

# Symbolic Weighted Language Models and Parsing over Infinite Alphabets

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## Abstract

We propose a framework for weighted parsing over infinite alphabets. It is based on language models called Symbolic Weighted Automata (**swA**) at the joint between Symbolic Automata (**sA**) and Weighted Automata (**wA**), as well as Transducers (**swT**) and Visibly Pushdown (**sw-VPA**) variants. Like **sA**, **swA** deal with large or infinite input alphabets, and like **wA**, they output a weight value in a semiring domain. The transitions of **swA** are labeled by functions from an infinite alphabet into the weight domain. This is unlike **sA** whose transitions are guarded by boolean predicates over symbols in an infinite alphabet and also unlike **wA** whose transitions are labeled by constant weight values, and who deal only with finite automata. We present some properties of **swA**, **swT** and **sw-VPA** models, that we use to define and solve a variant of parsing over infinite alphabets. We also briefly describe the application that motivated the introduction of these models: a parse-based approach to automated music transcription.

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## 1 Introduction

Parsing is the problem of structuring a linear representation on input (a finite word), according to a language model. Most of the context-free parsing approaches [15] assume a finite and reasonably small input alphabet. Such a restriction makes perfect sense in the context of NLP tasks such as constituency parsing, or of programming languages compilers or interpreters. Considering large or infinite alphabets can however be of practical interest, for instance, when dealing with large characters encodings such as UTF-16, *e.g.* for vulnerability detection in Web-applications [8], for the analyse (*e.g.* validation or filtering) of data streams or serialization of structured documents (with textual or numerical attributes) [26], or for processing timed execution traces [3]. The latter case is related to a study that motivated the present work: automated music transcription. In this problem, a music performance, represented symbolically in the form of a sequence of timed musical events, is converted into a score in Common Western Music Notation [14], structured according to nested grouping and metric strength of events. It can therefore be stated as a parsing problem [12], over an infinite alphabet of timed events.

Various extensions of language models for handling infinite alphabets have been studied. For instance, some automata with memory extensions allow restricted storage and comparison of input symbols, (see [26] for a survey), with pebbles for marking positions [25], registers [18], or the possibility to compute on subsequences with the same attribute values [2]. The automata at the core of model checkers compute on input symbols represented by large bitvectors [27] (sets of assignments of Boolean variables) and in practice, each transition accepts a set of such symbols (instead of an individual symbol), represented by Boolean

register: skip refs  
and details, add  
Mikolaj recent



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■ **Figure 1** Classes of Symbolic/Weighted Automata.  $\Sigma_{\text{fin}}$  is a finite alphabet,  $\Sigma_{\text{inf}}$  is a countable alphabet,  $\mathbb{B}$  is the Boolean algebra,  $\mathbb{S}$  is a commutative semiring,  $q \xrightarrow{\cdot} q'$  is a transition between states  $q$  and  $q'$ .

44 formula or Binary Decision Diagrams. Following a similar idea, in symbolic automata  
 45 (sA) [7, 8], the transitions are guarded by predicates over infinite alphabet domains. With  
 46 appropriate closure conditions on the sets of such predicates, all the good properties enjoyed  
 47 by automata over finite alphabets are preserved.

48 Other extensions of language models help in dealing with non-determinism, by the  
 49 computation of weight values. With an ambiguous grammar, there may exist several  
 50 derivations (*abstract syntax trees* – AST) yielding one input word. The association of one  
 51 weight value to each AST permits to select a best one (or  $n$  bests). This is roughly the  
 52 principle of *weighted parsing* approaches [13, 24, 23]. In *weighted language models*, like *e.g.*  
 53 probabilistic context-free grammars and weighted automata (wA) [11], a weight value is  
 54 associated to each transition rule, and the rule’s weights can be combined with a associative  
 55 product operator  $\otimes$  into the weight of an AST. A second operator  $\oplus$ , associative and  
 56 commutative, is moreover used to handle the ambiguity of the model, by summing the  
 57 weights of the possibly several (in general exponentially many) AST associated to a given  
 58 input word. Typically,  $\oplus$  will select the best of two weight values. The weight domain,  
 59 equipped with these two operators shall be, at minima, a *semiring* where  $\oplus$  can be extended  
 60 to infinite sums, such as the Viterbi semiring and the tropical min-plus algebra, see Figure 2.

