



Tribhuvan University
Institute of Science and Technology

A Comparative Evaluation of Buffer Replacement Algorithms LIRS-WSR and CCF-LRU for Flash Memory Based Systems

Dissertation

Submitted to

Central Department of Computer Science & Information Technology
Kirtipur, Kathmandu, Nepal

In partial fulfillment of the requirements
for the Masters Degree in Computer Science & Information Technology

By

Mahesh Kumar Yadav

Date: April 2, 2017



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Supervisor

Prof. Dr. Subarna Shakya

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Student's Declaration

I hereby declare that I am the only author of this work and that no sources other than the listed here have been used in this work.

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Supervisor's Recommendation

I hereby recommend that this dissertation prepared under my supervision by **Mr. Mahesh Kumar Yadav** entitled “**A Comparative Evaluation of Buffer Replacement Algorithms LIRS-WSR and CCF-LRU for Flash Memory Based Systems**” in partial fulfillment of the requirements for the degree of M.Sc. in Computer Science and Information Technology be processed for the evaluation.

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LETTER OF APPROVAL

We certify that we have read this dissertation and in our opinion, it is satisfactory in the scope and quality as a dissertation in the partial fulfillment of the requirement of Masters Degree in Computer Science and Information Technology.

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ABSTRACT

Flash memory has the characteristics of not-in-place update and asymmetric Input/output (I/O) costs among read, write, and erase operations, in which the cost of write/erase operations is much higher than that of read operation. Among different flash-aware buffer replacement algorithms Low Inter-reference Recency Set Write Sequence Reordering (LIRS-WSR) and Cold Clean First Least Recently Used (CCF-LRU) are two buffer replacement policies that can be suitable for flash-based systems. LIRS-WSR enhances Low Inter-reference Recency Set (LIRS) by reordering the writes of not-cold-dirty pages from the buffer cache to flash storage to focus on the reduction of the number of write/erase operations as well as preventing serious degradation of buffer hit ratio. CCF-LRU, which enhances the previous Clean First Least Recently Used (CF-LRU) and Least Recently Used Write Sequence Reordering (LRU-WSR) methods by differentiating clean pages into cold and hot ones, and evicting cold clean pages first and delaying the eviction of hot clean pages. The objective of this dissertation is mainly focused on evaluating the performance of LIRS-WSR and CCF-LRU buffer replacement algorithms. Finally, the comparative analysis based on quantitative analysis of those algorithms is performed based on the hit/miss rates and the number write counts.

Using the trace-driven simulation, when workload has high reference locality, LIRS-WSR has significantly superior performance than CCF-LRU in terms of both hit rate and write count. LIRS-WSR has higher hit rate up to 17.5% and minimizes write count up to 37.17% in comparison to CCF-LRU. For uniformly distributed workloads, the difference in hit rates, and write count of LIRS-WSR and CCF-LRU is comparatively small. LIRS-WSR outperforms CCF-LRU by increasing hit rate up to 3.8% and decreasing write count up to 9.23% in comparison to CCF-LRU in its worst case.

Keywords:

CCF-LRU, Flash memory, Hit Rate, LIRS, LIRS-WSR, Write Count

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LIST OF ABBREVIATIONS

2Q Two Queue

μs Microsecond

AD-LRU Adaptive Double Least Recently Used

ARC Adaptive Replacement Cache

CCF Cold Clean First

CAR Clock with Adaptive Replacement

CCF-LRU Cold Clean First Least Recently Used

CDCSIT Central Department of Computer Science and Information Technology

CFDC Clean First Dirty Clustered

CF-LRU Clean First Least Recently Used

CLOCK-Pro Clock with Pro

CPU Central Processing Unit

CRF Combined Recency and Frequency

DRAM Dynamic Random Access Memory

DOECE Department of Electronics and Computer Engineering

EELRU Early Eviction Least Recently Used

EEPROM Electrically Erasable Programmable Read Only Memory

FIFO First In First Out

GCLOCK Generalized CLOCK

GiB Giga Byte

GHz Giga Hertz

HIR High Inter-reference Recency

HIRS High Inter-reference Recency Set

I/O Input/output

IOE Institute of Engineering

IRR Inter-Reference Recency

kB Kilo Byte

LFU Least Frequently Used

LIR Low Inter-reference Recency

LIRS Low Inter-reference Recency Set

LIRS-WSR Low Inter-reference Recency Set Write Sequence Reordering

LRFU Least Recently Frequently Used

LRU Least Recently Used

LRU-WSR Least Recently Used Write Sequence Reordering

mA Milliampere

MRU Most Recently Used

ms Millisecond

NRU Not Recently Used

OLTP Online Transaction Processing

OPT or MIN OPTimum or MINimum

OS Operating System

PDA Personal Digital Assistant

RAM Random Access Memory

TU Tribhuvan University

USB Universal Serial Bus

WSR Write Sequence Reordering

Chapter 1

BACKGROUND AND INTRODUCTION

1.1 Background

1.1.1 Flash Memory

Flash memory is non-volatile, shock resistant, and power economic. With recent technology breakthroughs in both capacity and reliability, flash-memory storage systems are much more affordable than ever. As a result, flash memory is now among the top choices for storage media in embedded systems [1]. There are two major types of flash memory in the current market: NAND and NOR flash memory. NAND flash memory is mainly designed for data storage and NOR flash memory is for Electrically Erasable Programmable Read Only Memory (EEPROM) replacement [1]. Flash memory has become a powerful and cost-effective solid-state storage technology widely used in mobile electronics devices and other consumer applications. NAND Flash, which was designed with a very small cell size to enable a low cost-per-bit of stored data, has been used primarily as a high-density data storage medium for consumer devices such as digital still cameras and Universal Serial Bus (USB) solid-state disk drives. NOR Flash has typically been used for code storage and direct execution in portable electronics devices, such as cellular phones and Personal Digital Assistants (PDAs) [2]. However, several hardware limitations exist in a flash memory. Firstly, a data unit of erase operations is a block that is the set of fixed number of contiguous pages even if a data unit of read/write operations is a page. Secondly, it is impossible to re-write the page in-place in a flash memory. So, in order to update data of the page, a system should perform only one of the following:

1. Writing these data to a newly allocated page and invalidating the original page; or
2. Writing these data to the original page only after erasing the block containing that page.

Thirdly, the lifetime of a flash memory is shorter than the lifetime of a hard disk and a Dynamic Random Access Memory (DRAM). In other words, only a limited number of erase operations can be performed safely to each memory cell, typically, between 100,000 and 1,000,000 cycles. Finally, there exist differences among I/O latencies according to the kinds of I/O operations, i.e., read, write, and erase. The write operation is about 10 times slower than the read operation and the erase operation is about 20 times slower than the write operation [3, 4, 5].

Flash caching is needed for reducing flash I/O latencies. The traditional magnetic-disk-based buffering algorithms Least Recently Used (LRU) [6], LIRS [7], Adaptive Replacement Cache (ARC) [8] etc. focus on hit-ratio improvement alone, but not on write costs caused by the replacement process. So, their straight adoption would result in poor buffering performance and would demote the development of flash-based systems. The replacement policy should minimize the number of writes and erases operations on flash memory, and at the same time prevent the degradation of the hit ratio. Recently, CF-LRU [9], LIRS-WSR [3] and Adaptive Double Least Recently Used (AD-LRU) [10] were proposed as new buffering algorithms for flash-based systems. These new flash based buffer replacement policies consider not only buffer hit ratios but also replacement costs incurring when a dirty page has to be propagated to flash memory to make room for a requested page currently not in the buffer. These algorithms try to obtain the optimal I/O sequence from the given I/O sequence by discriminatively selecting the accesses according to the type of I/O operations. These algorithms favor to first evict clean pages from the buffer so that the number of writes incurring for replacements can be reduced.

Table 1.1: Characteristics of flash memory

Device	Current (mA)		Access time (4 kB)		
	Idle	Active	Read	Write	Erase
NOR	0.03	32	20 μ s	28 ms	1.2 sec
NAND	0.01	10	25 μ s	250 μ s	2 ms

Compared to the magnetic disk, flash memory has special properties. Firstly, flash memory has no latency associated with the mechanical head movement to locate the proper position to read or write data. Secondly, flash memory has asymmetric read and write operation characteristic

in terms of performance and energy consumption. Table 1.1 compares the access time and the energy consumption in flash memory when 4 kB data is read, written, or erased [1]. Thirdly, flash memory does not support in-place update; the write to the same page cannot be done before the page is erased. Thus, as the number of write operations increases so does the number of erase operations. If the erase operations are involved, the cost imbalance would be even worse. Finally, blocks of flash memory are worn out after the specified number of write/erase operations.

1.1.2 Performance Metrics

The off-line performance of buffer replacement algorithm is measured in terms of page fault count, hit rate and hit ratio, miss rate and miss ratio, and write count. When an accessed block of memory is currently mapped to the physical memory then hit occurs. If it doesn't map then miss occurs. The Higher hit rate of the algorithm exhibits higher performance. In the case of flash based system, a minimum number of write count is a measure for optimal cost algorithm. flash based system, a minimum number of write count is a measure for optimal cost algorithm.

1.1.2.1 Page Fault Counts

Page Fault is an interrupt generated when the processor references a page that is neither in cache nor in main memory. An efficient page replacement algorithm always produces less number of page faults. It can be computed by counting the occurrences of a number of page faults between some intervals of references.

1.1.2.2 Hit/Miss Rate

When the processor needs to read or write a location in main memory, it first checks whether that memory location is in the cache. This is accomplished by comparing the address of the memory location to all tags in the cache that might contain that address. If the processor finds that the memory location is in the cache, we say that a **cache hit** has occurred; otherwise, we speak of a **cache miss**.

Miss rate can be calculated by using the formula:

miss rate = 1-hit rate.

Hit ratio is calculated by subtracting miss ratio from 1.

Miss ratio (mr) is calculated by using the formula:

$$mr = 100 * ((\#pf - \#distinct) / (\#refs - \#distinct))$$

where #pf is a number of page faults, #distinct is the number of distinct pages referenced and #refs is the total number of referenced pages [11].

Hit ratio is also calculated by dividing a total number of hit counts by a total number of reference counts.

$$\text{Hit ratio} = \text{Total number of Hit Counts} / \text{Total number of Reference Counts}$$

To represent it as a percentage: $\text{Hit\%} = \text{Hit ratio} * 100$

1.1.2.3 Write Counts

Write count is a number of pages propagated to flash memory which can be calculated by counting the number of physical pages writes to flash memory and at the end of each test the dirty pages in the buffer are flushed to the flash memory to get exact write counts.

1.1.3 Program behavior

There are several factors that influence the performance of page replacement algorithm. The performance of page replacement algorithm relies on the pattern of pages that are referenced. The behavior of the program depends on the access pattern it references memory which further depends on working set and locality of reference.

1.1.3.1 Locality of Reference

The locality of reference, also known as the principle of locality. During the course of execution of program memory references tend to cluster forming certain locality. Locality varies on the basis of time and space. Temporal locality is based on time; it assumes that memory location referenced just now is likely to be referenced again in near future. Looping, subroutines, stacks, the variable used for counting and totaling etc. supports this assumption. Spatial locality is based on space, it assumes that once a memory is referenced there is a high chance of nearby memory location to be referenced again. Array traversal, sequential code execution, related

variable declaration nearby in source code supports this assumption. Hints of the locality are followed in any type memory reference sequence.

1.1.3.2 Memory Reference Patterns

Altogether three types of standard synthetic traces i.e. random traces, read-most traces and write-most traces are used in this dissertation.

1.1.3.2.1 Random Traces The page references having random read and write nature of pages are called random traces.

1.1.3.2.2 Write-most Traces The page references having most of the pages with write mode nature are called write-most traces.

1.1.3.2.3 Read-most Traces The page references having most of the pages with read mode nature are called read-most traces.

1.1.3.2.4 Zipf Traces Zipf trace has a referential locality 20/80 meaning that eighty percent of the references deal with the most active twenty percent of the pages.

1.2 Introduction

1.2.1 Problem Definition

Since the use of flash memory requires buffer replacement policies considering not only buffer hit ratios or miss ratios but also replacement costs incurring when a dirty page has to be propagated to flash memory, not in the buffer. There are many buffer replacement algorithms developed for flash-based systems. The traditional magnetic-disk-based buffering algorithms LRU [6], LIRS [7], ARC [8] etc. focus on hit-ratio improvement alone, but not on write costs caused by the replacement process. However, Flash memory has characteristics of out-of-place update and asymmetric I/O latencies for read-write and erase operations in the aspects of time and energy. So, the replacement algorithm with flash memory should consider not

only the hit count but also the replacement cost caused by selecting dirty victim pages. There are many buffer replacement algorithms developed for flash-based systems. The evaluation of these buffer replacement algorithms for flash-based systems in terms of hit rate and write counts is required to rate their performance. This dissertation work has mainly focused on the comparative evaluation of two algorithms: LIRS-WSR and CCF-LRU.

1.2.2 Objectives

The main objectives of this dissertation work are:

- To simulate the algorithm LIRS-WSR and CCF-LRU; and
- To perform a comparative analysis of LIRS-WSR and CCF-LRU buffer replacement algorithms for flash based systems in terms of hit rate and write count.

1.2.3 Motivation

Since the use of flash memory requires buffer replacement policies considering not only buffer hit ratios or miss ratios but also replacement costs incurring when a dirty page has to be propagated to flash memory, not in the buffer. As a consequence, a replacement policy should minimize the number of write/erase operations on flash memory and at the same time increase the hit ratio. Memory management is not only the burden of today's computing devices. It has been researched for decades. Whatever variety of storage devices found in today's market is the great achievement of computer science. But still, computer memory is the limited source which directly hampers the performance of computing system. Performance gain can be achieved by increasing the capacity of primary storage. The expectation of customer is to decrease cost price with sufficient working memory. Hence to fulfill this demand for manufacturing such device fewer materials are used and size of memory is being decreased. But rather than this technical view, it is not possible to gain performance without managing memory logically for its usability. Varieties of techniques had been tried for this achievement. Among such techniques, paging is the successful one. Page replacement algorithm is the main part of paging technique because deciding the victim page is a very tough job.

The emergence of single flash memory chip with several gigabytes capacity makes a strong tendency to replace the magnetic disk with flash memory for the secondary storage of mobile

computing devices. Most operating systems are customized for disk-based storage systems and their replacement policies only concern the number of cache hits. However, the operating systems that consider flash memory as secondary storage should consider different read and write cost of flash memory when they replace pages to reclaim free space. There are different buffer replacement algorithms proposed for flash-based storage systems. Some of them consider recency factor only, some consider cleanliness, some both of these factors and some consider recency, cleanliness, and frequency of page references as well.

1.3 Dissertation Organization

Background part of this dissertation work focuses the related basic terms and terminologies of the performance of buffer algorithms which are already mentioned. Some chapters are remaining which clarifies the topics fulfilling the objectives of this dissertation work.

Chapter 2 consists of literature review which briefly reviews the related topics. The literature review includes the summary of several traditional page replacement.

Chapter 3 consists of research methodology steps of the dissertation which includes details about data collection and memory references that show trace-driven input. Also, it includes details about the used programming tools, data structures, algorithm, and the flowchart to develop the simulator.

Chapter 4 consists results, analysis, and comparisons which include output results with several analyzing graphs which are tested for different workloads.

Chapter 5 consists of the conclusion of this whole dissertation work and the future work which shows guidelines for further research.

Chapter 2

LITERATURE REVIEW

2.1 Traditional Buffer Replacement Algorithms

2.1.1 OPT or MIN Page Replacement Algorithm

Various memory management techniques have been used from the beginning for the improvement of performance. Bélády [6] in 1966 developed optimal page replacement algorithm called OPTimum or MINimum (OPT or MIN). His algorithm depends on the principle of optimality which states “To obtain optimal performance the page to replace is the one that will not be used again for the furthest time into the future.” His optimal algorithm is not applicable for real implementation because our Operating System (OS) doesn’t know which pages will be used before execution. Hence it can be only simulated due to lack of future knowledge. It is used as a benchmark for measuring the effectiveness of other page replacement algorithms. OPT Replacement algorithm replaces the page that will not be used for the longest period of time by computing maximum forward distance. From the past experiences and research papers, the research on the page replacement algorithms are categorized into LRU based replacement algorithms and CLOCK based replacement algorithms [12].

