

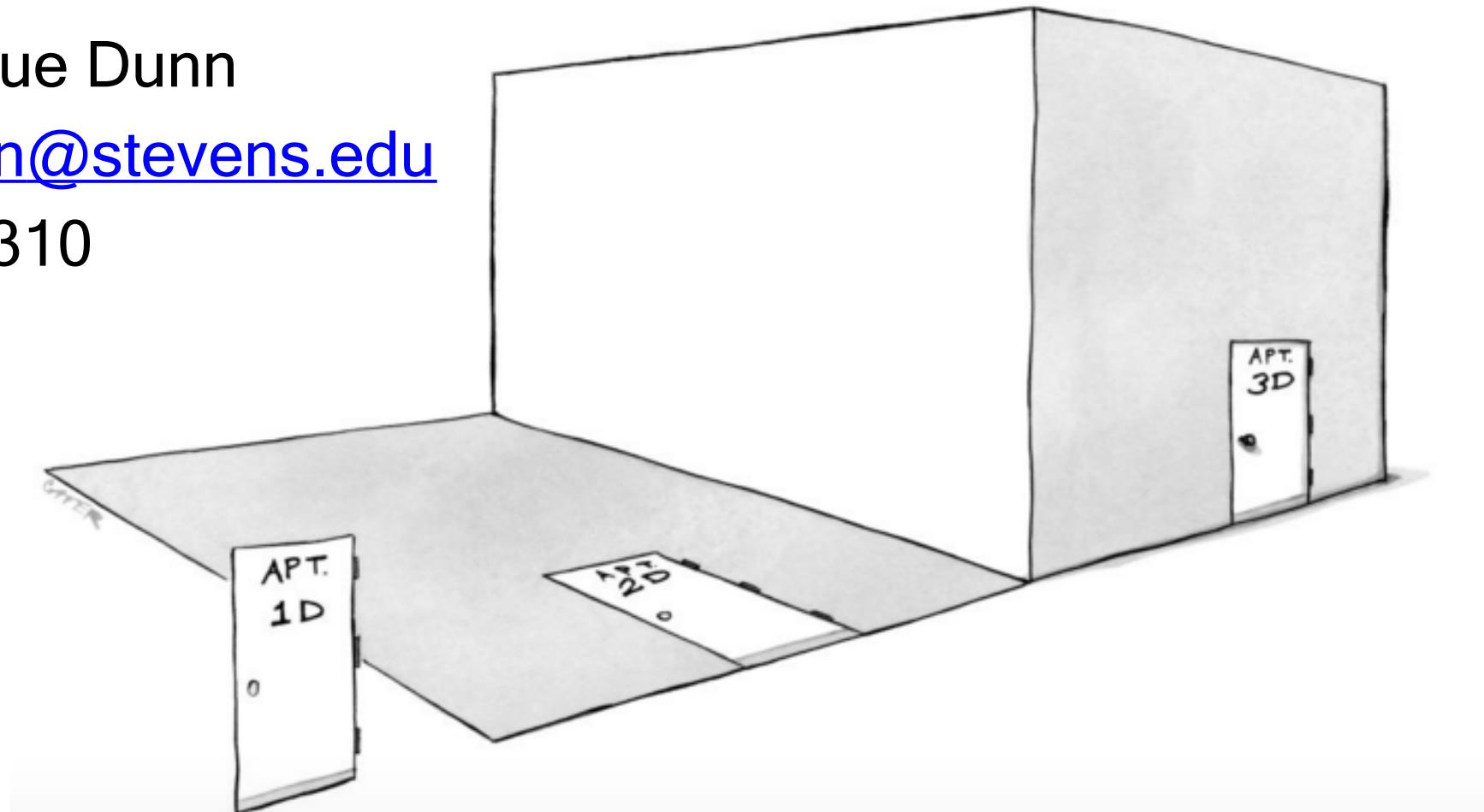
CS 532: 3D Computer Vision

Lecture 10

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Lieb 310



Homework # 4 available on Canvas

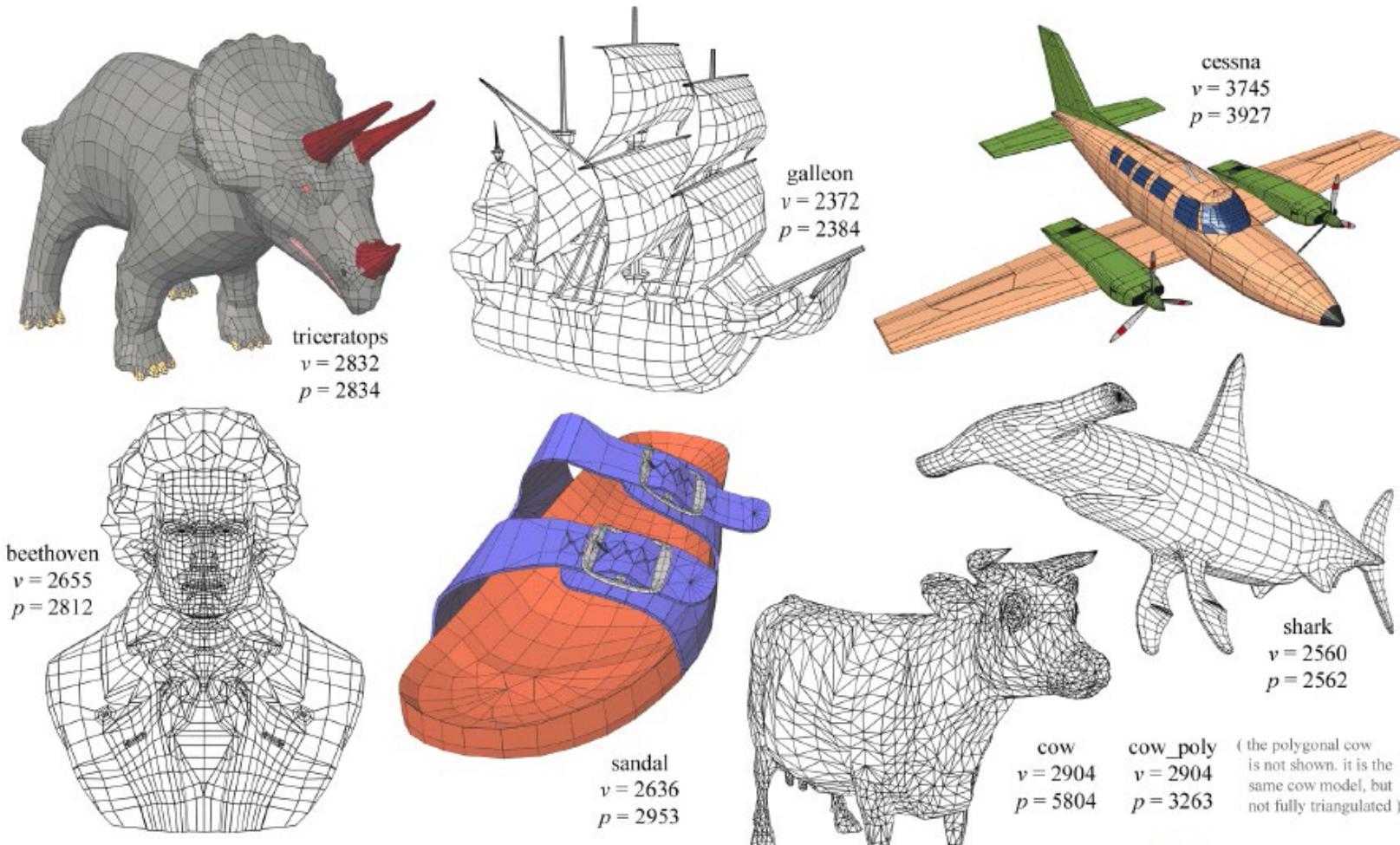
Due Nov 22

Lecture Outline

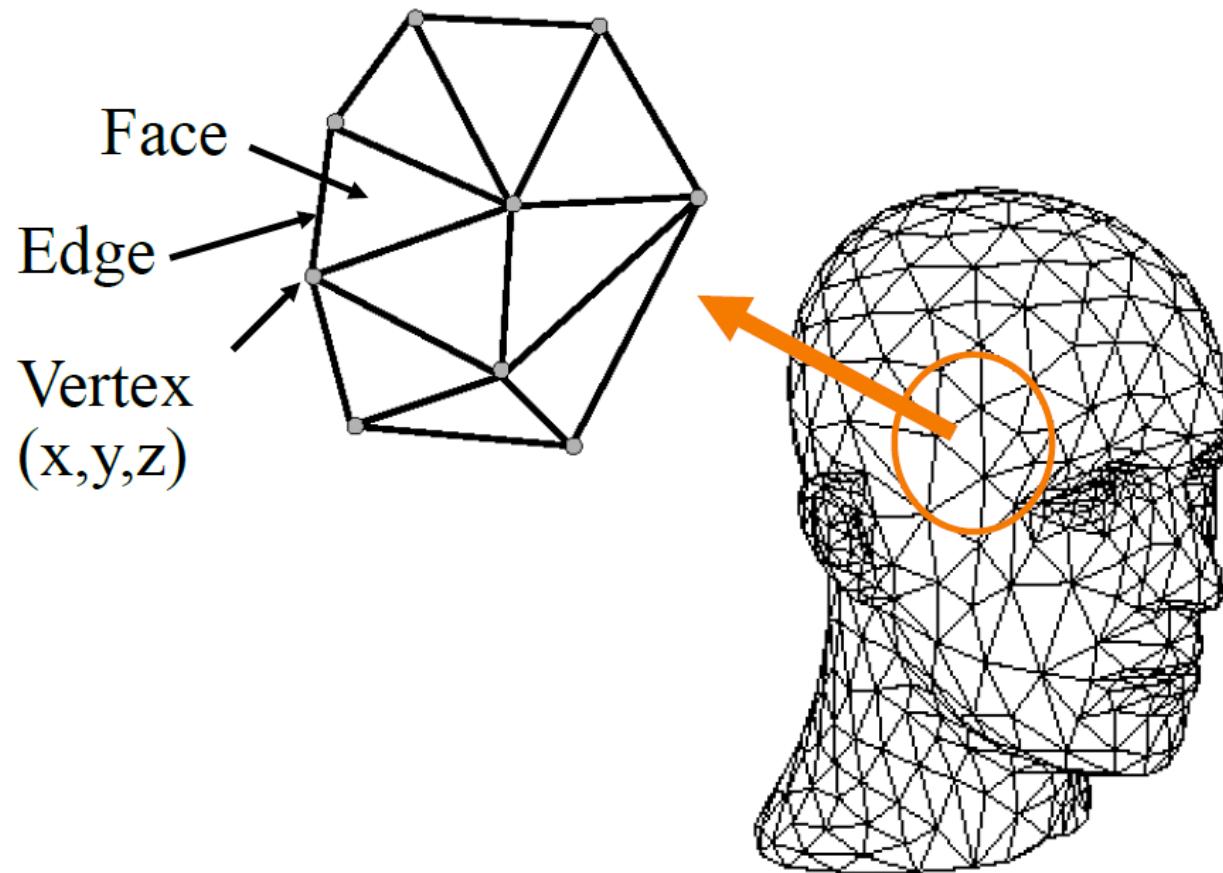
- Meshes
- Slides by:
 - S. Rusinkiewicz, T. Liu and V. Kim (Princeton University)
 - **Ching-Kuang Shene (Michigan Tech. Univ.)**
- David M. Mount, CMSC 754: Computational Geometry lecture notes, Department of Computer Science, University of Maryland, Spring 2012
 - Lecture 22

3D Polygonal Mesh

- Set of polygons representing a 2D surface embedded in 3D

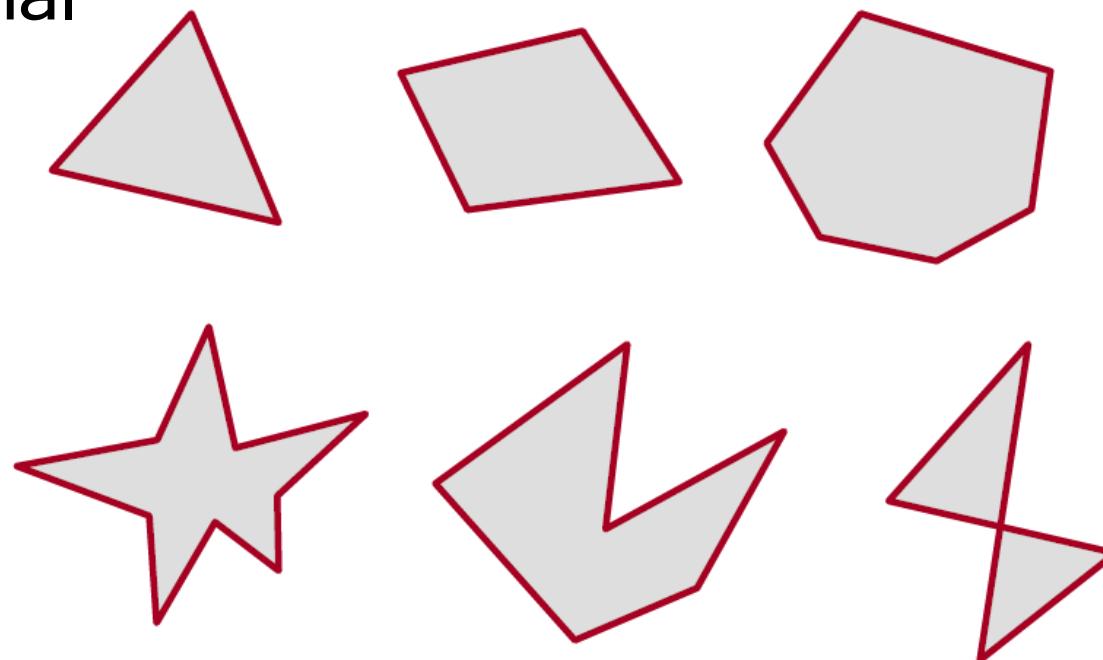


3D Polygonal Mesh



3D Polygon

- Region “inside” a sequence of coplanar points
- Points in counter-clockwise order
 - Define normal

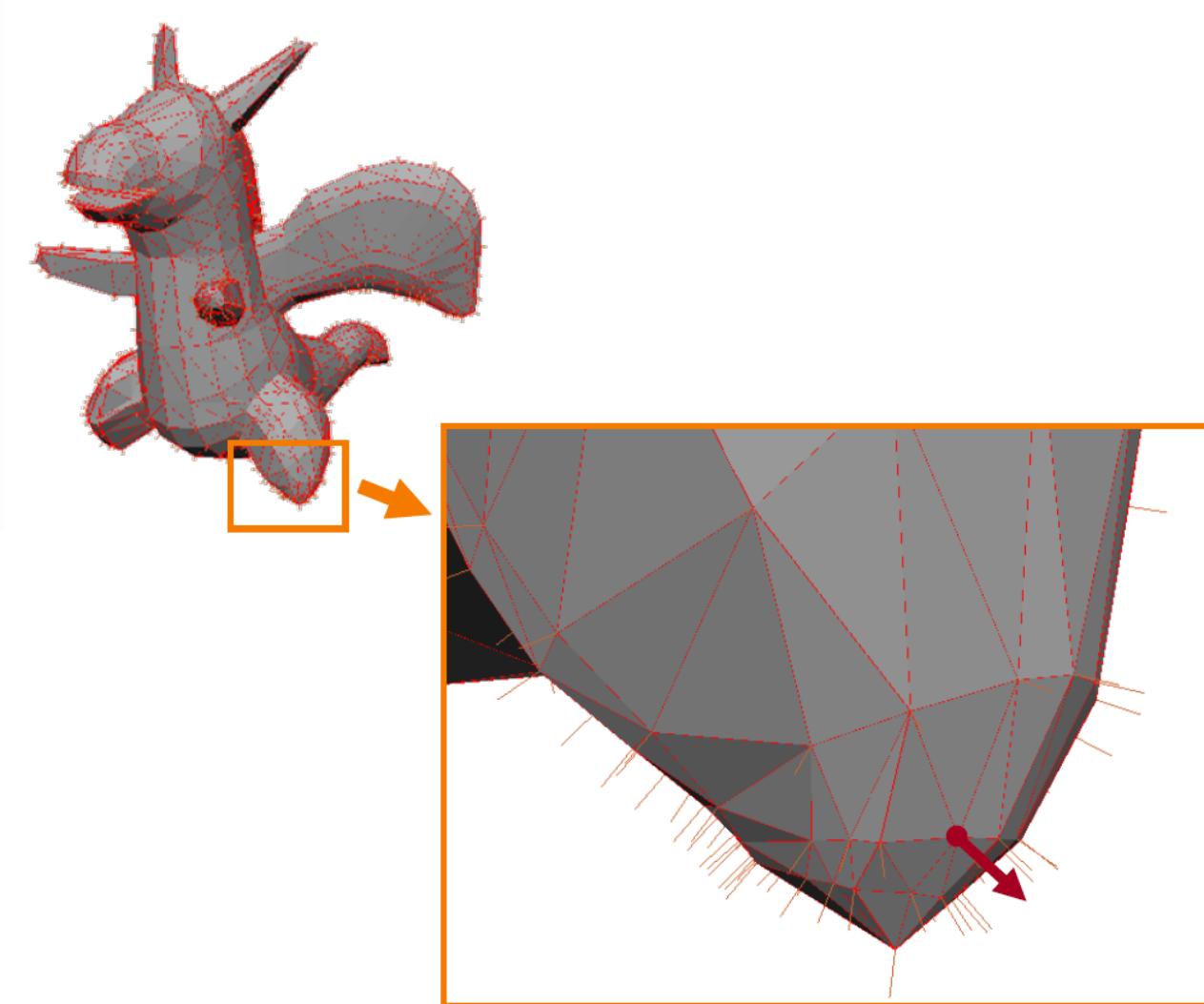


3D Polygonal Meshes

Why are they of interest?

