

# LECTURE 9

Computing an LL(1)  
Parsing Table

# REVIEW

Last lecture, we parsed an input string using our recursive descent parser. We then discussed how we could create an equivalent table-driven parser.

We ended the lecture with a discussion of how we might begin to build a parsing table.

Today, we'll create the parsing table from scratch and use it to parse an input string.

# LL(1) GRAMMARS

- A grammar whose parsing table has no multiply-defined entries is an LL(1) grammar.
- Uses one input symbol of look-ahead at each step to make a parsing decision.
- No ambiguous or left-recursive grammar can be LL(1).

# PARSING TABLE

The basic outline for creating a parsing table from a LL(1) grammar is the following:

- Compute the First sets of the non-terminals.
- Compute the Follow sets of the non-terminals.
- For each production  $N \rightarrow \omega$ ,
  - Add  $N \rightarrow \omega$  to  $M[N, t]$  for each  $t$  in  $\text{First}(\omega)$ .
  - If  $\text{First}(\omega)$  contains  $\epsilon$ , add  $N \rightarrow \omega$  to  $M[N, t]$  for each  $t$  in  $\text{Follow}(N)$ .
- All undefined entries represent a parsing error.

# FIRST SETS

Let's calculate our First sets. The rules are:

- If  $X$  is a terminal symbol,  $\text{First}(X) = X$ .
- If  $X$  is  $\epsilon$ , add  $\epsilon$  to  $\text{First}(X)$ .
- If  $X$  is a non-terminal, look at all productions where  $X$  is on left-hand side. Each production will be of the form:  $X \rightarrow Y_1 Y_2 \dots Y_k$  where  $Y$  is a nonterminal or terminal. Then:
  - Put  $\text{First}(Y_1) - \epsilon$  in  $\text{First}(X)$ .
  - If  $\epsilon$  is in  $\text{First}(Y_1)$ , then put  $\text{First}(Y_2) - \epsilon$  in  $\text{First}(X)$ .
  - If  $\epsilon$  is in  $\text{First}(Y_2)$ , then put  $\text{First}(Y_3) - \epsilon$  in  $\text{First}(X)$ .
  - ...
  - If  $\epsilon$  is in  $Y_1, Y_2, \dots, Y_k$ , then add  $\epsilon$  to  $\text{First}(X)$ .

# FIRST SETS

The very first thing we'll do is compute the First set of each of our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$First(factor) = ?$

# FIRST SETS

The very first thing we'll do is compute the First set of each of our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$First(factor) = \{ '(', int \}$

# FIRST SETS

The very first thing we'll do is compute the First set of each of our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$First(term) = ?$

$First(factor) = \{ '(', int \}$



# FIRST SETS

The very first thing we'll do is compute the First set of each of our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$First(term) = \{ '(', int \}$

$First(factor) = \{ '(', int \}$

# FIRST SETS

The very first thing we'll do is compute the First set of each of our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$First(expr) = \{ '(', int \}$

$First(term) = \{ '(', int \}$

$First(factor) = \{ '(', int \}$

# FIRST SETS

The very first thing we'll do is compute the First set of each of our non-terminals.

$program \rightarrow expr$

$First(program) = \{ '(', int \}$

$expr \rightarrow term\ expr\_tail$

$First(expr) = \{ '(', int \}$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$First(term) = \{ '(', int \}$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$First(factor) = \{ '(', int \}$

# FIRST SETS

The very first thing we'll do is compute the First set of each of our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$First(program) = \{ '(', int \}$

$First(expr) = \{ '(', int \}$

$First(expr\_tail) = \{ '+', \epsilon \}$

$First(term) = \{ '(', int \}$

$First(factor) = \{ '(', int \}$

# FIRST SETS

The very first thing we'll do is compute the First set of each of our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$First(program) = \{ '(', int \}$

$First(expr) = \{ '(', int \}$

$First(expr\_tail) = \{ '+', \epsilon \}$

$First(term) = \{ '(', int \}$

$First(term\_tail) = \{ '*', \epsilon \}$

$First(factor) = \{ '(', int \}$

# FOLLOW SETS

Now, we'll look at the Follow sets. The rules are:

- If  $N$  is the starting non-terminal, put  $\$$  in  $\text{Follow}(N)$ .
- If  $X \rightarrow \alpha N$ , where  $\alpha$  is some string of non-terminals and/or terminals, put  $\text{Follow}(X)$  in  $\text{Follow}(N)$ .
- If  $X \rightarrow \alpha N \beta$  where  $\alpha, \beta$  are some string of non-terminals and/or terminals, put  $\text{First}(\beta)$  in  $\text{Follow}(N)$ . If  $\text{First}(\beta)$  includes  $\epsilon$ , then put  $\text{Follow}(X)$  in  $\text{Follow}(N)$ .

