

Converting from Decimal to Binary

- convert 1693 to binary
- use a divisor of 2 to obtain the following sequence of quotients and remainders

dividend	quotient	remainder
1693	846	1
846	423	0
423	211	1
211	105	1
105	52	1
52	26	0
26	13	0
13	6	1
6	3	0
3	1	1
1	0	1

- read remainders in reverse order $1693_{10} = 11010011101_2$

Converting Between Hex and Binary

chart of values

decimal	hex	binary	decimal	hex	binary
0	0	0000	8	8	1000
1	1	0001	9	9	1001
2	2	0010	10	A	1010
3	3	0011	11	B	1011
4	4	0100	12	C	1100
5	5	0101	13	D	1101
6	6	0110	14	E	1110
7	7	0111	15	F	1111

to convert from binary to hex

- start at right of binary number
- convert each group of 4 digits into a hex value
- e.g., convert 11011101100_2 to hex

binary: 0110 1110 1100

hex: 6 E C

Octal

- $2^4 = 16$ and $2^3 = 8$
 - power = # of bits per hex/octal digit
- **Binary to Hex**
 - every 4 bits = 1 hex digit
- **Octal – base 8**
 - digits 0-7
- **Binary to Octal**
 - Every 3 bits = 1 octal digit

DEC	OCT	HEX	BIN	Notes
0	0	0	0	-
1	1	1	1	2^0
2	2	2	10	2^1
3	3	3	11	
4	4	4	100	2^2
5	5	5	101	
6	6	6	110	
7	7	7	111	
8	10	8	1000	2^3
9	11	9	1001	
10	12	A	1010	
11	13	B	1011	
12	14	C	1100	
13	15	D	1101	
14	16	E	1110	
15	17	F	1111	

Example Data Representations

C Data Type	Typical 32-bit	Typical 64-bit	x86-64
char	1	1	1
short	2	2	2
int	4	4	4
long	4	8	8
float	4	4	4
double	8	8	8
long double	—	—	10/16
pointer	4	8	8

Boolean Algebra

■ Developed by George Boole in 19th Century

- Algebraic representation of logic
 - Encode “True” as 1 and “False” as 0

And

- $A \& B = 1$ when both $A=1$ and $B=1$

&	0	1
0	0	0
1	0	1

Or

- $A | B = 1$ when either $A=1$ or $B=1$

	0	1
0	0	1
1	1	1

Not

- $\sim A = 1$ when $A=0$

\sim	
0	1
1	0

Exclusive-Or (Xor)

- $A \wedge B = 1$ when either $A=1$ or $B=1$, but not both

\wedge	0	1
0	0	1
1	1	0

General Boolean Algebras

■ Operate on Bit Vectors

- Operations applied bitwise

$$\begin{array}{rcl} \begin{array}{c} 01101001 \\ \& 01010101 \end{array} & \begin{array}{c} 01101001 \\ | \quad 01010101 \end{array} & \begin{array}{c} 01101001 \\ ^\wedge \quad 01010101 \end{array} \\ \hline \begin{array}{c} 01000001 \\ 01111101 \end{array} & \begin{array}{c} 01111101 \\ 00111100 \end{array} & \begin{array}{c} 10101010 \\ 00111100 \end{array} \end{array}$$

■ All of the Properties of Boolean Algebra Apply

Contrast: Logic Operations in C

■ Contrast to Bit-level Operators

- `&&`, `||`, `!`
 - View 0 as “false”
 - Anything nonzero as “true”
 - Always return 0 or 1
 - Early termination

■ Examples (char data type)

- `!0x41` → `0x00`
- `!0x00` → `0x01`
- `!!0x41` → `0x01`

- `0x69 && 0x55` → `0x01`
- `0x69 || 0x55` → `0x01`
- `p && *p` (avoids null pointer access)

Shift Operations

■ Left Shift: $x \ll y$

- Shift bit-vector x left y positions
 - Throw away extra bits on left
 - Fill with 0's on right

■ Right Shift: $x \gg y$

- Shift bit-vector x right y positions
 - Throw away extra bits on right
- Logical shift
 - Fill with 0's on left
- Arithmetic shift
 - Replicate most significant bit on left

■ Undefined Behavior

- Shift amount < 0 or \geq word size

Argument x	01100010
$\ll 3$	00010000
Log. $\gg 2$	00011000
Arith. $\gg 2$	00011000

Argument x	10100010
$\ll 3$	00010000
Log. $\gg 2$	00101000
Arith. $\gg 2$	11101000