- Simple, common representation
- Rendering with hardware support
- Output of many acquisition tools
- Input to many simulation/analysis tools

Surface Normals



Curvature

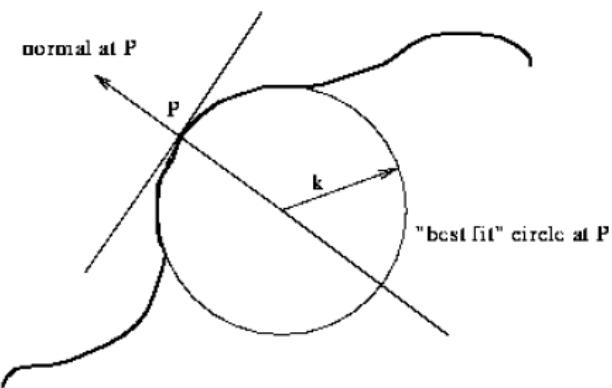
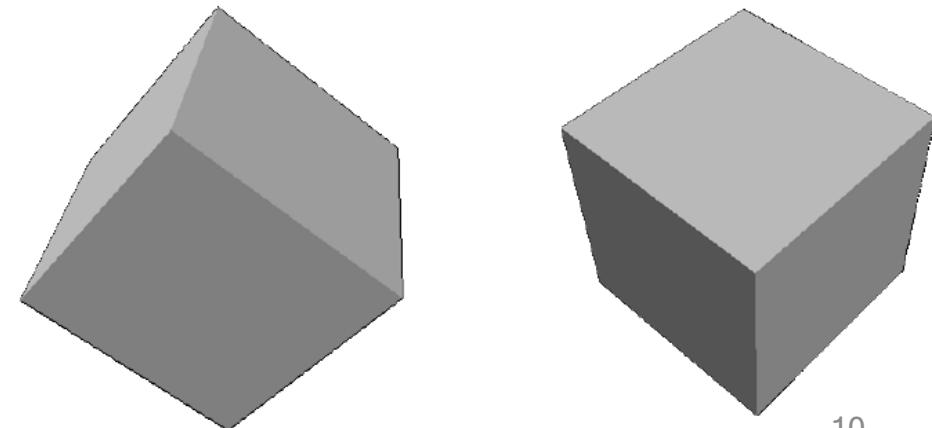
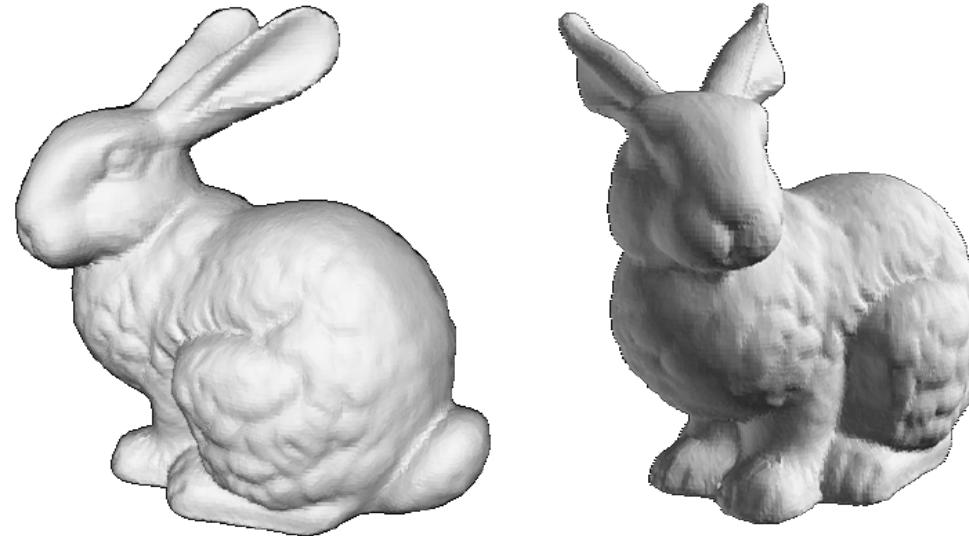


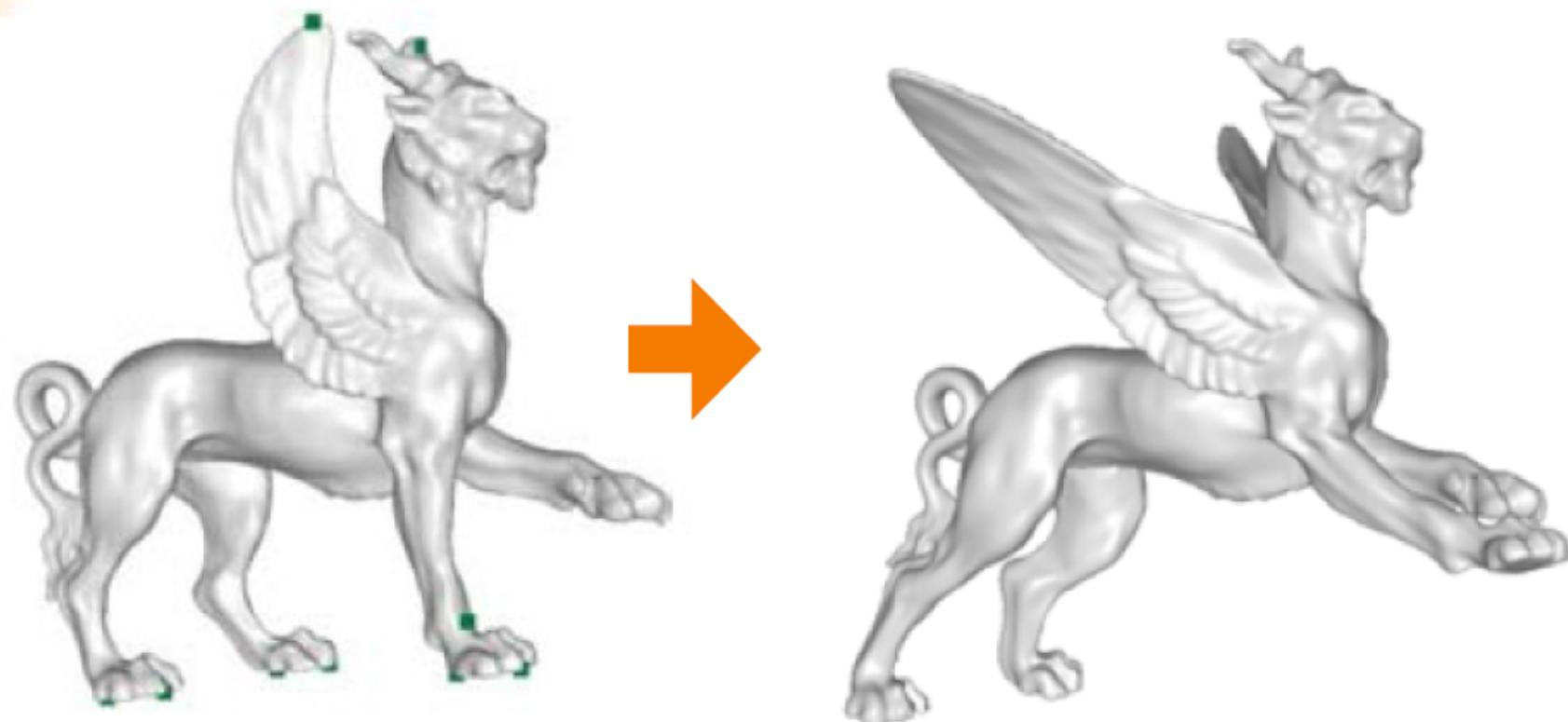
Figure 32: curvature of curve at P is $1/k$

Rigid Transformations

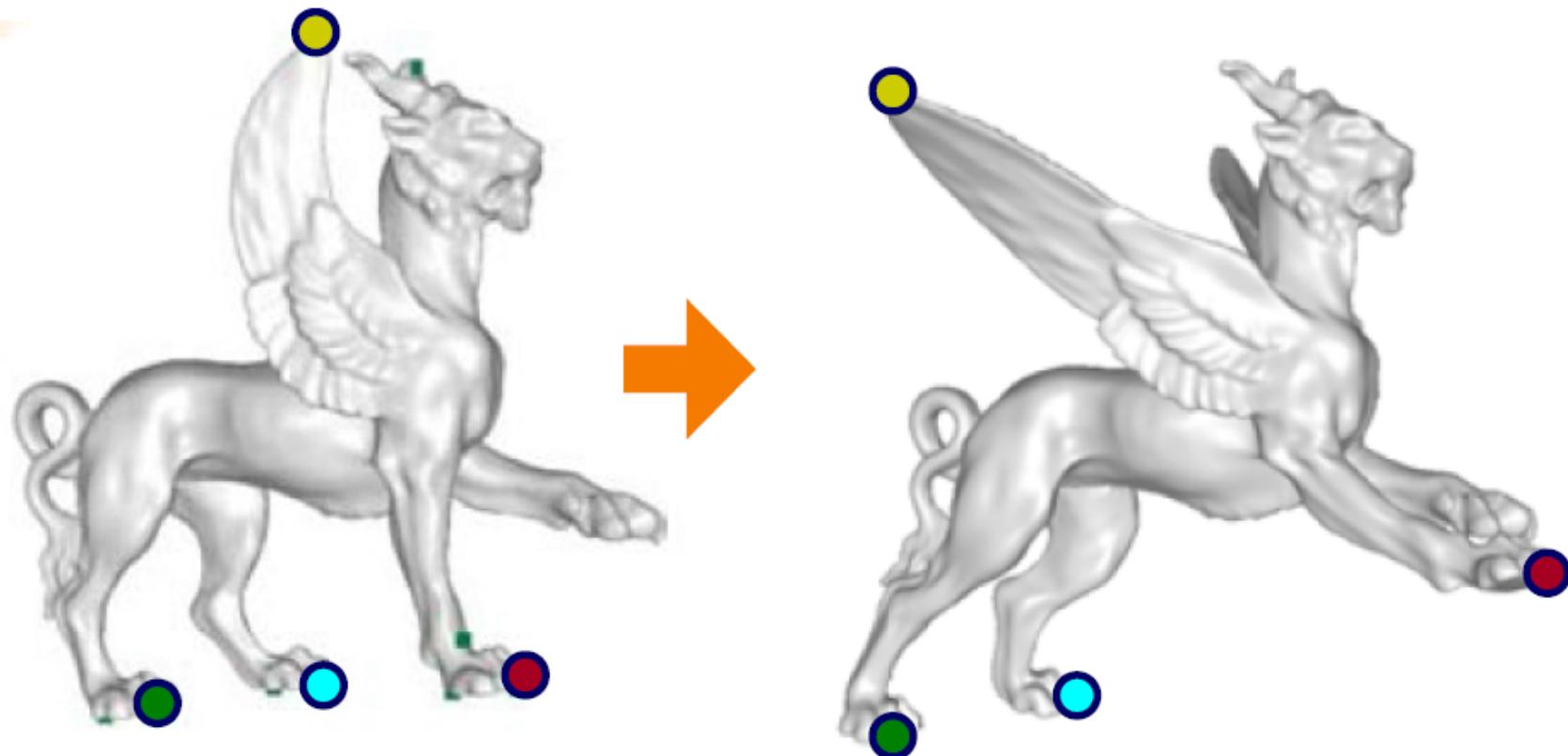
- Compare with implicit representations
 - level sets



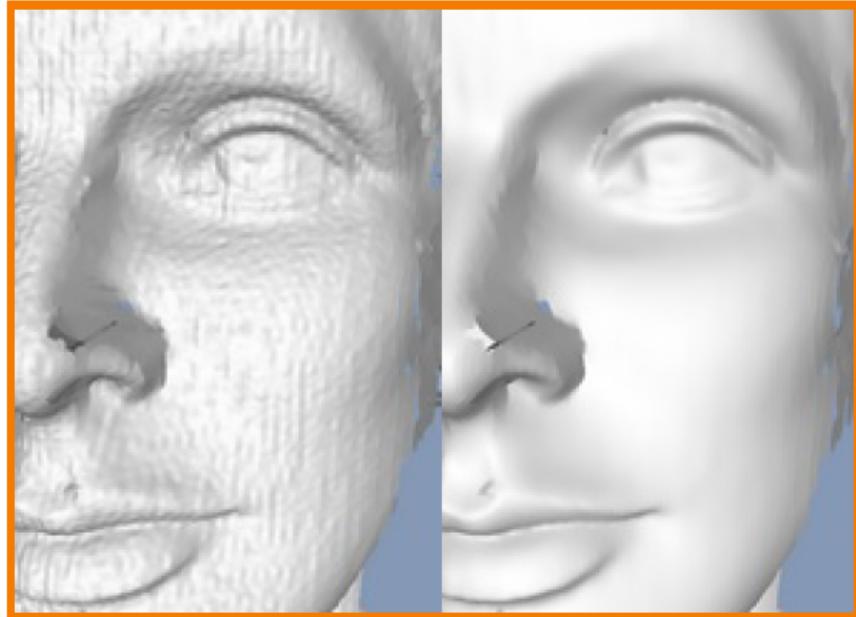
Deformations



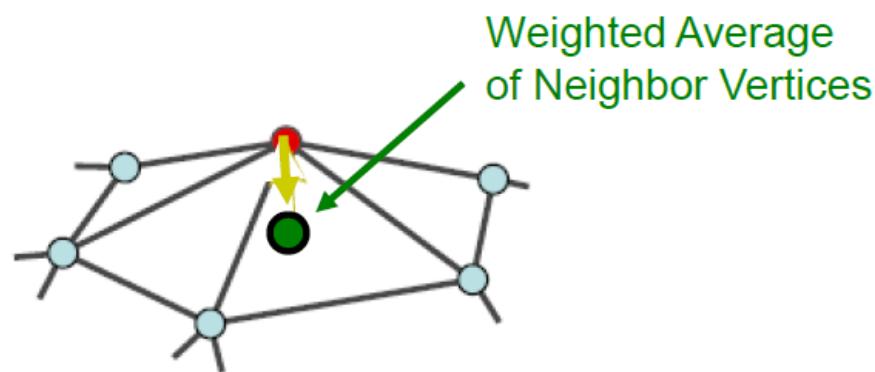
Deformations



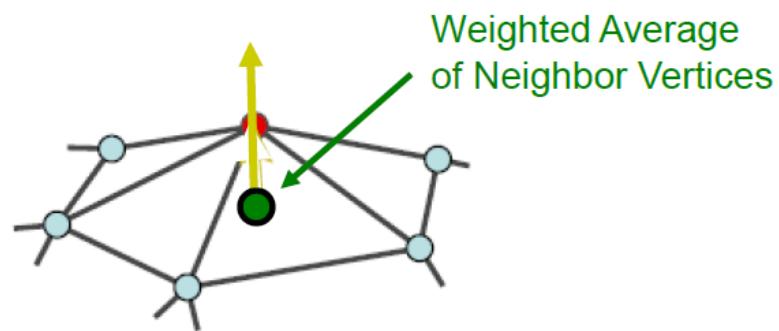
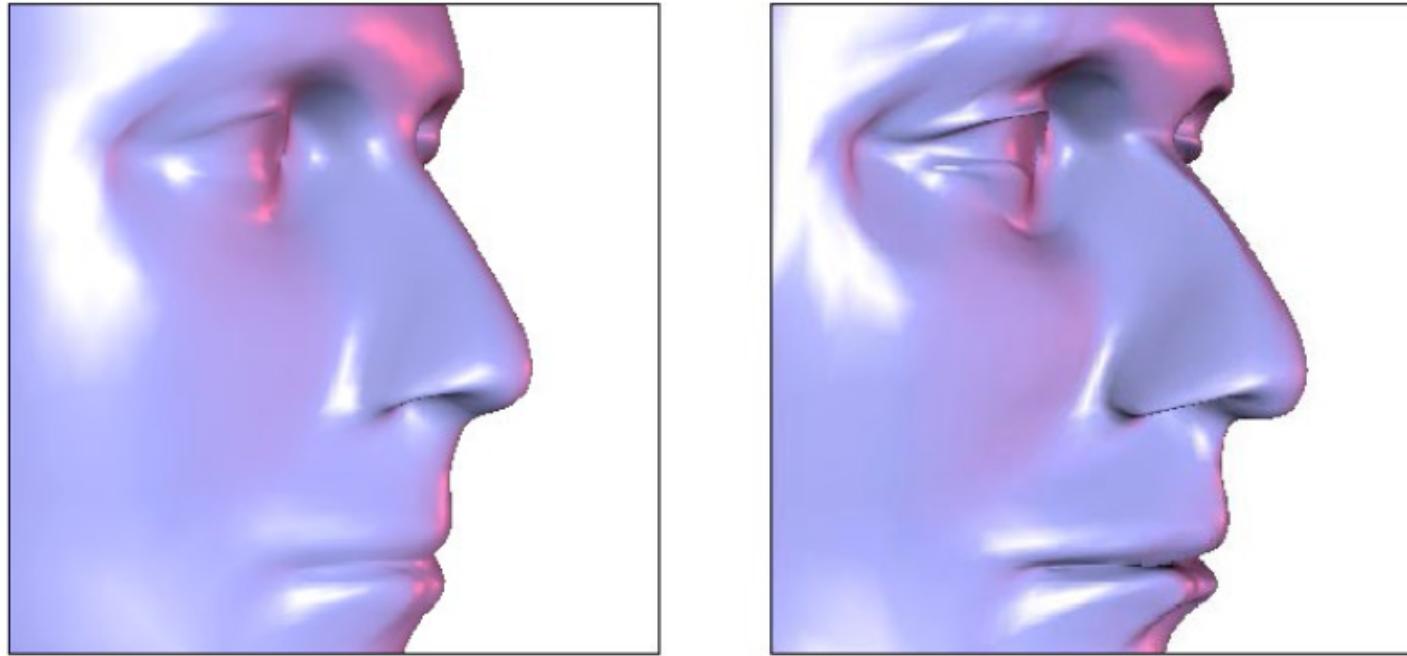
Smoothing



