
    counts.assign(1 << bits, 0);
    for (T &x : data) {
      T_key key = extract_key(x) - minimum;
      counts[(key >> bits_so_far) & ((1 << bits) - 1)]++;
    }
    int count_sum = 0;
    for (int &count : counts) {
      int current = count;
      count = count_sum;
      count_sum += current;
    }
    for (T &x : data) {
      T_key key = extract_key(x) - minimum;
      int key_section = (key >> bits_so_far) & ((1 << bits) -
          1);
      buffer[counts[key_section]++] = x;
    }
    swap(data, buffer);
    bits_so_far += bits;
  }
}
```

postfix-notation-solver.h
Description: Solves postfix (Reverse Polish) notation equation to solve prefix notation equation reverse e and change (i) and (ii)

Time: $\mathcal{O}(N)$

bf1f57, 31 lines

```cpp
template<typename T, typename P, typename F>
T postfixSolver(const vector<P> &e, const set<P> &ops, F ptot){
  vector<T> stk;
  for (auto cur: e)
    if(ops.count(cur)){
      T c;
      //operations:
      if(cur == "-"){
        T b = stk.back(); // (i) T a = stk.back();
        stk.pop_back();
        T a = stk.back(); //(ii) T b = stk.back();
        stk.pop_back();
        c = a - b;
      }
      else if(cur == "NOT"){
        T a = stk.back();
        stk.pop_back();
        c = ~a;
      }
      stk.push_back(c);
    } else
      stk.push_back(ptot(cur));
  return stk.back();
}
//example postfix:
vector<string> e = {"13", "14", "-", "NOT"};
int ans = postfixSolver<int>( e, {"-", "NOT"}, [](const string
    &s){ return stoi(s); } );
//example prefix:
vector<string> e = {"NOT", "-", "13", "14"};
reverse(e.begin(), e.end()); // DON'T FORGET!!!!!
int ans = postfixSolver<int>( e, {"-", "NOT"}, [](const string
    &s){ return stoi(s); } );
```

count-triangles.h
Description: Counts x, y >= 0 such that Ax + By <= C.

fbb812, 6 lines

```cpp
ll count_triangle(ll A, ll B, ll C) {
  if (C < 0) return 0;
  if (A > B) swap(A, B);
  ll p = C / B, k = B / A, d = (C - p * B) / A;
  return count_triangle(B - k * A, A, C - A * (k * p + d + 1))
      + (p + 1) * (d + 1) + k * p * (p + 1) / 2;
}
```

karatsuba.h
Description: Faster-than-naive convolution of two sequences: $c[x] = \sum a[i]b[x-i]$. Uses the identity $(aX + b)(cX + d) = acX^2 + bd + ((a + c)(b + d) - ac - bd)X$. Doesn't handle sequences of very different length welint. See also FFT, under the Numerical chapter.
Time: $\mathcal{O}(N^{1.6})$

37b858, 30 lines

```cpp
int size(int s) { return s > 1 ? 32-__builtin_clz(s-1) : 0; }

void karatsuba(lint *a, lint *b, lint *c, lint *t, int n) {
  int ca = 0, cb = 0;
  for(int i = 0; i < n; ++i) ca += !!a[i], cb += !!b[i];
  if (min(ca, cb) <= 1500/n) { // few numbers to multiply
    if (ca > cb) swap(a, b);
    for(int i = 0; i < n; ++i)
      if (a[i]) for(int j = 0; j < n; ++j) c[i+j] += a[i]*b[j];
  }
  else {
    int h = n >> 1;
    karatsuba(a, b, c, t, h); // a0*b0
    karatsuba(a+h, b+h, c+n, t, h); // a1*b1
    for(int i = 0; i < h; ++i) a[i] += a[i+h], b[i] += b[i+h];
    karatsuba(a, b, t, t+n, h); // (a0+a1)*(b0+b1)
    for(int i = 0; i < h; ++i) a[i] -= a[i+h], b[i] -= b[i+h];
    for(int i = 0; i < n; ++i) t[i] -= c[i]+c[i+n];
    for(int i = 0; i < n; ++i) c[i+h] += t[i], t[i] = 0;
  }
}