Encoding Integers

Unsigned

$$B2U(X) = \sum_{i=0}^{w-1} x_i \cdot 2^i$$

Two's Complement

$$B2T(X) = -x_{w-1} \cdot 2^{w-1} + \sum_{i=0}^{w-2} x_i \cdot 2^i$$

```
short int x = 15213;  
short int y = -15213;
```

Sign
Bit

■ C short 2 bytes long

	Decimal	Hex	Binary
x	15213	3B 6D	00111011 01101101
y	-15213	C4 93	11000100 10010011

■ Sign Bit

- For 2's complement, most significant bit indicates sign
 - 0 for nonnegative
 - 1 for negative

Two-complement Encoding Example (Cont.)

x =	15213:	00111011	01101101
y =	-15213:	11000100	10010011

Weight	15213		-15213	
1	1	1	1	1
2	0	0	1	2
4	1	4	0	0
8	1	8	0	0
16	0	0	1	16
32	1	32	0	0
64	1	64	0	0
128	0	0	1	128
256	1	256	0	0
512	1	512	0	0
1024	0	0	1	1024
2048	1	2048	0	0
4096	1	4096	0	0
8192	1	8192	0	0
16384	0	0	1	16384
-32768	0	0	1	-32768
Sum	15213		-15213	

Two's Complement – Simple Conversion

■ conversion using a two-step process

- reverse the bits of the positive representation
- add 1 to the result
- e.g.,

00001001 9

11110110 reverse all bits

11110111 add 1 = -9

■ only one representation for 0

00000000

$11111111 + 1 = 00000000$

■ one more negative number than positive number

Two's Complement – Alternate Conversion

■ alternate conversion using a two-step process

- reading from right to left, copy all values up to and including the first 1
- reverse the remainder of the bits
- e.g.,

00011100	28
11100100	-28

Numeric Ranges

■ Unsigned Values

- $UMin = 0$
000...0
- $UMax = 2^w - 1$
111...1

■ Two's Complement Values

- $TMin = -2^{w-1}$
100...0
- $TMax = 2^{w-1} - 1$
011...1

■ Other Values

- Minus 1
111...1

Values for $W = 16$

	Decimal	Hex	Binary
UMax	65535	FF FF	11111111 11111111
TMax	32767	7F FF	01111111 11111111
TMin	-32768	80 00	10000000 00000000
-1	-1	FF FF	11111111 11111111
0	0	00 00	00000000 00000000

Values for Different Word Sizes

	W			
	8	16	32	64
UMax	255	65,535	4,294,967,295	18,446,744,073,709,551,615
TMax	127	32,767	2,147,483,647	9,223,372,036,854,775,807
TMin	-128	-32,768	-2,147,483,648	-9,223,372,036,854,775,808

Observations

- | $TMin$ | = $TMax + 1$
 - Asymmetric range
- $UMax = 2 * TMax + 1$

C Programming

- #include <limits.h>
- Declares constants, e.g.,
 - ULONG_MAX
 - LONG_MAX
 - LONG_MIN
- Values are platform-specific

Unsigned & Signed Numeric Values

X	$B2U(X)$	$B2T(X)$
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	-8
1001	9	-7
1010	10	-6
1011	11	-5
1100	12	-4
1101	13	-3
1110	14	-2
1111	15	-1

■ Equivalence

- Same encodings for nonnegative values

■ Uniqueness

- Every bit pattern represents unique integer value
- Each representable integer has unique bit encoding

■ \Rightarrow Can Invert Mappings

- $U2B(x) = B2U^{-1}(x)$
 - Bit pattern for unsigned integer
- $T2B(x) = B2T^{-1}(x)$
 - Bit pattern for two's comp integer

Casting Surprises

■ Expression Evaluation

- If there is a mix of unsigned and signed in single expression,
signed values implicitly cast to unsigned
- Including comparison operations `<`, `>`, `==`, `<=`, `>=`
- Examples for $W = 32$: **TMIN = -2,147,483,648**, **TMAX = 2,147,483,647**

■ Constant ₁	Constant ₂	Relation	Evaluation
0	0U	<code>==</code>	unsigned
-1	0	<code><</code>	signed
-1	0U	<code>></code>	unsigned
2147483647	-2147483647-1	<code>></code>	signed
2147483647U	-2147483647-1	<code><</code>	unsigned
-1	-2	<code>></code>	signed
(unsigned)-1	-2	<code>></code>	unsigned
2147483647	2147483648U	<code><</code>	unsigned
2147483647	(int) 2147483648U	<code>></code>	signed