2.1.2 LRU Based Page Replacement

2.1.2.1 FIFO Page Replacement Algorithm

The simplest page-replacement algorithm is a First In First Out (FIFO) [6] algorithm. In this algorithm, operating system keeps track of all pages in the memory in a queue, oldest page

is in the front of the queue. When a page needs to be replaced the page in the front of the queue is selected for removal. Conceptually FIFO is a queue with limited size. Initially, the queue is filled by inserting page reference from the tail. When the queue is full new reference is inserted from tail and old reference is evicted from the head. FIFO is simple but suffers from Bélády's Anomaly, a strange situation in which page fault increased while increasing the number of the page frame. That is, with an increase in physical memory, FIFO can decrease page fault performance seemingly at random. Like random page replacement algorithm, FIFO still does not take advantage of locality trends. But it can be modified very easily.

2.1.2.2 LRU Page Replacement Algorithm

A good approximation to the optimal algorithm is based on the observation that pages that have been heavily used in the last few instructions will probably be heavily used again in the next few. Conversely, pages that have not been used for ages will probably remain unused for a long time. This idea suggests a realizable algorithm: when a page fault occurs, throw out the page that has been unused for the longest time. This strategy is called LRU paging [6]. This algorithm is purely based on recency of page references. Recency is evaluated by maintaining LRU stack that is a sorted list on the basis of virtual time which is the only factor for replacement. Thus LRU is simple but is not easy to implement without hardware support. It can adapt faster according to as program behavior. LRU like algorithm doesn't suffer from Bélády's Anomaly as FIFO. It gives a good approximation of optimal algorithm. Although LRU is theoretically realizable, it is not cheap. To fully implement LRU it is necessary to maintain a linked list of all pages in memory, with the most recently used page at head and least recently used page at the tail. The difficulty is that the list must be updated on every memory reference. Finding a page in the list, deleting it and then moving it to the head is a very time-consuming operation.

2.1.2.3 NRU Page Replacement Algorithm

Not Recently Used (NRU), sometimes known as the LRU, page replacement algorithm is an algorithm that favors keeping pages in memory that have been recently used. This algorithm works on the following principle: when a page is referenced, a referenced bit is set for that page, marking it as referenced. Similarly, when a page is modified, a modified bit is set. The setting of the bits is usually done by the hardware, although it is possible to do so on the software level

as well. Pages are categorized into four classes in NRU algorithm:

- Class 0 contains pages that are neither referenced nor modified;
- Class 1 contains pages that are modified but not referenced;
- Class 2 contains pages that are referenced but not modified; and
- Class 3 contains pages that are modified as well as referenced.

During page fault, NRU evicts any page from the lowest class [13].

2.1.2.4 LFU Page Replacement Algorithm

Least Frequently Used (LFU) [14] selects a victim page that has not been used often in the past. Instead of using a single recency factor as LRU, LFU defines additional information of frequency of use associated with each page. This frequency is calculated throughout the reference stream by maintaining counting information. Frequency count leads to the serious problem after a long duration of reference stream. Because when the locality changes, reaction to such certain change will be extremely slow. Assuming that a program either changes its set of active pages or terminates and is replaced by a completely different program, the frequency count will cause pages in the new locality to be immediately replaced since their frequency is much less than the pages associated with the previous program. Since the context has changed, the pages swapped out will most likely be needed again soon which leads to thrashing. One way to remedy this is to use a popular variant of LFU, which uses frequency counts of a page since it was last loaded rather than since the beginning of the page reference stream. Each time a page is loaded, its frequency counter is reset rather than being allowed to increase indefinitely throughout the execution of the program. LFU still tends to respond slowly to change in the locality.

2.1.2.5 EELRU Page Replacement Algorithm

Some algorithms use recency as history information like LRU and Most Recently Used (MRU). LRU is suitable for the good locality of reference whereas MRU is somewhat suitable for the weak locality of workloads. These two algorithms can be tuned to form adaptive algorithm

called Early Eviction Least Recently Used (EELRU) [15], which was proposed as an attempt to mix LRU and MRU, based only on the positions on the LRU queue that concentrates most of the memory references. This queue is only a representation of the main memory using the LRU model, ordered by the recency of each page. EELRU detects potential sequential access patterns analyzing the reuse of pages. One important feature of this algorithm is the detection of non-numerically adjacent sequential memory access patterns. Two tunable parameters used are early eviction point and late eviction point. LRU queue concentrates most of the memory references when it reaches late eviction point.

2.1.2.6 LRFU Page Replacement Algorithm

The LRU and LFU replacement policies are two extreme replacement policies. The LRU policy gives weight to only one reference for each block, that is, the most recent reference to the block while giving no weight to older ones representing one extreme, and the LFU gives equal weight to all references representing the other extreme. These extremes imply the existence of a spectrum between them. In [16], proposed such a spectrum which is called the Least Recently Frequently Used (LRFU) policy. The LRFU policy associates a value with each block. This value is called the Combined Recency and Frequency (CRF) value and quantifies the likelihood that the block will be referenced in the near future. The performance of the LRFU algorithm largely relies on a parameter called λ , which determines the relative weight of LRU or LFU and has to be adjusted according to the system configuration, even according to different workloads [17].

2.1.2.7 LRU-K Page Replacement Algorithm

The LRU policy takes into account the recency information while evicting pages, without considering the frequency. To consider the frequency information, LRU-K [18] was proposed which evicts pages with the largest backward K-distance. Backward K-distance $b_t(p, K)$ can be defined as the distance backward to the K^{th} most recent reference to page p where reference string is known up to time t (r_1, r_2, \dots, r_t). The value of parameter K can be taken as 1, 2 or 3. If $K=1$, it works as simple LRU algorithm. Highly increasing the value of K the overall performance of algorithm reduces. LRU-K can discriminate better between frequently referenced and infrequently referenced pages. Unlike the approach of manually tuning the assignment of

page pools to multiple buffer pools, LRU-K does not depend on any external hints. Unlike LFU and its variants, this algorithm copes well with temporally clustered patterns.

2.1.2.8 2Q Page Replacement Algorithm

Two Queue (2Q) [19] is a good buffering algorithm (giving a 5-10% improvement in hit rate over LRU for a wide variety of applications and buffer sizes, and never hurting), having constant time overhead, and requiring little or no tuning. It works well for the same intuitive reason that LRU/B works well: it bases buffer priority on sustained popularity rather than on a single access. 2Q algorithm quickly removes sequentially and cyclically referenced block with after a long interval. The algorithm uses special buffer queue $A1_{in}$ of size K_{in} , ghost buffer queue $A1_{out}$ of size K_{out} and the main buffer A_m . The special buffer contains all missed that is first time referenced block. Ghost buffer contains replaced blocks from the special buffer. Frequently accessed blocks are available in the main buffer. Hence victim blocks are always from the special buffer and main buffer.

2.1.2.9 LIRS Page Replacement Algorithm

Another important algorithm is LIRS which is already described in section 2.2.4. Its objective is to minimize the deficiencies presented by LRU using an additional criterion named Inter-Reference Recency (IRR) that represents the number of different pages accessed between the last two consecutive accesses to the same page. This means that LIRS does not replace the page that has not been referenced for the longest time, but it uses the access recency information to predict which pages have more probability to be accessed in a near future.

2.1.2.10 ARC Page Replacement Algorithm

The ARC [8] is an adaptive page replacement algorithm developed at the IBM Almaden Research Center. The algorithm keeps a track of both frequently used and recently used pages, along with some history data regarding eviction for both. It improves the LRU strategy by splitting the cache directory into two lists, T_1 and T_2 , for recently and frequently referenced entries. In turn, each of these is extended with a ghost list (B_1 or B_2) which is attached to the bottom of these two lists. These ghost lists act as score cards by keeping track of the history of recently evicted cache entries and the algorithm uses ghost hits to adapt to recent change in resource

usage. The ghost lists only contain meta-data (keys for the entries) and not the resource data itself, i.e. as an entry is evicted into a ghost list its data is discarded. The combined cache directory is organized in four LRU lists:

1. T_1 , for recent cache entries;
2. T_2 , for frequent entries, referenced at least twice;
3. B_1 , ghost entries recently evicted from the T_1 cache but are still tracked; and
4. B_2 , similar ghost entries, but evicted from T_2 .

T_1 and B_1 together are referred to as L_1 , a combined history of recent single references. Similarly, L_2 is the combination of T_2 and B_2 .

2.1.3 CLOCK Based Page Replacement Algorithm

2.1.3.1 CLOCK Page Replacement Algorithm

Frank Corbató (who later went on to win the ACM Turing Award) introduced CLOCK [20] as a one-bit approximation to LRU, and its performance characteristics are very similar to those of LRU. So all the performance disadvantages about LRU are also applied to CLOCK. In CLOCK, the memory spaces holding the pages can be regarded as a circular buffer. Here each page is associated with a bit, called reference bit, which is set by hardware whenever the page is accessed. When it is necessary to replace a page to service a page fault, the page pointed to by the hand is checked. If its reference bit is unset, the page is replaced. Otherwise, the algorithm resets its reference bit and keeps moving the hand to the next page [21].

When a page fault occurs, the page being pointed to by the hand is inspected.

- $R = 0$, the page is evicted, the new page is inserted into the clock in its place, and the hand is advanced one position; and
- $R = 1$, it is cleared and the hand is advanced to the next page. This process is repeated until a page is found with $R = 0$.

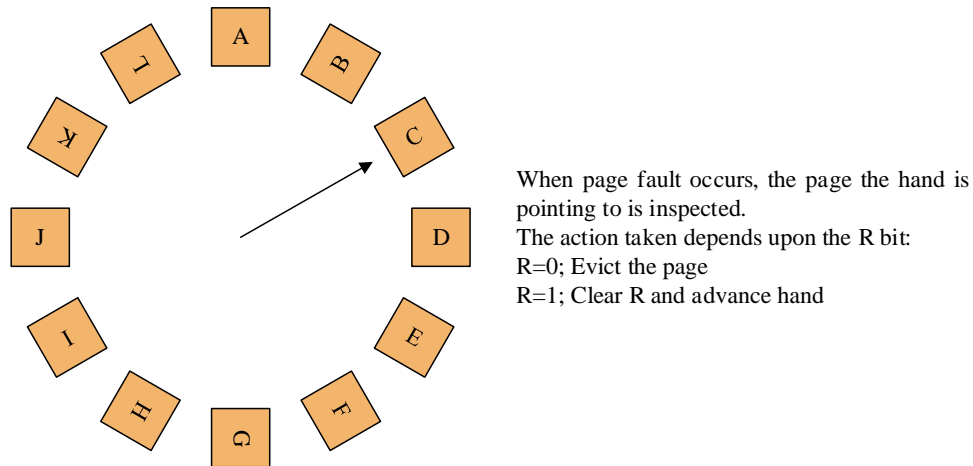


Figure 2.1: CLOCK page replacement algorithm

2.1.3.2 Clock-Pro

In Paper [21], proposed an improved CLOCK replacement policy, called Clock with Pro (CLOCK-Pro). It takes the same principle as that of LIRS-it uses reuse distance (called IRR in LIRS) rather than recency in its replacement decision. When a page is accessed, the reuse distance is the period of time in terms of the number of other distinct pages accessed since its last access. A page is categorized as a cold page if it has a large reuse distance or as a hot page if it has a small reuse distance. Although there is a reuse distance between any two consecutive references to a page, only the most current distance is relevant in the replacement decision.

2.1.3.3 CAR Page Replacement Algorithm

In paper [22], proposed a simple and elegant new algorithm, namely, Clock with Adaptive Replacement (CAR), that has several advantages over CLOCK:

1. It is scan-resistant;
2. It is self-tuning, and it adaptively and dynamically captures the recency and frequency features of a workload;
3. It uses essentially the same primitives as CLOCK, and, hence, is low-complexity and amenable to a high-concurrency implementation; and
4. It outperforms CLOCK across a wide range of cache sizes and workloads.

The algorithm CAR is inspired by the Adaptive Replacement Cache (ARC) algorithm, and inherits virtually all advantages of ARC including its high performance, but does not serialize cache hits behind a single global lock.

This algorithm uses two clocks T_1 & T_2 and two lists B_1 & B_2 . T_1 and T_2 contain cold pages and hot pages i.e. contain pages in the cache, while B_1 & B_2 maintain history information about the recently evicted pages from B_1 & B_2 respectively.

2.1.3.4 GCLOCK Page Replacement Algorithm

Paper [23] proposed a new algorithm, namely, Generalized CLOCK (GCLOCK). The GCLOCK buffer replacement policy uses a circular buffer and a weight associated with each page brought in the buffer to decide on which page to replace. Whenever a page is referenced, the associated count field is set to i . When a page fault occurs, a pointer that circles around this circular list of page frames is observed. If the count field pointed to zero, then the page is removed and the new page is placed in that frame. Otherwise, the count is decremented by 1, the pointer is advanced to the next count field and the process is repeated. When a new page is placed in the page frame, the count field is set to i if the page is to be referenced (demand fetch) and it is set to j if the page has been pre-paged and is not immediately referenced. This algorithm abbreviated by writing CLOCKP (j, i). The P indicates that this is a pre-paging algorithm (the pre-paging strategy has not been specified).

2.2 Buffer Replacement Algorithms for Flash-Based Systems

2.2.1 CF-LRU

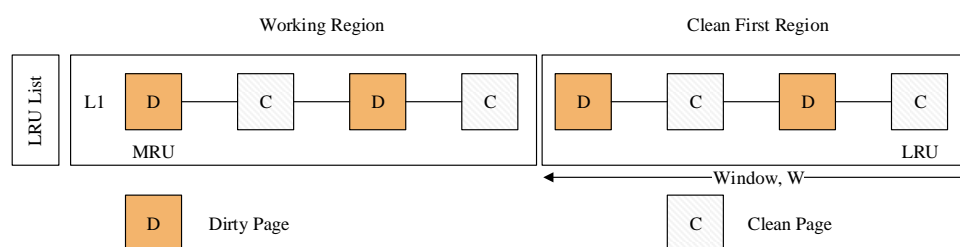


Figure 2.2: CF-LRU page replacement algorithm

CF-LRU [9] is the first algorithm designed for flash-based systems. CF-LRU tries to reduce the number of costly write operations and potential erase operations until the degradation of cache hit rate does not harm the performance. It modified the LRU policy by introducing a clean first window W , which starts from the LRU position and contains the least recently used $w*B$ pages, as shown in Figure 2.2, where B is the buffer size and w is the ratio of the window size to buffer size. When the victim is selected, CF-LRU first evicts least recently used clean pages in W . Hence it reduces the number of write operations because the clean page is not propagated to flash memory. If no clean page is found, then it behaves like LRU policy. CF-LRU has some problems such as clean-first window size is to be tuned to the current workload and cannot suit for differing workloads and it always replaces clean pages first, which causes the cold dirty pages residing in the buffer for a long time and, in turn, results in suboptimal hit ratio. The window size, W , can be tuned statically or dynamically. In this sense, CF-LRU is known as CF-LRU-static or CF-LRU-dynamic. In paper [9], CF-LRU-static and CF-LRU-dynamic have been compared with LRU policy for five different workloads. They found a result that CF-LRU static and dynamic reduces the replacement cost by 28.4% and 23.1% for swap system buffer cache and 26.2% and 23.5% for file system buffer cache with compared to LRU.