vector<lint> conv(vector<lint> a, vector<lint> b) {
  int sa = a.size(), sb = b.size(); if (!sa || !sb) return {};
  int n = 1<<size(max(sa,sb)); a.resize(n), b.resize(n);
  vector<lint> c(2*n), t(2*n);
  for(int i = 0; i < 2*n; ++i) t[i] = 0;
  karatsuba(&a[0], &b[0], &c[0], &t[0], n);
  c.resize(sa+sb-1); return c;
}
```

count-inversions.h
Description: Count the number of inversions to make an array sorted. Merge sort has another approach.
Time: $\mathcal{O}(nlog(n))$

"../data-structures/fenwick-tree.h"   7e4bc9, 6 lines

```cpp
FT<lint> bit(n);
lint inv = 0;
for (int i = n-1; i >= 0; --i) {
  inv += bit.query(values[i]); // careful with the interval
  bit.update(values[i], 1);    // [0, x) or [0, x] ?
}
```

histogram.h
Time: $\mathcal{O}(N)$

a77bf4, 17 lines

```cpp
int max_area(const vector<int>& height) {
  const int N = int(height.size());
  vector<int> L(N), R(N);
  for (int i = N-1; i >= 0; --i) {
    R[i] = i+1;
    while (R[i] < N && height[i] <= height[R[i]]) R[i] = R[R[i
        ]];
  }
  for (int i = 0; i < N; ++i) {
    L[i] = i-1;
    while (L[i] >= 0 && height[i] <= height[L[i]]) L[i] = L[L[i
        ]];
  }
  int area = 0;
  for (int i = 0; i < N; ++i)
    area = max(area, int64_t(R[i] - L[i] - 1) * heigh[i]);
  return area;
}
```

date-manipulation.h

088459, 44 lines

```cpp
string week_day_str[7] = {"Sunday", "Monday", "Tuesday", "
    Wednesday", "Thursday", "Friday", "Saturday"};
string month_str[13] = {"", "January", "February", "March", "
    April", "May", "June", "July", "August", "September", "
    October", "November", "December"};
map<string, int> week_day_int = {{"Sunday", 0}, {"Monday", 1},
    {"Tuesday", 2}, {"Wednesday", 3}, {"Thursday", 4}, {"
    Friday", 5}, {"Saturday", 6}};
map<string, int> month_int = {{"January", 1}, {"February", 2},
    {"March", 3}, {"April", 4}, {"May", 5}, {"June", 6}, {"
    July", 7}, {"August", 8}, {"September", 9}, {"October",
    10}, {"November", 11}, {"December", 12}};
int month[2][13] = {{0, 31, 28, 31, 30, 31, 30, 31, 31, 30, 31,
    30, 31}, {0, 31, 29, 31, 30, 31, 30, 31, 31, 30, 31, 30,
    31}};

/* O(1) - Checks if year y is a leap year. */
```

```cpp
bool leap_year(int y){
  return (y % 4 == 0 && y % 100 != 0) || y % 400 == 0;
}

/* O(1) - Increases the day by one. */
void update(int &d, int &m, int &y){
  if (d == month[leap_year(y)][m]){
    d = 1;
    if (m == 12) {
      m = 1;
      y++;
    }
    else m++;
  }
  else d++;
}

int intToDay(int jd) { return jd % 7; }
int dateToInt(int y, int m, int d) {
  return 1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
}
void intToDate(int jd, int &y, int &m, int &d) {
  int x, n, i, j;
  x = jd + 68569;
  n = 4 * x / 146097;
  x -= (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;
  x -= 1461 * i / 4 - 31;
  j = 80 * x / 2447;
  d = x - 2447 * j / 80;
  x = j / 11;
  m = j + 2 - 12 * x;
  y = 100 * (n - 49) + i + x;
}
```

n-queens.h
Description: NQueens

e97e9e, 43 lines