Thouis “Ray” Jones



Sharpen

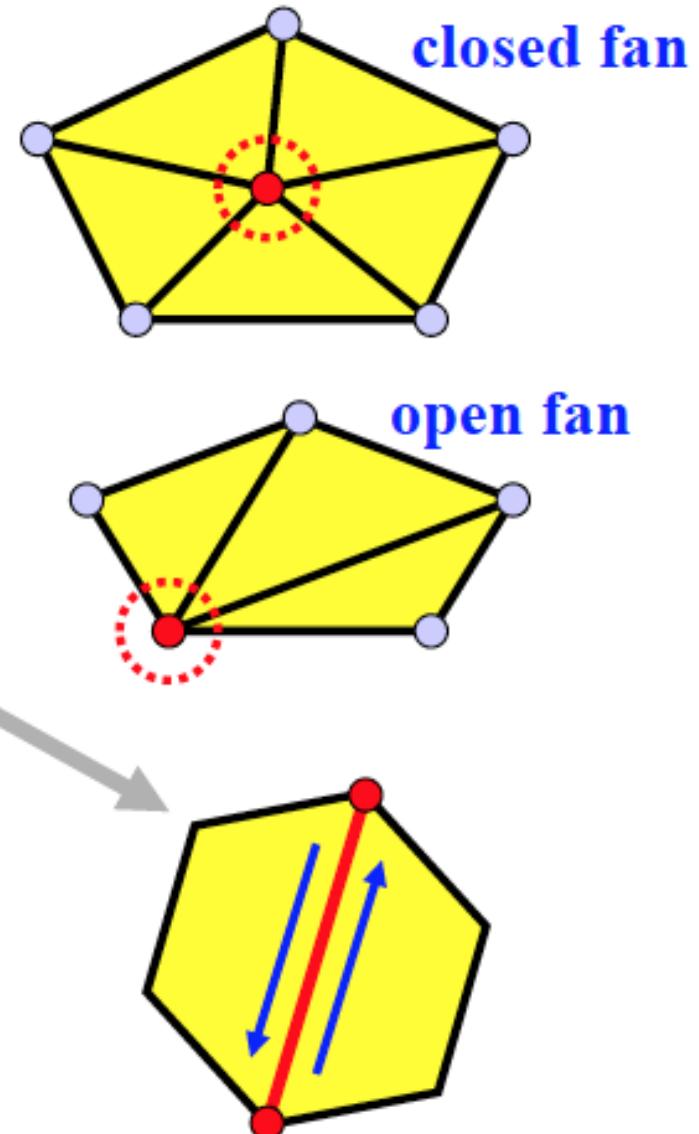


Definitions

- A **polygonal mesh** consists of three kinds of mesh elements: **vertices**, **edges**, and **faces**.
- The information describing the mesh elements are mesh **connectivity** and mesh **geometry**.
- The **mesh connectivity**, or topology, describes the incidence relations among mesh elements (e.g., adjacent vertices and edges of a face, etc).
- The **mesh geometry** specifies the position and other geometric characteristics of each vertex.

Definitions

- A polygonal mesh is a **manifold** if
 - Each edge is incident to only one or two faces, and
 - The faces incident to a vertex form a **closed** or an **open** fan.
- The orientation of a face is a cyclic ordering of the incident vertices.
- The orientation of a pair of adjacent faces is **compatible**, if the two vertices of the single common edge are in opposite order.
- A manifold mesh is **orientable** if any two adjacent faces have compatible orientation



Definitions

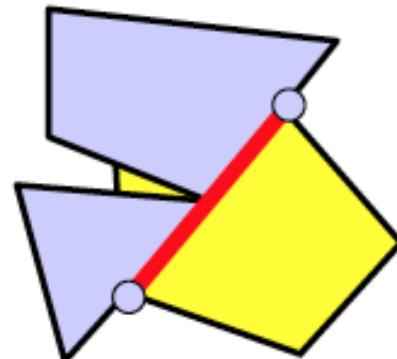
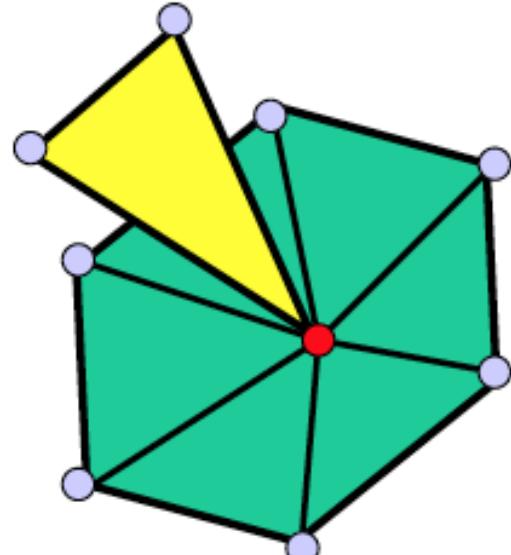
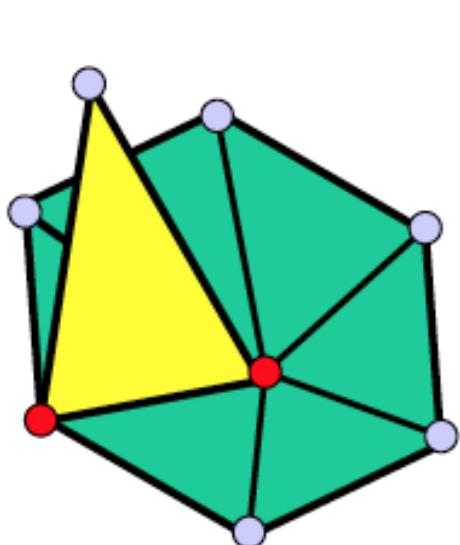
- Boundary edge: adjacent to exactly 1 face
- Regular edge: adjacent to exactly 2 faces
- Singular edge: adjacent to more than 2 faces
- Closed mesh: mesh with no boundary edges

Another Mesh Definition...

- A set of finite number of **closed polygons**
 - Intersection of inner polygonal areas is empty
 - Intersection of 2 polygons from is either empty, a point or an edge
 - Every edge belongs to at least one polygon
 - The set of all edges which belong only to one polygon are called edges of the polygonal mesh and are either empty or form a single closed polygon

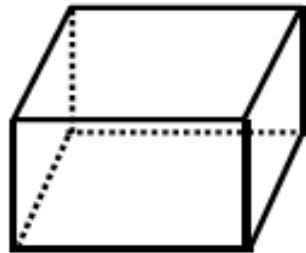
Non-Manifold Meshes

- Manifold Conditions:
 - Each edge is incident to only one or two faces,
 - The faces incident to a vertex form a closed or an open fan.
- The following examples are non-manifold meshes!

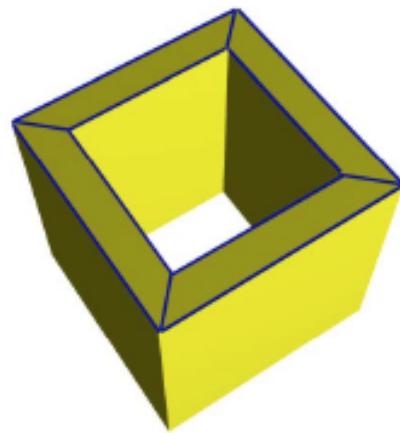


Euler-Poincaré Characteristic

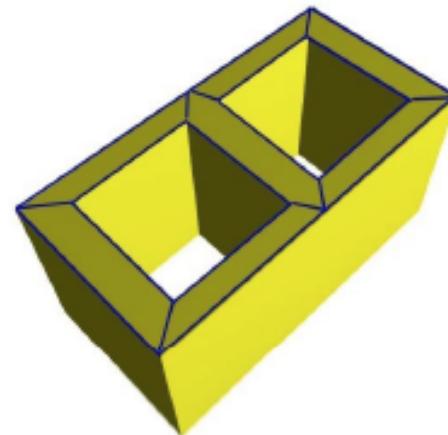
- Given a 2-manifold mesh \mathbf{M} without boundary, the Euler-Poincaré characteristic of \mathbf{M} is $\chi(\mathbf{M}) = V - E + F$, where V , E and F are the number of vertices, number of edges, and number of faces.



$$V=8, E=12, F=6$$
$$\chi(\mathbf{M}) = V - E + F = 2$$



$$V=16, E=32, F=16$$
$$\chi(\mathbf{M}) = V - E + F = 0$$

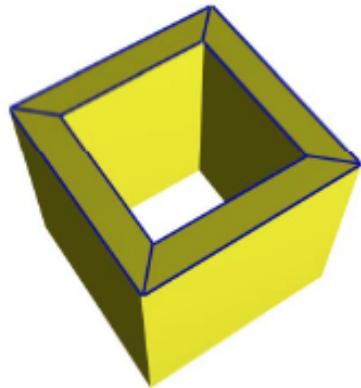


$$V=28, E=56, F=26$$
$$\chi(\mathbf{M}) = V - E + F = -2$$

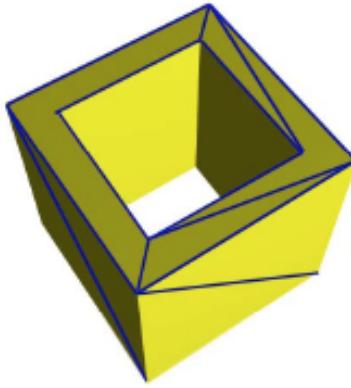
Euler-Poincaré Characteristic

- Euler-Poincaré characteristic $\chi(\mathbf{M}) = V-E+F$ is independent of tessellation.

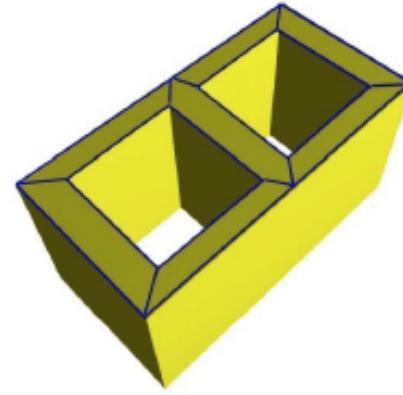
$$V=24, E=48, F=22$$
$$\chi(\mathbf{M}) = V-E+F = -2$$



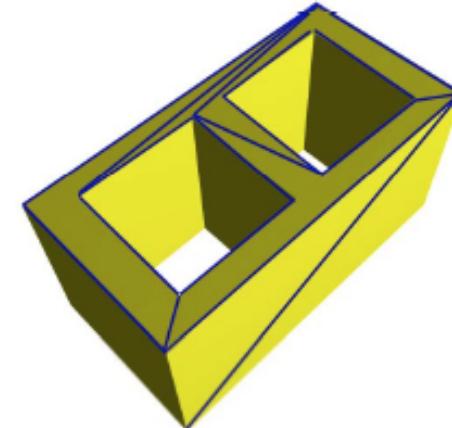
$$V=16, E=32, F=16$$
$$\chi(\mathbf{M}) = V-E+F=0$$



$$V=16, E=36, F=20$$
$$\chi(\mathbf{M}) = V-E+F=0$$

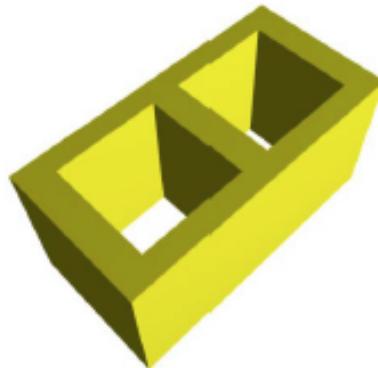


$$V=28, E=56, F=26$$
$$\chi(\mathbf{M}) = V-E+F=-2$$

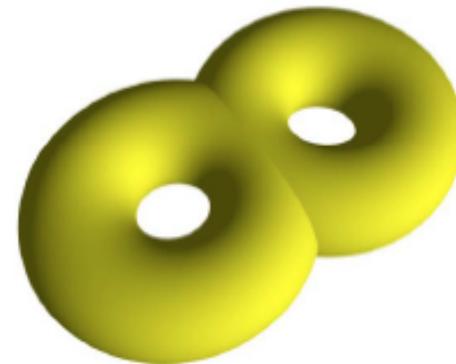


Homeomorphism

- Two 2-manifold meshes **A** and **B** are **homeomorphic** if their surfaces can be transformed to the other by twisting, squeezing, and stretching without cutting and gluing.
- Thus, boxes, spheres and ellipsoids are homeomorphic to each other.



is homeomorphic to



Homeomorphism

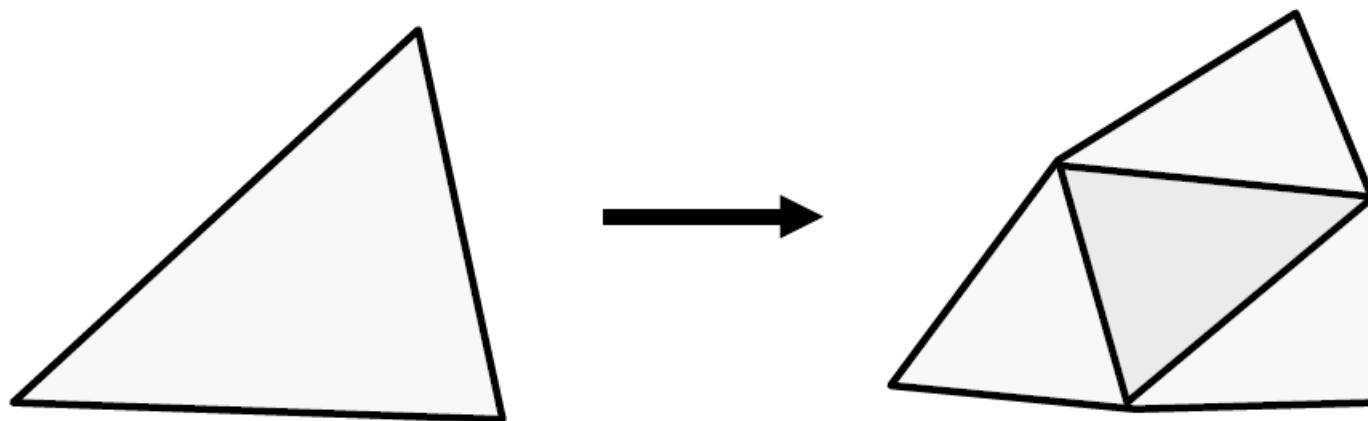
Two orientable 2-manifold meshes without boundary are *homeomorphic* if and only if they have the same Euler-Poincaré characteristic.