61 In this paper, we present a uniform framework for weighted parsing over infinite input  
 62 alphabets. It is based on *symbolic weighted* finite states language models (swM), gener-  
 63 alizing sA, with functions into an arbitrary semiring instead of Boolean guards, and wA,  
 64 by handling infinite alphabets, see Figure 1. In the transition rules of swM models, input  
 65 symbols appear as variables, and the weight associated to a transition rule is a function of  
 66 these variables. The models presented here are finite automata called symbolic-weighted  
 67 (swA), transducers (swT), and pushdown automata with a visibly restriction [1] (sw-VPA).  
 68 The latter model of automata computes on *nested words* [1], a structured form of words  
 69 parenthesized with markup symbols, corresponding to a linearization of trees. In the context  
 70 of parsing, they can represent (weighted) AST of CF grammars. More precisely, a sw-VPA  $A$   
 71 associates a weight value  $A(t)$  to a given a nested word  $t$ , which is the linearization of an  
 72 AST. On the other hand, a swT can define a distance  $T(s, t)$  between finite words  $s$  and  $t$  over  
 73 infinite alphabets. Then, the *SW-parsing* problem aims at finding  $t$  minimizing  $T(s, t) \otimes A(t)$   
 74 (*wrt* the ranking defined by  $\oplus$ ). The latter value is called the distance between  $s$  and  $A$   
 75 in [21]. Like weighted-parsing methods [13, 24, 23], our approach proceeds in two steps,

based on properties of the swM. The first step is an intersection (Bar-Hillel construction [15]) where, given a swT  $T$ , a sw-VPA  $A$ , and an input word  $s$ , a sw-VPA  $A_{T,s}$  is built, such that for all  $t$ ,  $A_{T,s}(t) = T(s, t) \otimes A(t)$ . In the second step, a best AST  $t$  is found by applying to  $A_{T,s}$  a best search algorithm similar to the shortest distance in graphs [20, 17].

The main contributions of the paper are: (i) the introduction of automata, swA, transducers, swT (Section 3), and visibly pushdown automata sw-VPA (Section 4), generalizing the corresponding classes of symbolic and weighted models, (ii) a polynomial best-search algorithm for sw-VPA, and (iii) a uniform framework (Section 5) for parsing over infinite alphabets, the keys to which are (iii.a) the swT-based definition of generic edit distances between input and output (yield) words, and (iii.b) the use, convenient in this context, of nested words, and sw-VPA, instead of syntax trees and grammars.

chap. intersection  
in [15]

expressiveness: VPA  
have restricted  
equality test. com-  
parable to pebble  
automata? → con-  
clusion

## 2 Preliminary Notions

### Semirings

We shall consider semirings for the weight values of our language models. A *semiring*  $\langle \mathbb{S}, \oplus, \otimes, 0, 1 \rangle$  is a structure with a domain  $\mathbb{S}$ , equipped with two associative binary operators  $\oplus$  and  $\otimes$ , with respective neutral elements  $0$  and  $1$ , and such that:

- $\oplus$  is commutative:  $\langle \mathbb{S}, \oplus, 0 \rangle$  is a commutative monoid and  $\langle \mathbb{S}, \otimes, 1 \rangle$  a monoid,
- $\otimes$  distributes over  $\oplus$ :  $\forall x, y, z \in \mathbb{S}$ ,  $x \otimes (y \oplus z) = (x \otimes y) \oplus (x \otimes z)$ , and  $(x \oplus y) \otimes z = (x \otimes z) \oplus (y \otimes z)$ ,
- $0$  is absorbing for  $\otimes$ :  $\forall x \in \mathbb{S}$ ,  $0 \otimes x = x \otimes 0 = 0$ .