```cpp
int ans;
bitset<30> rw, ld, rd; //2*MAX_N -1
bitset<30> iniqueens; //2*MAX_N -1
vector<int> col;
void init(int n){
  ans=0;
  rw.reset();
  ld.reset();
  rd.reset();
  col.assign(n,-1);
}
void init(int n, vector<pair<int,int>> initial_queens){
  //it does NOT check if initial queens are at valid positions
  init(n);
  iniqueens.reset();
  for(pair<int,int> pos: initial_queens){
    int r=pos.first, c= pos.second;
    rw[r] = ld[r-c+n-1] = rd[r+c]=true;
    col[c]=r;
    iniqueens[c] = true;
  }
}

void backtracking(int c, int n){
  if(c==n){
    ans++;
    for(int r:col) cout<<r+1<<" ";
    cout<<"\n";
    return;
```

```cpp
  }
  else if(iniqueens[c]){
    backtracking(c+1,n);
  }
  else for(int r=0;r<n;r++){
    if(!rw[r] && !ld[r-c+n-1] && !rd[r+c]){
      // if(board[r][c]!=blocked && !rw[r] && !ld[r-c+n-1] && !
          rd[r+c]){ // if there are blocked possitions
      rw[r] = ld[r-c+n-1] = rd[r+c]=true;
      col[c]=r;
      backtracking(c+1,n);
      col[c]=-1;
      rw[r] = ld[r-c+n-1] = rd[r+c]=false;
    }
  }
}
```

sudoku-solver.h

6be906, 41 lines

```cpp
int N,m; // N = n*n, m = n; where n equal number of rows or
    columns
array<array<int, 10>, 10> grid;
struct SudokuSolver {
  bool UsedInRow(int row,int num){
    for(int col = 0; col < N; ++col)
      if(grid[row][col] == num) return true;
    return false;
  }
  bool UsedInCol(int col,int num){
    for(int row = 0; row < N; ++row)
      if(grid[row][col] == num) return true;
    return false;
  }
  bool UsedInBox(int row_0,int col_0,int num){
    for(int row = 0; row < m; ++row)
      for(int col = 0; col < m; ++col)
        if(grid[row+row_0][col+col_0] == num) return true;
    return false;
  }
  bool isSafe(int row,int col,int num){
    return !UsedInRow(row,num) && !UsedInCol(col,num) && !
        UsedInBox(row-row%m,col-col%m,num);
  }
  bool find(int &row,int &col){
    for(row = 0; row < N; ++row)
      for(col = 0; col < N; ++col)
        if(grid[row][col] == 0) return true;
    return false;
  }
  bool Solve(){
    int row, col;
    if(!find(row,col)) return true;
    for(int num = 1; num <= N; ++num){
      if(isSafe(row,col,num)){
        grid[row][col] = num;
        if(Solve()) return true;
        grid[row][col] = 0;
      }
    }
    return false;
  }
};
```

floyd-cycle.h
Description: Detect loop in a list. Consider using mod template to avoid overflow.
Time: $\mathcal{O}(n)$

b456ab, 9 lines

```cpp
template<class F> pair<int,int> find(int x0, F f) {
```

```cpp
  int t = f(x0), h = f(t), mu = 0, lam = 1;
  while (t != h) t = f(t), h = f(f(h));
  h = x0;
  while (t != h) t = f(t), h = f(h), ++mu;
  h = f(t);
  while (t != h) h = f(h), ++lam;
  return {mu, lam};
}
```

10.3   Dynamic programming

divide-and-conquer-dp.h
Description: Given $a[i] = \min_{lo(i) \leq k < hi(i)}(f(i,k))$ where the (minimal) optimal k increases with i, computes $a[i]$ for $i = L..R - 1$.
Time: $\mathcal{O}((N + (hi - lo)) \log N)$

c7f2d1, 29 lines