Low-level Operations

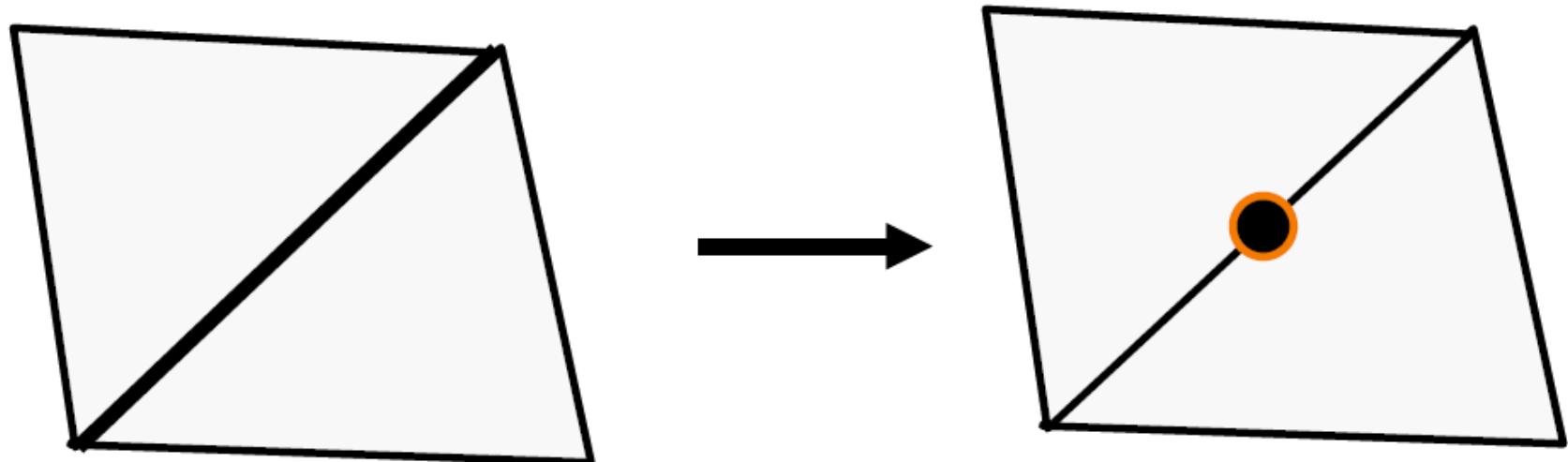
- Subdivide face
- Subdivide edge
- Collapse edge
- Merge vertices
- Remove vertex

Subdivide Face

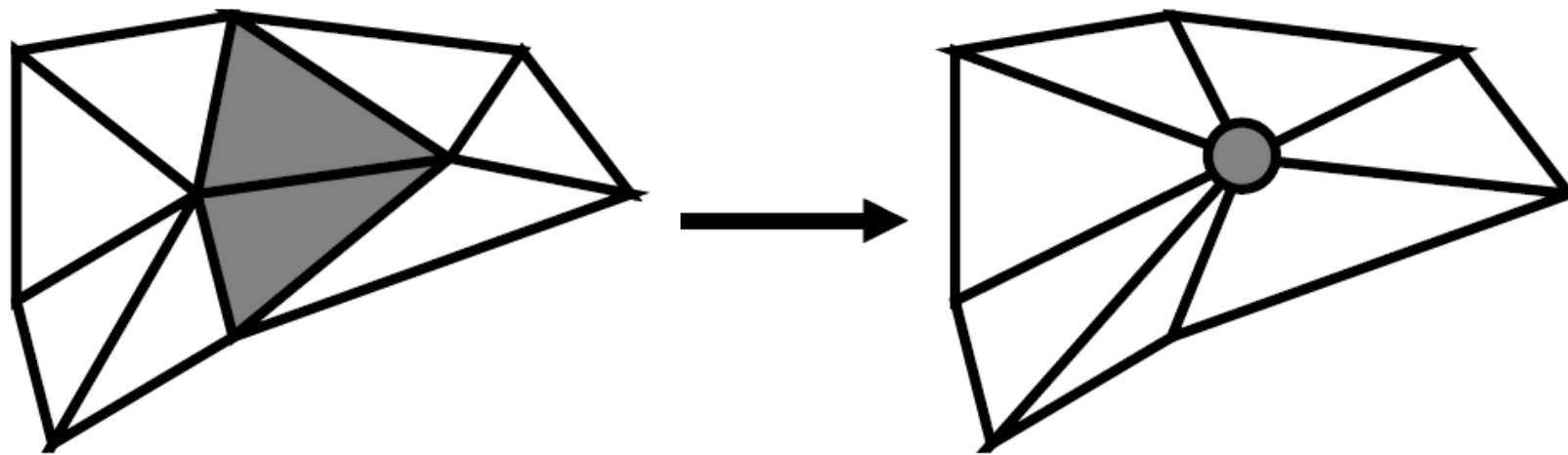
- How should we split current triangle?



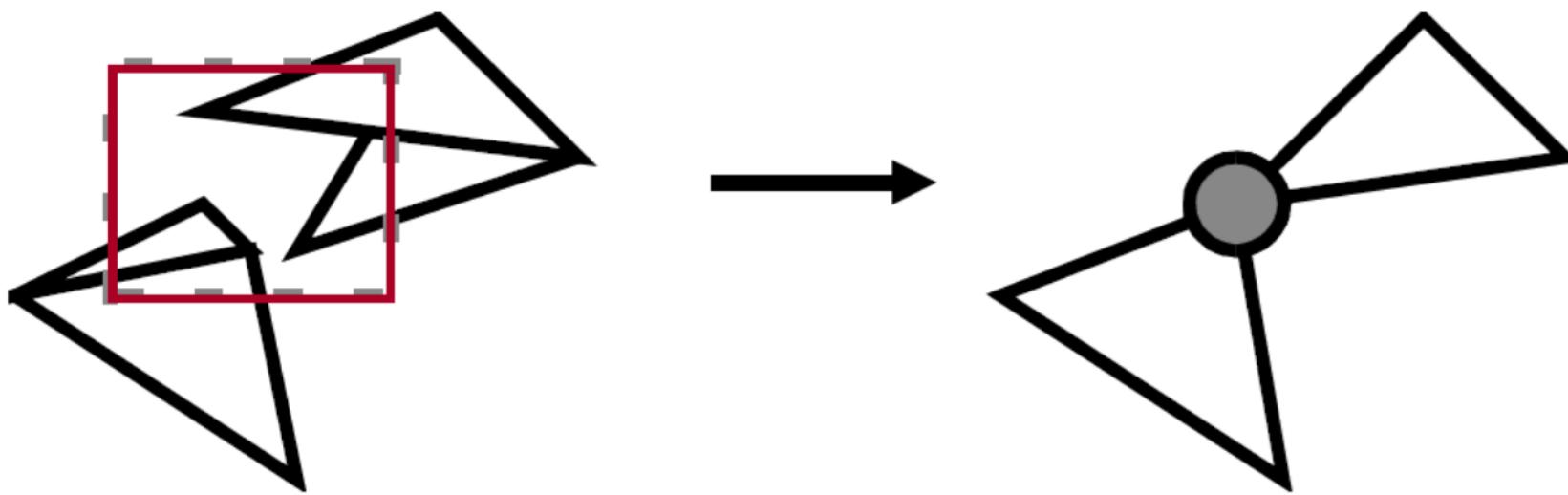
Subdivide Edge



Collapse Edge



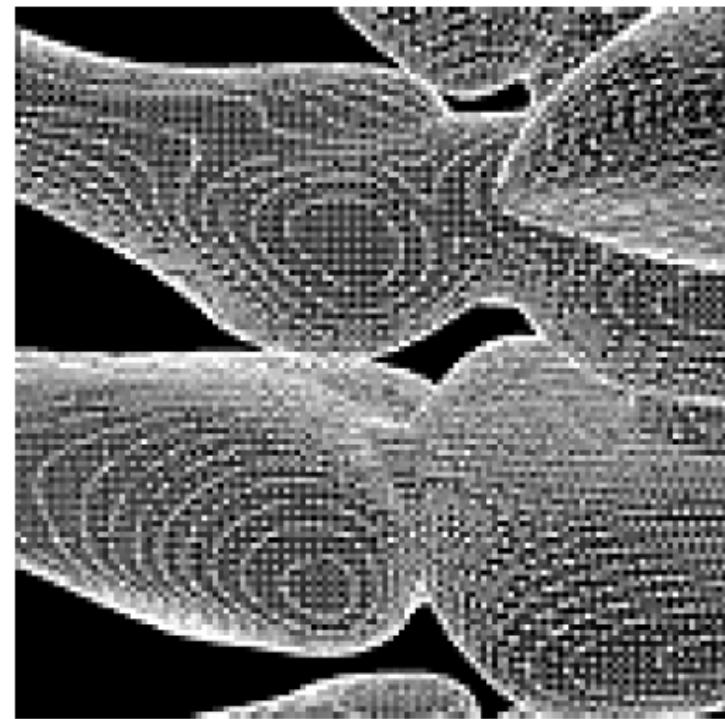
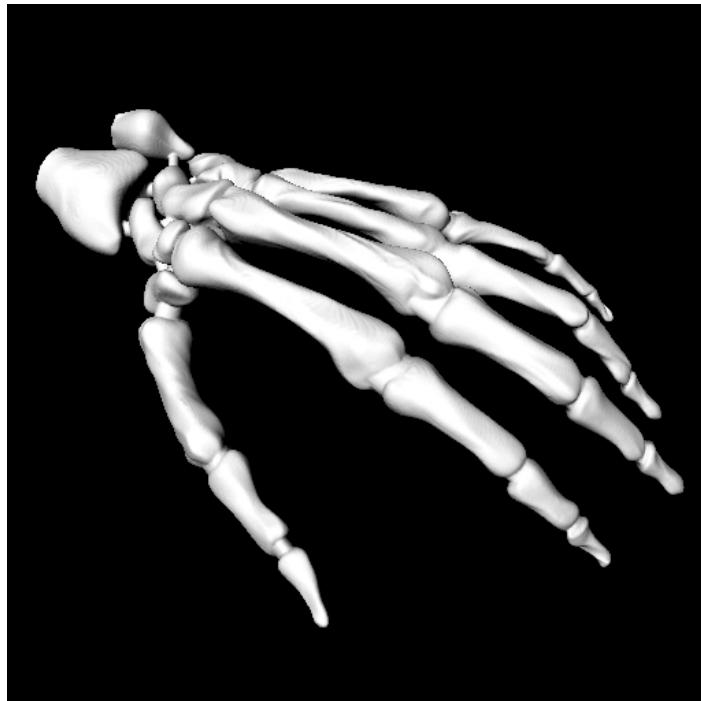
Merge Vertices



Polygonal Mesh Representation

Important properties of mesh representation

- Efficient traversal of topology
- Efficient use of memory
- Efficient updates

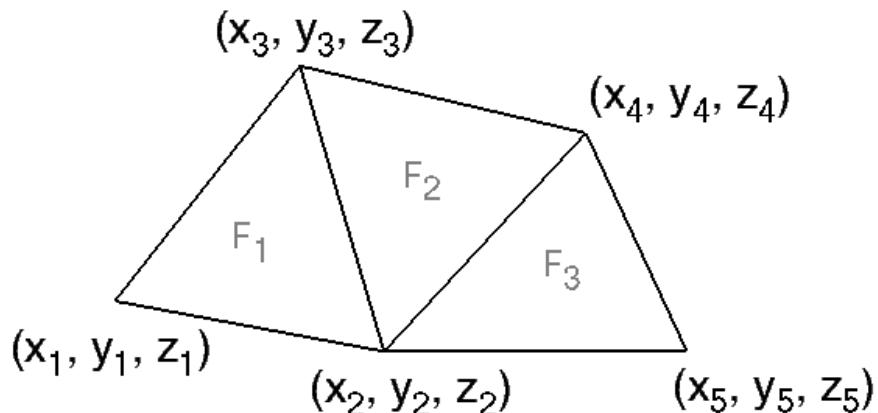


Possible Data Structures

- List of independent faces
- Vertex and face tables
- Adjacency lists
- Winged edge
- Half edge
- etc.

Independent Faces

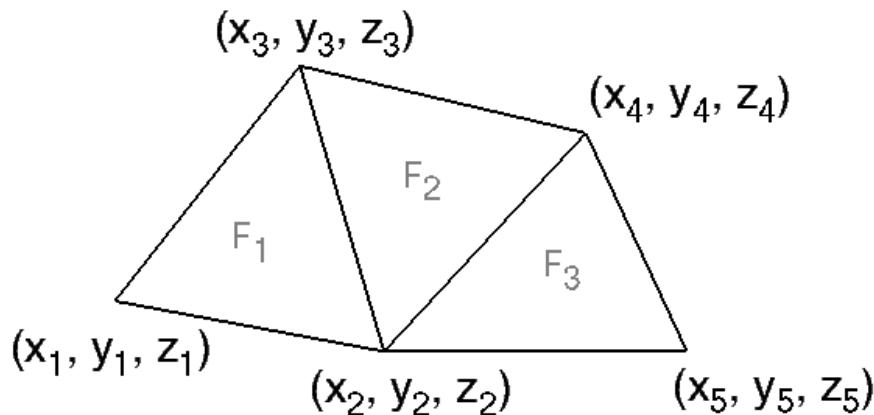
- A.k.a triangle soup
- Each face lists vertex coordinates
 - Redundant vertices
 - No adjacency information



FACE TABLE			
F ₁	(x ₁ , y ₁ , z ₁)	(x ₂ , y ₂ , z ₂)	(x ₃ , y ₃ , z ₃)
F ₂	(x ₂ , y ₂ , z ₂)	(x ₄ , y ₄ , z ₄)	(x ₃ , y ₃ , z ₃)
F ₃	(x ₂ , y ₂ , z ₂)	(x ₅ , y ₅ , z ₅)	(x ₄ , y ₄ , z ₄)

Vertex and Face Tables

- Each face lists vertex references
 - Shared vertices
 - Still no adjacency information

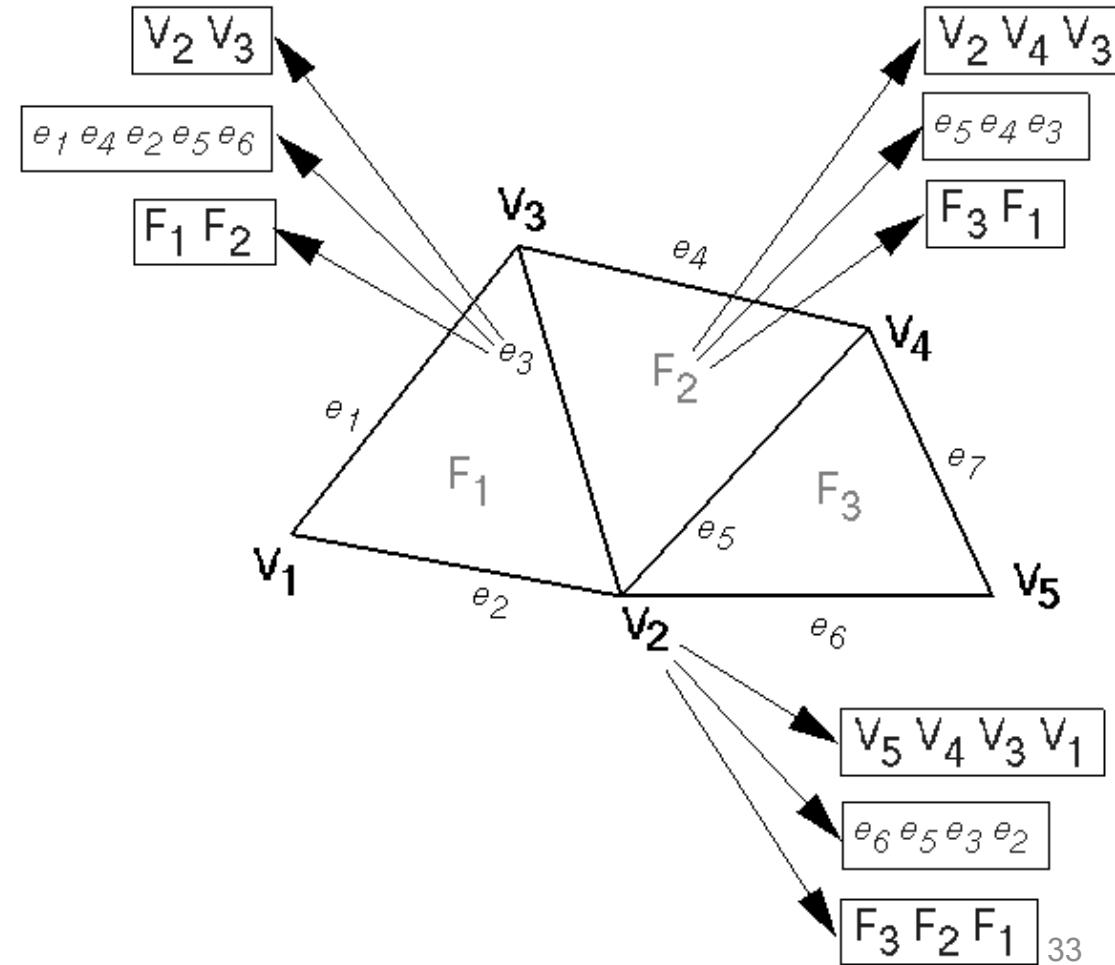


VERTEX TABLE			
V ₁	X ₁	Y ₁	Z ₁
V ₂	X ₂	Y ₂	Z ₂
V ₃	X ₃	Y ₃	Z ₃
V ₄	X ₄	Y ₄	Z ₄
V ₅	X ₅	Y ₅	Z ₅

FACE TABLE			
F ₁	V ₁	V ₂	V ₃
F ₂	V ₂	V ₄	V ₃
F ₃	V ₂	V ₅	V ₄

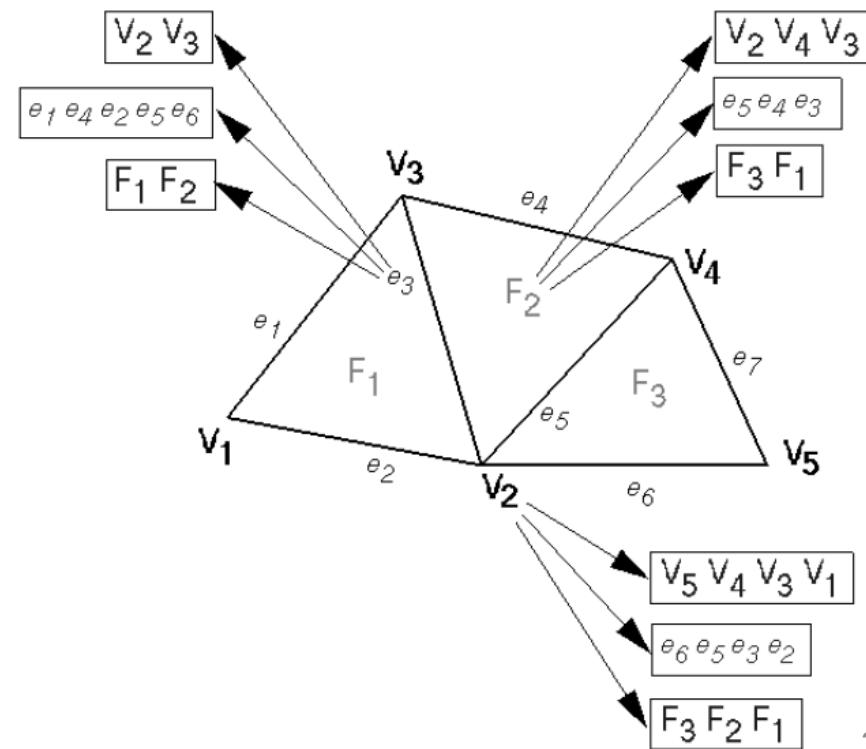
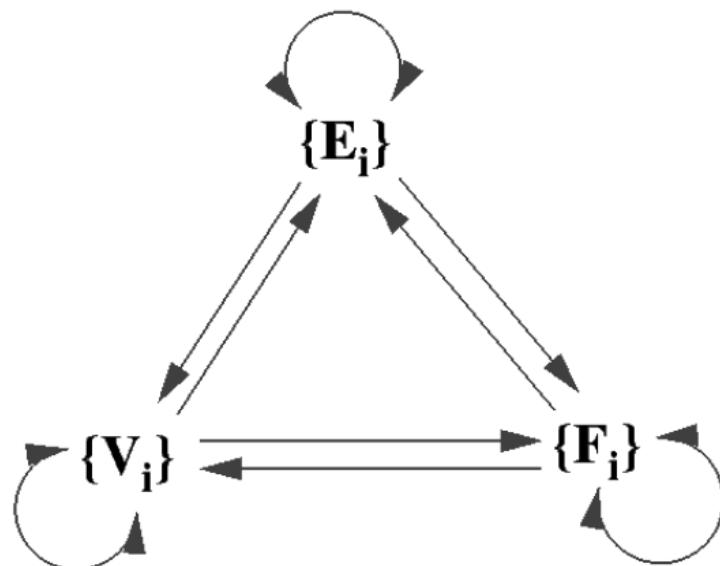
Adjacency Lists