Intuitively, in the models presented in this paper,  $\oplus$  selects an optimal value from two given values, in order to handle non-determinism, and  $\otimes$  combines two values into a single value, in a chaining of transitions.

A semiring  $\mathbb{S}$  is *commutative* if  $\otimes$  is commutative. It is *idempotent* if for each  $x \in \text{dom}(\mathbb{S})$ ,  $x \oplus x = x$ . Every idempotent semiring  $\mathbb{S}$  induces a partial ordering  $\leq_\oplus$  called the *natural ordering* of  $\mathbb{S}$  [20] and defined, by: for all  $x$  and  $y$ ,  $x \leq_\oplus y$  iff  $x \oplus y = x$ . The natural ordering is sometimes defined in the opposite direction [10]; We follow here the direction that coincides with the usual ordering on the Tropical semiring *min-plus* (Figure 2). An idempotent semiring  $\mathbb{S}$  is called *total* if it  $\leq_\oplus$  is total *i.e.* when for all  $x, y \in \mathbb{S}$ , either  $x \oplus y = x$  or  $x \oplus y = y$ .

► **Lemma 1** (Monotony, [20]). *Let  $\langle \mathbb{S}, \oplus, \otimes, 0, 1 \rangle$  be an idempotent semiring. For all  $x, y, z \in \mathbb{S}$ , if  $x \leq_\oplus y$  then  $x \oplus z \leq_\oplus y \oplus z$ ,  $x \otimes z \leq_\oplus y \otimes z$  and  $z \otimes x \leq_\oplus z \otimes y$ .*

When the property of Lemma 1 holds,  $\mathbb{S}$  is called *monotonic*. Another important semiring property in the context of optimization is superiority [16], which corresponds to the *non-negative weights* condition in shortest-path algorithms [9]. Intuitively, it means that combining elements with  $\otimes$  always increase their weight. Formally, it is defined as the property (i) below.

► **Lemma 2** (Superiority, Boundedness). *Let  $\langle \mathbb{S}, \oplus, \otimes, 0, 1 \rangle$  be an idempotent semiring. The two following statements are equivalent:*

- i. *for all  $x, y \in \mathbb{S}$ ,  $x \leq_\oplus x \otimes y$  and  $y \leq_\oplus x \otimes y$*
- ii. *for all  $x \in \mathbb{S}$ ,  $1 \oplus x = 1$ .*

**Proof.** (ii)  $\Rightarrow$  (i) :  $x \oplus (x \otimes y) = x \otimes (1 \oplus y) = x$ , by distributivity of  $\otimes$  over  $\oplus$ . Hence  $x \leq_\oplus x \otimes y$ . Similarly,  $y \oplus (x \otimes y) = (1 \oplus x) \otimes y = y$ , hence  $y \leq_\oplus x \otimes y$ . (i)  $\Rightarrow$  (ii) : by the second inequality of (i), with  $y = 1$ ,  $1 \leq_\oplus x \otimes 1 = x$ , *i.e.*, by definition of  $\leq_\oplus$ ,  $1 \oplus x = 1$ . ◀

The results are es-  
tablished for a gen-  
eral class of semir-  
ings. They can be  
instantiated for con-  
crete cases

	domain	$\oplus$	$\otimes$	$\mathbb{0}$	$\mathbb{1}$
Boolean	$\{\perp, \top\}$	$\vee$	$\wedge$	$\perp$	$\top$
Counting	$\mathbb{N}$	$+$	$\times$	$0$	$1$
Viterbi	$[0, 1] \subset \mathbb{R}$	$\max$	$\times$	$0$	$1$
Tropical min-plus	$\mathbb{R}_+ \cup \{\infty\}$	$\min$	$+$	$\infty$	$0$

■ **Figure 2** Some commutative, bounded, total and complete semirings.

