

# COMP90054 – Week 7 tutorial

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## Best-Supporters and Relaxed plan heuristics

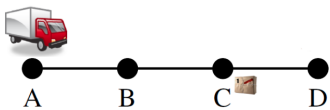
### Best-supporter functions $bs$

#### Definition

- Best-supporter function  $bs$ : For every fact  $p \in (F \setminus s)$ , returns an action that can achieve  $p$  with the cheapest cost within relaxation
- In this subject we use the two following closed well-founded best-supporter functions
  - $bs_s^{\text{add}}(p) := \underset{a \in A, p \in \text{add}_a}{\text{argmin}} \left( c(a) + h^{\text{add}}(s, \text{pre}_a) \right)$
  - $bs_s^{\text{max}}(p) := \underset{a \in A, p \in \text{add}_a}{\text{argmin}} \left( c(a) + h^{\text{max}}(s, \text{pre}_a) \right)$

#### Example

• “Logistics”:



■ Initial state  $I: t(A), p(C)$ .  
 ■ Goal  $G: t(A), p(D)$ .  
 ■ Actions  $A: dr(X, Y), lo(X), ul(X)$ .

	$t(A)$	$t(B)$	$t(C)$	$t(D)$	$p(T)$	$p(A)$	$p(B)$	$p(C)$	$p(D)$
$h^{\text{add}}$	0	1	2	3	3	4	5	0	7
$bs_s^{\text{add}}$	-	$dr(A, B)$	$dr(B, C)$	$dr(C, D)$	$lo(C)$	$ul(A)$	$ul(B)$	-	$ul(D)$

- For  $p = t(B)$ ,  $bs_s^{\text{add}}(p) = \underset{a \in \{dr(A, B), dr(C, B)\}}{\text{argmin}} \left( c(a) + h^{\text{add}}(s, \text{pre}_a) \right)$ 
  - $a = dr(A, B) \rightarrow 1 + h^{\text{add}}(t(A)) = 1$
  - $a = dr(C, B) \rightarrow 1 + h^{\text{add}}(t(C)) = 3$

Relaxed plan heuristics  $h^{FF}$ 

## Idea

1. Compute a best-supporter function  $bs$
2. Extract a relaxed plan  $RPlan$  by applying  $bs$  to each  $g \in G$  and collecting all actions
3. Use the relaxed plan heuristics  $h^{FF} = \sum_{a \in RPlan} c(a)$  to estimate  $h^*$

## Algorithm

Relaxed Plan Extraction for state  $s$  and best-supporter function  $bs$ 

```

Open := G \ s; Closed := ∅; RPlan := ∅
while Open ≠ ∅ do:
  select g ∈ Open
  Open := Open \ {g}; Closed := Closed ∪ {g};
  RPlan := RPlan ∪ {bs(g)}; Open := Open ∪ (prebs(g) \ (s ∪ Closed))
endwhile
return RPlan

```

- Requires a closed well-founded best-supporter function to make sense

Properties of  $h^{FF}$ 

- $h^{FF}$  is pessimistic:  $h^{FF} \geq h^+$  so  $\exists \Pi, s$  such that  $h^{FF}(s) > h^*(s)$
- $h^{FF}$  agrees with  $h^+$  on  $\infty$ :  $h^{FF}(s) = \infty \Leftrightarrow h^+(s) = \infty$
- $h^{FF}$  does not over-count like  $h^{add}$

## Example



- Initial state  $I: t(A), p(C)$ .
- Goal  $G: t(A), p(D)$ .
- Actions  $A: dr(X, Y), lo(X), ul(X)$ .

- Logistics":

	$t(A)$	$t(B)$	$t(C)$	$t(D)$	$p(T)$	$p(A)$	$p(B)$	$p(C)$	$p(D)$
$bs_s^{add}$	-	$dr(A, B)$	$dr(B, C)$	$dr(C, D)$	$lo(C)$	$ul(A)$	$ul(B)$	-	$ul(D)$

- Relaxed plan extraction

g	Open	Closed	RPlan
-	{p(D)}	{}	{}
p(D)	{t(D), p(T)}	{p(D)}	{ul(D)}
t(D)	{p(T), t(C)}	{p(D), t(D)}	{dr(C, D), ul(D)}
t(C)	{p(T), t(B)}	{p(D), t(D), t(C)}	{dr(B, C), dr(C, D), ul(D)}
t(B)	{p(T)}	{p(D), t(D), t(C), t(B)}	{dr(A, B), dr(B, C), dr(C, D), ul(D)}
p(T)	{}	{p(D), t(D), t(C), t(B), p(T)}	{lo(C), dr(A, B), dr(B, C), dr(C, D), ul(D)}

- Relaxed plan heuristics  $h^{FF} = 1 + 1 + 1 + 1 + 1 = 5$

## Width-based Planning

### Iterated Width $IW$

- Uses the idea of novelty  $w$ :  $w(s)$  is the size of the smallest subset of atoms in  $s$  that is true for the first time in the search
- $IW(k)$  is a breadth-first search that prunes generated states  $s$  with  $w(s) > k$
- $IW$  is an iterative/sequence of calls to  $IW(k)$  for  $k = 0, 1, 2, \dots$  over problem  $P$  until problem solved or  $i$  exceeds number of variables in problem

### Properties of $IW$

- The order of search node expansion can affect which nodes are pruned
- $IW(k)$  expands at most  $O(n^k)$  states

### Example

Left->right order

- $IW(1)$ : {1,2,3}
- $IW(2)$ : {1,2,3,4}
- $IW(3)$ : {1,2,3,4}

Right->left order

- $IW(1)$ : {1,3}
- $IW(2)$ : {1,3,2,5}
- $IW(3)$ : {1,3,2,5,4}