2.2.2 CFDC

Clean First Dirty Clustered (CFDC) [24] improves the efficiency of the buffer manager by flushing pages in a clustered fashion, based on the observation that flash writes with strong spatial locality can be served by flash disks more efficiently than random writes. It manages

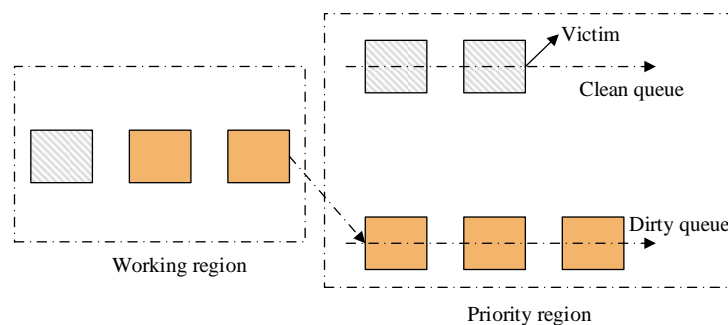


Figure 2.3: CFDC page replacement algorithm

the buffer in two regions: the working region W for keeping hot pages that are frequently revisited and the priority region P responsible for optimizing replacement costs by assigning varying priorities to pages. A parameter λ , called priority window, determines the size ratio

of P relative to the total buffer. Therefore, if the buffer has B pages, then P contains λ pages and the remaining $(1-\lambda)*B$ pages are managed in W . Various conventional replacement policies can be used to maintain high hit ratios in W and, therefore, prevent hot pages from entering P . CFDC improves the efficiency of buffer manager by flushing pages in clustered fashion based on the observation that flash writes with strong spatial locality can be served by flash disks more efficiently than random writes. In paper [24], CFDC has been compared with LRU and CF-LRU for different four workloads in the database engine. The results show CFDC outperforms both competing policies, with a performance gain between 14% and 41% over CF-LRU, which, in turn, is only slightly better than LRU with a maximum performance gain of 6%.

2.2.3 LRU-WSR

LRU-WSR [25] is a flash-aware algorithm based on LRU and Second Chance [20], using only a single list as an auxiliary data structure. The idea is to evict clean and cold-dirty pages and keep the hot-dirty pages in the buffer as long as possible. When a victim page is needed, it starts searching from the LRU end of the list. If a clean page is found, it will be returned immediately (LRU and clean-first strategy). If a dirty page marked as “cold” is found, it will also be returned; otherwise, it will be marked “cold” (Second Chance), moved to the MRU end of the list and the search continues. Although LRU-WSR considers the hot/cold property of dirty pages, which is not tackled by CF-LRU, it has a high dependency on the write locality of workloads. It shows low performance in case of low write locality, which may cause dirty pages to be quickly evicted. In paper [25], LRU-WSR has been compared with LRU, CF-LRU algorithms for different workloads collected from PostgreSQL, GCC, Viewperf, and Cscope. LRU-WSR has been found 1.4 times faster than LRU. In most of the cases, LRU-WSR has higher hit ratio and lower write count than others.

2.2.4 LIRS-WSR

LIRS-WSR [3] algorithm is designed for a buffer cache of the flash memory based storage system. The objective of LIRS-WSR is reducing the number of flushes of dirty pages from the buffer into flash memory when page replacement occurs, To achieve this objective, it uses the strategy: delaying eviction of the page which is dirty and has high access frequency as possible. It enhances an existing LIRS buffer replacement algorithm with add-on buffer replacement

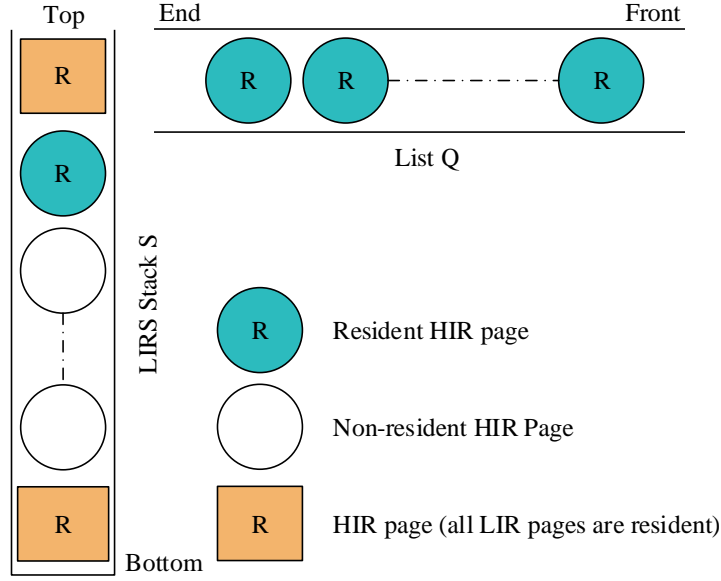


Figure 2.4: Two Lists of the LIRS algorithm

strategy, namely Write Sequence Reordering (WSR). WSR reorders writing not-cold dirty pages from the buffer cache to the disk to reduce the number of write operations while preventing excessive degradation of the hit ratio. The LIRS [7] algorithm uses history information of data accesses in the form of two metrics - the IRR and the Recency. The IRR of a data block refers to the number of other distinct blocks accessed between the last two consecutive accesses of the data block in question while recency refers to the number of other distinct blocks accessed between the last reference to the current time.

LIRS algorithm uses two sets of pages based on IRR. Set of pages with low IRR value is taken as a hot block and called LIRS. Set of pages with high IRR value is taken as a cold block and called High Inter-reference Recency Set (HIRS). Blocks that can be most probably used in future are taken as hot blocks whereas blocks that may not be used in near future are taken as cold blocks. Hence, High Inter-reference Recency (HIR) blocks are always replaced and Low Inter-reference Recency (LIR) blocks are never replaced. LIRS always selects HIR page with the largest recency as a victim for replacement. WSR policy is developed to adapt LIRS with flash memory [3]. The basic scheme of WSR is following:

1. Use cold-detection algorithm to judge whether the page is cold or not; and
2. Delays flushing dirty pages which are not regarded as cold.

LIRS-WSR is implemented using 2 lists: LIR stack S which stores all LIR pages as well as HIR pages regardless of the residence status—some of them are resident and others are not

(actually, only their meta-data are stored in the list)—and HIR list Q that stores HIR resident pages. The operations on these two data structures are same that of LIRS. Every page has additional status either cold or not-cold. Initially, all pages are cold, this cold flag is cleared if the pages are referenced again when they are in stack S or queue Q. If a page is introduced to the buffer for a write request for the first time, it becomes a dirty page and enters the top of the stack S as an LIR page. Every time when Stack's bottom is moved to HIR Q, WSR policy is applied. That is, if bottom LIR page is dirty and not cold, then it's cold flag is set and moved to the head of Stack, otherwise, it is moved to the head of HIR Q. All other operations like pruning, switching between LIR and HIR pages are same as that of LIRS.

2.2.5 CCF-LRU

An efficient new buffer replacement algorithm for flash memory based storage systems called CCF-LRU [26]. The goal of CCF-LRU is to improve the overall I/O performance by focusing on reducing the write count incurring in the replacement process. In order to accomplish this goal, it tries to first evict clean pages with low access frequencies. If there are no such clean pages, it will evict the dirty pages with low access frequencies instead of the clean pages with high access frequencies. Using the cold-detection mechanism of LRU-WSR, pages in the buffer list can be classified into the following four groups, namely cold clean page, hot clean page, cold dirty page, and hot dirty page. A cold flag attached with each page is used to distinguish cold pages from hot pages. The CCF-LRU algorithm maintains two LRU lists, which are called

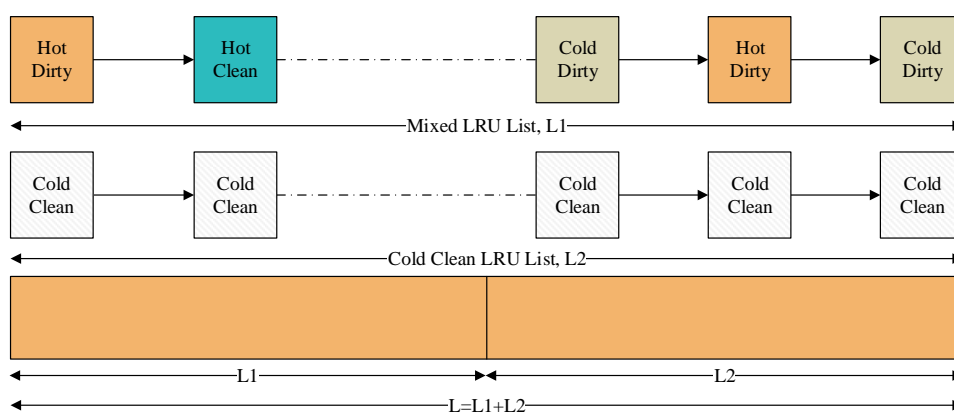


Figure 2.5: Mixed LRU list and cold clean LRU list in CCF-LRU

mixed LRU list and cold clean LRU list. The mixed LRU list contains L1 pages and is used to maintain hot clean pages and dirty pages regardless of the status of its cold flag and the cold

clean LRU list with the size of L2 is only for cold clean pages. If the buffer contains a total of L pages, the sizes of the two LRU lists are both from 0 to L. Moreover, the sum of L1 and L2 is L. The first referenced pages are regarded as cold by default, each of which is inserted into the cold clean LRU list with a cold flag. When the page in the cold clean LRU list is referenced again or becomes dirty, it will be moved from the cold clean LRU list to the MRU position in the mixed LRU list. When the page in the mixed LRU list is referenced, it will be moved to the MRU position of the mixed LRU list. The CCF-LRU selects a victim page by the following rules in order:

1. If the cold clean LRU list is not empty, the LRU page in the cold clean LRU list is selected as the victim; and
2. If the cold clean LRU list is empty, the LRU page in the mixed LRU list is chosen as the victim candidate. If the candidate is a cold dirty page, it is selected as the victim. If the candidate is a hot dirty page, it is labeled as cold and moved to the MRU position of the mixed LRU list. If the candidate is a hot clean page, it is set to cold and moved from the mixed LRU list to the MRU position of the cold clean LRU list and we continue to check the LRU position in the mixed LRU list. If there is no victim found after traversing the mixed LRU list, it needs to call the CCF-LRU algorithm one more time to select a victim.

2.2.6 AD-LRU

AD-LRU algorithm [10] is buffer replacement algorithm for flash-based systems which focuses on reducing the write costs of the buffer replacement algorithm while keeping a high hit ratio. It tries to integrate the properties: recency, frequency, and cleanness of pages into the buffer replacement policy. AD-LRU has two LRU queues: Cold LRU queue and Hot LRU queue, to capture the concept of recency and frequency of the page references, among which Cold LRU queue stores the pages referenced only once and Hot LRU queue maintains the pages that are referenced at least twice. The sizes of these two LRU queues are dynamically adjusted according to changes in reference patterns. When a page is first referenced, it is put in the head of cold LRU queue. The pages move from cold LRU queue to head of hot LRU queue when it is referenced again and when a page in hot LRU queue is selected as a victim, it is demoted to head of cold LRU queue. During the eviction procedure, least recently used a clean page from cold LRU queue is selected as a victim.

Chapter 3

RESEARCH METHODOLOGY

The main purpose of research is to discover answers to the questions through the applications of scientific procedures. Research is a careful study performed to find out new things in a systematic way. In a scientific method of research at the first problem is formulated then output information is generated from collected input data and output is analyzed and finally, the result is generalized [27]. This dissertation work is truly scientific and flows in the same way. Page replacement algorithm is one of the major strategies to manage memory efficiently. The main exploration of this dissertation focuses on LIRS-WSR and CCF-LRU algorithms developed to address flash memory characteristics in memory management. Out of different types of research methodologies, this dissertation is based on the trace-driven simulation approach. All the data collected which are traces of page references. Output information gathered is analyzed in a quantitative approach. Finally, the conclusion is drawn with the help of analyzed data.

3.1 Data Collection

Data are the sources of information. Hence data should be collected very carefully. In this dissertation work, four types of synthetic traces [10] have been used in the simulation experiment, i.e., random trace, read-most trace (e.g., of decision support systems), write-most trace (e.g., of Online Transaction Processing (OLTP) systems) and Zipf trace as Workload 1, Workload 2, Workload 3, and Workload 4 respectively. These data are real memory traces. Workload represents the different locality of memory reference patterns that are generated during execution of the process in real OS. There is total 100,000 pages reference in each of the first three traces,

which are restricted to a set of pages whose numbers range from 0 to 49,999. The total number of page references in the Zipf trace is set to 500000 in order to obtain a good approximation, while the page numbers still fall in $[0, 49999]$. Zipf trace has a referential locality “20/80” meaning that eighty percent of the references deal with the most active twenty percent of the pages. The Sample of Workload 1, Workload 2, Workload 3, and Workload 4 are in appendix A.1, appendix A.2, appendix A.3, and appendix A.4 respectively. Tables 3.1, 3.2, 3.3, and 3.4 show the details concerning these workloads.

Table 3.1: Trace for Random Access

Attributes	Value
Total I/O references	100,000
Total distinct references	43247
R/W ratio	50%/50%
Reference patterns	Uniform

Table 3.2: Trace for Read-most Access

Attributes	Value
Total I/O references	100,000
Total distinct references	43212
R/W ratio	90%/10%
Reference patterns	Uniform

Table 3.3: Trace for Write-most Access

Attributes	Value
Total I/O references	100,000
Total distinct references	43182
R/W ratio	10%/90%
Reference patterns	Uniform

Table 3.4: Trace for Zipf

Attributes	Value
Total I/O references	500,000
Total distinct references	47023
R/W ratio	50%/50%
Reference locality	20%/80%

3.2 Development Methodology and Tools

The simulator is built by using an incremental approach. The LRU stack automatically maintains recency factor. Information of recently referenced block is available on top of the stack and the oldest in the bottom of the stack. Every time when the block is accessed it is kept on top of the stack. LIRS-WSR and CCF-LRU algorithms are also implemented by using doubly link list and list interface. The algorithms have been implemented in C++ programming language using Microsoft® Visual Studio and intellij idea on Intel® Core™ i5-4210U Central Processing Unit (CPU) @ 1.7 GHz with 4 GiB Random Access Memory (RAM) Microsoft® Windows 10 1607, 64 bit OS.

3.3 LIRS-WSR

The LIRS [7] algorithm can be implemented using 2 lists: LIR stack S which stores all LIR pages as well as HIR pages regardless of the residence status some of them are resident and others are not (actually, only their meta-data are stored in the list) and HIR list Q that stores HIR resident pages. The sum of the size of HIRS and size of LIRS equals to the size of the cache. HIR block that may be resident or non-resident can be promoted to LIR block. At the same time to maintain the LIRS and HIRS size, oldest LIR block must be demoted to HIR-resident block. Then one of the resident HIR blocks becomes the victim. The promotion/demotion

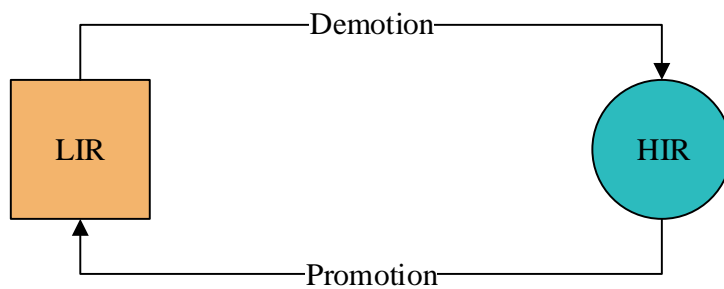


Figure 3.1: General LIR vs. HIR Transition Diagram

policy is shown in the Figures 3.1, 3.2, and 3.3 show the specific promotion/demotion policy among LIR, resident HIR and non-resident HIR, so as to maintain partition size. Every page has additional status either cold or not-cold. Its main purpose is to maintain recency value. As we move toward bottom recency factor increases. Bottom most one is always LIR block, which is the oldest block having higher recency factor and topmost one is the recent block having

recency factor equals to zero. Each stack node contains information about reference block.