```cpp
struct DP { // Modify at will:
  vector<int>a, freq;
  vector<ll>old, cur;
  ll cnt; int lcur, rcur;
  DP(const vector<int>&_a, int n): a(_a), freq(n), old(n+1,
      linf), cur(n+1, linf), cnt(0), lcur(0), rcur(0){}
  int lo(int ind) { return 0; }
  int hi(int ind) { return ind; }
  void add(int k, int c){ cnt += freq[a[k]]++; }
  void del(int k, int c){ cnt -= --freq[a[k]]; }
  ll C(int l, int r){
    while(lcur > l) add(--lcur, 0);
    while(rcur < r) add(rcur++, 1);
    while(lcur < l) del(lcur++, 0);
    while(rcur > r) del(--rcur, 1);
    return cnt;
  }
  ll f(int ind, int k) { return old[k] + C(k, ind); }
  void store(int ind, int k, ll v) { cur[ind] = v; }
  void rec(int L, int R, int LO, int HI) {
    if (L >= R) return;
    int mid = (L + R) >> 1;
    pair<ll, int> best(LLONG_MAX, LO);
    for(int k = max(LO,lo(mid)); k <= min(HI,hi(mid)); ++k)
      best = min(best, make_pair(f(mid, k), k));
    store(mid, best.second, best.first);
    rec(L, mid, LO, best.second);
    rec(mid+1, R, best.second, HI);
  }
};
```

knuth-dp.h
Description: When doing DP on intervals: $a[i][j] = \min_{i<k<j}(a[i][k] + a[k][j]) + f(i,j)$, where the (minimal) optimal k increases with both i and j, one can solve intervals in increasing order of length, and search $k = p[i][j]$ for $a[i][j]$ only between $p[i][j-1]$ and $p[i+1][j]$. This is known as Knuth DP. Sufficient criteria for this are if $f(b,c) \leq f(a,d)$ and $f(a,c) + f(b,d) \leq f(a,d) + f(b,c)$ for all $a \leq b \leq c \leq d$. Consider also: LineContainer (ch. Data structures), monotone queues, ternary search.
Time: $\mathcal{O}(N^2)$

cht-dp.h
Description: Transforms dp of the form (or similar) $dp[i] = \min_{j<i}(dp[j] + b[j]*a[i])$. Time goes from $O(n^2)$ to $O(n \log n)$, if using online line container, or $O(n)$ if lines are inserted in order of slope and queried in order of x. To apply try to find a way to write the factor inside minimization as a linear function of a value related to i. Everything else related to j will become constant.

edit-distance.h

Description: Find the minimum numbers of edits required to convert string s into t. Only insertion, removal and replace operations are allowed.

101083, 32 lines

```cpp
int edit_dist(string &s, string &t) {
    const int n = int(s.size()), m = int(t.size());
    vector<vector<int>> dp(n+1, vector<int>(m+1, n+m+2));
    vector<vector<int>> prv(n+1, vector<int>(m+1, 0));
    dp[0][0] = 0;
    for (int i = 0; i <= n; i++) {
        for (int j = 0; j <= m; j++) {
            if (i < n) { // remove
                int cnd = dp[i][j] + 1;
                if (cnd < dp[i+1][j]) {
                    dp[i+1][j] = cnd;
                    prv[i+1][j] = 1;
                }
            }
            if (j < m) { // insert
                int cnd = dp[i][j] + 1;
                if (cnd < dp[i][j+1]) {
                    dp[i][j+1] = cnd;
                    prv[i][j+1] = 2;
                }
            }
            if (i < n && j < m) { // modify
                int cnd = dp[i][j] + (s[i] != t[j]);
                if (cnd < dp[i+1][j+1]) {
                    dp[i+1][j+1] = cnd;
                    prv[i+1][j+1] = 3;
                }
            }
        }
    }
    return dp[n][m];
}
```

LIS.h

Description: Compute indices for the longest increasing subsequence.
Time: $\mathcal{O}(N \log N)$

e962da, 16 lines

```cpp
template<class I> vector<int> lis(const vector<I>& S) {
    if (S.empty()) return {};
    vector<int> prev(S.size());
    using p = pair<I, int>; vector<p> res;
    for(int i = 0; i < (int)S.size(); i++) {
        // change 0 -> i for longest non-decreasing subsequence
        auto it = lower_bound(res.begin(), res.end(), p {S[i], 0});
        if (it == res.end()) res.emplace_back(), it = res.end()-1;
        *it = {S[i], i};
        prev[i] = it == res.begin() ? 0 : (it-1)->second;
    }
    int L = res.size(), cur = res.back().second;
    vector<int> ans(L);
    while (L--) ans[L] = cur, cur = prev[cur];
    return ans;
}
```

digit-dp.h

Description: Compute how many # between 1 and N have K distinct digits in the base L without leading zeros;
Usage: auto hex_to_dec = [&](char c) -> int {
return ('A' <= c && c <= 'F' ? (10 + c - 'A') : (c - '0'));
};
digit_dp<modnum<int(1e9) + 7>, hex_to_dec>(N, K);
Time: $\mathcal{O}(NK)$

8138af, 26 lines

