#### CS 143 HW 6

- 1.  $\{A,B,C,F\}$ + =  $\{A,B,C,D,E,F\}$ . Because  $\{A,B,C,F\}$ + contains a key, we can say that this decomposition is lossless.
- 2. AC -> B, BC -> A
- 3. (a) If we have a cyclic functional dependency, then we can determine that it's one-to-one.
- (b) If it is acyclic FD, then it can be many-to-one.
- 4. (a)  $\{E\}$  + =  $\{A, B, C, D, E\}$ . Because of this, E is a key.
- (b)  $\{B, C\} + = \{B, C, D, E, A\}$ . Because of this and that BC is minimal, BC contains a key.

5. 
$$\{A\} + = \{A, B, C, D, E\} \mid \{C\} + = \{C, E\} \mid \{B\} + = \{D\}$$

Because these non-trivial FD's does not equal R, we can see that this is not in BCNF.

R1(A, B, C, D, E)

## R2(A, F) - in BCNF

Now using FD's for R1:

$${A}+={A, B, C, D, E} | {C}+={C, E}$$

Since  $\{C\}$ + doesn't equal R1, we split R1 into:

### R3(C, E) - in BCNF

R4(A, B, C, D)

Using FD's for R4:

$${A}+={A, B, C, D} \mid {B}+={B, D}$$

Since  $\{B\}$ + doesn't equal R4, we split R4 into:

#### **R5(B,D)**

#### **R6(A,B,C)**

This should be in BCNF using R2, R3, R5, R6.

6. We also know that the following are tuples in R:

7.  $\{AB\}$ + =  $\{A, B, E\}$ 

Split R into:

R1 (A, B, E)

R2 (A, B, C, D, F)

Split R1 into:

# R3(A, B) - in 4NF

## R4(B, E) - in 4NF

Split R2 using the MVD A ->-> B:

R5(A, B)

# R6(A, C, D, F)

Now this is in 4NF using R3, R4, R5, R6.