# FOLLOW SETS

Now let's calculate the Follow sets for our non-terminals.

$program \rightarrow expr$

$Follow(program) = \{\$ \}$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

# FOLLOW SETS

Now let's calculate the Follow sets for our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$Follow(program) = \{\$ \}$

$Follow(expr) = \{\$, '\}'$



# FOLLOW SETS

Now let's calculate the Follow sets for our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$Follow(program) = \{\$ \}$

$Follow(expr) = \{\$, '\}'$

$Follow(expr\_tail) = \{\$, '\}'$

# FOLLOW SETS

Now let's calculate the Follow sets for our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$Follow(program) = \{\$, \epsilon\}$

$Follow(expr) = \{\$, '\epsilon'\}$

$Follow(expr\_tail) = \{\$, '\epsilon'\}$

$Follow(term) = \{'+', \$, '\epsilon'\}$

# FOLLOW SETS

Now let's calculate the Follow sets for our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$Follow(program) = \{\$ \}$

$Follow(expr) = \{\$, '\}'$

$Follow(expr\_tail) = \{\$, '\}'$

$Follow(term) = \{'+', \$, '\}'$

$Follow(term\_tail) = \{'+', \$, '\}'$

# FOLLOW SETS

Now let's calculate the Follow sets for our non-terminals.

$program \rightarrow expr$

$expr \rightarrow term\ expr\_tail$

$expr\_tail \rightarrow +\ term\ expr\_tail \mid \epsilon$

$term \rightarrow factor\ term\_tail$

$term\_tail \rightarrow *\ factor\ term\_tail \mid \epsilon$

$factor \rightarrow (expr) \mid int$

$Follow(program) = \{\$, \epsilon\}$

$Follow(expr) = \{\$, '\epsilon'\}$

$Follow(expr\_tail) = \{\$, '\epsilon'\}$

$Follow(term) = \{'+', \$, '\epsilon'\}$

$Follow(term\_tail) = \{'+', \$, '\epsilon'\}$

$Follow(factor) = \{'*', '+', \$, '\epsilon'\}$

# FIRST AND FOLLOW SETS

So, now we have our First and Follow sets. From here, we can construct our parsing table.

$program \rightarrow expr$

$First(program) = \{ '(', int \}$

$Follow(program) = \{ \$ \}$

$expr \rightarrow term \ expr\_tail$

$First(expr) = \{ '(', int \}$

$Follow(expr) = \{ \$, ' \}$

$expr\_tail \rightarrow + \ term \ expr\_tail \mid \epsilon$

$First(expr\_tail) = \{ '+', \epsilon \}$

$Follow(expr\_tail) = \{ \$, ' \}$

$term \rightarrow factor \ term\_tail$

$First(term) = \{ '(', int \}$

$Follow(term) = \{ '+', \$, ' \}$

$term\_tail \rightarrow * \ factor \ term\_tail \mid \epsilon$

$First(term\_tail) = \{ '*', \epsilon \}$

$Follow(term\_tail) = \{ '+', \$, ' \}$

$factor \rightarrow (expr) \mid int$

$First(factor) = \{ '(', int \}$

$Follow(factor) = \{ '*', '+', \$, ' \}$

# CONSTRUCTING AN LL(1) PARSING TABLE

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

$First(program) = \{(' , int)\}$   
 $First(expr) = \{(' , int)\}$   
 $First(expr\_tail) = \{+', \epsilon\}$   
 $First(term) = \{(' , int)\}$   
 $First(term\_tail) = \{*', \epsilon\}$   
 $First(factor) = \{(' , int)\}$

$Follow(program) = \{\$\}$   
 $Follow(expr) = \{\$, '\}'$   
 $Follow(expr\_tail) = \{\$, '\}'$   
 $Follow(term) = \{+', \$, '\}'$   
 $Follow(term\_tail) = \{+', \$, '\}'$   
 $Follow(factor) = \{*', '+', \$, '\}'$

Let's start with program, production number 1. What is  $First(expr)$ ?

