

How is the Consensus Reached in Ethereum?

@juinc

Outline

- Overview
- Section 10: Blocktree to Blockchain
- Section 11: Block Finalisation
- Appendix J: ETHASH

Review

Please refer to my previous talk:

- [Let's Make a Blockchain \(Basic\)](#)
- [How is Your Contract Executed in Ethereum? \(EVM\)](#)

Overview

- Consensus: general agreement on the World State
- The consensus can be reached through different schemes:
 - Proof of Work
 - Proof of Stake
 - ...
- The process of applying or wrapping a block is called finalization

Section 10: Blocktree to Blockchain

- The canonical blockchain is a path from root to leaf through the entire block tree
- Define Total Difficulty:

$$B_t \equiv B'_t + B_d$$

$$B' \equiv P(B_H)$$

- Choose the path that has the most computation done upon it, or, the most difficult path

Section 11: Block Finalization

Peer (Non-miner) v.s. Miner

- Step 1: Validate (Determine) Ommers
- Step 2: Validate (Determine) Transactions
- Step 3: Apply Rewards
- Step 4: Validate (Determine) State and Nonce

Recall: State Transition Function

$$\sigma_{t+1} \equiv \Pi(\sigma_t, B)$$

$$B \equiv (\dots, (T_0, T_1, \dots))$$

$$\Pi(\sigma, B) \equiv \Omega(B, \Upsilon(\Upsilon(\sigma, T_0), T_1) \dots)$$

where

$$\sigma' = \Upsilon(\sigma, T)$$

(explained in [my last talk](#))

Step 1: Validate (Determine) Ommers

$$\|B_U\| \leq 2 \bigwedge_{U \in B_U} V(U) \wedge k(U, P(B_H)_H, 6)$$

where

$$k(U, H, n) \equiv \begin{cases} false & \text{if } n = 0 \\ s(U, H) \vee k(U, P(H)_H, n - 1) & \text{otherwise} \end{cases}$$

and

$$s(U, H) \equiv (P(H) = P(U) \wedge H \neq U \wedge U \notin B(H)_U)$$

Step 2: Validate (Determine) Transactions

$$B_{H_g} = \ell(\mathbf{R})_u$$

Step 3: Apply Rewards

$$\begin{aligned}\Omega(B, \sigma) &\equiv \sigma' : \sigma' = \sigma \text{ except:} \\ \sigma'[B_{Hc}]_b &= \sigma[B_{Hc}]_b + (1 + \frac{\|B_U\|}{32})R_b \\ \forall U \in B_U : \\ \sigma'[U_c]_b &= \sigma[U_c]_b + (1 + \frac{1}{8}(U_i - B_{Hi}))R_b\end{aligned}$$

where

$$R_b = 5 \times 10^{18}$$

Step 4: Validate (Determine) State and Nonce

A. Validate

$$\Pi(\boldsymbol{\sigma}, B) \equiv \Omega(B, \ell(\mathbf{R})_{\boldsymbol{\sigma}})$$

where

$$\mathbf{R}[n]_{\boldsymbol{\sigma}} = \begin{cases} \Gamma(B) & \text{if } n < 0 \\ \Upsilon(\mathbf{R}[n-1]_{\boldsymbol{\sigma}}, B_{\mathbf{T}}[n]) & \text{otherwise} \end{cases}$$

Step 4: Validate (Determine) State and Nonce

A. Validate (Cont.)

and

$$\mathbf{R}[n]_u = \begin{cases} 0 & \text{if } n < 0 \\ \Upsilon^g(\mathbf{R}[n-1]_\sigma, B_{\mathbf{T}}[n]) \\ \quad + \mathbf{R}[n-1]_u & \text{otherwise} \end{cases}$$

and

$$\mathbf{R}[n]_1 = \Upsilon^1(\mathbf{R}[n-1]_\sigma, B_{\mathbf{T}}[n])$$

Step 4: Validate (Determine) State and Nonce

A. Validate (Cont.)

and

$$\Gamma(B) \equiv \begin{cases} \sigma_0 & \text{if } P(B_H) = \emptyset \\ \sigma_i : \text{TRIE}(L_S(\sigma_i)) = P(B_H)_{H_r} & \text{otherwise} \end{cases}$$

and

$$\sigma' = \Upsilon(\sigma, T)$$

Step 4: Validate (Determine) State and Nonce

B. Determine

Block Transition Function

$$\Phi(B) \equiv B' : B' = B^* \quad \text{except:}$$

$$B'_n = n : x \leq \frac{2^{256}}{H_d}$$

$$B'_m = m \quad \text{with } (x, m) = \text{PoW}(B_{\mathcal{H}}^*, n, \mathbf{d})$$

$$B^* \equiv B \quad \text{except: } B_r^* = r(\Pi(\Gamma(B), B))$$

Step 4: Validate (Determine) State and Nonce

B. Determine (Cont.)

Proof-of-work Function (Ethash)

$$\text{PoW}(H_{\mathcal{H}}, H_n, \mathbf{d}) = \{\mathbf{m}_c(\text{KEC}(\text{RLP}(L_H(H_{\mathcal{H}}))), H_n, \mathbf{d}), \text{KEC}(\mathbf{s}_h(\text{KEC}(\text{RLP}(L_H(H_{\mathcal{H}}))), H_n) + \mathbf{m}_c(\text{KEC}(\text{RLP}(L_H(H_{\mathcal{H}}))), H_n, \mathbf{d}))\}$$

Appendix J: Ethash

Why Ethash?

- ASIC-resistance
- Light Client Verifiability

ref: [Ethash](#)

Appendix J: Ethash

1. Determine Cache Size and Dataset (DAG) Size
2. Determine Seed Hash
3. Determine Cache
 - Calculate Initial Cache from Seed and Cache Size
 - Calculate Cache from RandMemoHash (RMH) Algorithm
4. Determine Dataset (DAG) from Cache
5. Determine Mix from Dataset (DAG), Header, and Nonce
6. Compare Mix to Target, repeat 5 until Mix is smaller than Target (PoW)

Conclusion

- By looking into block finalization process, we can understand more on the whole picture of ethereum
- Implementation details might be different from what yellow paper describes