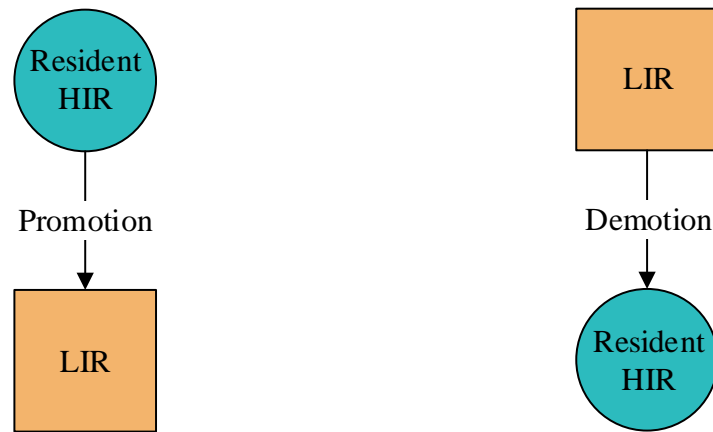


Figure 3.2: Specific LIR vs. Resident HIR Transition Diagram

Here information of every page reference is not available in stack S due to the major event stack pruning. Some information is also available in queue Q and some outdated information is also left in Stack. Queue Q contains collection referenced pages that are resident HIR blocks available in the cache. Hence the size of HIR cache partition determines the size of Queue Q. Initially, all pages are cold, this cold flag is cleared if the pages are referenced again when

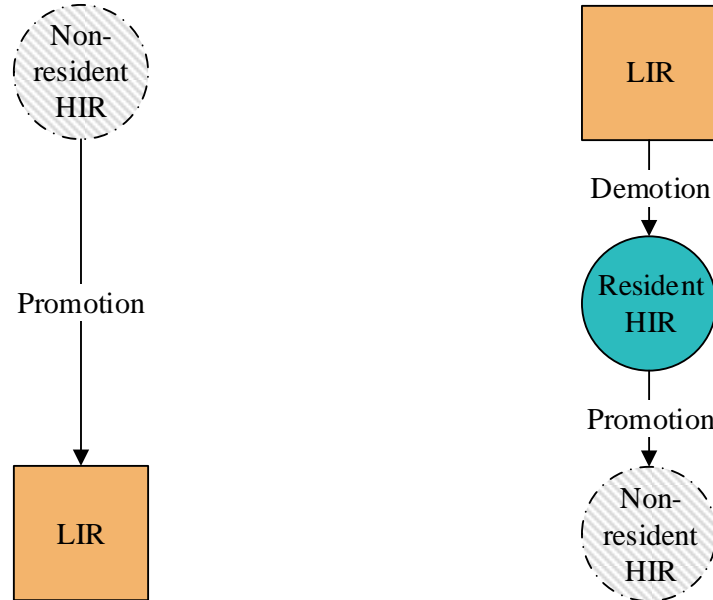


Figure 3.3: LIR vs. Non-Resident HIR Transition Diagram

they are in stack S or queue Q. The block in the Queue can be removed from anywhere if it is promoted to LIR. Comparing IRR and recency value is automatically done by the use of Q which increases performance. If a page is introduced to the buffer for a write request for the first time, it becomes a dirty page and enters the top of the stack S as an LIR page. Every time

when Stack's bottom is moved to HIR Q, WSR policy is applied. That is, if bottom LIR page is dirty and not cold, then it's cold flag is set and moved to the head of Stack, otherwise it is moved to the head of HIR Q. That is, only clean or dirty cold pages are moved to the head of HIR Q from the bottom of S. Stack pruning operation is performed on every move of operation performed on bottom LIR page of S.

3.3.1 Stack Pruning Function

The major function stack pruning is conducted during the status change. The Bold assumption of the algorithm is that bottom of stack S is always LIR block. While changing status, the page in the bottom of the stack S is demoted to HIR resident for that it is kept in queue Q. At that time next LIR bottom is chosen which is nearer from the bottom of stack S and all other HIR bottom are removed one by one. Information of thus removed HIRs is available in queue Q if it is a resident. Stack pruning is also conducted if the accessed block P is the bottom LIR because the recent block is always moved to the top of stack S. Stack pruning decreases the size of the stack hence the stack doesn't keep track of outdated references.

3.3.2 Data Structure

The LIRS-WSR algorithm can be implemented by using two LRU lists LIRS stack and HIR queue. Each node in the LIR list and HIR Q are implemented as a doubly linked list.

3.3.3 Algorithm

Algorithm 3.1 LIRS-WSR

INPUT: Pages

OUTPUT: Miss rate, hit rate etc.

```
1: Start
2: Read new page
3: if page is in Stack S then
4:   if If page is LIR page then
5:     Page hit LIR page
6:     Clear cold flag, i.e. cold=0
7:     Move the page to the head of Stack, S
8:     if If the page is at bottom of S then
9:       Prune the Stack, S
10:    end if
11:  else if Page is HIR then
12:    if Page is resident HIR then
13:      Page hit, HIR resident page
14:      Move the page to the head of Stack, making it LIR
15:      Clear cold flag i.e. cold=0
16:      Remove the page from HIR list Q
17:      while Stack's bottom is dirty and not cold do
18:        Move bottom page of Stack, S to the head of Stack, S
19:        Set cold=1 for this page
20:        Prune the Stack
21:      end while
22:      Move stack bottom to head of the HIR list, making it HIR
23:      Prune the Stack
24:    else
25:      Page is non-resident HIR
26:      Page miss
27:      Remove Tail of HIR Q
28:      Move it to the head of Stack S, making it LIR
29:      Clear cold flag of the page
30:      while Stack's bottom is dirty and not cold do
31:        Set the cold flag of bottom page, i.e. Cold=1
32:        Move the page to the head of Stack S
33:        Prune the Stack
34:      end while
35:      Move the bottom of Stack to head of the HIR queue Q
36:      Prune the Stack
37:    end if
38:  end if
39: else if Page is in the HIR Q then
40:   Page hit in HIR Queue
41:   Move to the head of HIR, Q
42:   Add to the head Stack S
```

Algorithm 3.1 LIRS-WSR (cont.)

```
43: else
44:   Page miss occurs
45:   if Free memory is available then
46:     if Free memory is larger than HIR_Limit then
47:       Add page to the head of Stack S
48:       Make it LIR page.
49:       Decrease free memory by one
50:     else
51:       if Page is write then
52:         Add the page to the head of Stack S
53:         Make it LIR
54:         while Stack's bottom is dirty and not cold do
55:           Move the bottom of Stack S to the head of Stack S
56:           Set cold=1 for this page
57:           Prune the Stack
58:         end while
59:         Move Stack bottom to head of the HIR list
60:         Make this page resident HIR
61:         Prune the Stack
62:         Decrease free memory by one
63:       else
64:         Page is read
65:         Add the page to head of queue Q
66:         Add the page to the head of Stack
67:         Decrease free memory by one
68:       end if
69:     end if
70:   else
71:     Memory is full
72:     Remove tail of HIR queue Q
73:     if Page is write then
74:       Add the page to the head of Stack S
75:       Make it LIR
76:       while Stack's bottom is dirty and not cold do
77:         Move the bottom of Stack S to the head of Stack S
78:         Set cold=1 for this page
79:         Prune the Stack
80:       end while
81:       Move Stack bottom to head of the HIR list
82:       Make this page resident HIR
83:       Prune the Stack
84:     else
85:       Page is read
86:       Add the page to head of Stack S
87:       Add the page to the head of queue Q
88:     end if
89:   end if
90: end if
91: Stop
```

3.3.4 Flowchart

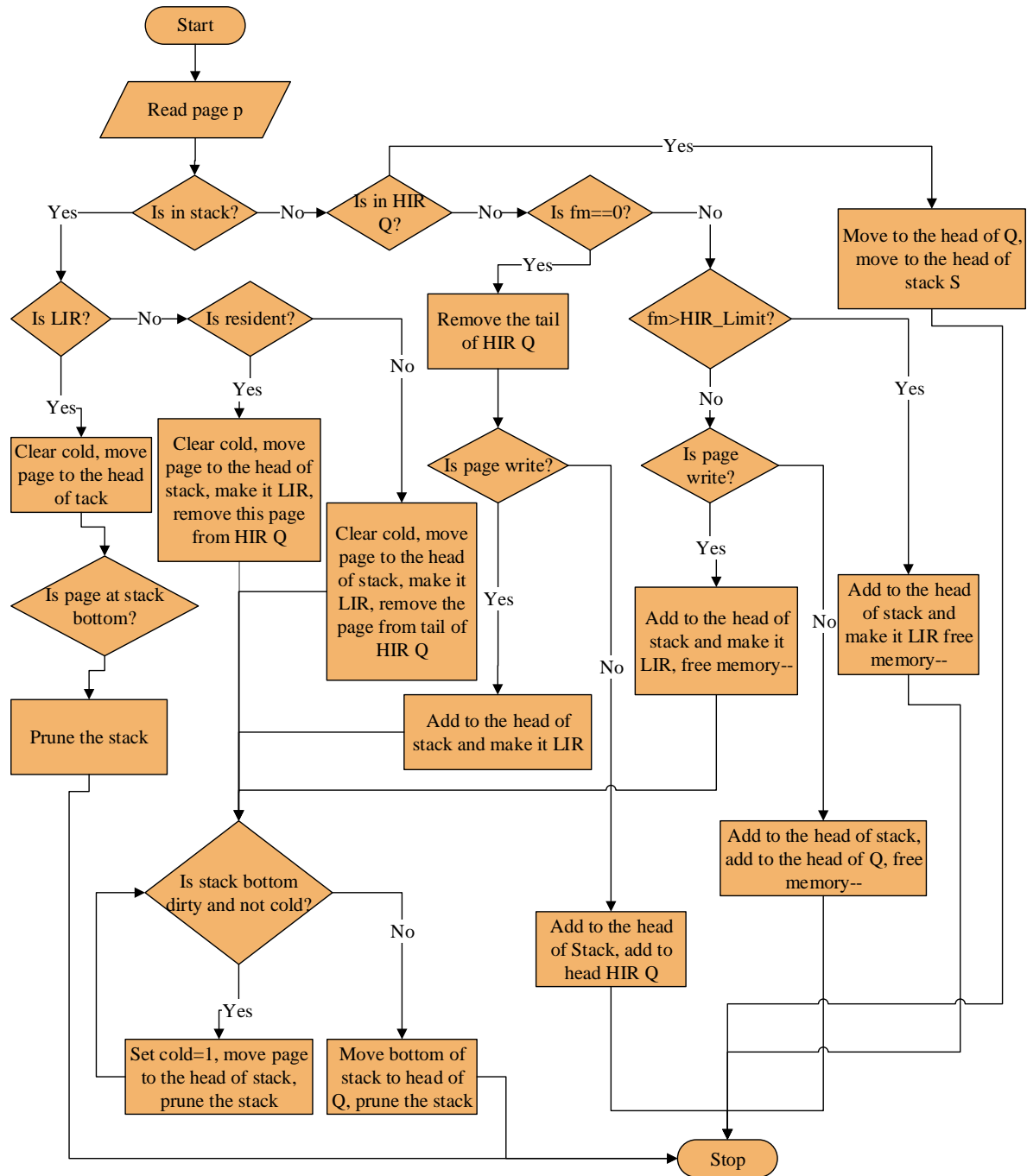


Figure 3.4: Flowchart of LIRS-WSR Algorithm

3.3.5 Tracing

Size of LIRS: 3

Size of HIRS: 1

Cache Size: 3+1=3

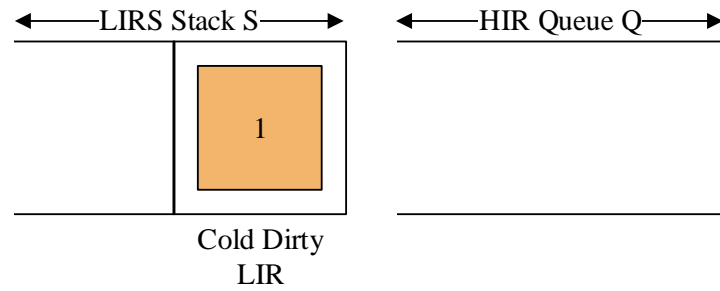
Input References: 1,1 1,2 0,3 0,1 1,4 1,3 0,5 0,2 1,3 where 1 = write, 0 = read

Total Number of References: 9, Number of Distinct References: 5

Another page status: cold/hot, dirty/clean

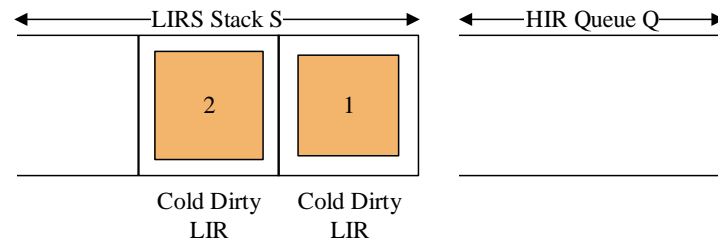


Figure 3.5: Symbols of LIRS-WSR tracing



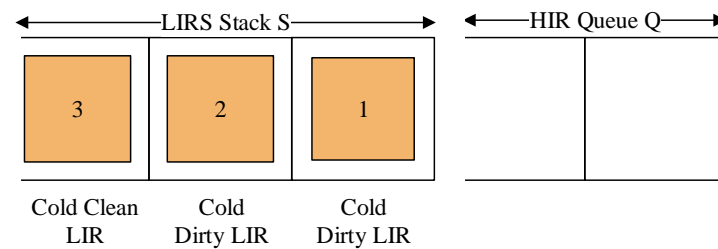
(a) State at Virtual Time 1

Accessing Page 1, 1 page miss, page fault occurred



(b) State at Virtual Time 2

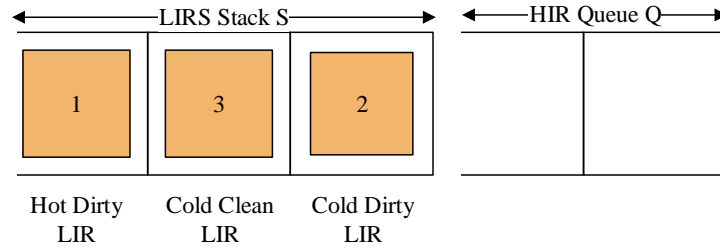
Accessing Page 1, 2 page miss, page fault occurred



(c) State at Virtual Time 3

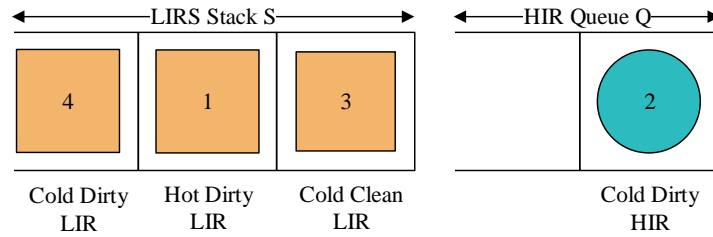
Accessing Page 0, 3 page miss, page fault occurred

Figure 3.6: Tracing of LIRS-WSR



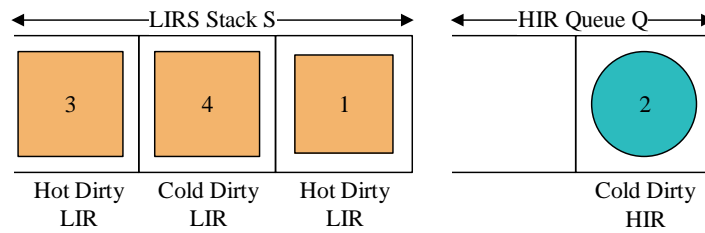
(d) State at Virtual Time 4

Accessing Page 0, 1 page hit occurred, Cleared cold



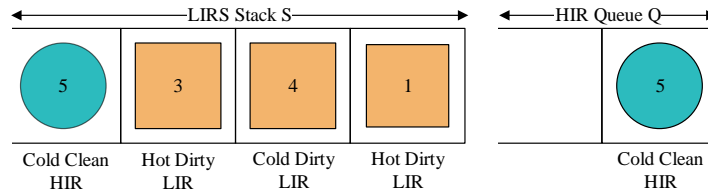
(e) State at Virtual Time 5

Accessing Page 1, 4 page miss, added to the head of the stack, 2 is demoted to HIR Q, page fault



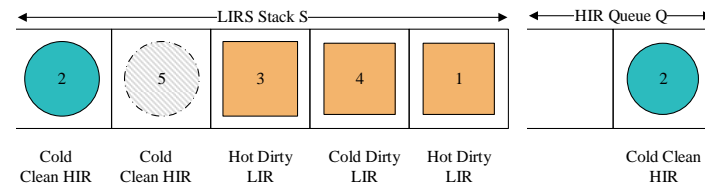
(f) State at Virtual Time 6

Accessing Page 1, 3 Page hit, moved to the head of the stack, Cleared cold



(g) State at Virtual Time 7

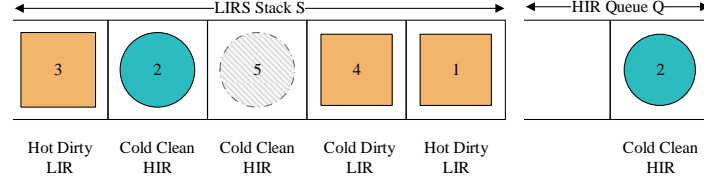
Accessing Page 0, 5 page miss, page fault, added to the head of stack and also in Q head, 2 is replaced, write count=1



(h) State at Virtual Time 8

Accessing Page 0, 2 page miss, page fault occurred, added to the head of stack and also in Q head

Figure 3.6: Tracing of LIRS-WSR (cont.)