N	(	int	*	+	)	\$
program						
expr						
expr_tail						
term						
term_tail						
factor						

# CONSTRUCTING AN LL(1) PARSING TABLE

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

$First(program) = \{ '(', int \}$   
 $First(expr) = \{ '(', int \}$   
 $First(expr\_tail) = \{ '+', \epsilon \}$   
 $First(term) = \{ '(', int \}$   
 $First(term\_tail) = \{ '*', \epsilon \}$   
 $First(factor) = \{ '(', int \}$

$Follow(program) = \{ \$ \}$   
 $Follow(expr) = \{ \$, ') \}$   
 $Follow(expr\_tail) = \{ \$, ') \}$   
 $Follow(term) = \{ '+', \$, ') \}$   
 $Follow(term\_tail) = \{ '+', \$, ') \}$   
 $Follow(factor) = \{ '*', '+', \$, ') \}$

Note: You are not looking at First() of the left side of the production. You are calculating First() of the right side!

N	(	int	*	+	)	\$
program	(1)	(1)				
expr						
expr_tail						
term						
term_tail						
factor						

# CONSTRUCTING AN LL(1) PARSING TABLE

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

$First(program) = \{ '(', int \}$   
 $First(expr) = \{ '(', int \}$   
 $First(expr\_tail) = \{ '+', \epsilon \}$   
 $First(term) = \{ '(', int \}$   
 $First(term\_tail) = \{ '*', \epsilon \}$   
 $First(factor) = \{ '(', int \}$

$Follow(program) = \{ \$ \}$   
 $Follow(expr) = \{ \$, ') \}$   
 $Follow(expr\_tail) = \{ \$, ') \}$   
 $Follow(term) = \{ '+', \$, ') \}$   
 $Follow(term\_tail) = \{ '+', \$, ') \}$   
 $Follow(factor) = \{ '*', '+', \$, ') \}$

N	(	int	*	+	)	\$
program	(1)	(1)				
expr	(2)	(2)				
expr_tail						
term						
term_tail						
factor						



# CONSTRUCTING AN LL(1) PARSING TABLE

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

$First(program) = \{ '(', int \}$   
 $First(expr) = \{ '(', int \}$   
 $First(expr\_tail) = \{ '+', \epsilon \}$   
 $First(term) = \{ '(', int \}$   
 $First(term\_tail) = \{ '*', \epsilon \}$   
 $First(factor) = \{ '(', int \}$

$Follow(program) = \{ \$ \}$   
 $Follow(expr) = \{ \$, ') \}$   
 $Follow(expr\_tail) = \{ \$, ') \}$   
 $Follow(term) = \{ '+', \$, ') \}$   
 $Follow(term\_tail) = \{ '+', \$, ') \}$   
 $Follow(factor) = \{ '*', '+', \$, ') \}$

N	(	int	*	+	)	\$
program	(1)	(1)				
expr	(2)	(2)				
expr_tail				(3)	(4)	(4)
term						
term_tail						
factor						

# CONSTRUCTING AN LL(1) PARSING TABLE

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

$First(program) = \{ '(', int \}$   
 $First(expr) = \{ '(', int \}$   
 $First(expr\_tail) = \{ '+', \epsilon \}$   
 $First(term) = \{ '(', int \}$   
 $First(term\_tail) = \{ '*', \epsilon \}$   
 $First(factor) = \{ '(', int \}$

$Follow(program) = \{ \$ \}$   
 $Follow(expr) = \{ \$, ')\}$   
 $Follow(expr\_tail) = \{ \$, ')\}$   
 $Follow(term) = \{ '+', \$, ')\}$   
 $Follow(term\_tail) = \{ '+', \$, ')\}$   
 $Follow(factor) = \{ '*', '+', \$, ')\}$

N	(	int	*	+	)	\$
program	(1)	(1)				
expr	(2)	(2)				
expr_tail				(3)	(4)	(4)
term	(5)	(5)				
term_tail						
factor						

