MapReduce

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Review

Common Big Data Challenges

Volume:

- Too much data for processing → reduce processing time by using distributed & parallel computing
 - Performing word counting on 300 articles simultaneously
- Solution: Distributed & Parallel Computing

Parallel Computing

- Make use of parallel architectures (e.g. multi-core, multi-processor, clusters of machines, etc.) to improve computing performance
 - Using multiple cores to speedup video processing (Facetime HD!)
 - Using multiple processors in web servers
 - Using GPUs to increase rendering frame rate in games
 - Using clusters to run simulations
- Have a long history in High Performance Computing
 - Similar to streaming, getting popular in the Big Data era

Python's Core HOF

- map(): applies a function over an iterable to produce a new iterable
- map(lambda x: int(x)+1, ['0', '1']) -> [1, 2]
- filter(): return only values that satisfy a predicate in the new iterable
- filter(lambda x: x<1, [0, 1]) -> [0]
- reduce(): accumulate a sequence of values from left to right
- reduce(lambda x, y: x+y, xrange(5), 100) -> 110
- sorted(): return a new sorted list using a comparator function
 sorted(xrange(5))
 sorted(xrange(5), cmp=lambda x,y:y-x) -> [4,3,2,1,0]

Outline

- MapReduce
- Designing MapReduce Algorithms

MapReduce

- By definition, big data is too large to handle by conventional means.
 Sooner or later, you just can't scale up anymore
 - machines cannot grow too large...
- But we can add more computers!
- MapReduce paradigm allows us to scale out even with commodity hardware
 - first proposed in the big data world by Google Jeffrey Dean and Sanjay Ghemawat.
 - MapReduce: Simplified Data Processing on Large Clusters. 2004.

- A programming paradigm (for big data processing)
 - Data is split into distributable chunks
 - Perform a series of transformations on those chunks
 - Transformations are run in parallel on the chunks (data parallelism)
- MapReduce is scalable by adding more machines to process chunks
 - scaling out on commodity hardware
- The foundation for Hadoop, which is an implementation of MapReduce

- A programming paradigm that process data in 2 phases/operations: map(), and then reduce()
 - instead of map(), filter(), reduce(), sort() like in the previous class, we are limited to only map() and reduce()
- In a nutshell that's it:
 - Provide a data collection (separable records)
 - Apply a user-defined map function on each data record
 - Then reduce the mapped output with another user-defined function

Fixed

pipeline

MapReduce works on (key,value) pairs

```
INPUT: list of key-value pairs of (k1,v1)
MAP: (k1,v1) \rightarrow [list of (k2,v2)]
SHUFFLE: combine (k2,v2) \rightarrow (k2, [list of v2])
REDUCE: (k2, [list of v2]) \rightarrow (k3,v3)
OUTPUT: list of (k3,v3)
```

