What is Cedille$_{\text{core}}$?

- Independent implementation of CDLE (fact check?)
- Under 1000 lines of code, in Haskell
- Austere sublanguage of Cedille
- Cedille sanity check

Cedille vs. Cedille$_{\text{core}}$

$(\lambda\ x.\ t)$

$(\lambda\ x : T.\ t)$

$[\ t_1,\ t_2\ ]$

δ

Type Checking

$(\rho\ e\ @\ x.\ T_x\ -\ t)$

ϕ β

$[\ t_1,\ t_2\ @\ x.\ T_x\ ]$

Scope Shadowing

ε

$(\beta<u_T>\{u_t\})$

$(\beta\{u_t\})$

Type Synthesis

ς

χ

θ $(\rho<3\ 4\ 6>\ e\ -\ t)$

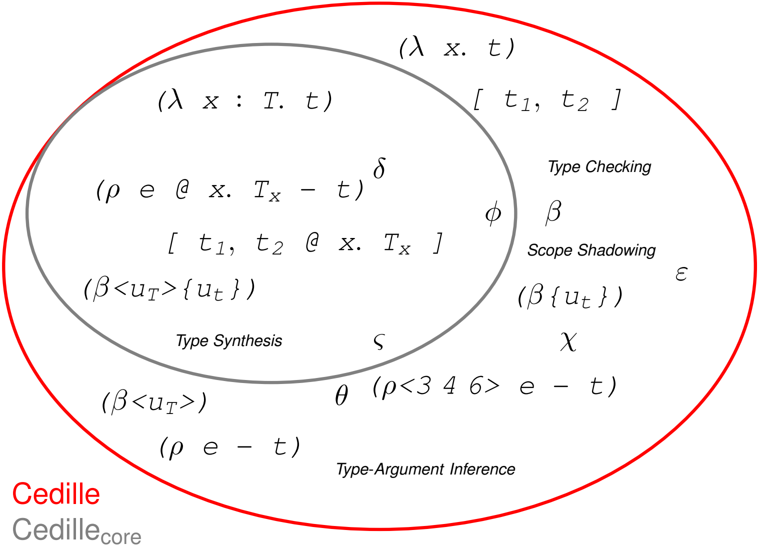
$(\beta<u_T>)$

$(\rho\ e\ -\ t)$

Type-Argument Inference

Cedille

Cedille vs. Cedille$_{\text{core}}$



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- Expand module parameters and import arguments