# CONSTRUCTING AN LL(1) PARSING TABLE

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

$First(program) = \{ '(', int \}$   
 $First(expr) = \{ '(', int \}$   
 $First(expr\_tail) = \{ '+', \epsilon \}$   
 $First(term) = \{ '(', int \}$   
 $First(term\_tail) = \{ '*', \epsilon \}$   
 $First(factor) = \{ '(', int \}$

$Follow(program) = \{ \$ \}$   
 $Follow(expr) = \{ \$, ']' \}$   
 $Follow(expr\_tail) = \{ \$, ']' \}$   
 $Follow(term) = \{ '+', \$, ']' \}$   
 $Follow(term\_tail) = \{ '+', \$, ']' \}$   
 $Follow(factor) = \{ '*', '+', \$, ']' \}$

N	(	int	*	+	)	\$
program	(1)	(1)				
expr	(2)	(2)				
expr_tail				(3)	(4)	(4)
term	(5)	(5)				
term_tail			(6)	(7)	(7)	(7)
factor						

# CONSTRUCTING AN LL(1) PARSING TABLE

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

$First(program) = \{ '(', int \}$   
 $First(expr) = \{ '(', int \}$   
 $First(expr\_tail) = \{ '+', \epsilon \}$   
 $First(term) = \{ '(', int \}$   
 $First(term\_tail) = \{ '*', \epsilon \}$   
 $First(factor) = \{ '(', int \}$

$Follow(program) = \{ \$ \}$   
 $Follow(expr) = \{ \$, ') \}$   
 $Follow(expr\_tail) = \{ \$, ') \}$   
 $Follow(term) = \{ '+', \$, ') \}$   
 $Follow(term\_tail) = \{ '+', \$, ') \}$   
 $Follow(factor) = \{ '*', '+', \$, ') \}$

N	(	int	*	+	)	\$
program	(1)	(1)				
expr	(2)	(2)				
expr_tail				(3)	(4)	(4)
term	(5)	(5)				
term_tail			(6)	(7)	(7)	(7)
factor	(8)	(9)				

# CONSTRUCTING AN LL(1) PARSING TABLE

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

$First(program) = \{ '(', int \}$   
 $First(expr) = \{ '(', int \}$   
 $First(expr\_tail) = \{ '+', \epsilon \}$   
 $First(term) = \{ '(', int \}$   
 $First(term\_tail) = \{ '*', \epsilon \}$   
 $First(factor) = \{ '(', int \}$

$Follow(program) = \{ \$ \}$   
 $Follow(expr) = \{ \$, ') \}$   
 $Follow(expr\_tail) = \{ \$, ') \}$   
 $Follow(term) = \{ '+', \$, ') \}$   
 $Follow(term\_tail) = \{ '+', \$, ') \}$   
 $Follow(factor) = \{ '*', '+', \$, ') \}$

N	(	int	*	+	)	\$
program	(1)	(1)	-	-	-	-
expr	(2)	(2)	-	-	-	-
expr_tail	-	-	-	(3)	(4)	(4)
term	(5)	(5)	-	-	-	-
term_tail	-	-	(6)	(7)	(7)	(7)
factor	(8)	(9)	-	-	-	-

# PARSING A STRING

Let's pick a sample string from our grammar to parse.

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

$3 * ( 1 + 2 ) * 4$



int   '\*'   '('   int   '+'   int   ')'   '\*'   int   '\$'

# PARSING A STRING

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

N	(	int	*	+	)	\$
program	(1)	(1)	-	-	-	-
expr	(2)	(2)	-	-	-	-
expr_tail	-	-	-	(3)	(4)	(4)
term	(5)	(5)	-	-	-	-
term_tail	-	-	(6)	(7)	(7)	(7)
factor	(8)	(9)	-	-	-	-

Stack (bottom  $\rightarrow$  top)

\$ program  
 \$ expr  
 \$ expr\_tail term  
 \$ expr\_tail term\_tail factor  
 \$ expr\_tail term\_tail int  
 \$ expr\_tail term\_tail  
 \$ expr\_tail term\_tail factor \*  
 \$ expr\_tail term\_tail factor

Input tokens

int \* ( int + int ) \* int \$  
 int \* ( int + int ) \* int \$  
 int \* ( int + int ) \* int \$  
 int \* ( int + int ) \* int \$  
 int \* ( int + int ) \* int \$  
 \* ( int + int ) \* int \$  
 \* ( int + int ) \* int \$  
 ( int + int ) \* int \$