- Store all vertex, edge and face adjacencies
 - Efficient adjacency traversal
 - Extra storage requirements



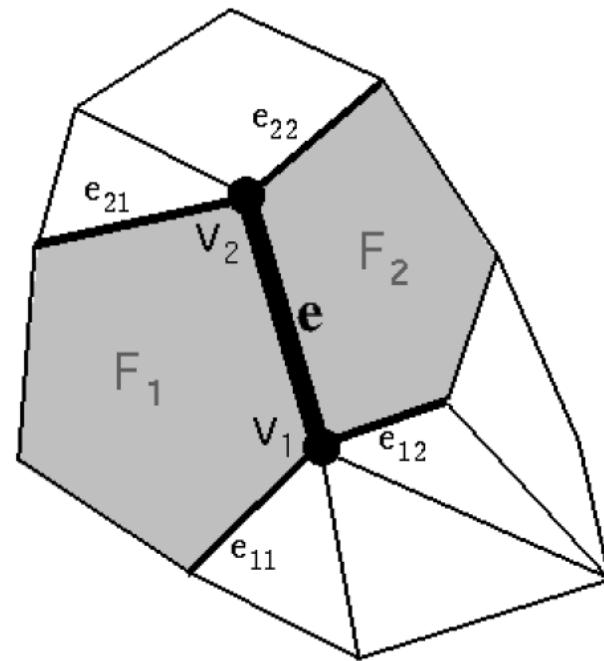
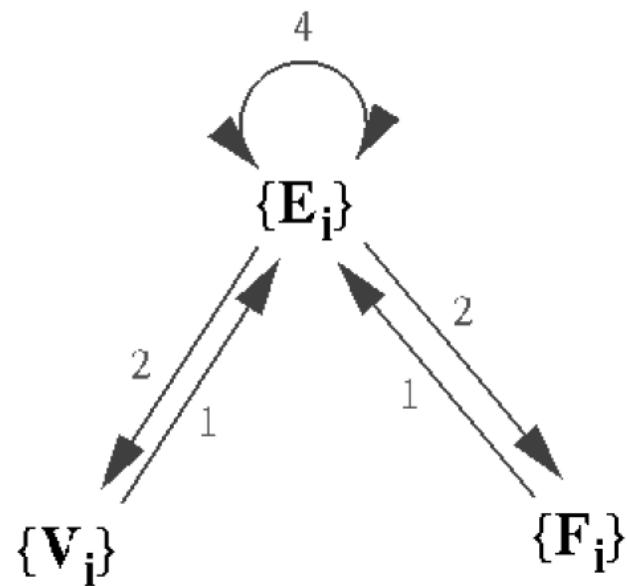
Partial Adjacency Lists

- Can we store only some adjacency relationships and derive others?

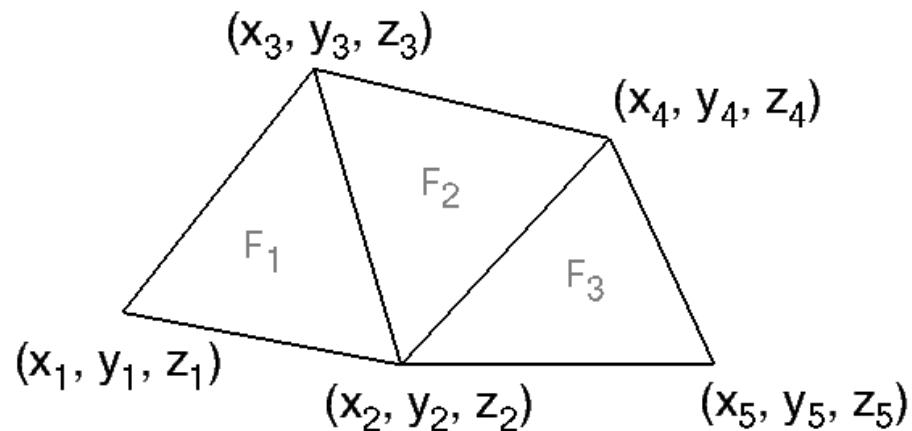


Winged Edge

- Adjacency encoded in edges
 - All adjacencies in $O(1)$ time
- Little extra storage (fixed records)
- Arbitrary polygons



Winged Edge



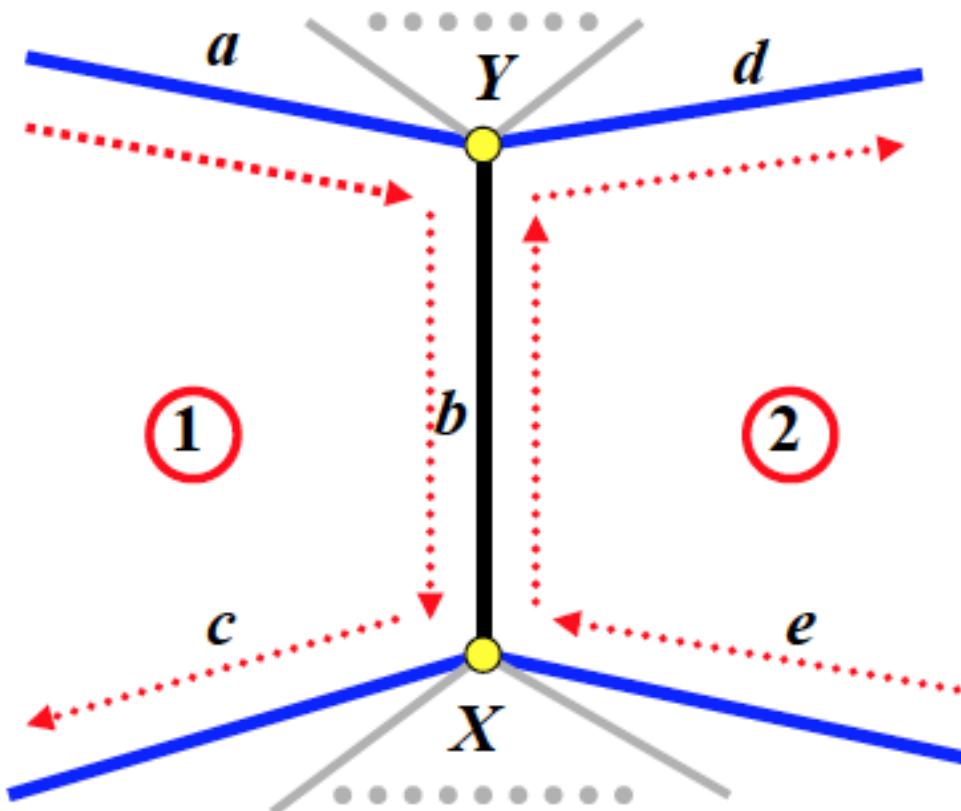
VERTEX TABLE				
V ₁	X ₁	Y ₁	Z ₁	e ₁
V ₂	X ₂	Y ₂	Z ₂	e ₆
V ₃	X ₃	Y ₃	Z ₃	e ₃
V ₄	X ₄	Y ₄	Z ₄	e ₅
V ₅	X ₅	Y ₅	Z ₅	e ₆

EDGE TABLE					11	12	21	22
e ₁	V ₁	V ₃	F ₁	e ₂	e ₂	e ₄	e ₃	
e ₂	V ₁	V ₂	F ₁	e ₁	e ₁	e ₃	e ₆	
e ₃	V ₂	V ₃	F ₁	F ₂	e ₂	e ₅	e ₁	e ₄
e ₄	V ₃	V ₄	F ₂	F ₂	e ₁	e ₃	e ₇	e ₅
e ₅	V ₂	V ₄	F ₂	F ₃	e ₃	e ₆	e ₄	e ₇
e ₆	V ₂	V ₅	F ₃		e ₅	e ₂	e ₇	e ₇
e ₇	V ₄	V ₅		F ₃	e ₄	e ₅	e ₆	e ₆

FACE TABLE	
F ₁	e ₁
F ₂	e ₃
F ₃	e ₅

Winged Edge

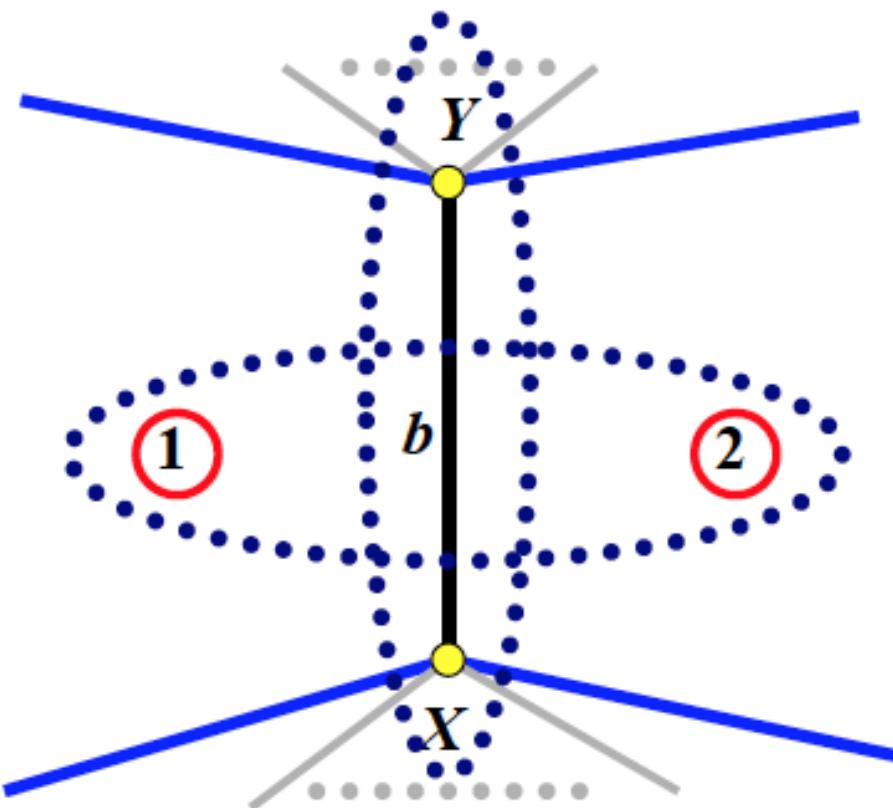
If all faces are oriented clock-wise, each edge has *eight* pieces of incident information.



- Given edge: $b=XY$
- Incident faces: 1 and 2
- Pred. & succ. edges of 1
- Pred. & succ. edges of 2
- The *wings* of edge $b=XY$ are faces 1 and 2.
- Edge b is a *winged-edge*

Winged Edge

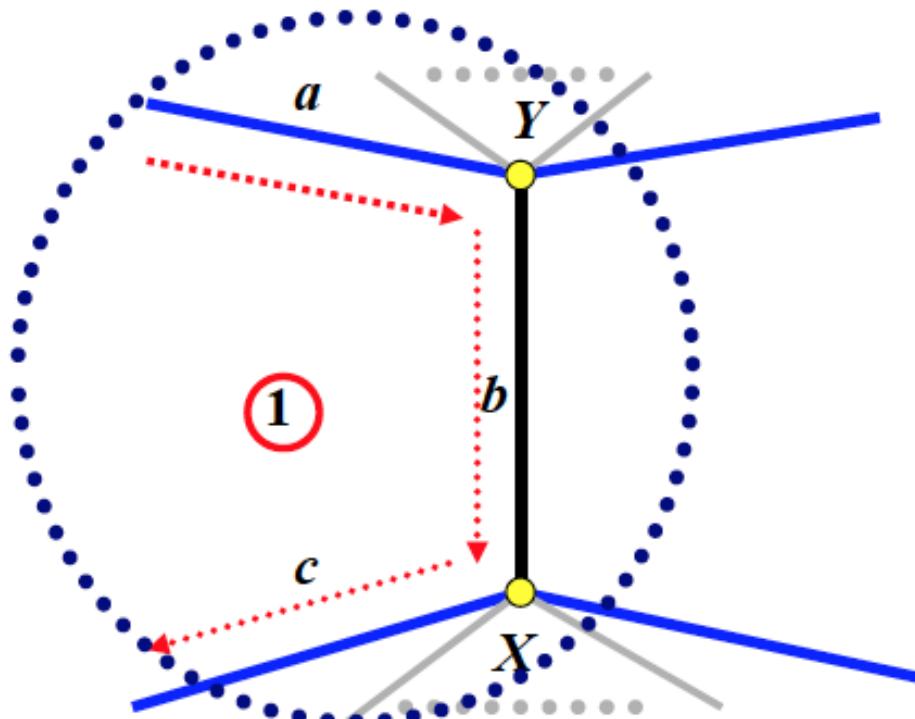
If all faces are oriented clock-wise, each edge has *eight* pieces of incident information.



- The first *four* pieces:
 - The *two* vertices of *b*: *X* and *Y*
 - The *two* incident faces: *1* and *2*

Winged Edge

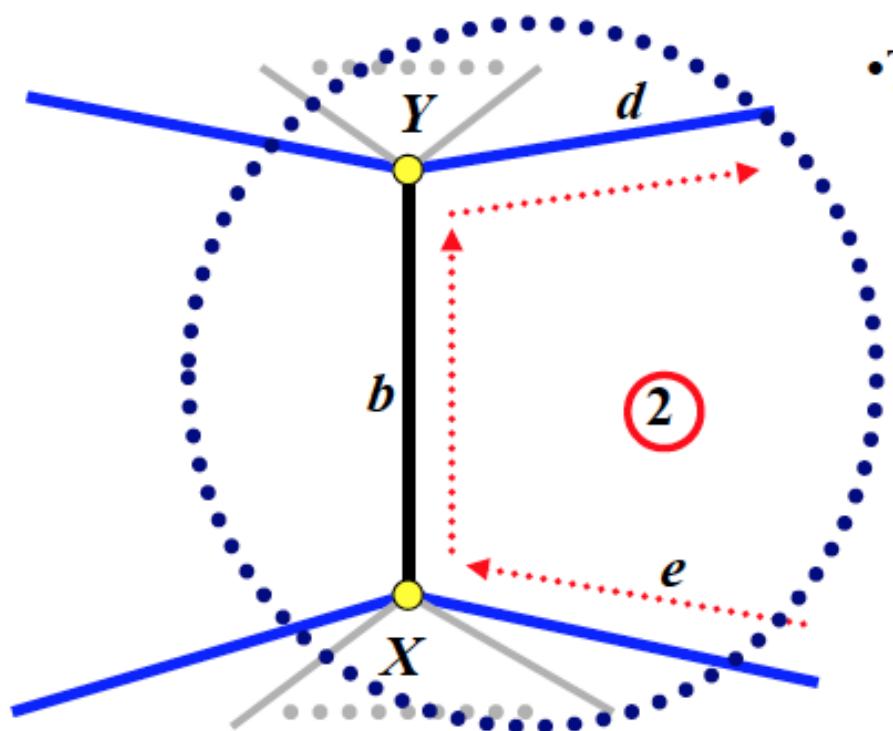
If all faces are oriented clock-wise, each edge has *eight* pieces of incident information.