```cpp
template<typename T, class F> T digit_dp(const string& S, int K
    , F& L) {
    const int base = 16, len = int(S.size());
```

```cpp
    vector<bool> w(base);
    vector<vector<T>> dp(len + 1, vector<T>(base + 2));
    int cnt = 0;
    for (int d = 0; d < len; ++d) {
        // adding new digit to numbers with prefix < s
        for (int x = 0; x <= base; ++x) {
            dp[d + 1][x] += dp[d][x] * x;
            dp[d + 1][x + 1] += dp[d][x] * (base - x);
        }
        // adding strings whith prefix only 0's and last digit != 0
        if (d) dp[d + 1][1] += (base - 1);
        // adding prefix equal to s and last digit < s, first digit
        //     cannot be 0
        for (int x = 0; x < L(S[d]); ++x) {
            if (d == 0 && x == 0) continue;
            if (w[x]) dp[d + 1][cnt] += 1;
            else dp[d + 1][cnt + 1] += 1;
        }
        // marking if the last digit appears in the prefix of s
        if (w[L(S[d])] == false)
            w[L(S[d])] = true, cnt++;
    }
    // adding string k
    dp[len][cnt] += 1; return dp[len][K];
}
```

LCS.h

Description: Finds the longest common subsequence.
Memory: $\mathcal{O}(nm)$
Time: $\mathcal{O}(nm)$ where n and m are the lengths of the sequences.

463080, 14 lines

```cpp
template<class T> T lcs(const T &X, const T &Y) {
    int a = X.size(), b = Y.size();
    vector<vector<int>> dp(a+1, vector<int>(b+1));
    for(int i = 1; i <= a; ++i) for(int j = 1; j <= b; j++)
        dp[i][j] = X[i-1]==Y[j-1] ? dp[i-1][j-1]+1 :
        max(dp[i][j-1],dp[i-1][j]);
    int len = dp[a][b];
    T ans(len, 0);
    while (a && b)
        if (X[a-1] == Y[b-1]) ans[--len] = X[--a], --b;
        else if (dp[a][b-1] > dp[a-1][b]) --b;
        else --a;
    return ans;
}
```

knapsack-unbounded.h

Description: Knapsack problem but now take the same item multiple items is allowed.
Time: $\mathcal{O}(N \log N)$

1acfe0, 10 lines

```cpp
int knapsack(vector<int> &v, vector<int> &w, int total) {
    vector<int> dp(total+1, -1);
    int result = 0; dp[0] = 0;
    for (int i = 0; i <= total; ++i) for (int j = 0; j < n; ++j)
        if (w[j] <= i && dp[i - w[j]] >= 0)
            dp[i] = max(dp[i], dp[i - w[j]] + v[j]);
    int result = 0;
    for (int i = 0; i <= total; ++i) result = max(result, dp[i]);
    return result;
}
```

knapsack-bounded.h

Description: You are given n types of items, each items has a weight and a quantity. Is possible to fill a knapsack with capacity X using any subset of items?
Time: $\mathcal{O}(W \cdot N)$

b24799, 11 lines

```cpp
auto solve(vi weight, vi cnt, int X) {
```

```cpp
    vector<int> dp(X+1, 0);
    for (int i = 0; i < N; ++i)
        for (int j = X-weight[i]; j >= 0; --j) {
            if (!dp[j]) continue;
            int k = cnt[i], s = j + weight[i];
            while (k > 0 && s <= X && !dp[s])
                dp[s] = 1, --k, s += weight[i];
        }
    return dp[X];
}
```

knapsack-bounded-costs.h

Description: You are given N types of items, you have $cnt[i]$ items of i-th type, and each item of i-th type $weight[i]$ and $cost[i]$. What is the maximal cost you can get by picking some items weighing exactly X in total?
Time: $\mathcal{O}(N \cdot W)$

"../data-structures/monotonic-queue.h"   03d171, 15 lines