In [16], when the property (i) holds,  $\mathbb{S}$  is called *superior wrt* the ordering  $\leq_\oplus$ . We have seen in the proof of Lemma 2 that it implies that  $\mathbb{1} \leq_\oplus x$  for all  $x \in \mathbb{S}$ . Similarly, by the first inequality of (i) with  $y = \mathbb{0}$ ,  $x \leq_\oplus x \otimes \mathbb{0} = \mathbb{0}$ . Hence, in a superior semiring, it holds that for all  $x \in \mathbb{S}$ ,  $\mathbb{1} \leq_\oplus x \leq_\oplus \mathbb{0}$ . Intuitively, from an optimization point of view, it means that  $\mathbb{1}$  is the best value, and  $\mathbb{0}$  the worst. In [20],  $\mathbb{S}$  with the property (ii) of Lemma 2 is called *bounded* – we shall use this term in the rest of the paper. It implies that, when looking for a best path in a graph whose edges are weighted by values of  $\mathbb{S}$ , the loops can be safely avoided (because, for all  $x \in \mathbb{S}$  and  $n \geq 1$ ,  $x \oplus x^n = x \otimes (\mathbb{1} \oplus x^{n-1}) = x$ ).

► **Lemma 3.** *Every bounded semiring is idempotent.*

**Proof.** By boundedness,  $\mathbb{1} \oplus \mathbb{1} = \mathbb{1}$ , and idempotency follows by multiplying both sides by  $x$  and distributing. ◀

We shall need below infinite sums with  $\oplus$ . A semiring  $\mathbb{S}$  is called *complete* [11] if it has an operation  $\bigoplus_{i \in I} x_i$  for every family  $(x_i)_{i \in I}$  of elements of  $\text{dom}(\mathbb{S})$  over an index set  $I \subset \mathbb{N}$ , such that:

i. *infinite sums extend finite sums:*

$$\bigoplus_{i \in \emptyset} x_i = \mathbb{0}, \quad \forall j \in \mathbb{N}, \bigoplus_{i \in \{j\}} x_i = x_j, \quad \forall j, k \in \mathbb{N}, j \neq k, \bigoplus_{i \in \{j, k\}} x_i = x_j \oplus x_k,$$

ii. *associativity and commutativity:*

$$\text{for all } I \subseteq \mathbb{N} \text{ and all partition } (I_j)_{j \in J} \text{ of } I, \bigoplus_{j \in J} \bigoplus_{i \in I_j} x_i = \bigoplus_{i \in I} x_i,$$

iii. *distributivity of product over infinite sum:*

$$\text{for all } I \subseteq \mathbb{N}, \bigoplus_{i \in I} (x \otimes y_i) = x \otimes \bigoplus_{i \in I} y_i, \text{ and } \bigoplus_{i \in I} (x_i \otimes y) = (\bigoplus_{i \in I} x_i) \otimes y.$$

## Label Theory

We shall now define the functions labeling the transitions of SW automata and transducers, generalizing the Boolean algebras of [7] from Boolean to other semiring domains. We consider *alphabets*, which are countable sets of symbols denoted  $\Sigma, \Delta, \dots$ . Given a semiring  $\langle \mathbb{S}, \oplus, \otimes, \mathbb{0}, \mathbb{1} \rangle$ , a *label theory* over  $\mathbb{S}$  is a set  $\bar{\Phi}$  of recursively enumerable sets denoted  $\Phi_\Sigma$ , containing unary functions of type  $\Sigma \rightarrow \mathbb{S}$ , or  $\Phi_{\Sigma, \Delta}$ , containing binary functions  $\Sigma \times \Delta \rightarrow \mathbb{S}$ , and such that:

- for all  $\Phi_{\Sigma, \Delta} \in \bar{\Phi}$ , we have  $\Phi_\Sigma \in \bar{\Phi}$  and  $\Phi_\Delta \in \bar{\Phi}$
- every  $\Phi_\Sigma \in \bar{\Phi}$  contains all the constant functions from  $\Sigma$  into  $\mathbb{S}$ ,
- for all  $\alpha \in \mathbb{S}$  and  $\phi \in \Phi_\Sigma$ ,  $\alpha \otimes \phi : x \mapsto \alpha \otimes \phi(x)$ , and  $\phi \otimes \alpha : x \mapsto \phi(x) \otimes \alpha$  belong to  $\Phi_\Sigma$ , and similarly for  $\oplus$  and for  $\Phi_{\Sigma, \Delta}$
- for all  $\phi, \phi' \in \Phi_\Sigma$ ,  $\phi \otimes \phi' : x \mapsto \phi(x) \otimes \phi'(x)$  belongs to  $\Phi_\Sigma$
- for all  $\eta, \eta' \in \Phi_{\Sigma, \Delta}$ ,  $\eta \otimes \eta' : x, y \mapsto \eta(x, y) \otimes \eta'(x, y)$  belongs to  $\Phi_{\Sigma, \Delta}$

- 154 – for all  $\phi \in \Phi_\Sigma$  and  $\eta \in \Phi_{\Sigma,\Delta}$ ,  $\phi \otimes_1 \eta : x, y \mapsto \phi(x) \otimes \eta(x, y)$  and
- 155  $\eta \otimes_1 \phi : x, y \mapsto \eta(x, y) \otimes \phi(x)$  belong to  $\Phi_{\Sigma,\Delta}$
- 156 – for all  $\psi \in \Phi_\Delta$  and  $\eta \in \Phi_{\Sigma,\Delta}$ ,  $\psi \otimes_2 \eta : x, y \mapsto \psi(y) \otimes \eta(x, y)$  and
- 157  $\eta \otimes_2 \psi : x, y \mapsto \eta(x, y) \otimes \psi(y)$  belong to  $\Phi_{\Sigma,\Delta}$
- 158 – similar closures hold for  $\oplus$ .

partial application is needed?

160 In what follows, we might omit the subscripts in  $\otimes_1$ ,  $\otimes_2$ ,  $\oplus_1$ ,  $\oplus_2$  when there is no ambiguity,  
 161 and keep them only for the special case  $\Sigma = \Delta$ , i.e.  $\eta \in \Phi_{\Sigma,\Sigma}$ . When the semiring  $\mathbb{S}$  is  
 162 complete, let us consider the following operators on the functions of a label theory.

$$\begin{aligned} \bigoplus_\Sigma : \Phi_\Sigma &\rightarrow \mathbb{S}, \phi \mapsto \bigoplus_{a \in \Sigma} \phi(a) \\ \bigoplus_\Sigma^1 : \Phi_{\Sigma,\Delta} &\rightarrow \Phi_\Delta, \eta \mapsto (y \mapsto \bigoplus_{a \in \Sigma} \eta(a, y)) \quad \bigoplus_\Delta^2 : \Phi_{\Sigma,\Delta} \rightarrow \Phi_\Sigma, \eta \mapsto (x \mapsto \bigoplus_{b \in \Delta} \eta(x, b)) \end{aligned}$$

164 Similarly as for the above product and sum of functions, the superscripts in  $\bigoplus_\Sigma^1$  and  $\bigoplus_\Sigma^2$   
 165 shall be reserved to the ambiguous case of  $\Phi_{\Sigma,\Sigma}$ , in order to distinguish between the first  
 166 and the second argument.

167 ► **Definition 4.** A label theory  $\bar{\Phi}$  is complete when its underlying semiring  $\mathbb{S}$  is complete,  
 168 and for all  $\Phi_{\Sigma,\Delta} \in \bar{\Phi}$  and all  $\eta \in \Phi_{\Sigma,\Delta}$ ,  $\bigoplus_\Sigma \eta \in \Phi_\Delta$  and  $\bigoplus_\Delta \eta \in \Phi_\Sigma$ .

notion of diagram of functions akin BDD for transitions in practice

170 The following facts are immediate.

mv appendix?