- The pred. and succ. edges of b with respect to face 1: a and c

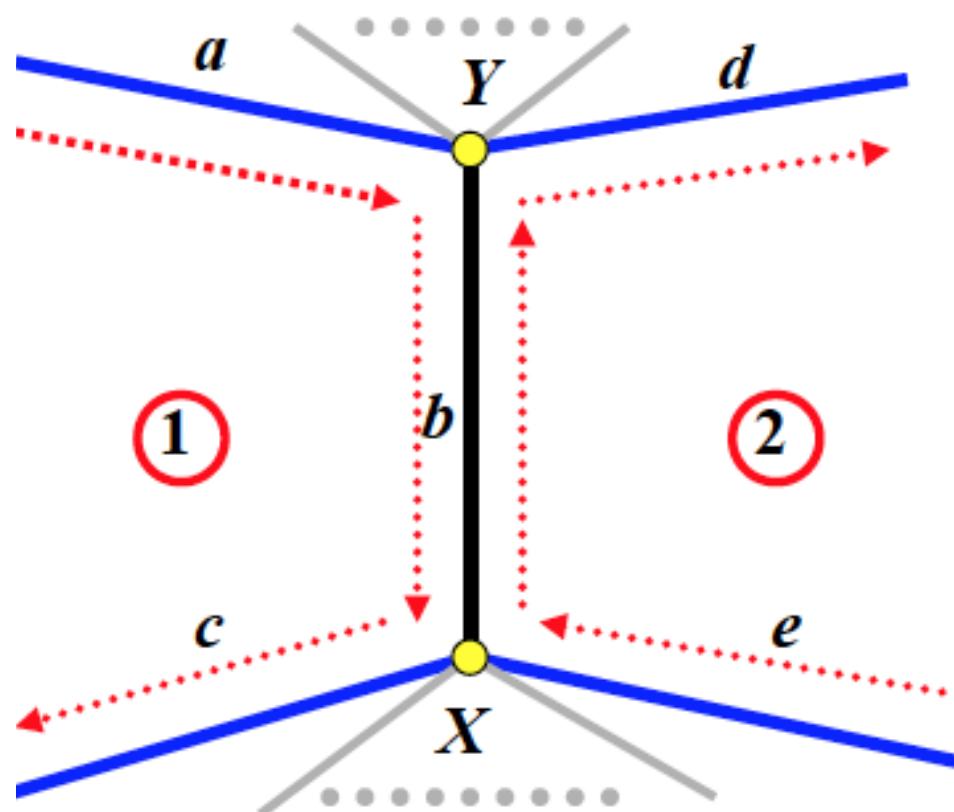
Winged Edge

If all faces are oriented clock-wise, each edge has *eight* pieces of incident information.



- The pred. and succ. edges of *b* with respect to face 2: *e* and *d*

Left vs. Right Faces

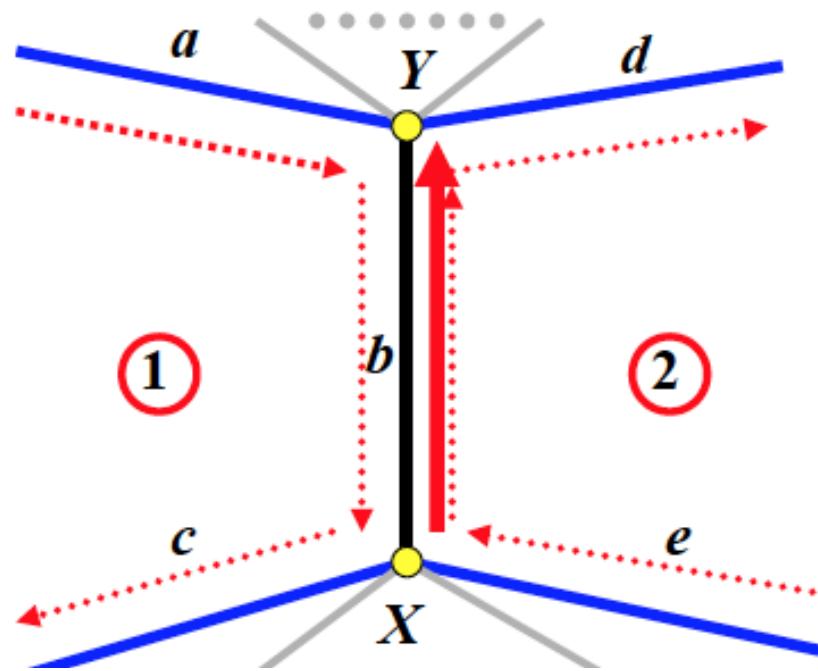


- Which one is left, 1 or 2?
- Choose a direction for edge b , say from X to Y or from Y to X .
- Going from the start vertex to the end vertex, we know which face is the left one!
- If the start vertex is X , the *left* face is face 1.

Winged Edge Info

Information for Edge b (from X to Y)

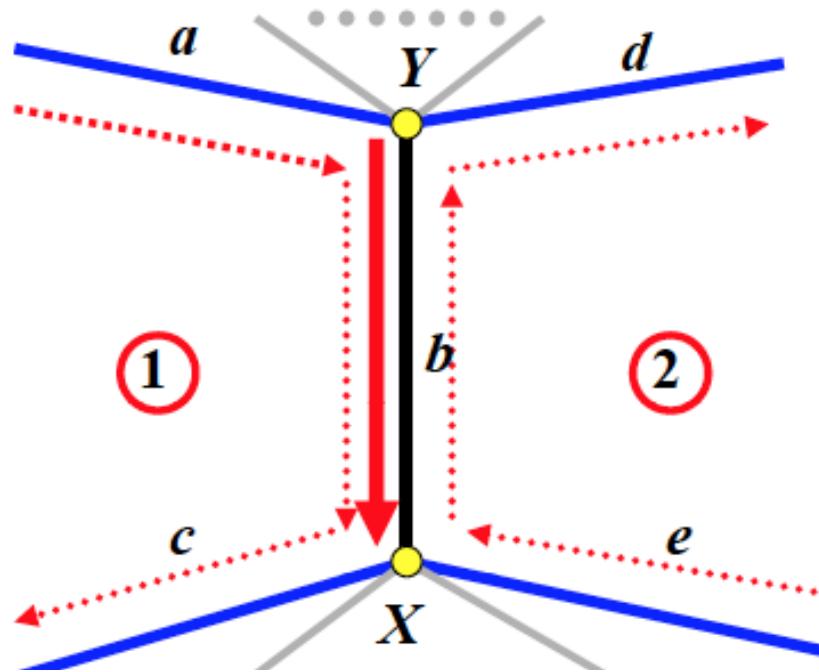
Start Vertex	End Vertex	Left Face	Right Face	Left Pred.	Left Succ.	Right Pred.	Right Succ.
X	Y	1	2	a	c	e	d



Winged Edge Info

Information for Edge b (from Y to X)

Start Vertex	End Vertex	Left Face	Right Face	Left Pred.	Left Succ.	Right Pred.	Right Succ.
Y	X	2	1	e	d	a	c

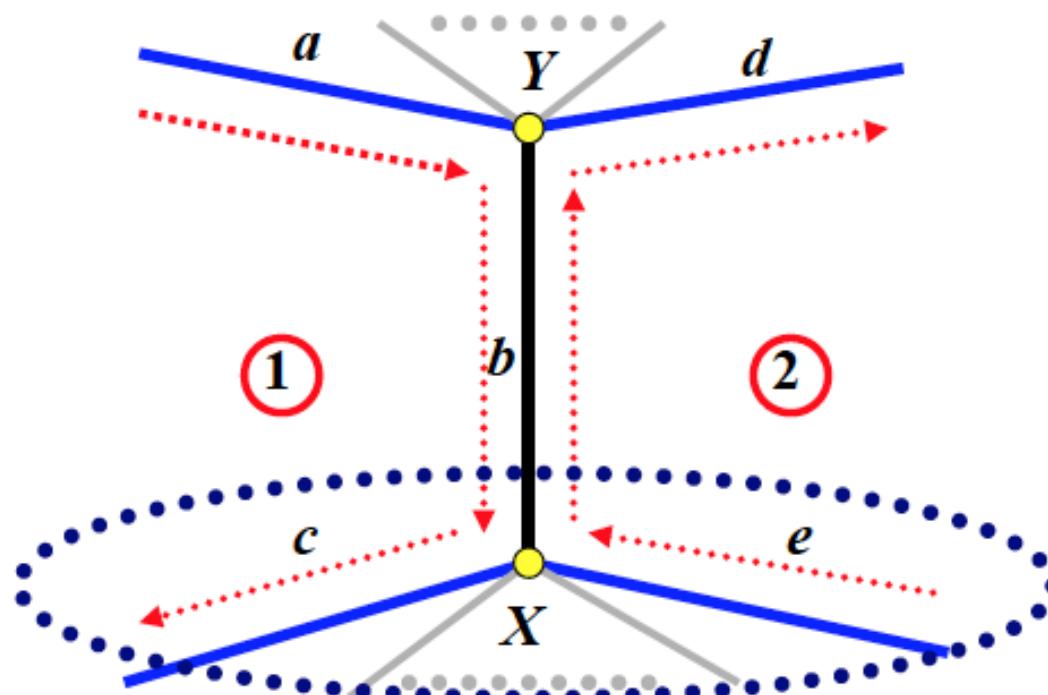


Winged Edged Data Structure

- The winged-edge data structure has three tables: **edge** table, **vertex** table, and **face** table.
- Each edge has one row in the **edge** table.
 - Each row contains the eight pieces of information of that edge.
- Each vertex has one entry in the **vertex** table.
 - Each entry has a pointer to an incident edge (in the edge table) of that vertex.
- Each face has one entry in the **face** table.
 - Each entry has a pointer to an incident edge (in the edge table) of that face.

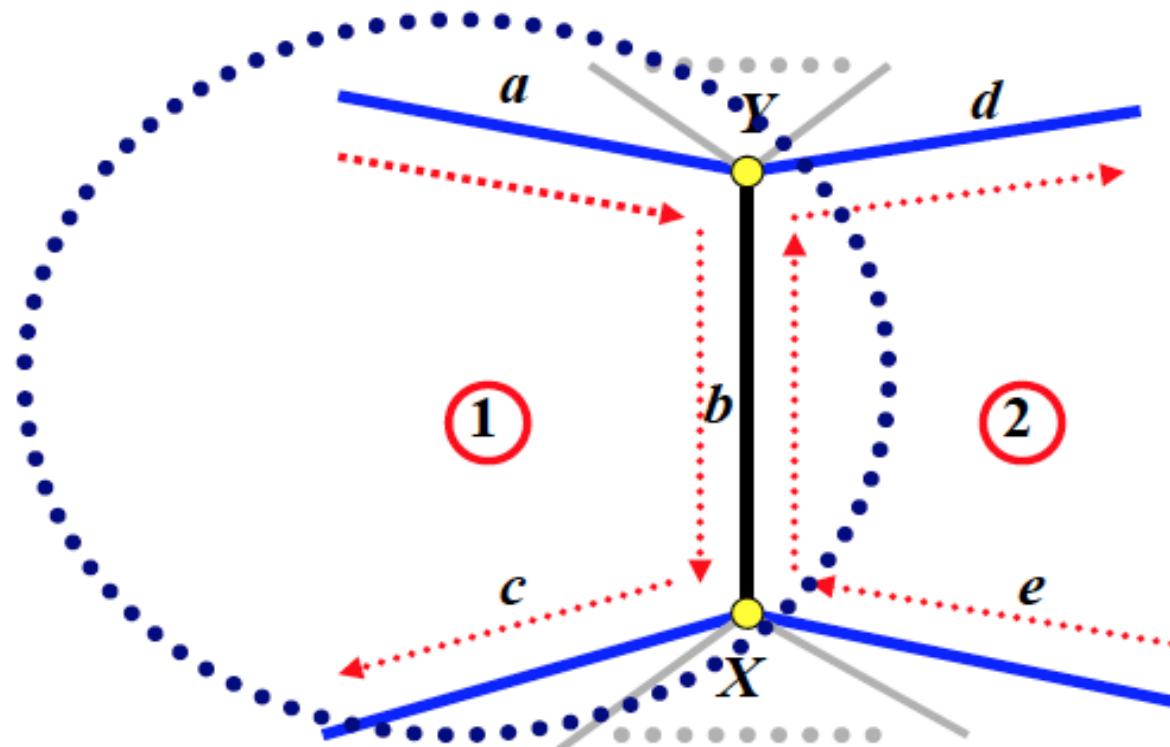
Winged Edged Data Structure

The vertex table entry for vertex X may be the edge table entry of edges c, b, e , or any other incident edge.



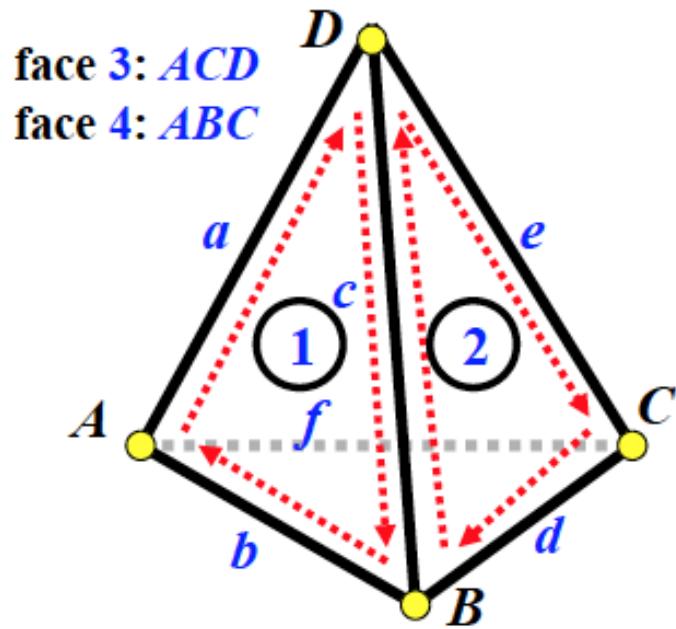
Winged Edged Data Structure

The face table entry for face 1 may be the edge table entry of edges *a*, *b*, *c*, or any other incident edge.



Winged Edged Data Structure

- The following tetrahedron has four vertices **A**, **B**, **C** and **D**, six edges **a**, **b**, **c**, **d**, **e**, **f**, and four faces **1**, **2**, **3** and **4**.



Vertex Table

Vertex	Edge
<i>A</i>	<i>a</i>
<i>B</i>	<i>b</i>
<i>C</i>	<i>d</i>
<i>D</i>	<i>a</i>

Face Table

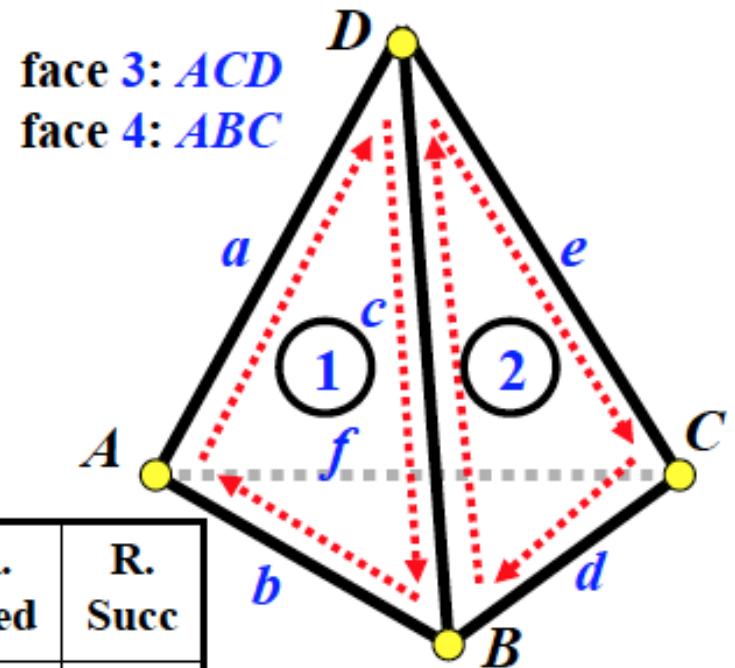
Face	Edge
1	<i>a</i>
2	<i>c</i>
3	<i>a</i>
4	<i>b</i>

Winged Edged Data Structure

- The following tetrahedron has four vertices A, B, C and D, six edges a, b, c, d, e, f, and four faces 1, 2, 3 and 4.

Edge Table

Edge	Start Vtx	End Vtx	L. Face	R. Face	L. Pred	L. Succ	R. Pred	R. Succ
a	A	D	3	1	e	f	b	c
b	A	B	1	4	c	a	f	d
c	B	D	1	2	a	b	d	e
d	B	C	2	4	e	c	b	f
e	C	D	2	3	c	d	f	a
f	A	C	4	3	d	b	a	e

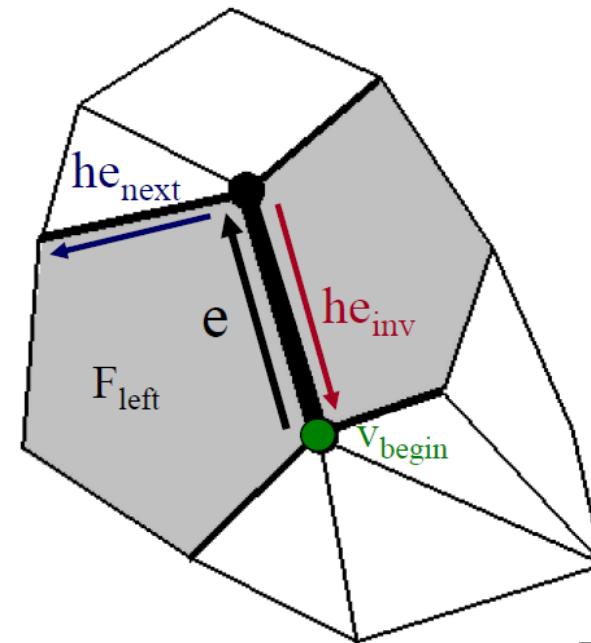


Winged Edged Data Structure

- The winged-edge data structure permits a program to answer many topological inquiries very efficiently.
- If (1) V , E and F are the numbers of vertices, edges, and faces and (2) each entry in the table uses one memory unit, the vertex table, edge, table, and face table require V , $8E$ and F memory units, respectively.