Production Used

(1)  
 (2)  
 (5)  
 (9)  
  
 (6)

# PARSING A STRING

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

N	(	int	*	+	)	\$
program	(1)	(1)	-	-	-	-
expr	(2)	(2)	-	-	-	-
expr_tail	-	-	-	(3)	(4)	(4)
term	(5)	(5)	-	-	-	-
term_tail	-	-	(6)	(7)	(7)	(7)
factor	(8)	(9)	-	-	-	-

Stack (bottom  $\rightarrow$  top)

\$ expr\_tail term\_tail factor  
\$ expr\_tail term\_tail ) expr (  
\$ expr\_tail term\_tail ) expr  
\$ expr\_tail term\_tail ) expr\_tail term  
\$ expr\_tail term\_tail ) expr\_tail term\_tail factor  
\$ expr\_tail term\_tail ) expr\_tail term\_tail int  
\$ expr\_tail term\_tail ) expr\_tail term\_tail  
\$ expr\_tail term\_tail ) expr\_tail  
\$ expr\_tail term\_tail ) expr\_tail term '+'

Input tokens

( int + int ) \* int \$  
( int + int ) \* int \$  
int + int ) \* int \$  
int + int ) \* int \$  
int + int ) \* int \$  
int + int ) \* int \$  
+ int ) \* int \$  
+ int ) \* int \$  
+ int ) \* int \$

Production Used

(8)  
  
  
(2)  
(5)  
(9)  
  
(7)  
(3)



# PARSING A STRING

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

N	(	int	*	+	)	\$
program	(1)	(1)	-	-	-	-
expr	(2)	(2)	-	-	-	-
expr_tail	-	-	-	(3)	(4)	(4)
term	(5)	(5)	-	-	-	-
term_tail	-	-	(6)	(7)	(7)	(7)
factor	(8)	(9)	-	-	-	-

Stack (bottom  $\rightarrow$  top)

\$ expr\_tail term\_tail ) expr\_tail term '+'  
\$ expr\_tail term\_tail ) expr\_tail term  
\$ expr\_tail term\_tail ) expr\_tail term\_tail factor  
\$ expr\_tail term\_tail ) expr\_tail term\_tail int  
\$ expr\_tail term\_tail ) expr\_tail term\_tail  
\$ expr\_tail term\_tail ) expr\_tail  
\$ expr\_tail term\_tail )  
\$ expr\_tail term\_tail

Input tokens

+ int ) \* int \$  
int ) \* int \$  
int ) \* int \$  
int ) \* int \$  
 ) \* int \$  
 ) \* int \$  
 ) \* int \$  
 \* int \$

Production Used

(3)  
  
(5)  
(9)  
  
(7)  
(4)

# PARSING A STRING

1.  $program \rightarrow expr$
2.  $expr \rightarrow term\ expr\_tail$
3.  $expr\_tail \rightarrow +\ term\ expr\_tail$
4.  $expr\_tail \rightarrow \epsilon$
5.  $term \rightarrow factor\ term\_tail$
6.  $term\_tail \rightarrow *\ factor\ term\_tail$
7.  $term\_tail \rightarrow \epsilon$
8.  $factor \rightarrow (expr)$
9.  $factor \rightarrow int$

N	(	int	*	+	)	\$
program	(1)	(1)	-	-	-	-
expr	(2)	(2)	-	-	-	-
expr_tail	-	-	-	(3)	(4)	(4)
term	(5)	(5)	-	-	-	-
term_tail	-	-	(6)	(7)	(7)	(7)
factor	(8)	(9)	-	-	-	-

Stack (bottom  $\rightarrow$  top)

\$ expr\_tail term\_tail  
 \$ expr\_tail term\_tail factor \*  
 \$ expr\_tail term\_tail factor  
 \$ expr\_tail term\_tail int  
 \$ expr\_tail term\_tail  
 \$ expr\_tail  
 \$

Accept!

Input tokens

\* int \$  
 \* int \$  
 int \$  
 int \$  
 \$  
 \$  
 \$

Production Used

(6)  
  
 (9)  
  
 (7)  
 (4)

# NEXT LECTURE

Review of Lexical and Syntax Analysis, then we'll begin Semantic Analysis.