Binary Multiplication

■ 8-bit binary multiplication

$$\begin{array}{r} 01010101 \\ \times 00011001 \\ \hline 01010101 \\ 01010101 \\ \hline 100001001101 \end{array}$$

truncated: 0100 1101

Power-of-2 Multiply with Shift

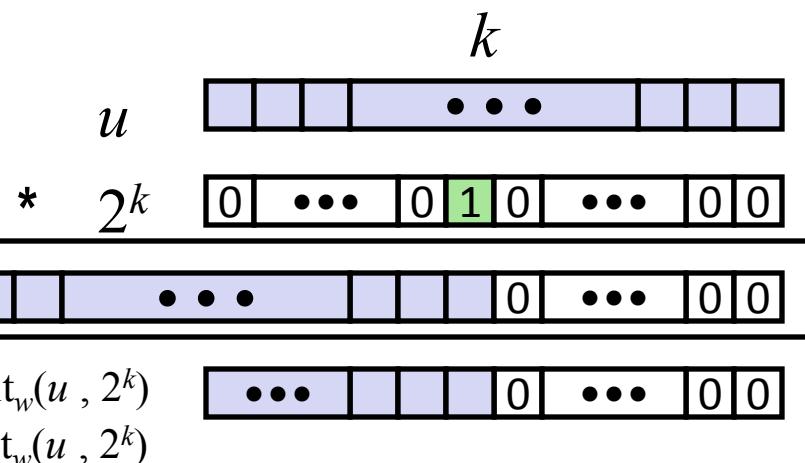
Operation

- $u \ll k$ gives $u * 2^k$
- Both signed and unsigned

Operands: w bits

True Product: $w+k$ bits

Discard k bits: w bits



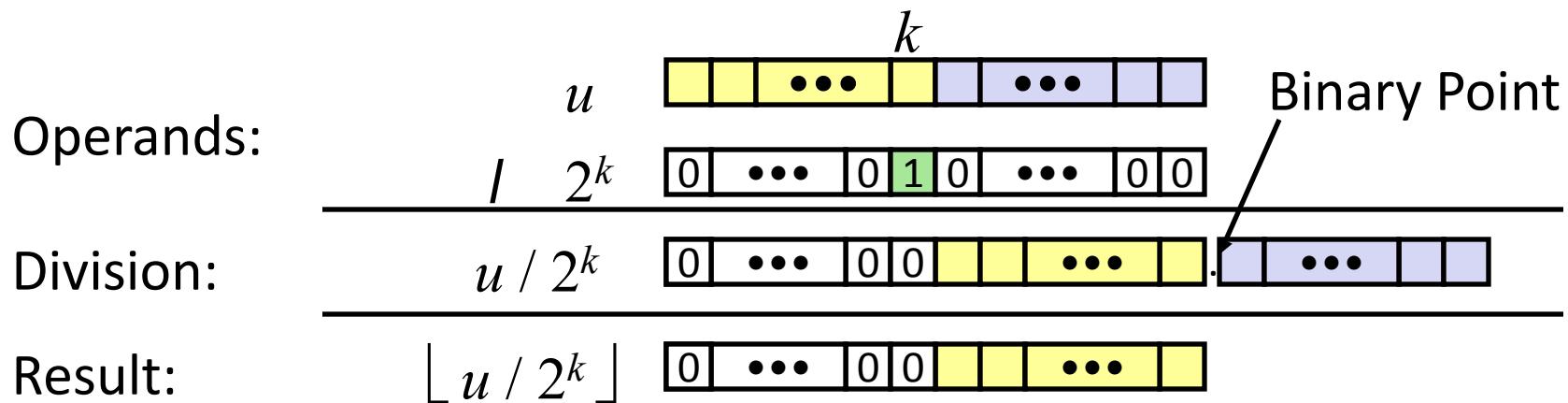
Examples

- $u \ll 3 == u * 8$
- $(u \ll 5) - (u \ll 3) == u * 24$
- Most machines shift and add faster than multiply
 - Compiler generates this code automatically

Unsigned Power-of-2 Divide with Shift

Quotient of Unsigned by Power of 2

- $u \gg k$ gives $\lfloor u / 2^k \rfloor$
- Uses logical shift



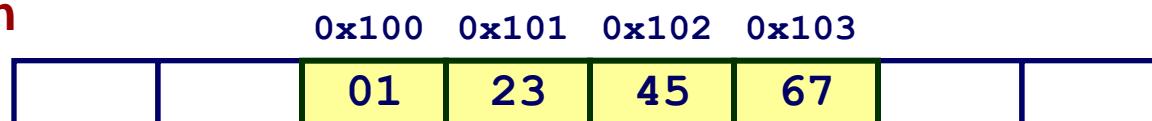
	Division	Computed	Hex	Binary
x	15213	15213	3B 6D	00111011 01101101
x >> 1	7606.5	7606	1D B6	00011101 10110110
x >> 4	950.8125	950	03 B6	00000011 10110110
x >> 8	59.4257813	59	00 3B	00000000 00111011