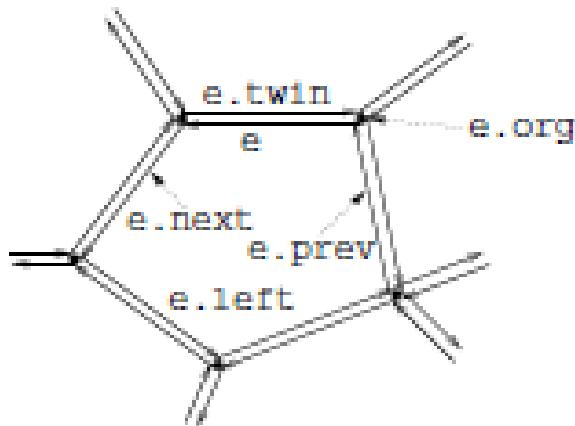
Half Edge

- Adjacency encoded in edges
 - All adjacencies in $O(1)$ time
 - Little extra storage (fixed records)
 - Arbitrary polygons
- Similar to winged-edge,
except adjacency
encoded in half-edges



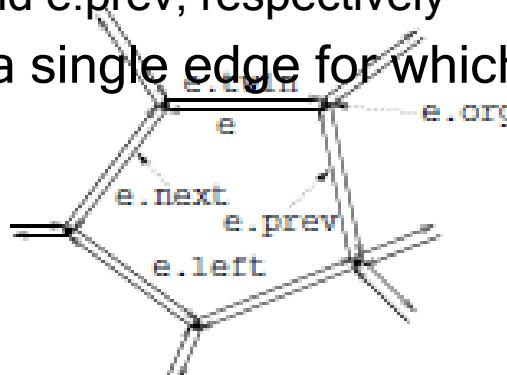
Half Edge

- Each undirected edge represented by two *directed* half edges
 - Unambiguously defines left and right
- Assume that there are no holes in faces



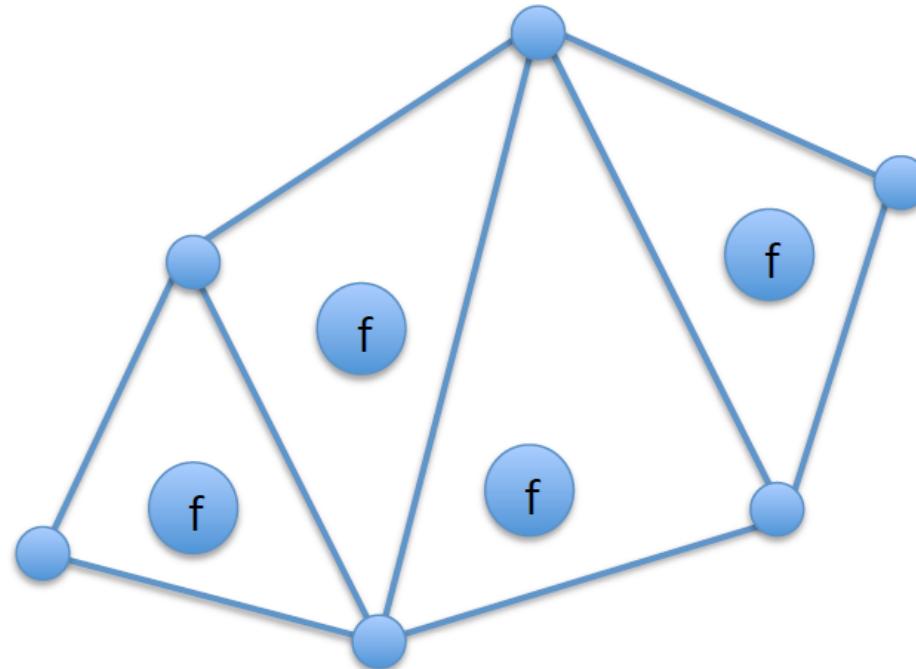
Half Edge

- Each vertex stores:
 - its coordinates
 - a pointer `v.inc_edge` to any directed edge that has vertex as its origin
- Each directed edge is **associated** with:
 - a pointer to the oppositely directed edge, called its twin
 - an origin and destination vertex
 - two faces, one to its left and one to its right.
- We only **store**:
 - a pointer to the origin vertex `e.org` (`e.dest` can be accessed as `e.twin.org`)
 - a pointer to the face to the left of the edge `e.left` (we can access the face to the right from the twin edge)
 - pointers to the next and previous directed edges in counterclockwise order about the incident face, `e.next` and `e.prev`, respectively
- Each face `f` stores a pointer to a single edge for which this face is the incident face, `f.inc_edge`



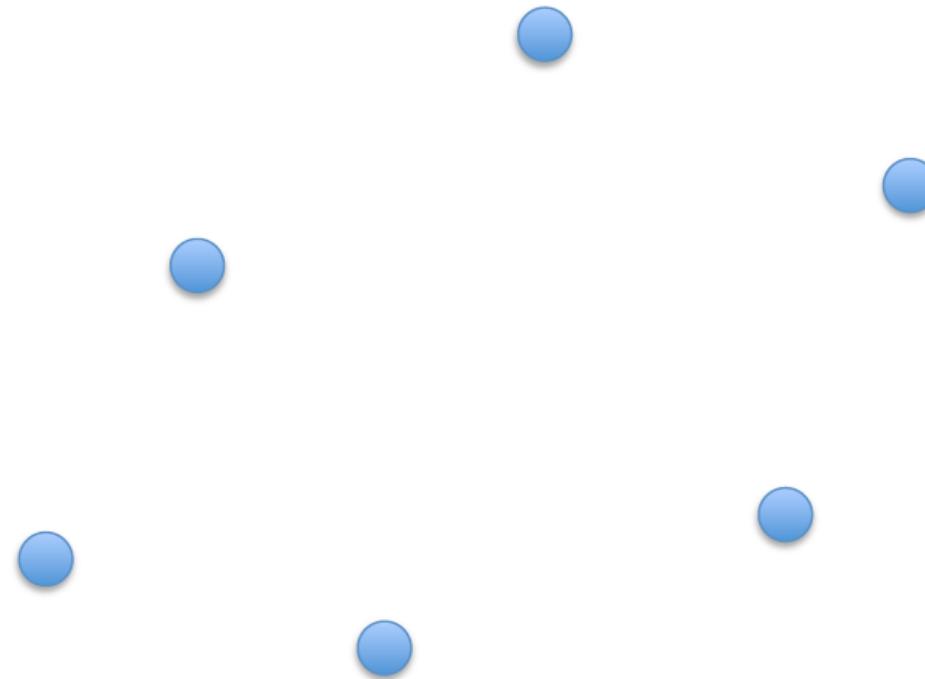
How to Load a Shape

- From file with vertices and triangles



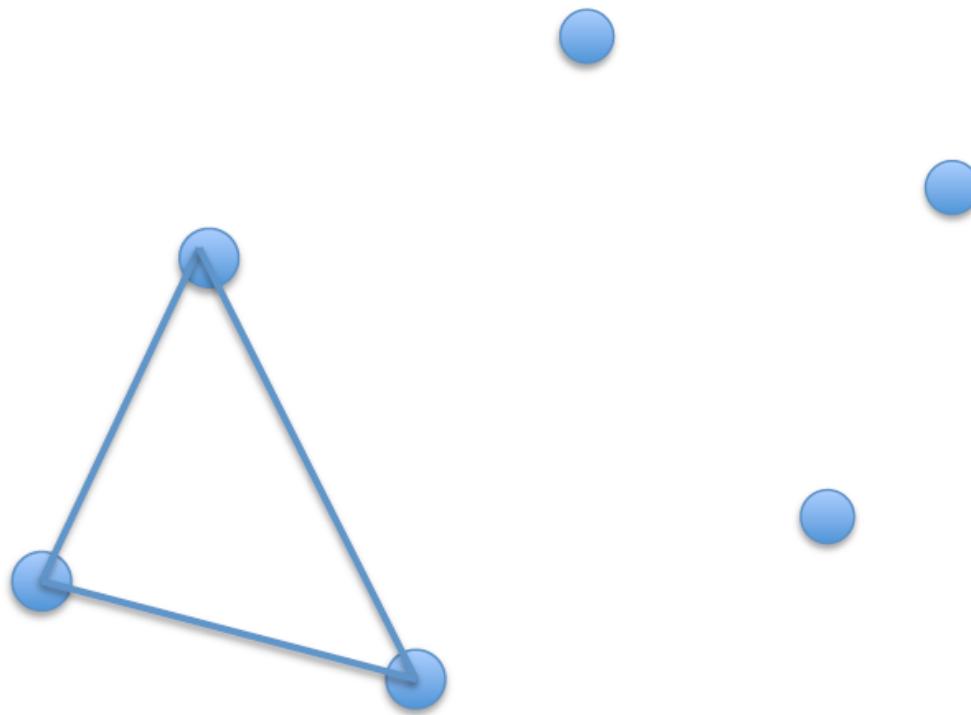
How to Load a Shape

- Add vertex coordinates
to list



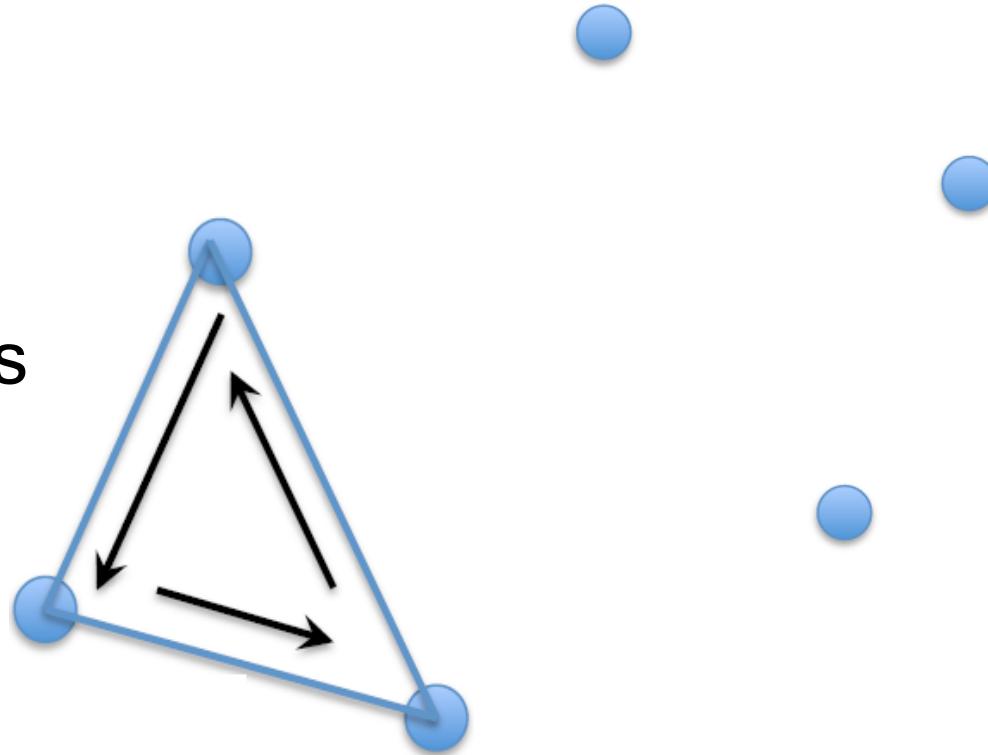
How to Load a Shape

- Add vertex coordinates to list
- Add half-edges with faces



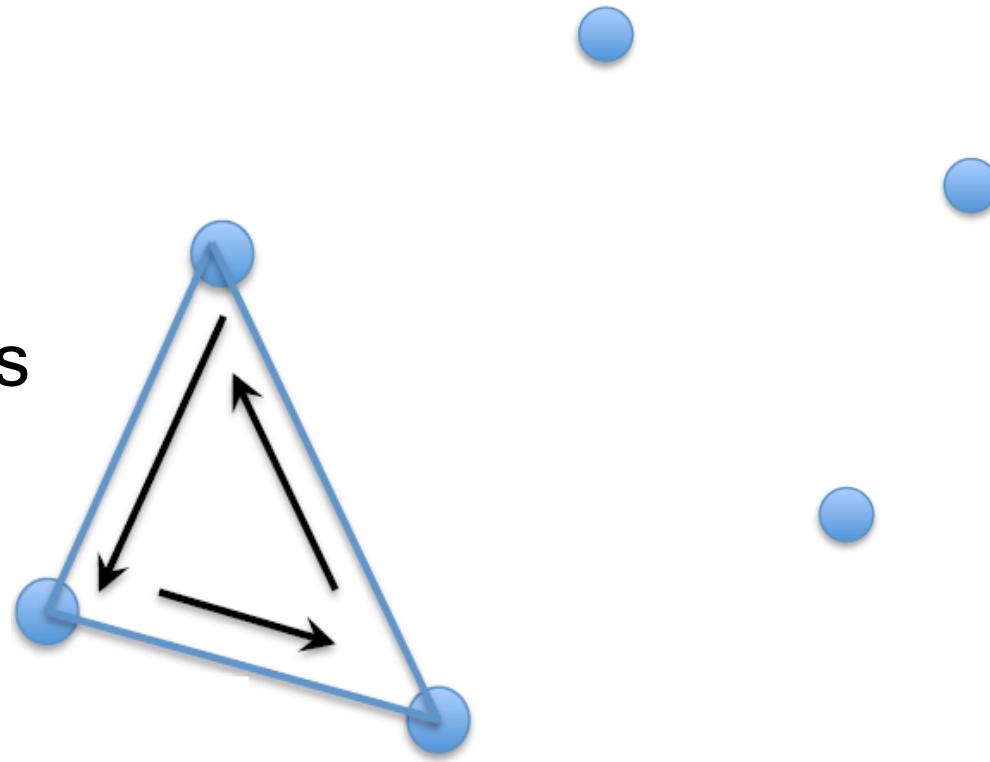
How to Load a Shape

- Add vertex coordinates to list
- Add half-edges with faces
 - Inner half-edges are sufficient



How to Load a Shape

- Add vertex coordinates to list
- Add half-edges with faces
 - Inner half-edges are sufficient
 - Update vertex pointers to half-edges

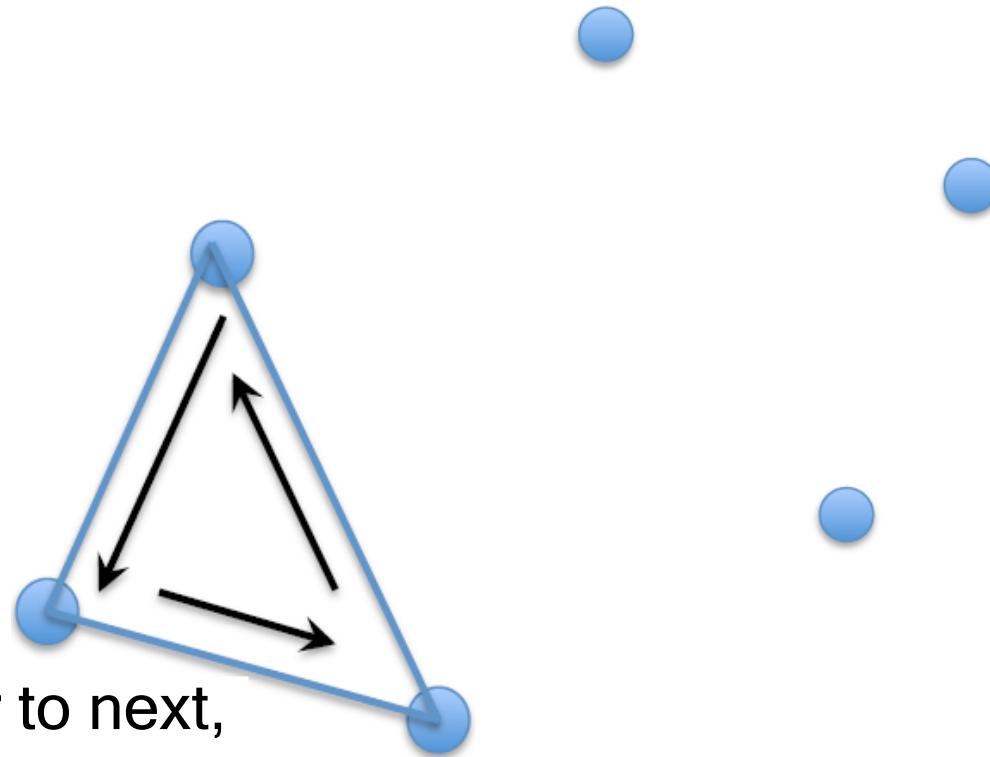


How to Load a Shape

- Add vertex coordinates to list

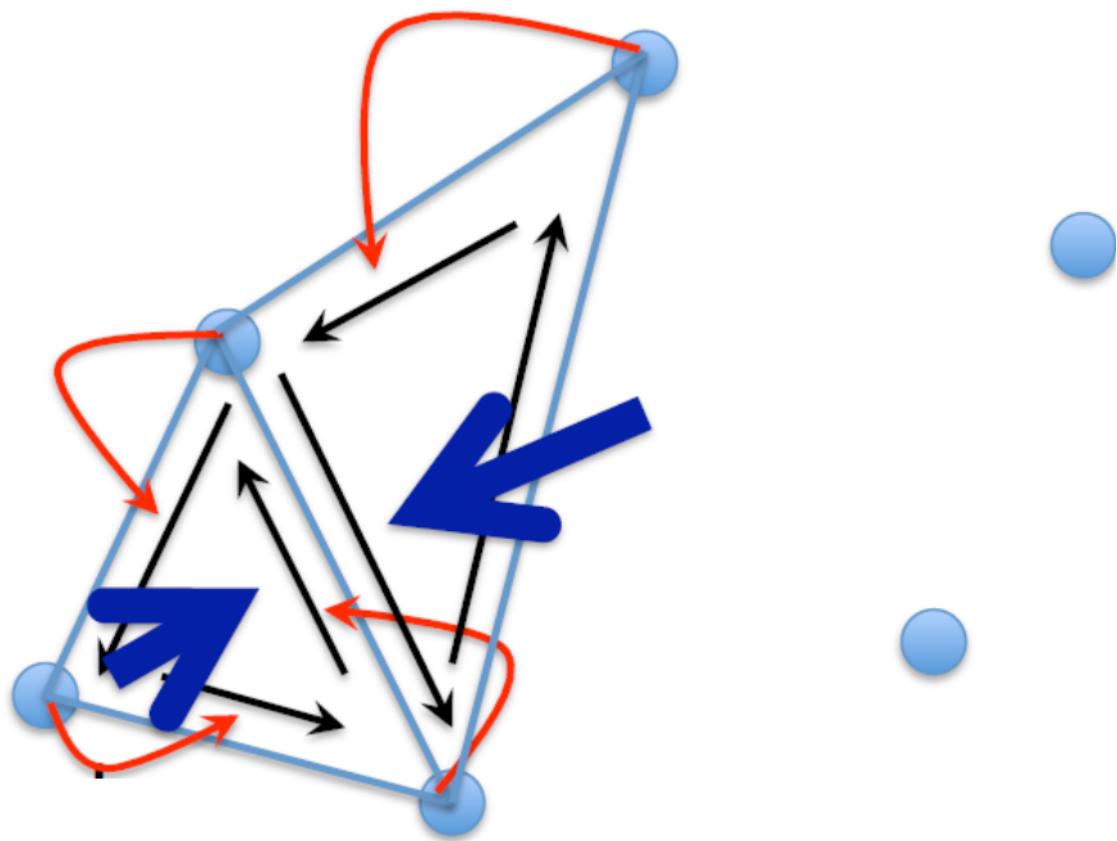
- Add half-edges with faces

- Inner half-edges are sufficient
 - Update vertex pointers to half-edges
 - Half-edges: pointer to next, pointer to face
 - Faces: pointer to one of the inner half-edges