171 ► **Lemma 5.** For  $\bar{\Phi}$  complete  $\alpha \in \mathbb{S}$ ,  $\phi, \phi' \in \Phi_\Sigma$ ,  $\psi \in \Phi_\Delta$ , and  $\eta \in \Phi_{\Sigma,\Delta}$ :

- 172 i.  $\bigoplus_\Sigma \bigoplus_\Delta^2 \eta = \bigoplus_\Delta \bigoplus_\Sigma^1 \eta$
- 173 ii.  $\alpha \otimes \bigoplus_\Sigma \phi = \bigoplus_\Sigma (\alpha \otimes \phi)$  and  $(\bigoplus_\Sigma \phi) \otimes \alpha = \bigoplus_\Sigma (\phi \otimes \alpha)$ , and similarly for  $\oplus$
- 174 iii.  $(\bigoplus_\Sigma \phi) \oplus (\bigoplus_\Sigma \phi') = \bigoplus_\Sigma (\phi \oplus \phi')$  and  $(\bigoplus_\Sigma \phi) \otimes (\bigoplus_\Sigma \phi') = \bigoplus_\Sigma (\phi \otimes \phi')$
- 175 iv.  $(\bigoplus_\Delta^2 \eta) \oplus (\bigoplus_\Delta^2 \eta') = \bigoplus_\Delta^2 (\eta \oplus \eta')$ , and  $(\bigoplus_\Delta^2 \eta) \otimes (\bigoplus_\Delta^2 \eta') = \bigoplus_\Delta^2 (\eta \otimes \eta')$
- 176 v.  $\phi \otimes (\bigoplus_\Delta^2 \eta) = \bigoplus_\Delta (\phi \otimes_1 \eta)$ , and  $(\bigoplus_\Delta^2 \eta) \otimes \phi = \bigoplus_\Delta (\eta \otimes_1 \phi)$ , and similarly for  $\oplus$
- 177 vi.  $\psi \otimes (\bigoplus_\Sigma^1 \eta) = \bigoplus_\Sigma (\psi \otimes_2 \eta)$ , and  $(\bigoplus_\Sigma^1 \eta) \otimes \psi = \bigoplus_\Sigma (\eta \otimes_2 \psi)$ , and similarly for  $\oplus$

178 Intuitively, the operators  $\bigoplus_\Sigma$  return global minimum, wrt  $\leq_\oplus$ , of functions of  $\bar{\Phi}$ . A label  
 179 theory is called *effective* when for all  $\phi \in \Phi_\Sigma$  and  $\eta \in \Phi_{\Sigma,\Delta}$ ,  $\bigoplus_\Sigma \phi$ ,  $\bigoplus_\Sigma \eta$ , and  $\bigoplus_\Delta \eta$  can be  
 180 effectively computed from  $\phi$  and  $\eta$ .

precise/restrict complexity

### 3 SW Automata and Transducers

182 We follow the approach of [21] for the computation of distances, between words and languages,  
 183 using weighted transducers, and extend it to infinite alphabets. The models introduced in  
 184 this section generalize weighted automata and transducers [11] by labeling each transition  
 185 with a weight function (instead of a simple weight value), that takes the input and output  
 186 symbols as parameters. These functions are similar to the guards of symbolic automata [7, 8],  
 187 but they can return values in a generic semiring, whereas the latter guards are restricted to  
 188 the Boolean semiring.

189 Let  $\bar{\mathbb{S}}$  be a commutative semiring,  $\Sigma$  and  $\Delta$  be alphabets called respectively *input* and *output*,  
 190 and  $\bar{\Phi}$  be a label theory over  $\mathbb{S}$  containing  $\Phi_\Sigma$ ,  $\Phi_\Delta$ ,  $\Phi_{\Sigma,\Delta}$ .