(i) State at Virtual Time 9
Accessing Page 1, 3 page hit

Figure 3.6: Tracing of LIRS-WSR (cont.)
Total page faults = 6, Total page hits = 3 and Write counts = 1

3.4 CCF-LRU

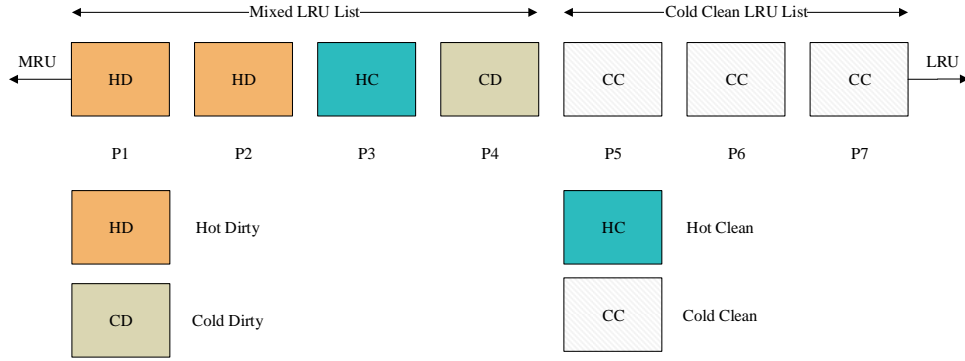


Figure 3.7: An example of CCF-LRU

CCF-LRU [26] enhances the CF-LRU and LRU-WSR algorithms by further differentiating clean pages into cold and hot ones. It uses the cold-detection mechanism of LRU-WSR to classify the buffered pages into four classes, including cold clean pages, hot clean pages, cold dirty pages, and hot dirty pages. As illustrated in Figure 3.7, the cold clean pages are stored in a cold clean LRU list while other types of pages are held in the mixed LRU list. When a replacement occurs, CCF-LRU always evicts cold clean pages first. If there are no cold clean pages, cold dirty pages are preferentially chosen as the victim and the eviction of hot dirty pages are delayed as long as possible. The relation of four classified buffer pages can be defined as follows:

$$C_{CC} < C_{CD} < C_{HC} < C_{HD} \quad (3.1)$$

Here, the C_{CC} is the cost of the evicting cold clean page, the C_{CD} is the cost of evicting cold dirty page, the C_{HC} is the cost of evicting hot clean page and the C_{HD} is the cost of evicting hot dirty page.

From the expression 3.1, we must reduce the number of evicting hot dirty page and hot clean

page as far as possible. If we want to improve the overall I/O performance for flash memory based storage system. However, the LRU-WSR keeps cold clean pages in the buffer until they are evicted, which not only lowers the hit ratio but also makes cold dirty page and hot clean page evicted more quickly. Besides, it always evicts clean pages without considering their access frequencies, which may lower the hit ratio. The main idea of the Cold Clean First (CCF) strategy is described as follows:

1. It uses the cold-detection algorithm to judge whether the page is cold or hot;
2. It evicts cold clean page preferentially as possible, especially for clean page accessed only once recently; and
3. If there is no cold clean page mentioned above, evicting cold dirty page instead of the hot clean page.

3.4.1 Data Structure

The CCF-LRU algorithm can be implemented by using two LRU lists, Mixed LRU and Cold Clean LRU list. These both lists have MRU and LRU and these both lists are implemented as list interface.

3.4.2 Flowchart

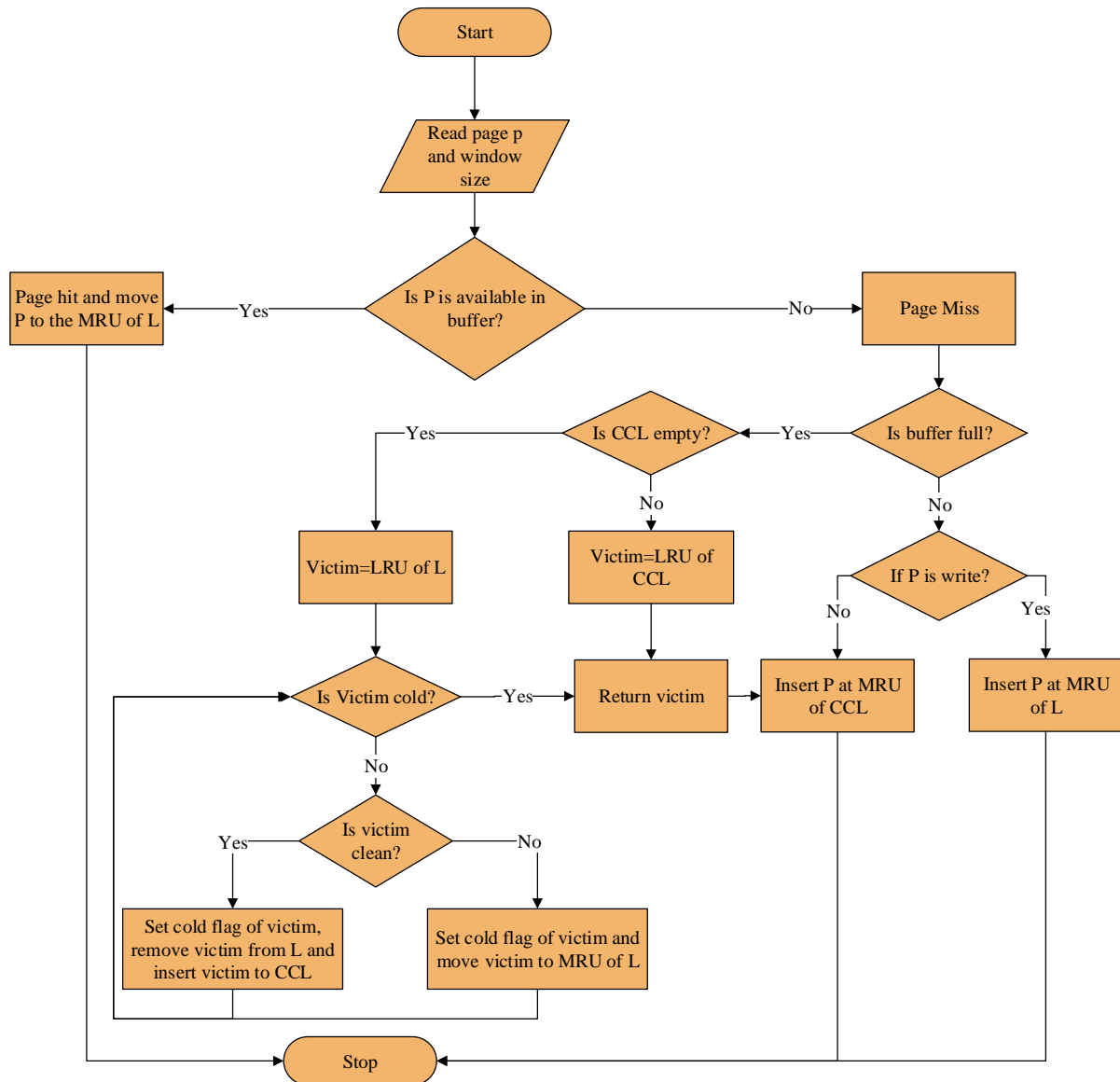


Figure 3.8: Flowchart of CCF-LRU Algorithm

3.4.3 Tracing

Size of L: 2

Size of CCL: 2

Cache Size: 2+2=4

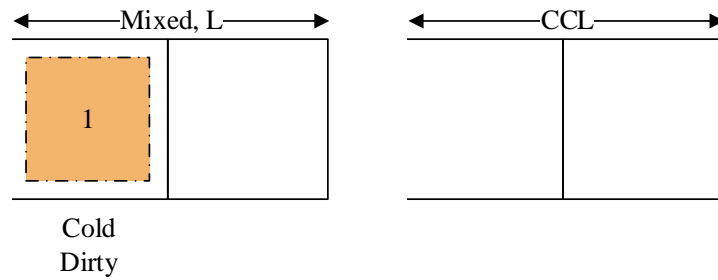
Input References: 1,1 1,2 0,3 0,1 1,4 1,3 0,5 0,2 1,3 where 1 = write, 0 = read

Total Number of References: 9, Number of Distinct References: 5

Another page status: cold/hot, dirty/clean

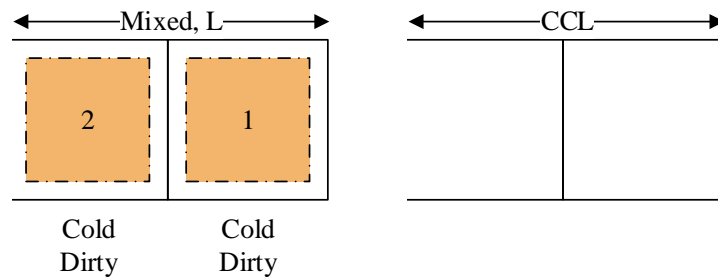


Figure 3.9: Symbols of CCF-LRU tracing



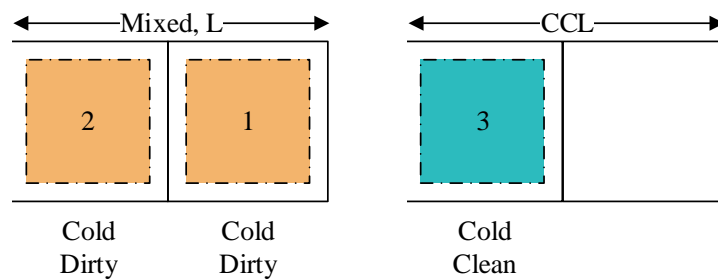
(a) State at Virtual Time 1

Accessing Page 1, 1 page miss, page fault occurred, inserted at MRU of Mixed



(b) State at Virtual Time 2

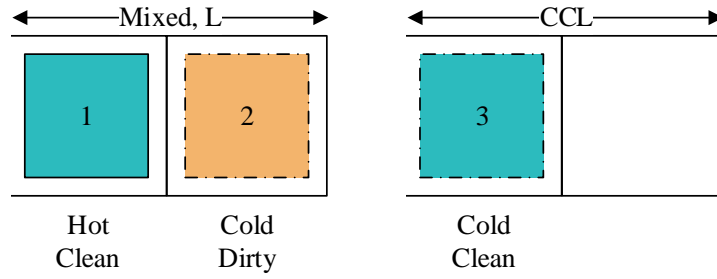
Accessing Page 1, 2 page miss, page fault occurred, inserted at MRU of Mixed



(c) State at Virtual Time 3

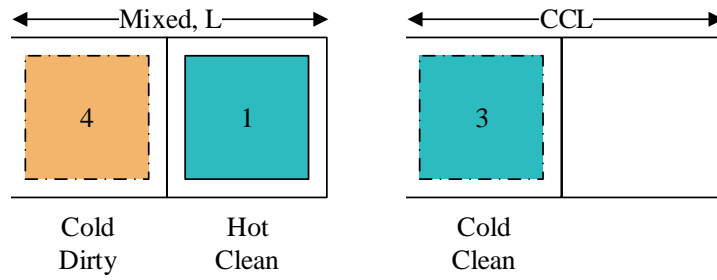
Accessing Page 0, 3 page miss, page fault occurred, inserted at MRU of CCL

Figure 3.10: Tracing of CCF-LRU



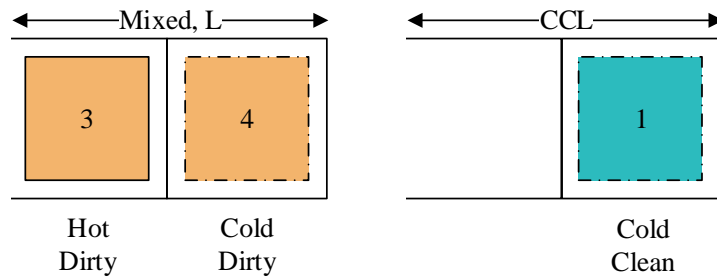
(d) State at Virtual Time 4

Accessing Page 0, 1 page hit in Mixed, Moved to MRU, Cleared cold



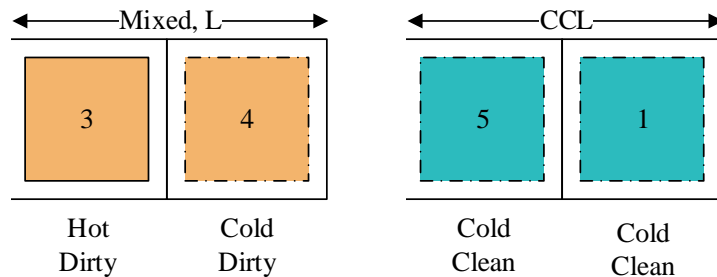
(e) State at Virtual Time 5

Accessing Page 1, 4 page miss, 4 is added to the MRU of the Mixed, 2 is replaced, write count=1, page fault occurred



(f) State at Virtual Time 6

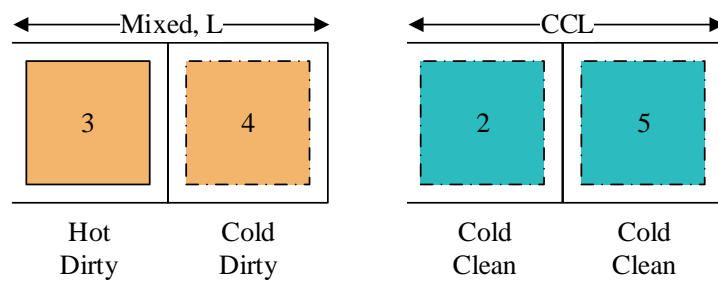
Accessing Page 1, 3 Page hit in CCL, Cleared cold and moved to MRU of L



(g) State at Virtual Time 7

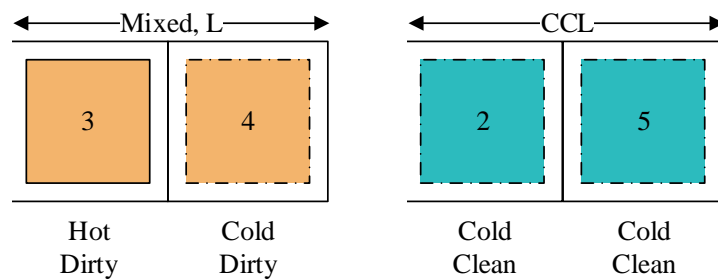
Accessing Page 0, 5 page miss, page fault, inserted at MRU of CCL

Figure 3.10: Tracing of CCF-LRU (cont.)