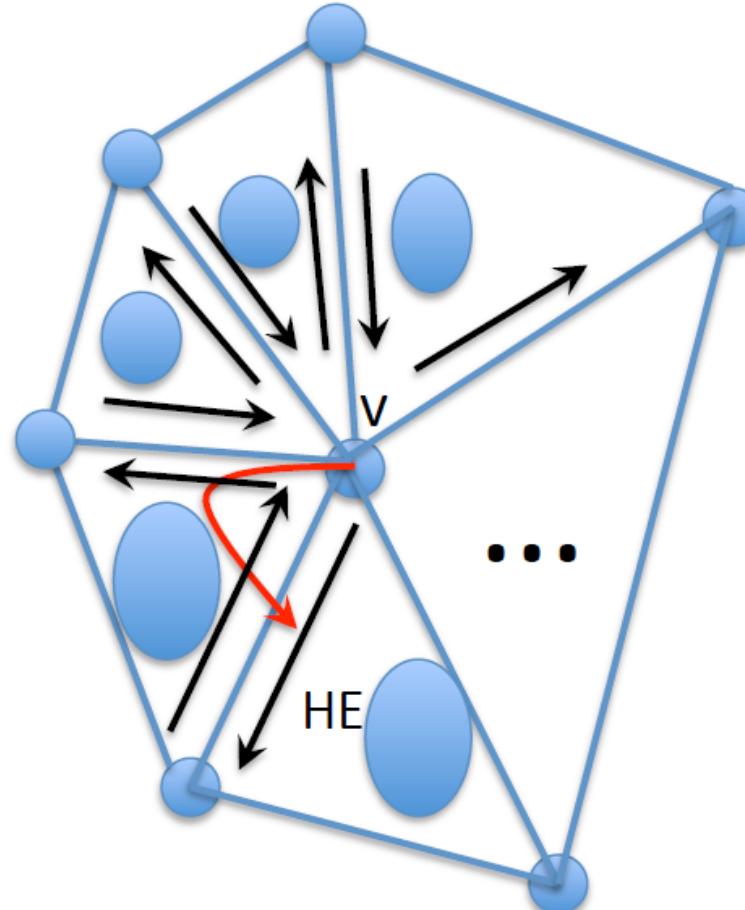
How to Load a Shape

- Continue adding incrementally



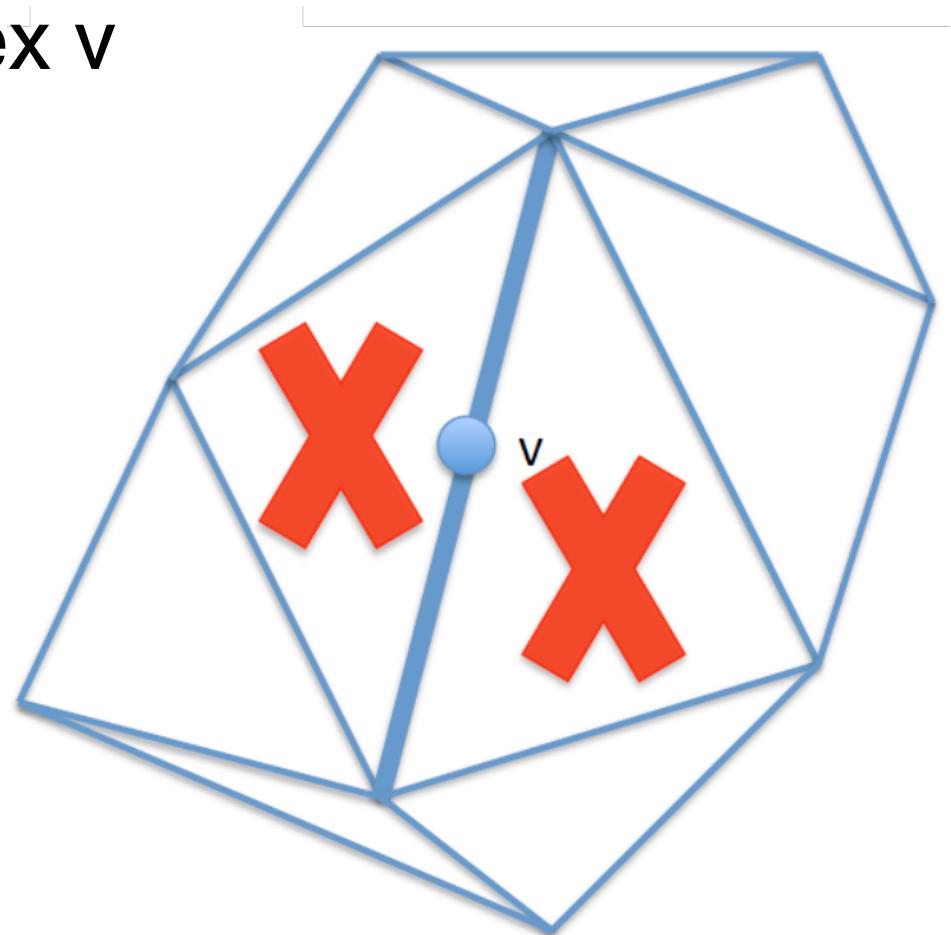
Finding Adjacent Faces

- Check all outgoing half edges
 - V points to a half edge HE
 - $\text{ADD_FACE}(HE)$
 - Iterate:
 - $X=HE.\text{twin}$
 - $Y=X.\text{next}$
 - $\text{ADD_FACE}(Y)$
 - $HE:=Y$



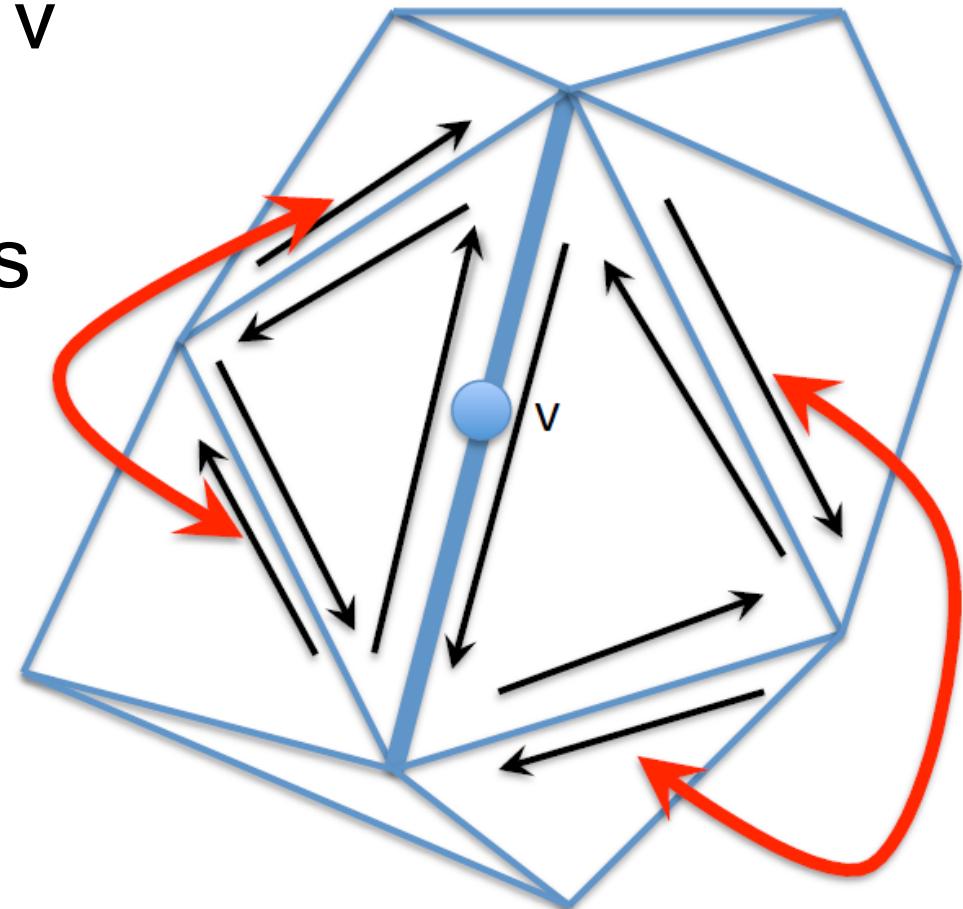
Collapsing an Edge

- Create a new vertex v
- Remove faces



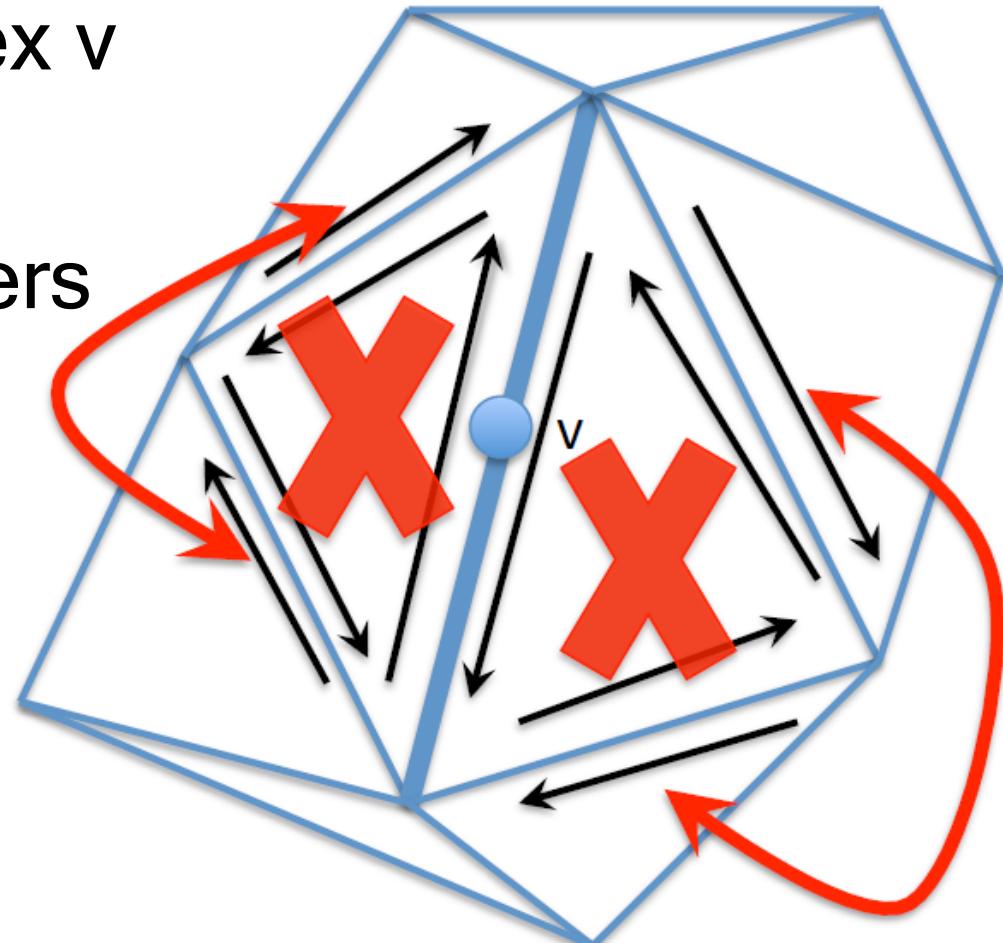
Collapsing an Edge

- Create a new vertex v
- Remove faces
- Change twin pointers



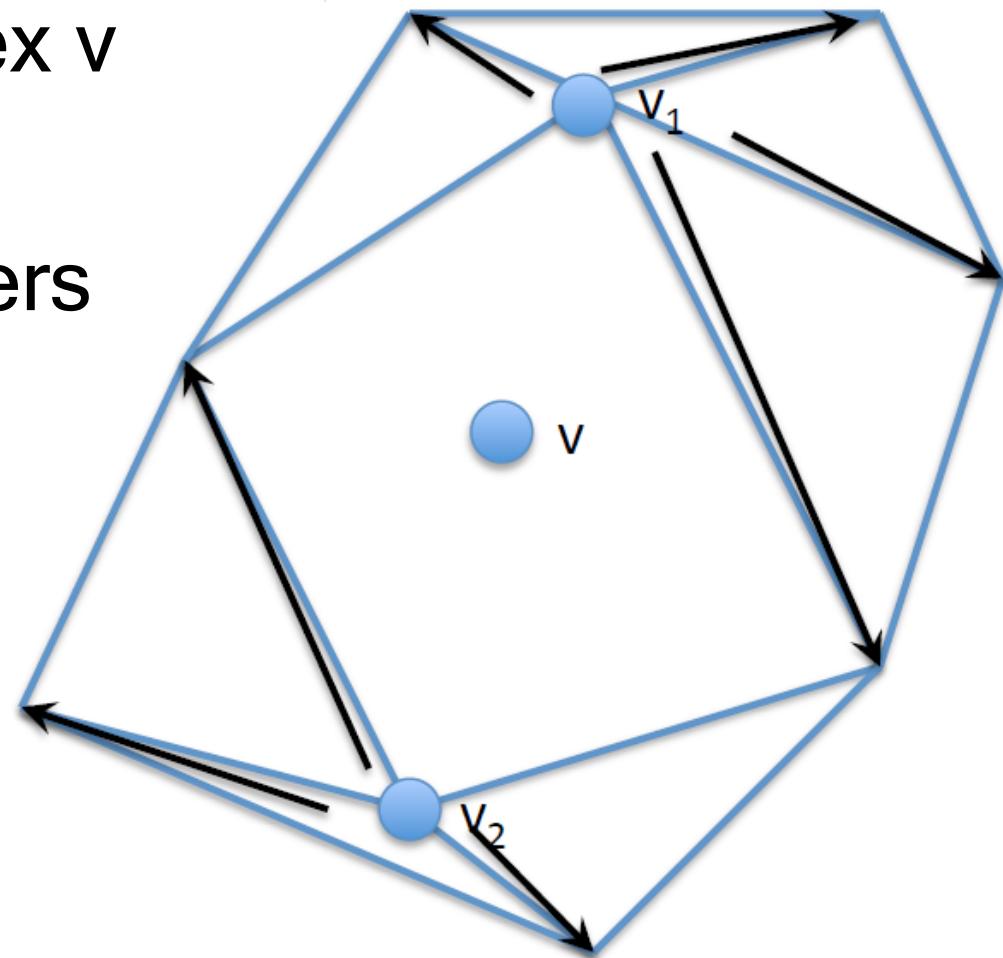
Collapsing an Edge

- Create a new vertex v
- Remove faces
- Change twin pointers
- Remove edges



Collapsing an Edge

- Create a new vertex v
- Remove faces
- Change twin pointers
- Remove edges
- Change pointers from half-edges to v_1 and v_2



Collapsing an Edge

- Create a new vertex v
- Remove faces
- Change twin pointers
- Remove edges
- Change pointers from half-edges to v_1 and v_2
- Remove v_1 and v_2
- Pick an outgoing edge for v