191 ► **Definition 6.** A symbolic-weighted transducer (*swT*) over  $\Sigma$ ,  $\Delta$ ,  $\mathbb{S}$  and  $\bar{\Phi}$  is a tuple  
 192  $T = \langle Q, \text{in}, \bar{w}, \text{out} \rangle$ , where  $Q$  is a finite set of states,  $\text{in} : Q \rightarrow \mathbb{S}$  (respectively  $\text{out} : Q \rightarrow \mathbb{S}$ )  
 193 are functions defining the weight for entering (respectively leaving) computation in a state,  
 194 and  $\bar{w}$  is a triplet of transition functions  $w_{10} : Q \times Q \rightarrow \Phi_\Sigma$ ,  $w_{01} : Q \times Q \rightarrow \Phi_\Delta$ , and  
 195  $w_{11} : Q \times Q \rightarrow \Phi_{\Sigma,\Delta}$ .

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We call *number of transitions* of  $T$  the number of pairs of states  $q, q' \in Q$  such that  $w_{10}$  or  $w_{01}$  or  $w_{11}$  is not the constant  $\mathbb{0}$ . For convenience, we shall sometimes present transitions as functions of  $Q \times (\Sigma \cup \{\varepsilon\}) \times (\Delta \cup \{\varepsilon\}) \times Q \rightarrow \mathbb{S}$ , overloading the function names, such that, for all  $q, q' \in Q$ ,  $a \in \Sigma$ ,  $b \in \Delta$ ,

$$\begin{aligned} w_{10}(q, a, \varepsilon, q') &= \phi(a) & \text{where } \phi = w_{10}(q, q') \in \Phi_\Sigma, \\ w_{01}(q, \varepsilon, b, q') &= \psi(b) & \text{where } \psi = w_{01}(q, q') \in \Phi_\Delta, \\ w_{11}(q, a, b, q') &= \eta(a, b) & \text{where } \eta = w_{11}(q, q') \in \Phi_{\Sigma, \Delta}. \end{aligned}$$

The  $\text{swT } T$  computes on pairs of words  $\langle s, t \rangle \in \Sigma^* \times \Delta^*$ ,  $s$  and  $t$ , being respectively called *input* and *output* word. More precisely,  $T$  defines a mapping from  $\Sigma^* \times \Delta^*$  into  $\mathbb{S}$ , based on an intermediate function  $\text{weight}_T$  defined recursively, for every states  $q, q' \in Q$ , and every strings  $\langle s, t \rangle \in \Sigma^* \times \Delta^* \setminus \{\langle \varepsilon, \varepsilon \rangle\}$  considered as concatenation of the symbol  $a \in \Sigma$  (resp.  $b \in \Delta$ ) with a word  $u \in \Sigma^*$  (resp.  $v \in \Delta^*$ ), by :

$$\begin{aligned} \text{weight}_T(q, \varepsilon, \varepsilon, q') &= \mathbb{1} \quad \text{if } q = q' \text{ and } \mathbb{0} \text{ otherwise} \\ \text{weight}_T(q, s, t, q') &= \bigoplus_{\substack{q'' \in Q \\ s=au, a \in \Sigma}} w_{10}(q, a, \varepsilon, q'') \otimes \text{weight}_T(q'', u, t, q') \\ &\quad \oplus \bigoplus_{\substack{q'' \in Q \\ t=bv, b \in \Delta}} w_{01}(q, \varepsilon, b, q'') \otimes \text{weight}_T(q'', s, v, q') \\ &\quad \oplus \bigoplus_{\substack{q'' \in Q \\ s=au, t=bv}} w_{11}(q, a, b, q'') \otimes \text{weight}_T(q'', u, v, q') \end{aligned} \tag{1}$$

We recall that, by convention (Section 2), an empty sum with  $\bigoplus$  is equal to  $\mathbb{0}$ . Intuitively, using a transition  $w_{ij}(q, a, b, q')$  means for  $T$ : when reading respectively  $a$  and  $b$  at the current positions in the input and output words, increment the current position in the input word if and only if  $i = 1$ , and in the output word iff  $j = 1$ , and change state from  $q$  to  $q'$ . When  $a = \varepsilon$  (resp.  $b = \varepsilon$ ), the current symbol in the input (resp. output) is not read. Since  $\mathbb{0}$  is absorbing for  $\otimes$  in  $\mathbb{S}$ , one term  $w_{ij}(q, a, b, q'')$  equal to  $\mathbb{0}$  in the above expression will be ignored in the sum, meaning that there is no possible transition from state  $q$  into state  $q'$  while reading  $a$  and  $b$ . This is analogous to the case of a transition's guard not satisfied by  $\langle a, b \rangle$  for symbolic transducers.