(h) State at Virtual Time 8

Accessing Page 0, 2 page miss, page fault occurred, , inserted at MRU of CCL



(i) State at Virtual Time 9

Accessing Page 1, 3 page hit in Mixed

Figure 3.10: Tracing of CCF-LRU (cont.)

Total page faults = 6, Total page hits = 3 and Write counts = 1

3.4.4 Algorithm

Algorithm 3.2 CCF-LRU

INPUT: Pages

OUTPUT: Miss rate, hit rate etc.

```
1: Start
2: Read new page
3: if Page is in CCL then
4:   Page hit in CCL
5:   if Page is write then
6:     Increase Write Request by 1
7:     Set clean flag false i.e. clean=0
8:   end if
9:   Clear cold flag, i.e. cold=0
10:  Move the page to the MRU of L
11: else if Page is in L then
12:   Page hit in L
13:   if Page is write then
14:     Increase Write Request by 1
15:     Set clean flag false i.e. clean=0
16:   else
17:     Set clean flag true i.e. clean=1
18:   end if
19:   Clear cold flag of page, i.e. cold=0
20:   Move page to the MRU position in L
21: else
22:   Page miss occurs
23:   if Page is read then
24:     if Free memory available in CCL then
25:       Insert page to MRU of CCL
26:       Set cold flag true i.e. cold=1
27:       Set clean flag true i.e. clean=1
28:     else
29:       Remove page from LRU of CCL
30:       Insert page to MRU of CCL
31:       Set cold flag true i.e. cold=1
32:       Set clean flag true i.e. clean=1
33:     end if
34:   else
35:     Increase Write Request by 1
36:     if Free memory available in L then
37:       Insert page to MRU of L
38:       Set cold flag i.e. cold=1
39:       Set clean flag false i.e. clean=0
40:     else
41:       for Each page from Tail to head of L do
42:         Victim=the LRU page in L
```

Algorithm 3.2 CCF-LRU (cont.)

```
43:         if Cold flag of victim is set then
44:             | Return victim
45:         else if Victim is clean page then
46:             | Set the cold flag of victim
47:             | Remove victim from L
48:             | Insert victim into CCL's MRU position
49:         else
50:             | Set cold flag of victim
51:             | Move victim to the MRU position in L
52:             | CALL AGAIN CCF-LRU
53:         end if
54:     end for
55:     if If victim is not clean then
56:         | Increase Write Count by 1
57:     end if
58:     Remove victim
59:     Insert page into MRU of L
60:     Set cold flag i.e. cold=1
61:     Set clean flag false i.e. clean=0
62: end if
63: end if
64: end if
65: Stop
```

Chapter 4

RESULTS, ANALYSIS AND COMPARISONS

4.1 Testing

Each workload is tested in LIRS-WSR and CCF-LRU simulator by varying the cache size from 512 to 18432. In the case of LIRS-WSR algorithms HIR, LIR partition is maintained as 1% and 99% of cache size. For CCF-LRU parameter WindowsSize is set 0.5 of the cache size for all Workloads.

4.1.1 Test Result of Workload 1 (Trace with Random Access)

No. of References = 100000, No. of Distinct Reference = 43247

No. of Write References = 49974

Table 4.1: Test Result of Workload 1

Buffer Size	LIRS-WSR				CCF-LRU			
	Page Fault	Miss Rate	Hit Rate	Write Count	Page Fault	Miss Rate	Hit Rate	Write Count
512	98955	98.2	1.8	48950	99011	98.3	1.7	49248
1024	97926	96.3	3.7	47959	97939	96.4	3.6	48450
2048	95847	92.7	7.3	46008	95893	92.8	7.2	46987
4096	91892	85.7	14.3	42538	91880	86.7	14.3	44273
6144	87951	78.8	21.2	39344	87931	78.7	21.3	41700
8192	84103	72.0	28.0	36482	84236	72.2	27.8	39439
9216	82253	68.7	31.3	35180	82440	69.1	30.9	38352
10240	80421	65.5	34.5	33950	80628	65.9	34.1	37263
12288	76800	59.1	40.9	31694	77083	59.6	40.4	35176
14336	73358	53.1	46.9	29706	73771	53.8	46.2	33178
16384	70076	47.3	52.7	27996	70666	48.3	51.7	31405
18432	66834	41.6	58.4	26446	67581	42.9	57.1	29604

4.1.2 Test Result of Workload 2 (Trace with Read-Most Access)

No. of References = 100000, No. of Distinct Reference = 43212

No. of Write References = 9919

Table 4.2: Test Result of Workload 2

Buffer Size	LIRS-WSR				CCF-LRU			
	Page Fault	Miss Rate	Hit Rate	Write Count	Page Fault	Miss Rate	Hit Rate	Write Count
512	98948	98.1	1.9	9078	98956	98.2	1.8	9197
1024	97954	96.4	3.6	8505	97907	96.3	3.7	8620
2048	95943	92.9	7.1	7680	95961	92.9	7.1	7701
4096	91877	85.7	14.3	6547	92169	86.2	13.8	6457
6144	87870	78.6	21.4	5832	88399	79.6	20.4	5599
8192	84144	72.1	27.9	5339	84936	73.5	26.5	5021
9216	82330	68.9	31.1	5135	83147	70.3	29.7	4754
10240	80491	65.6	34.4	4973	81429	67.3	32.7	4539
12288	76953	59.4	40.6	4706	78165	61.5	38.5	4138
14336	73402	53.2	46.8	4527	75029	56.0	44.0	3803
16384	70137	47.4	52.6	4403	71930	50.6	49.4	3525
18432	68828	41.6	58.4	4334	68975	45.4	54.6	3284

4.1.3 Test Result of Workload 3 (Trace with Write-Most Access)

No. of References = 100000, No. of Distinct Reference = 43182

No. of Write References = 89145

Table 4.3: Test Result of Workload 3

Buffer Size	LIRS-WSR				CCF-LRU			
	Page Fault	Miss Rate	Hit Rate	Write Count	Page Fault	Miss Rate	Hit Rate	Write Count
512	98933	98.1	1.9	88084	98960	98.2	1.8	88356
1024	97859	96.2	3.8	87018	97936	96.4	3.6	87574
2048	95831	92.7	7.3	84999	95954	92.9	7.1	86014
4096	91903	85.7	14.3	81111	92218	86.3	13.7	82995
6144	88066	79.0	21.0	77380	88825	80.3	19.7	80043
8192	84211	72.2	27.8	73621	86050	75.4	24.6	77176
9216	82390	69.0	31.0	71869	84939	73.5	26.5	75730
10240	80578	65.8	34.2	70126	83978	71.8	28.2	74265
12288	77008	59.5	40.5	66720	82126	68.5	31.5	71420
14336	73484	53.3	46.7	63398	80381	65.5	34.5	68677
16384	70197	47.5	52.5	60300	78573	62.3	37.7	65882
18432	66856	41.7	58.3	57298	76801	59.2	40.8	63125

4.1.4 Test Result of Workload 4 (Zipf Trace)

No. of References = 500000, No. of Distinct References = 47023

No. of Write References = 244790

Table 4.4: Test Result of Workload 4

Buffer Size	LIRS-WSR				CCF-LRU			
	Page Fault	Miss Rate	Hit Rate	Write Count	Page Fault	Miss Rate	Hit Rate	Write Count
512	361423	69.4	30.6	163145	356228	68.3	31.7	176669
1024	327682	62.0	38.0	144910	321647	60.6	39.4	160260
2048	288224	53.2	46.8	123920	282157	52.9	48.1	141871
4096	242417	43.1	56.9	100009	237893	42.1	57.9	121455
6144	212645	36.6	63.4	84817	209385	35.9	64.1	108352
8192	189960	31.6	68.4	73661	187861	31.1	68.9	98409
9216	180187	29.4	70.6	68958	178954	29.1	70.9	94140
10240	171186	27.4	72.6	64718	170853	27.3	72.7	90437
12288	155797	24.0	76.0	57657	156432	24.1	75.9	83595
14336	142185	21.0	79.0	51627	143880	21.4	78.6	77632
16384	130416	18.4	81.6	46630	132570	18.9	81.1	72044
18432	119561	16.0	84.0	42206	122647	16.7	83.3	67197

4.2 Analysis and Comparisons

All the results obtained from simulation is analyzed by drawing different graphs. Hit rate and write count are used as criteria for analyzing their goodness.

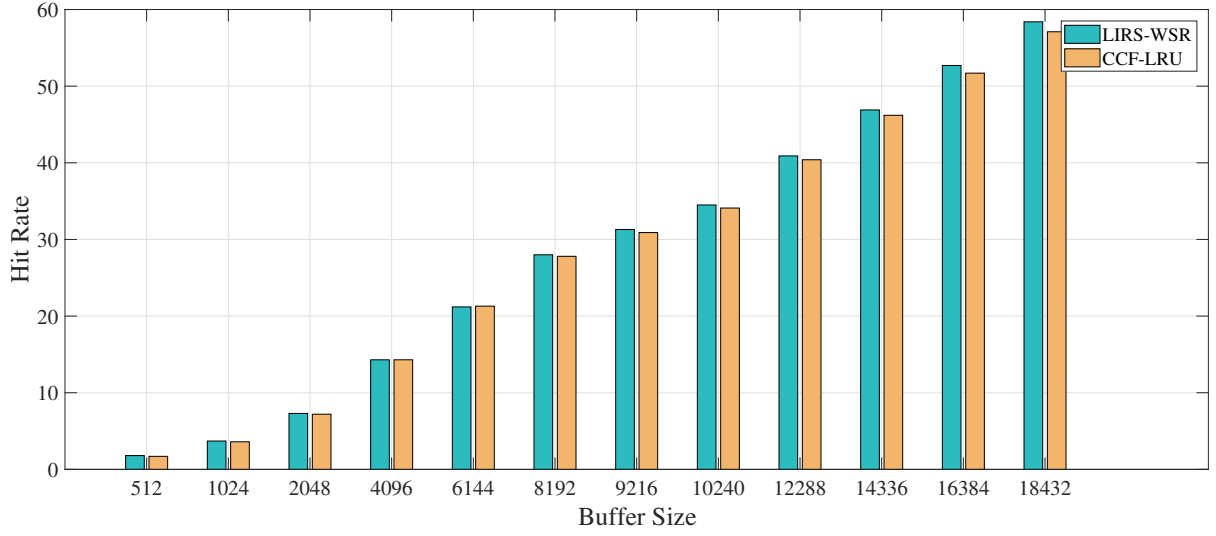


Figure 4.1: Graph of Hit Rate for Workload 1

4.2.1 Hit Rate Analysis

The graph of Figures 4.1, 4.2, 4.3, and 4.4 show that the LIRS-WSR algorithm is better than the CCF-LRU algorithm since LIRS-WSR has higher hit rate for different cache sizes. In Figure 4.1 and Figure 4.2 for workloads random and read-most access type traces the hit rates of algorithms are not much different because of the uniform distribution of page references. In

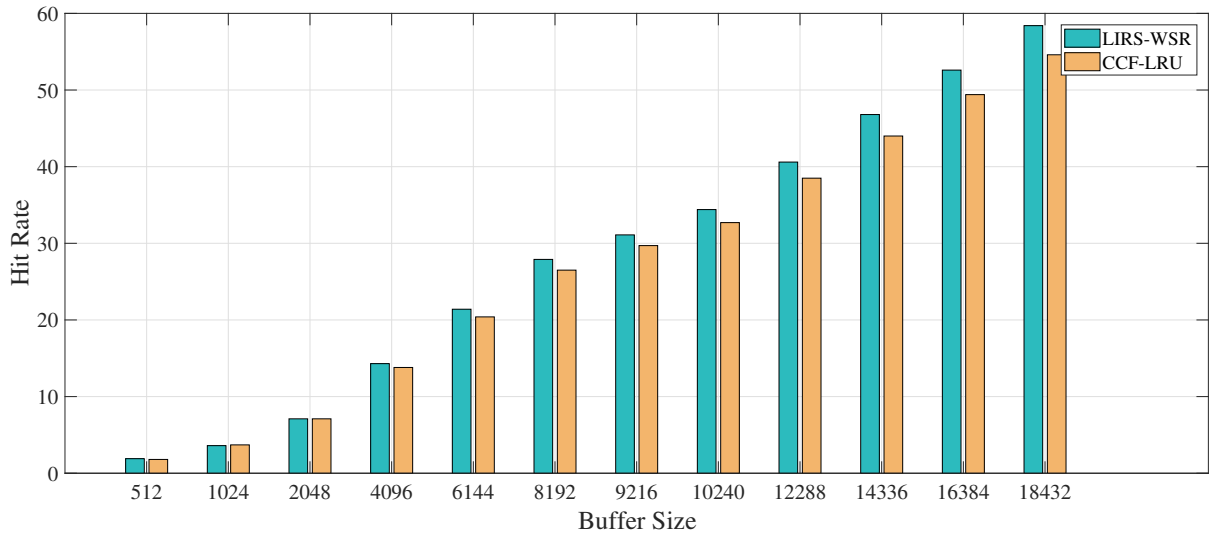


Figure 4.2: Graph of Hit Rate for Workload 2

these workloads, there is not a clear distinction between hot and cold pages as reference locality is not high. Despite the nature of page references in these workloads, LIRS-WSR has better hit rate as it adapts the changes in page references pattern dynamically. LIRS-WSR always treats pages with the write request as LIR pages or hot pages. The hit rate increases with large

buffer size because buffer can hold more pages that increase page hit. The graph of Figure

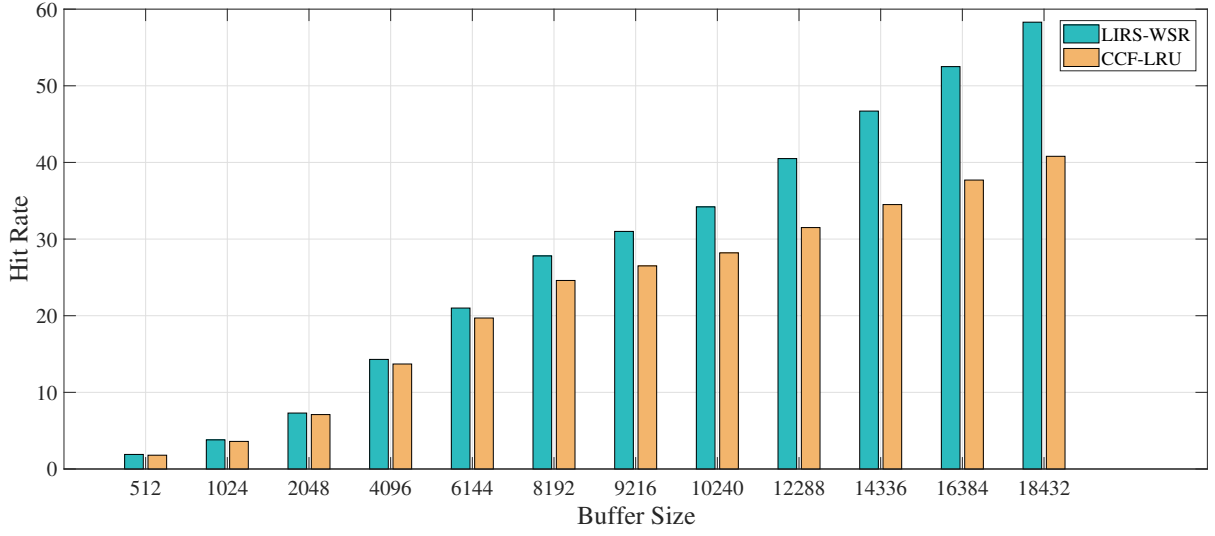


Figure 4.3: Graph of Hit Rate for Workload 3

4.3 shows a significant difference in hit rate. This is due to high reference locality of page references in write-most trace where 90% of pages are referenced as write. LIRS-WSR adapts more effectively to distinguish hot and cold pages in this workload. So, the hit rate is much higher than that of CCF-LRU. As a maximum value, LIRS-WSR has up to 17.5% higher hit