Byte Ordering Example

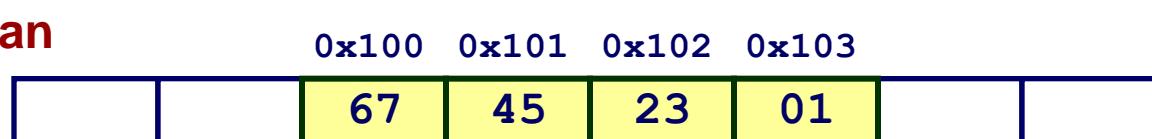
Example

- Variable x has 4-byte value of 0x01234567
- Address given by &x is 0x100

BigEndian



Little Endian



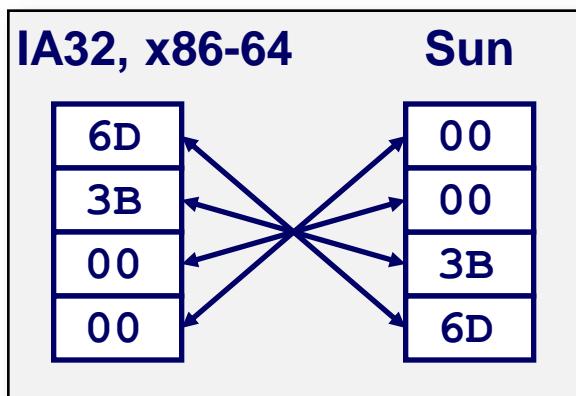
Representing Integers

Decimal: 15213

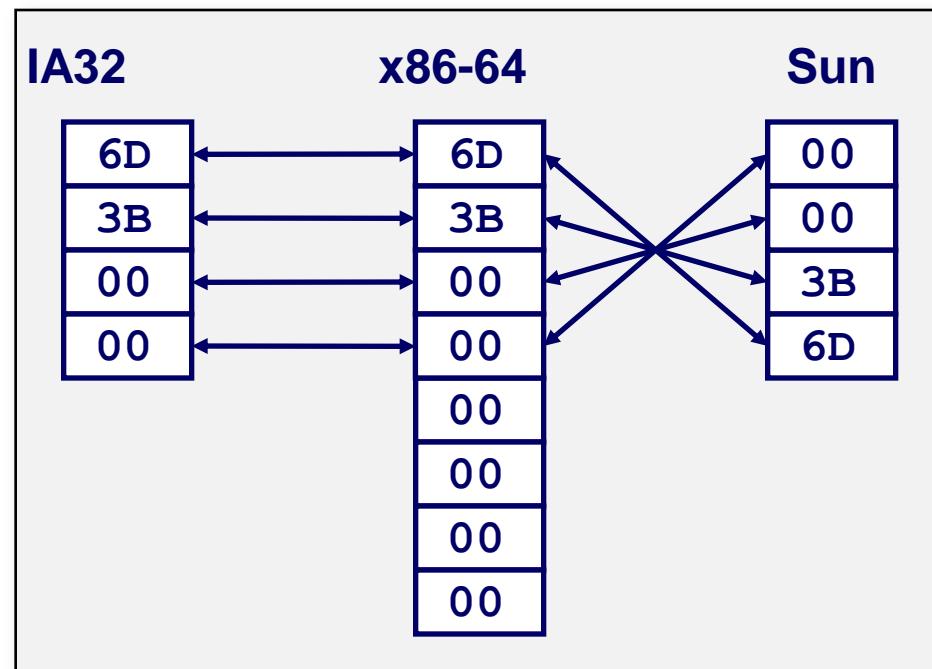
Binary: 0011 1011 0110 1101

Hex: 3 B 6 D

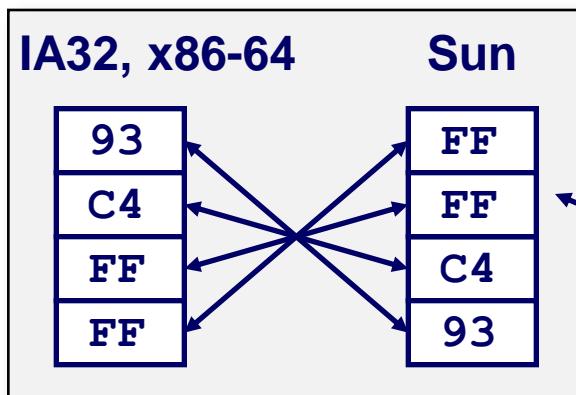
```
int A = 15213;
```



```
long int C = 15213;
```



```
int B = -15213;
```



Two's complement representation