

# Project Four

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## LAB 9

### CRAFTING A COMPILER EXERCISES:

EXERCISE: 5.5

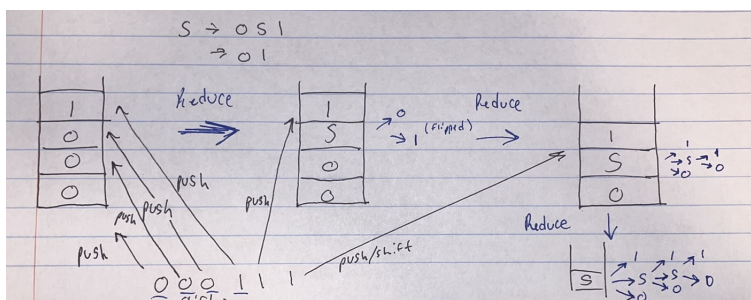
Transform the following grammar into LL(1) form using the techniques presented in Section 5.5:

```
1 DeclList -> DeclList ; Decl
2 -> Decl
3 Decl -> IdList : Type
4 IdList -> IdList , id
5 -> id
6 Type -> ScalarType
7 -> array ( ScalarTypeList ) of Type
8 ScalarType -> id
9 -> Bound . . Bound
10 Bound -> Sign intconstant
11 -> id
12 Sign -> +
13 -> -
14 -> λ
15 ScalarTypelist -> ScalarTypeList , ScalarType
16 -> ScalarType
```

1. DeclList	=>	Decl DeclList *
2. DeclList *	=>	DeclList
3.	=>	$\epsilon$
4. Decl	=>	IdList Type
5. IdList	=>	id IdList *
6. IdList *	=>	IdList
7.	=>	$\epsilon$
8. Type	=>	ScalarType
9.	=>	array (ScalarTypeList) of Type
10. ScalarType	=>	Bound *
11. Bound	=>	Sign intConstant
12.	=>	Bound *
13. Bound *	=>	id Bound **
14. Bound **	=>	id
15.	=>	$\epsilon$
16. Sign	=>	+
17.	=>	-
18.	=>	$\lambda$
19. ScalarTypeList	=>	ScalarType ScalarTypeList *
20. ScalarTypeList *	=>	ScalarTypeList
21.	=>	$\epsilon$

EXERCISE: 4.5.3

a) The input 000111 according to the grammar of Exercise 4.5.1.



$$S \rightarrow SS +$$

$$\rightarrow SS *$$

$$\rightarrow a$$

The diagram illustrates the LR(0) item sets and transitions for the grammar  $S \rightarrow SS +$ ,  $S \rightarrow SS *$ . The transitions are labeled with grammar symbols ( $a$ ,  $+$ ,  $*$ ) and the action "reduce".

The diagram shows the sequence of reductions for the expression  $a + a * a + a + a$ . The stack of states, the stack of symbols, and the stack of operators are shown at each step.

The sequence of reductions is as follows:

- Initial state:  $S$
- Shift  $a$ :  $S \rightarrow a$
- Shift  $+$ :  $S \rightarrow a +$
- Shift  $a$ :  $S \rightarrow a + a$
- Shift  $*$ :  $S \rightarrow a + a *$
- Shift  $a$ :  $S \rightarrow a + a * a$
- Shift  $+$ :  $S \rightarrow a + a * a +$
- Shift  $a$ :  $S \rightarrow a + a * a + a$
- Shift  $+$ :  $S \rightarrow a + a * a + a +$
- Shift  $a$ :  $S \rightarrow a + a * a + a + a$
- Reduce  $S \rightarrow SS +$ :  $S \rightarrow (a + a * a + a) + a$
- Reduce  $S \rightarrow SS *$ :  $S \rightarrow (a + a * a + a) * a$
- Reduce  $S \rightarrow SS +$ :  $S \rightarrow ((a + a * a + a) * a) + a$
- Reduce  $S \rightarrow SS *$ :  $S \rightarrow ((a + a * a + a) * a) * a$
- Reduce  $S \rightarrow SS +$ :  $S \rightarrow (((a + a * a + a) * a) * a) + a$
- Reduce  $S \rightarrow SS *$ :  $S \rightarrow (((a + a * a + a) * a) * a) * a$
- Reduce  $S \rightarrow SS +$ :  $S \rightarrow ((((a + a * a + a) * a) * a) * a) + a$
- Reduce  $S \rightarrow SS *$ :  $S \rightarrow ((((((a + a * a + a) * a) * a) * a) * a) * a) * a$

# PROJECT 4

## CODE GENERATION

In the Compilation process, Code Generation is the fourth and final step. In this step, we accept the Abstract Syntax Tree and Symbol Table produced by the Semantic Analyzer and make the 6502 Op Codes needed to execute our programs in a specific Operating System.

To test my Code Generator I used the test cases below:

### TEST CASE 1

```
1  {
2    int a
3    a = 1
4    {
5      int a
6      a = 2
7      print(a)
8    }
9    string b
10   b = "drake"
11   print("collin")
12   print(a)
13   print(b)
14 }$
```

### TEST CASE 2

```
1  {
2    int a
3    boolean b
4    a = 2+3+4
5    b = (2 == 2)
6    print(a)
7    print(b)
8    print(2+2)
9    print((2 == 3))
10 }$
```

### TEST CASE 3

```
1  {
2    int a
3    a = 5
4    if(true == false){
5      a = 7
6      print(a)
7    }
8  }$
```

#### TEST CASE 4

```
1      {
2          int a
3          a = 5
4          if (2==2){
5              a = 7
6              print(a)
7              if (2==3){
8                  print(a)
9              }
10             string b
11             b = "you win"
12             print(b)
13         }
14         print(true)
15     }$
```

#### TEST CASE 5

```
1      {
2          int x
3          x = 9
4          print(x)
5          while (x == 9){
6              x = 1+x
7              print(x)
8          }
9      }$
```

## APPENDIX

### BRANCHTABLEVARIABLE.JAVA

```
1  /*
2     Branch table variable file
3     Creates objects for the code generation branch table
4  */
5
6  //The branch table variable class!
7  public class branchTableVariable {
8      String temp;    //jump to variable
9      int distance;   //how many bytes to jump
10
11     //Static table variable constructor -- initializes all variables
12     public branchTableVariable(String temp, int distance){
13         this.temp = temp;
14         this.distance = distance;
15     }
16 }
```

### STATICTABLEVARIABLE.JAVA

```
1  /*
2     Static table variable file
3     Creates objects for the code generation variable table
4  */
5
6  //The static table variable object class!
7  public class staticTableVariable {
8      String tempAddress; //store the accumulator in a temp location/address TX
9      String var; //what variable is this?
10     int scope; //keeps track of the scope the variable is located in
11     int offset; //number of offset following code section of the array
12                 (location of variable stored)
13
14     //Static table variable constructor -- initializes all variables
15     public staticTableVariable(String temp, String var, int scope, int offset){
16         this.tempAddress = temp;
17         this.var = var;
18         this.scope = scope;
19         this.offset = offset;
20     }
21 }
```

## REFERENCES

### LINKS

Below are the resources I have used to create simple, readable, and beautiful code.

- 6502 op codes: [labouseur.com](http://labouseur.com)
- output checking: [Arnell Compiler](#)