executed in parallel

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User-defined functions
```

Input

- Data must be separable into records
 - Lines of text
 - Rows of tables
 - CSV: yes JSON, XML: no
- Key/value pairs
 - Key = line number, record index
 - Value = text string, row data
- Keys are mostly ignored in many cases (e.g. just passing text lines)

Map Phase

- Transform each input record using a user-defined function
 - \circ (k1,v1) \rightarrow [(k2,v2),...] could be an empty list
- one-to-many relationship
 - vs. one-to-one of Python's map() higher order function
 - Python's filter() is part of the Map phase
- Mainly process by streaming data chunks
 - [(k1,v1)] → [(k2,v2),...]
 - o generators → generators

Shuffle and Sort Phase

- For each key k2, collect all v2 from outputs of all map processes shuffling into (k2, [v2])
 - Done by distributed partitioning and local sort
- Outputs are (k2, [v2]) being distributed to "reducers"
- [(k2,v2)] → (k2, [v2])
- **(k2,v2)** are intermediate key/value pairs

Reduce Phase

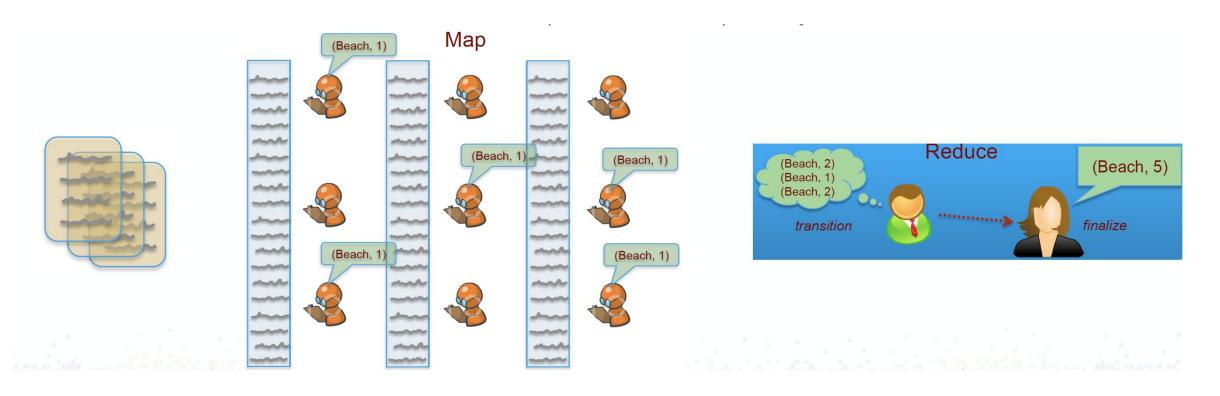
- Transform the output of the Shuffle phase using a user-defined function
 - $\circ (k2,[v2]) \rightarrow (k3, v3)$
- Equivalent to reduceByKey() ~ groupByKey() x reduce()
 - vs. Python's reduce()
- Also process in streaming data chunks
 - [(k2, [v2])] → [(k3, v3)]
- This phase is optional, can be omitted.

The Classic Example: WordCount

- The "Hello World" of MapReduce
- Given a set of documents, compute the frequency of each word
- INPUT (k1, v1): (document: line number, text line)
- OUTPUT (k3, v3): (word, frequency)
- INTERMEDIATE (k2, v2): (word, 1)
 - Transform each line to a list of words and their frequencies (=1)
 - Combine all tuples by words, and add all the frequency together

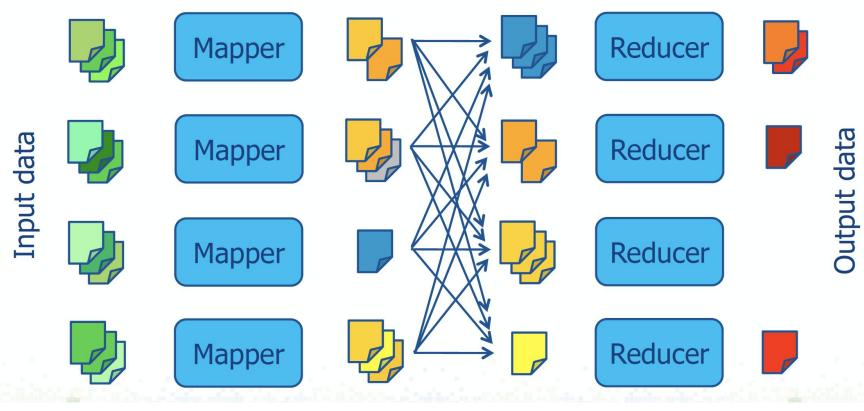
The Classic Example: WordCount

- The "Hello World" of MapReduce
- Given a set of documents, compute the frequency of each word



MapReduce Dataflow

Intermediate (key,value) pairs



The Shuffle

When to use MapReduce?

- Data parallelism problems ("embarrassingly parallel")
 - Word count
 - Distributed grep
 - Reverse index
 - Document OCR (Optical character recognition)
- Data must be splittable into chunks and records

Designing MapReduce Algorithms

Designing MapReduce Algorithms

- Key decision: What should be done by map, and what by reduce?
 - map can do something to each individual key-value pair, but it can't look at other key-value pairs
- Example: Filtering out key-value pairs we don't need
 - map can emit more than one intermediate key-value pair for each incoming key-value pair
- Example: Incoming data is text, map produces (word,1) for each word

Designing MapReduce Algorithms

- Key decision: What should be done by map, and what by reduce?
 - reduce can aggregate data; it can look at multiple values, as long as map has mapped them to the same (intermediate) key
 - Example: Count the number of words, add up the total cost, ...
- Need to get the intermediate format right!
 - If reduce needs to look at several values together, map must emit them using the same key!

Filtering with MapReduce

- Goal: find lines/files/tuples with a particular characteristic
- Examples:
 - grep Web logs for requests to a specific domain
 - find in the Web logs the hostnames accessed by a particular IP
 - locate all the files that contain the words 'Apple' and 'Jobs'
- Generally: map does most of the work, utilizing one-to-many relationships to discard unqualified records. reduce may simply be the identity

Aggregation with MapReduce

- Goal: compute the maximum, the sum, the average, ..., over a set of values
- Examples:
 - Count the number of requests to a website
 - Find the most popular domain
 - Average the number of requests per page per Web site
- Often: map may be simple or the identity, reduce does the aggregation

Combiners

- Certain functions can be decomposed into partial steps:
 - Can take counts of two sub-partitions, sum them up to get a complete count for the partition
 - Can take maxes of two sub-partitions, max them to get a complete max for the partition
- Multiple map jobs on the same machine may write to the same reduce key



Common Mistakes to Avoid

- Mapper and reducer should be stateless
- Don't use static variables after map +
- reduce return, they should remember
- nothing about the processed data!
- Reason: No guarantees about which
- key-value pairs will be processed by
- which workers!

```
HashMap h = new HashMap();
map(key, value) {
  if (h.contains(key)) {
    h.add(key,value);
    emit(key, "A");
  }
}
```

Wrong!

Common Mistakes to Avoid

- Mapper must not map too much data to the same key
- In particular, don't map everything to the same key!! Otherwise the reduce worker will be overwhelmed!
- It's okay if some reduce workers have more work than others

```
map(key, value) {
  emit("FOO", key + " " + value);
}
Wrong!
```