```cpp
auto solve(vi weight, vi cost, vi cnt, int X) {
    vector<int> dp(X+1, 0); int N = int(weight.size());
    vector<max_monotonic_queue<int>> M(X+1);
    for (int i = 0; i < N; ++i) {
        for (int j = 0; j < min(X+1, weight[i]); ++j)
            M[j] = max_monotonic_queue<int>();
        for (int j = 0; j <= X; ++j) {
            auto& que = M[j % weight[i]];
            if (que.size() > cnt[i]) que.pop();
            que.add(cost[i]);
            que.push(dp[j]);
            dp[j] = que.get_val();
        }
    } return dp[X];
}
```

knapsack-bitset.h

Description: Find first value greater than m that cannot be formed by the sums of numbers from v.

25166e, 11 lines

```cpp
bitset<int(1e7)> dp, dp1;
int knapsack(vector<int> &items, int n, int m) {
    dp[0] = dp1[0] = true;
    for (int i = 0; i < n; ++i) {
        dp1 <<= items[i];
        dp |= dp1;
        dp1 = dp;
    }
    dp.flip();
    return dp._Find_next(m);
}
```

two-max-equal-sum.h

Description: Two maximum equal sum disjoint subsets, s[i] = 0 if v[i] wasn't selected, s[i] = 1 if v[i] is in the first subset and s[i] = 2 if v[i] is in the second subset
Time: $\mathcal{O}(n * S)$

056620, 15 lines

```cpp
auto twoMaxEqualSumDS(const vector<int> &v){
    int sum=accumulate(v.begin(), v.end(), 0), n=int(v.size());
    vector<int> old(2*sum + 1, INT_MIN/2), dp(2*sum + 1), s(n);
    vector<vector<int>> rec(n, vector<int>(2*sum + 1));
    int i; old[sum] = 0;
    for(i = 0; i < n; ++i, swap(old, dp))
        for(int a, b, d = v[i]; d <= 2*sum - v[i]; d++){
            dp[d] = max(old[d], a = old[d - v[i]] + v[i]);
            dp[d] = max(dp[d], b = old[d + v[i]]);
            rec[i][d] = dp[d] == a ? 1 : dp[d] == b ? 2 : 0;
        }
    for(int j = i-1, d = sum; j >= 0; --j)
        d+=(s[j] = rec[j][d]) ? s[j] == 2 ? v[j] : - v[j] : 0;
    return make_pair(old[sum], s);
}
```

```
}
```

max-zero-submatrix.h
Description: Computes the area of the largest submatrix that contains only 0s
Time: $\mathcal{O}(NM)$

<div align="right">d7bff2, 18 lines</div>