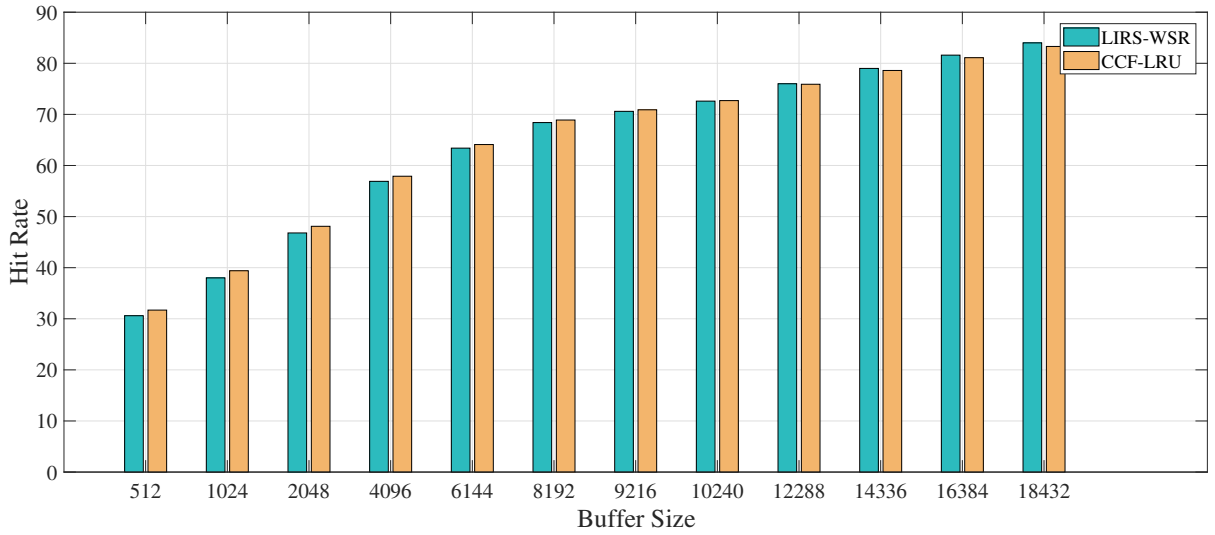


Figure 4.4: Graph of Hit Rate for Workload 4

rate than CCF-LRU algorithm for workload 3 (write-most trace) and in workload 4, the hit rate of LIRS-WSR increases for larger buffer. The hit rate of CCF-LRU is larger when buffer size is smaller. As minimum, LIRS-WSR has upto 0.7% higher hit rate than CCF-LRU for larger buffer. LIRS-WSR has higher hit rate than CCF-LRU in the case of high reference locality workloads.

4.2.2 Write Count Analysis

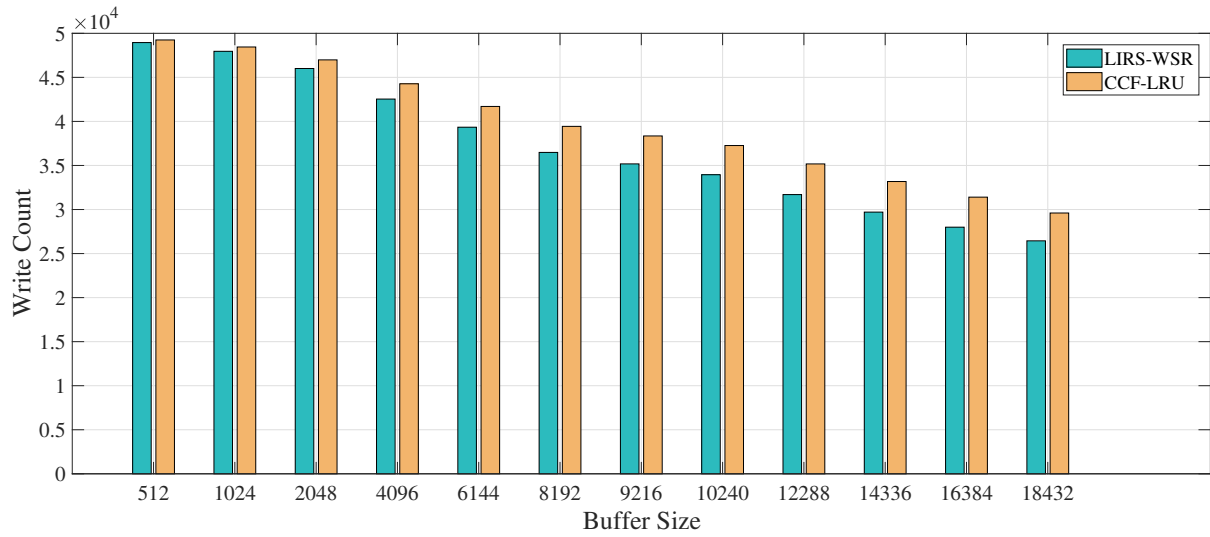


Figure 4.5: Graph of Write Count for Workload 1

The graph in Figures 4.5, 4.6, 4.7, and 4.8 show the number of pages propagated to flash memory. The number of pages flushed to flash memory is write count. The number is obtained by counting the eviction of page references with write request during page replacement event. At the end of simulation all the dirty pages in the buffer are also added to the count to get the exact write count. From the above graphs, it is clear that LIRS-WSR has smaller write count for the three workloads(random, write-most and zipf) used in the simulation. Workload 1 has 50%

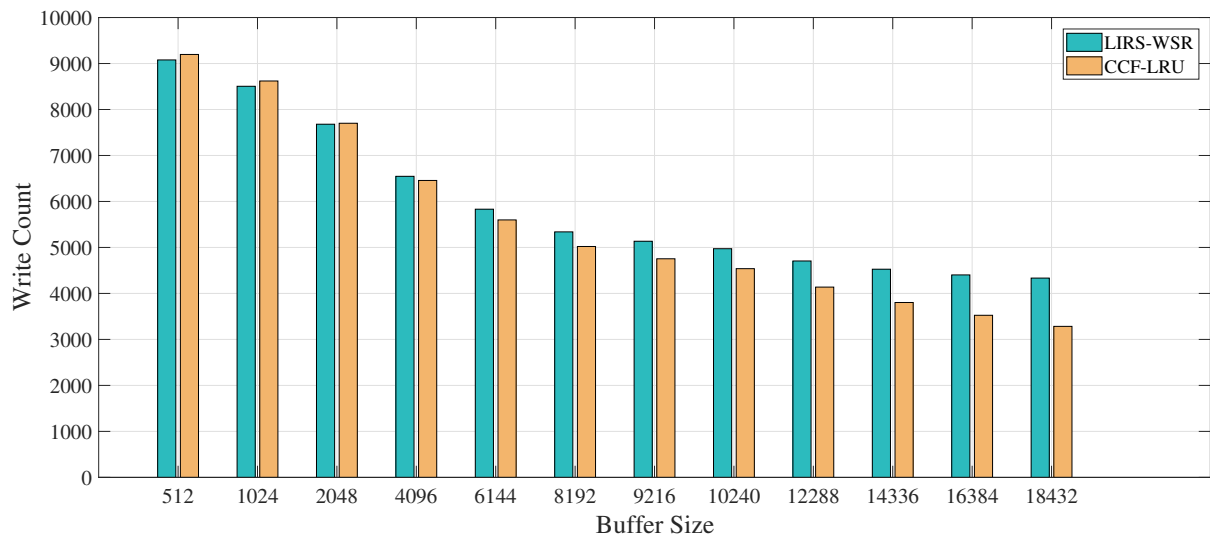


Figure 4.6: Graph of Write Count for Workload 2

pages as write pages which are uniformly distributed. For smaller buffer size, CCF-LRU has slightly higher write count than LIRS-WSR. This is because there is only half buffer (Mixed

LRU) list holds write pages, dirty pages are evicted faster in CCF-LRU than LIRS-WSR. In LIRS-WSR dirty pages are kept in LIRS Stack as LIR pages and due to WSR policy, they are kept for longer time. Workload 2 has read-most access pattern in page references with 10%

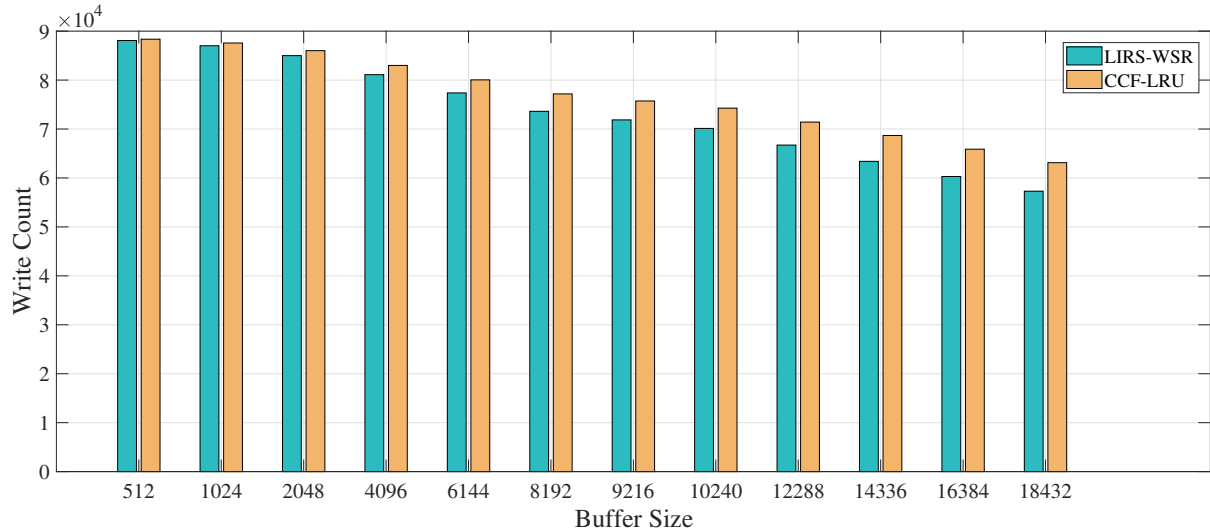


Figure 4.7: Graph of Write Count for Workload 3

writes and 90% reads. For workload 2, for smaller buffer, both algorithm has slightly difference in write count and LIRS-WSR outperforms. As buffer increases, CCF-LRU outperforms LIRS-WSR for read-most trace. Workload 3 has 90% writes and only 10% reads in page ref-

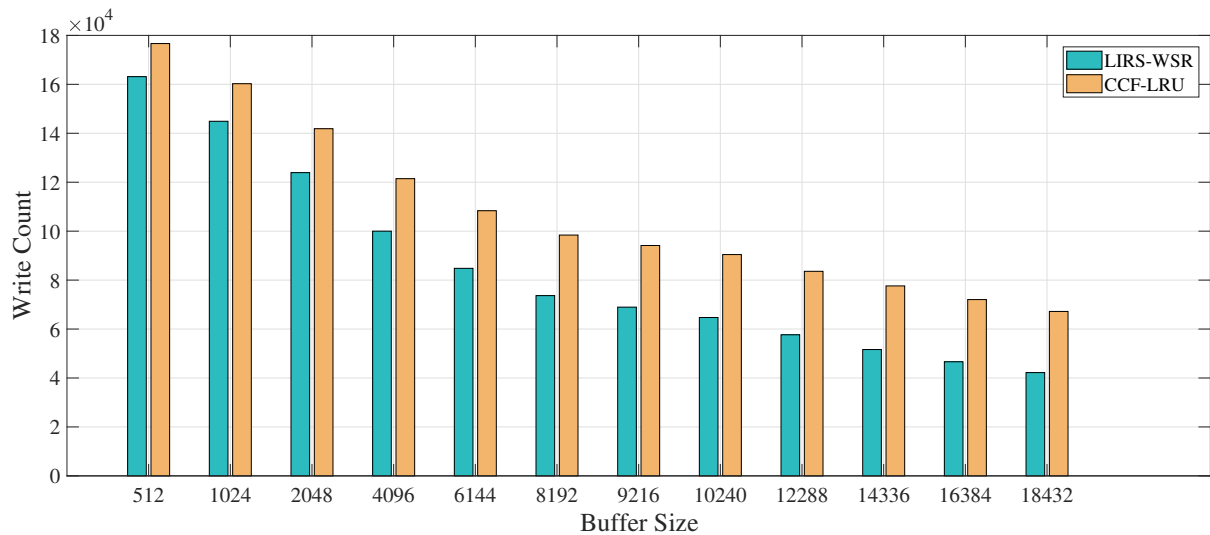


Figure 4.8: Graph of Write Count for Workload 4

erences. In this case, LIRS-WSR works better for all buffer sizes. There is a large number of write pages in trace, CCF-LRU cannot accommodate all dirty pages in Mixed LIRS set and evicts dirty pages as well continuously. LIRS-WSR adapts reference pattern so it has better output.

Workload 4 zipf trace has 50%/50% read and write references but has high reference locality of 80% page references are references to 20% of pages. LIRS-WSR adapts changes in reference pattern and locality. So it has much less write count than CCF-LRU for all buffer sizes.

For write-most trace (Workload 3) write count is decreased up to 9.23% by LIRS-WSR with a comparison to CCF-LRU algorithm.

For zipf trace (workload 4) LIRS-WSR decreased write count up to 37.17% with a comparison to CCF-LRU algorithm. This is the largest gap between the value of write count of these two algorithms. Thus, LIRS-WSR minimizes write counts significantly when reference locality is high.

Chapter 5

CONCLUSION

5.1 Conclusion

Flash memory has become an alternative to the magnetic disks, which brings new challenges to the traditional disk-based system. To efficiently support the characteristics of flash storage devices, traditional buffering approaches need to be revised to take into account the imbalance I/O property of flash memory. LIRS-WSR uses delayed eviction strategy when the dirty page is to be replaced. It uses recency and little bit frequency information in replacement policy. CCF-LRU captures frequency and recency of page references by using two LRU queues to classify all the buffer pages into a Mixed and a cold clean set.

From the simulation of these two algorithms for varying buffer size, it is found that the LIRS-WSR outperforms CCF-LRU for hit rate and write count. Especially, when workload has high reference locality, LIRS-WSR has significantly superior performance than CCF-LRU in terms of both hit rate and write count. This is because of LIRS-WSR's good adaptive technique to handle changes in reference patterns. But in case of workload 2 (read-most trace) CCF-LRU performs better than LIRS-WSR for write counts.

For uniformly distributed workloads, the difference in hit rates of CCF-LRU and LIRS-WSR is comparatively small. LIRS-WSR leads CCF-LRU in hit rate by a value up to 3.8%. For high reference locality workloads LIRS-WSR has significantly higher hit rate up to 17.5% in comparison to CCF-LRU.

The CCF-LRU may perform better in write count for uniformly distributed locality workloads when the buffer size is highly smaller in comparison to the size of workload as it treats all

write pages as hot pages and delays eviction. For larger buffer size and high write reference locality workloads, LIRS-WSR outperforms CCF-LRU in write count for workloads random, write-most and zipf. It seems that in the case of uniformly distributed write-most access type workload, write count is decreased up to 9.23% by LIRS-WSR, but for high reference locality workload LIRS-WSR minimizes write count up to 37.17% with a comparison to CCF-LRU algorithm. Thus, LIRS-WSR minimizes write counts significantly when reference locality is high.

5.2 Limitations

In this work, the size of HIR block is chosen 1% of total buffer size in LIRS algorithm and in CCF-LRU algorithm Mixed LRU list and Cold Clean LRU list are divided into two half of buffer size. The dynamic approach can be used to self-tune these parameters. Further research can be done to find the optimal value of these parameters for different workloads. In addition to this, in this work, only four different memory traces have been used for simulation purpose. Three of which are of uniform reference patterns and last one is with “20/80” reference locality. These are the limitations of this work. This work can be further extended by using a variety of real memory trace with different reference locality.

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Appendix A

Sample Input Traces

A.1 Random Input Trace

1,8575 0,17754 0,33289 0,3838 0,19942 1,25113 1,35145 1,1939 0,40780 0,12831 0,31724
1,37162 1,861 1,35912 0,39216 1,10863 0,15454 0,32425 0,42141 1,34769 0,29923 1,3050
0,4043 0,39113 1,11686 1,25837 0,4941 0,7882 0,39262 1,32631 0,36490 1,11934 1,8851
1,16962 0,37665 1,23980 0,41727 0,15074 1,19029 1,1750 0,49554 1,18797 1,6747 0,31276
1,786 1,42798 0,30971 0,42594 1,49503 0,23075 1,8717 1,13521 1,988 0,22467 1,12586 1,45284
1,39329 0,45058 0,14795 0,21120 0,7786 1,43211 0,47655 0,42213 0,919 0,603 0,4844 0,44923
1,29324 0,26292 1,31526 1,38097 1,39819 0,30117 1,14208 1,27844 1,8361 1,16455 1,5699
0,10670 1,1066 0,9039 1,6477 1,41170 0,23504 1,32354 0,14280 1,36795 1,8732 1,46002
1,4880 1,5637 1,21680 1,3496 1,3220 0,13282 1,42670 0,11669 0,2716 1,49749 0,12437 0,42550
0,27038 0,26790 1,44095 1,25674 0,4498 1,32206 0,33123 0,9846 0,46190 0,20089 0,14060
1,28875 1,16434 1,12575 1,47687 1,41433 0,16610 0,3411 1,23633 1,17429 0,49681 1,25625
1,34155 0,33804 0,21089 0,16647 1,3104 0,3843 1,7142 0,30193 0,12695 1,28453 0,9115
0,25532 0,47722 1,47868 1,49752 0,6476 1,41825 1,7631 0,14127 0,29127 1,12805 0,48855
1,33911 0,41079 1,25483 1,39430 0,1037 0,3297 0,16599 1,36036 0,15578 0,10091 1,25578
0,23037 0,24073 1,16386 0,15490 1,1048 0,19682 0,8798 0,26493 0,48889 1,7791 1,35987
1,16638 1,45825 1,38057 0,30566 0,48228 0,38949 1,47502 0,26137 1,22920 1,32430 1,7944
1,35589 1,40867 1,47773 0,46838 1,44616 1,39286 1,39175 1,20480 1,22293 0,20389 0,23900
0,18555 0,46427 0,8516 1,49886 0,10679 0,9400 1,24467 0,3709 0,411 0,25540 0,22153 1,3954
0,23179 0,9759 1,33020 0,2711 1,42697 1,34063 1,22716 1,23599 0,25436 1,22036 ...

A.2 Read-most Input Trace

0,47138 0,8885 0,46509 0,30725 1,15160 0,2460 0,9807 0,46791 1,5087 0,11237 0,22932
0,37902 0,6713 0,34922 0,4119 0,42689 0,25737 0,39402 0,9355 0,10606 0,641 0,27320 0,38193
0,21972 0,42518 0,10783 0,28314 1,1900 0,13867 0,39219 0,46605 1,38017 0,46494 0,23527
0,38630 1,21176 0,293 0,12907 0,39277 0,40610 0,7266 1,41366 0,30769 1,8749 1,10029
0,1320 0,46614 0,41918 0,26128 0,41673 0,19547 0,48693 0,37972 0,38947 0,15954 0,3438
0,18472 0,16481 0,6566 0,9291 0,43502 1,33032 0,3183 0,19948 0,6053 1,38512 0,46694
0,33131 0,29974 0,19584 0,49468 1,24278 0,17376 0,46130 0,4161 0,3133 0,45468 0,35567
0,36470 0,24196 1,34021 0,39449 0,18771 0,19982 0,26021 0,17350 0,44669 0,11232 0,2877
0,14913 0,26197 0,37578 0,44932 0,27057 0,8577 0,21545 1,19614 0,26010 0,31719 0,21978
0,9246 0,32690 0,35125 0,29523 0,34981 1,3135 1,2971 0,1054 0,15836 0,29720 0,39483
0,42668 0,23341 1,7058 1,37083 0,5836 0,39234 0,30664 0,47423 0,48384 0,49832 0,47732
0,6181 0,28049 0,20673 0,14815 1,16584 0,35416 0,15178 1,22743 0,37824 0,20809 0,43815
0,7992 1,22767 1,981 0,6349 0,22302 0,1909 0,37810 0,24271 0,27349 0,21940 0,11289 0,3186
0,14000 0,38546 1,20359 0,34039 0,3939 0,3492 0,44098 0,2151 0,17422 0,30562 1,24662
0,23074 0,26344 1,31895 0,6416 0,48410 0,15522 0,14390 0,34163 0,13073 0,19750 0,985
0,48011 1,18012 0,11608 0,14481 0,34997 0,22648 0,26672 0,15980 0,49335 1,34079 0,11814
0,31534 0,20259 0,11874 0,45185 1,20792 0,39186 0,18681 0,24097 0,8582 0,26107 1,11335
0,33248 0,31662 0,47539 0,2856 0,41237 0,19933 0,10902 1,6574 0,14599 0,39656 0,15879
0,9645 0,32760 0,11311 0,14258 0,38921 0,47086 0,24615 0,36799 1,23373 0,30556 0,6997
0,5647 0,22385 0,14890 0,5537 0,11311 0,18829 0,9608 0,44776 0,35106 0,21597 0,18245
0,25921 0,19819 0,41022 0,2924 0,33953 0,9818 0,21029 0,1955 0,26130 0,48683 0,16144
0,42243 0,1071 0,48155 0,17289 0,24699 1,19033 0,18424 0,39192 0,45975 0,949 0,25811
0,35775 0,28294 0,7946 0,2748 0,36907 0,5078 0,27022 1,14669 0,37419 0,12382 0,8955
0,43073 0,4139 0,37292 0,31386 0,15131 0,44501 0,40518 0,6139 0,49892 0,22521 0,9057
1,43638 0,45879 0,30391 0,14690 0,25367 0,10125 0,24894 0,41810 0,49555 0,38776 0,16140
0,49637 0,36102 0,13534 0,4838 1,33623 0,19639 0,33611 1,38969 0,34042 0,32887 0,13925
0,12100 0,10997 0,8528 0,11794 0,23601 0,15213 0,29736 0,47737 0,15336 0,16109 0,10809
0,36945 0,49102 0,40775 1,2132 0,12292 1,28002 0,29787 0,12657 0,27496 0,11586 0,35950
0,19189 0,7309 0,16707 0,45708 0,43469 0,32897 0,21864 0,18648 0,36112 0,29233 0,18148
0,37425 0,21023 0,8947 0,30022 0,43937 0,18352 0,32213 0,26617 0,6472 0,9465 ...

A.3 Write-most Input Trace

1,12527 1,1216 1,698 1,35286 1,39722 1,25887 1,45028 1,47558 1,44966 1,10018 1,41052
0,8011 1,42731 1,9714 1,39263 1,40196 1,6269 1,39623 1,33031 1,1853 1,29107 1,5242 1,1010
1,28122 1,35606 1,27792 1,19845 1,24155 1,20899 1,37819 1,27592 1,1272 1,2536 1,35733
1,33645 1,37360 1,13287 1,35073 1,24973 1,31865 1,7424 1,5993 1,8751 1,2237 1,39556
1,8440 1,35811 1,25015 1,42880 1,12603 1,8230 1,45262 1,10924 0,40802 1,24112 1,38237
1,31304 1,5412 1,43801 1,29898 1,10638 1,47683 1,4487 1,44810 1,8571 1,9911 1,33896
0,35169 1,35950 1,9344 1,2859 1,32483 1,2158 1,46525 1,32777 1,20380 0,25035 1,5188
1,6797 0,24879 1,8889 1,19975 1,8644 1,8494 1,17945 1,5175 1,29078 1,36322 1,46605 0,3722
1,23254 1,35573 1,44707 1,16353 1,23944 1,24724 0,40235 1,9453 1,33001 1,23185 1,19468
1,4818 1,18662 1,14189 0,1378 1,16011 1,18092 1,36090 1,37183 1,4364 1,33538 1,41008
1,19253 1,34763 1,21453 1,5052 1,38178 1,39783 1,33887 1,46310 1,2396 1,41563 1,18490
1,18554 1,46076 1,3812 1,46712 1,22442 0,15937 1,38230 1,45473 1,6945 1,24479 1,9632
1,21724 1,12421 1,20451 1,35388 0,980 1,4486 1,47436 1,44968 1,42560 1,34505 1,42484
1,8868 1,13237 1,45460 1,40381 1,46871 1,18937 1,1389 1,22092 1,20688 1,30869 0,45818
1,47306 1,3497 1,1803 1,6096 1,24012 1,43783 1,7630 1,24744 1,47367 1,42187 1,43951
1,21302 1,26076 1,12092 0,38106 1,21666 1,45645 1,12638 1,5712 0,14779 1,33647 1,29306
1,20191 1,33315 1,26443 1,11996 1,28139 1,18374 1,24340 1,26206 1,6606 1,1590 1,16723
1,48509 1,29078 1,36414 1,5498 0,24528 1,43092 1,11633 1,27217 1,10035 1,5380 1,2269
1,41075 1,7928 1,8105 1,3437 1,22547 1,45582 0,8817 1,38670 1,20172 1,30414 1,47214
1,19627 1,26446 1,40787 1,39687 1,3454 1,37369 1,30931 1,33101 1,18169 1,22790 1,11904
0,47052 1,3672 1,42585 1,9384 1,5275 1,13720 1,19348 1,49136 1,20843 1,19068 1,25883
1,16481 1,27189 0,29307 1,16008 1,45273 0,9839 1,38955 1,48500 1,48560 1,47897 1,37830
1,39217 1,9133 1,18904 0,10499 1,48972 0,42043 1,45152 1,1636 1,12524 1,39143 1,37057
1,9006 1,47238 1,45840 1,5534 1,45368 0,28865 1,6060 1,41228 1,31789 1,48175 1,22391
1,23196 1,34069 0,27033 1,11358 1,21846 1,38558 1,36046 1,4791 1,26938 1,20824 1,4823
1,48716 1,44135 1,28505 1,49252 1,44939 1,36081 1,29232 0,30656 0,47723 1,48222 1,35146
1,878 1,18288 1,8098 1,31077 1,8318 0,21097 0,7152 0,13565 1,46677 1,1957 1,31401 1,39787
1,27588 1,17227 1,31164 1,47753 1,12432 1,2839 1,47863 1,26882 1,6630 1,21134 1,19651
1,27453 1,14355 1,10102 1,29343 0,7942 1,1493 1,28572 0,38982 1,9057 1,15971 1,890 1,41953
1,49738 1,23491 1,31693 1,33812 1,32832 1,9872 0,9447 1,3797 0,32651 1,40169 ...

A.4 Zipf Input Trace

1,8550 0,3609 1,654 1,17913 0,145 0,2550 1,5970 0,2461 1,33806 0,17 0,1 0,17 0,1,370 0,159
0,10290 0,54 0,4 1,40078 0,481 0,14300 0,1 0,16 0,18 1,1167 0,7 1,27473 0,47 0,127 0,286
0,35 1,1 0,63 0,15 0,17 0,574 1,1815 0,173 0,6 0,9172 0,5565 1,69 1,7723 0,39491 0,2020
1,1 0,16 1,17217 1,3717 1,3294 1,31 1,40143 0,49198 0,15221 0,191 0,49491 1,2842 1,2797
0,25825 0,7165 1,40 0,3 1,47 0,6 0,8631 0,7375 0,9649 0,3530 0,21 1,508 1,8 0,16 1,20349
1,4506 1,279 1,111 0,1472 0,2768 0,36002 0,168 0,631 0,50 0,44415 1,800 0,1847 0,1353
0,115 0,28497 1,2611 0,697 1,1728 0,1 1,32 0,57 1,358 0,522 1,4 0,612 0,2599 1,2 1,5 0,47719
1,8 1,889 0,345 0,1136 0,242 1,10958 0,1178 0,17 1,3 1,20063 0,1992 0,4 1,7485 1,1406
0,168 1,87 1,602 0,1638 0,265 0,15042 1,42 0,16832 0,12 0,49373 0,2 1,2880 0,28 1,761 0,2
1,412 0,30 0,40 0,1244 0,146 1,9 0,30606 0,3 1,2 1,327 0,27188 0,29109 0,4 0,22156 0,145
1,11837 1,12173 1,41 0,49 0,11 1,1256 0,13969 1,18099 1,202 0,9684 0,31846 0,1303 1,5233
1,109 0,35427 0,29287 0,2080 1,74 0,303 1,34 1,4342 1,172 1,46482 0,4 1,30884 0,10994 0,1
1,36777 1,22311 0,4502 1,104 0,2 0,2258 0,26534 0,10 1,1693 1,15000 1,890 0,14941 1,5200
1,89 0,610 0,790 1,24567 0,2514 0,9 1,55 0,39162 1,1007 1,117 1,1203 1,6 1,34659 1,13
0,21797 1,6 0,123 1,2306 0,431 0,763 0,205 0,25 0,5509 0,663 1,920 0,1340 0,36399 1,179
0,17631 0,2 0,96 1,624 0,549 0,1383 1,3269 0,114 1,26 1,11056 0,37000 1,8808 1,39 1,10
0,1006 1,3 0,12883 1,592 1,4360 0,13 0,2387 1,4 0,183 0,1789 0,6300 1,207 1,33724 1,3747
0,23521 0,67 0,2285 0,8867 1,3101 0,33604 1,4 0,99 0,21 1,4084 1,105 1,19222 1,24966
0,18158 0,45453 0,329 1,13 1,625 0,33 0,45 0,2427 1,14065 0,37 1,213 1,14 1,42281 0,7631
0,275 0,204 1,9059 0,39832 0,9638 0,26 0,10692 1,22052 0,29487 1,1265 1,8160 1,10 1,232
0,30695 1,696 1,296 0,7309 0,3 0,45299 0,52 0,1 0,5550 0,16061 0,2100 0,7392 0,6111 0,28824
1,810 1,15492 1,534 0,21235 1,858 0,2 0,4657 1,1 1,4492 0,35641 0,12254 0,6213 1,9804
0,49110 1,31345 1,3 0,137 0,11385 1,2 1,2387 1,9581 0,1 0,9 1,19562 0,2436 1,14530 1,25
0,35087 1,650 0,5731 0,1481 0,437 1,97 0,8280 1,10875 1,422 0,791 1,9719 0,1 0,23043
0,22546 0,292 1,2 1,5320 0,5 1,10 1,83 0,3253 0,1128 1,6808 0,39912 1,5172 0,175 0,26391
0,679 0,3 1,27741 1,24082 1,3 0,692 0,8297 0,80 0,843 0,26940 0,2092 1,37 0,2655 1,38
1,2622 0,3115 1,23585 1,25069 1,1705 0,11 1,427 1,230 0,14530 0,788 0,27979 1,44389
1,2713 1,93 0,1 0,292 1,42296 0,1307 0,484 0,115 0,8 0,389 0,7576 1,36924 0,180 0,3128
0,27230 0,2 0,28 0,475 1,83 0,36531 1,8931 0,9726 1,2378 1,7889 1,1504 1,481 1,31 0,256
1,10779 1,6 0,64 0,17928 1,1031 0,3837 1,5605 0,145 0,3127 1,42369 1,548 1,12 0,5 ...