```cpp
const int MAXN = 100, MAXM = 100;
array<array<int, MAXN>, MAXM> A, H;
int solve(int N, int M) {
  stack<int, vector<int>> s; int ret = 0;
  for (int j = 0; j < M; j++) for (int i = N - 1; i >= 0; i--)
      H[i][j] = A[i][j] ? 0 : 1 + (i == N - 1 ? 0 : H[i + 1][j
      ]);
  for (int i = 0; i < N; i++) {
    for (int j = 0; j < M; j++) {
      int minInd = j;
      while (!s.empty() && H[i][s.top()] >= H[i][j]) {
        ret = max(ret, (j - s.top()) * (H[i][s.top()]));
        minInd = s.top(); s.pop(); H[i][minInd] = H[i][j];
      }
      s.push(minInd);
    }
    while (!s.empty()) ret = max(ret, (M - s.top()) * H[i][s.
        top()]); s.pop();
  }
  return ret;
}
```

10.4   Debugging tricks

- `signal(SIGSEGV, [](int) { _Exit(0); });` converts segfaults into Wrong Answers. Similarly one can catch SIGABRT (assertion failures) and SIGFPE (zero divisions). _GLIBCXX_DEBUG violations generate SIGABRT (or SIGSEGV on gcc 5.4.0 apparently).

- `feenableexcept(29);` kills the program on NaNs (1), 0-divs (4), infinities (8) and denormals (16).

10.5   Optimization tricks

10.5.1   Bit hacks

- `x & -x` is the least bit in `x`.

- `for (int x = m; x; ) { --x &= m; ... }` loops over all subset masks of `m` (except `m` itself).

- `c = x&-x, r = x+c; (((r^x) >> 2)/c) | r` is the next number after `x` with the same number of bits set.

- `rep(b,0,K) rep(i,0,(1 << K)) if (i & 1 << b) D[i] += D[i^(1 << b)];` computes all sums of subsets.

10.5.2   Pragmas

- `#pragma GCC optimize ("Ofast")` will make GCC auto-vectorize for loops and optimizes floating points better (assumes associativity and turns off denormals).

- `#pragma GCC target ("avx,avx2")` can double performance of vectorized code, but causes crashes on old machines.

- `#pragma GCC optimize ("trapv")` kills the program on integer overflows (but is really slow).

fast-input.h
Description: Returns an integer. Usage requires your program to pipe in input from file. Can replace calls to `gc()` with `getchar_unlocked()` if extra speed isn't necessary (60% slowdown).
Usage: `./a.out < input.txt`
Time: About 5x as fast as cin/scanf.

<div align="right">b31afb, 18 lines</div>

```cpp
struct GC {
  char buf[1 << 16];
  size_t bc = 0, be = 0;
  char operator()() {
    if (bc >= be) {
      buf[0] = 0, bc = 0;
      be = fread(buf, 1, sizeof(buf), stdin);
    }
    return buf[bc++]; // returns 0 on EOF
  }
} gc;
int readInt() {
  int a, c;
  while ((a = gc()) < 40);
  if (a == '-') return -readInt();
  while ((c = gc()) >= 48) a = a * 10 + c - 480;
  return a - 48;
}
```

bump-allocator.h
Description: When you need to dynamically allocate many objects and don't care about freeing them. "new X" otherwise has an overhead of something like 0.05us + 16 bytes per allocation.

<div align="right">745db2, 8 lines</div>

```cpp
// Either globally or in a single class:
static char buf[450 << 20];
void* operator new(size_t s) {
  static size_t i = sizeof buf;
  assert(s < i);
  return (void*)&buf[i -= s];
}
void operator delete(void*) {}
```

bump-allocator-stl.h
Description: BumpAllocator for STL containers.
Usage: `vector<vector<int, small<int>>> ed(N);`

<div align="right">bb66d4, 14 lines</div>

```cpp
char buf[450 << 20] alignas(16);
size_t buf_ind = sizeof buf;

template<class T> struct small {
  typedef T value_type;
  small() {}
  template<class U> small(const U&) {}
  T* allocate(size_t n) {
    buf_ind -= n * sizeof(T);
    buf_ind &= 0 - alignof(T);
    return (T*)(buf + buf_ind);
  }
  void deallocate(T*, size_t) {}
};
```

hashmap.h
Description: Faster/better hash maps, taken from CF

<div align="right">09a72f, 17 lines</div>

```cpp
#include <bits/extc++.h>
struct splitmix64_hash {
  static uint64_t splitmix64(uint64_t x) {
    x += 0x9e3779b97f4a7c15;
```

```cpp
    x = (x^(x >> 30)) * 0xbf58476d1ce4e5b9;
    x = (x^(x >> 27)) * 0x94d049bb133111eb;
    return x^(x >> 31);
  }
  size_t operator()(uint64_t x) const {
    static const uint64_t FIXED_RANDOM = std::chrono::
        steady_clock::now().time_since_epoch().count();
    return splitmix64(x + FIXED_RANDOM);
  }
};
template <typename K, typename V, typename Hash =
    splitmix64_hash>
using hash_map = __gnu_pbds::gp_hash_table<K, V, Hash>;
template <typename K, typename Hash = splitmix64_hash>
using hash_set = hash_map<K, __gnu_pbds::null_type, Hash>;
```

unrolling.h

<div align="right">520e76, 5 lines</div>

```cpp
#define F {...; ++i;}
int i = from;
while (i&3 && i < to) F // for alignment, if needed
while (i + 4 <= to) { F F F F }
while (i < to) F
```

fast-mod.h
Description: Compute $a\%b$ about 4 times faster than usual, where b is constant but not known at compile time. Fails for $b = 1$.

<div align="right">c977c5, 10 lines</div>

```cpp
typedef unsigned long long ull;
typedef __uint128_t L;
struct FastMod {
  ull b, m;
  FastMod(ull b) : b(b), m(ull((L(1) << 64) / b)) {}
  ull reduce(ull a) {
    ull q = (ull)((L(m) * a) >> 64), r = a - q * b;
    return r >= b ? r - b : r;
  }
};
```

custom-comparator.h

<div align="right">1e3970, 6 lines</div>

```cpp
auto cmp = [](const kind_t& a, const kind_t& b) {
    return a.func() < b.func();
};
set<kind_t, decltype(cmp)> my_set(cmp);
map<kind_t, int, decltype(cmp)> my_map(cmp);
priority_queue<kind_t, vector<kind_t>, decltype(cmp)> my_pq(cmp
    );
```

10.6   Bit Twiddling Hack

hacks.h

<div align="right">829b7d, 21 lines</div>

```cpp
// Iterate on non-empty submasks of a bitmask.
for (int s = m; s > 0; s = (m & (s - 1)))
// Iterate on non-zero bits of a bitset B.
for (int j = B._Find_next(0); j < MAXV; j = B._Find_next(j))
ll next_perm(ll v) { // compute next perm i.e.
    ll t = v | (v-1);  // 00111,01011,01101,10011 ...
    return (t + 1) | (((~t & -~t) - 1)>>(__builtin_ctz(v) + 1));
}
template<typename F>  // All subsets of size k of {0..N-1}
void iterate_k_subset(ll N, ll k, F f){
    ll mask = (1ll << k) - 1;
    while (!(mask & 1ll<<N)) { f(mask);
      ll t = mask | (mask-1);
      mask = (t+1)|(((~t & -~t)-1)>>(__builtin_ctzll(mask)+1));
    }
}
template<typename F> // All subsets of set
```

```
void iterate_mask_subset(ll set, F f){ ll mask = set;
  do  f(mask), mask = (mask-1) & set;
  while (mask != set);
}
```

bitset.h
Description: Some bitset functions
b9f55a, 17 lines

```
int main() {
  bitset<100> bt;
  cin >> bt;
  cout << bt[0] << "\n";
  cout << bt.count() << "\n"; // number of bits set
  cout << (~bt).none() << "\n"; // return true if has no bits
      set
  cout << (~bt).any() << "\n"; // return true if has any bit
      set
  cout << (~bt).all() << "\n"; // retun true if has all bits
      set
  cout << bt._Find_first() << "\n"; // return first set bit
  cout << bt._Find_next(10) << "\n";// returns first set bit
      after index i
  cout << bt.flip() << '\n'; // flip the bitset
  cout << bt.test(3) << '\n'; // test if the ith bit of bt is
      set
  cout << bt.reset(3) << '\n'; // reset the ith bit
  cout << bt.set() << '\n'; // turn all bits on
  cout << bt.set(4, 1) << '\n'; // set the 4th bit to value 1
  cout << bt << "\n";
}
```

10.7   Random Numbers

random-numbers.h
Description: An example on the usage of generator and distribution. Use shuffle instead of random shuffle.
b96f2b, 8 lines

```
mt19937 rng(random_device{}());
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().
    count());
shuffle(permutation.begin(), permutation.end(), rng);
uniform_int_distribution<int> uid(1, 100); //[1, 100]
unsigned xrand() {
  static unsigned x = 314159265, y = 358979323, z = 846264338,
      w = 327950288;
  unsigned t = x ^ x << 11; x = y; y = z; z = w; return w = w ^
      w >> 19 ^ t ^ t >> 8;
}
```