The expression (1) can be seen as a stateful definition of an edit-distance between a word  $s \in \Sigma^*$  and a word  $t \in \Delta^*$ , see also [22]. Intuitively,  $w_{10}(q, a, \varepsilon, r)$  is the cost of the deletion of the symbol  $a \in \Sigma$  in  $s$ ,  $w_{01}(q, \varepsilon, b, r)$  is the cost of the insertion of  $b \in \Delta$  in  $t$ , and  $w_{11}(q, a, b, r)$  is the cost of the substitution of  $a \in \Sigma$  by  $b \in \Delta$ . The cost of a sequence of such operations transforming  $s$  into  $t$ , is the product, with  $\otimes$ , of the individual costs of the operations involved; and the distance between  $s$  and  $t$  is the sum, with  $\oplus$ , of all possible products. Formally, the weight associated by  $T$  to  $\langle s, t \rangle \in \Sigma^* \times \Delta^*$  is:

$$T(s, t) = \bigoplus_{q, q' \in Q} \text{in}(q) \otimes \text{weight}_T(q, s, t, q') \otimes \text{out}(q') \tag{2}$$

► **Example 7.** In Common Western Music Notation [14], several symbols may be used to represent one single sounding event. For instance, several notes can be combined with a tie,

unique → similar

similar → single

like in  $\text{♩}$ , and one note can be augmented by half its duration with a dot like in  $\text{♩.}$ . These notations are perceived equivalent when played, as their duration is equal, yet the notation is different. We thus want to be able to recognize all the possible notations of equivalent durations of notes, when comparing a written score and a music performance.

We propose a small weighted transducer model that computes a distance between an input sequence of sounding events (music "performance") to an output sequence of written events (music "score"). Let us consider the tropical (*min-plus*) semiring  $\mathbb{S}$  of Figure 2 and let  $\Sigma = \mathbb{R}_+$  be an input alphabet of event dates and  $\Delta = \{\mathbf{e}, -\} \times \mathbb{R}_+$  be an output alphabet of symbols with timestamps. A symbol  $\langle \mathbf{e}, d \rangle \in \Delta$  represents an event starting at date  $d$ , and  $\langle -, d \rangle$  is a continuation of the previous event.

We consider a **swT** with two states  $q_0$  and  $q_1$  whose purpose is to compare a recorded performance  $s \in \Sigma^*$  with a notated music sheet  $t \in \Delta^*$ . One timestamp  $d_i \in \Sigma$  may correspond to one notated event  $\langle \mathbf{e}, d'_i \rangle \in \Delta$ , in which case the weight value computed by the **swT** is the time distance between both (see transitions  $w_{11}$  below). If  $\langle \mathbf{e}, d'_i \rangle$  is followed by continuations  $\langle -, d'_{i+1} \rangle \dots$ , they are just skipped with no cost (transitions  $w_{01}$  or weight 1).

$$\begin{array}{llll} w_{11}(q_0, d, \langle \mathbf{e}, d' \rangle, q_0) & = & |d' - d| & w_{11}(q_1, d, \langle \mathbf{e}, d' \rangle, q_0) & = & |d' - d| \\ w_{01}(q_0, \varepsilon, \langle -, d' \rangle, q_0) & = & 1 & w_{01}(q_1, \varepsilon, \langle -, d' \rangle, q_0) & = & 1 \\ w_{10}(q_0, d, \varepsilon, q_1) & = & \alpha & & & \end{array}$$

We also must be able to take performing errors into account, while still being able to compare with the score, since a performer could, for example, play an unwritten extra note. This is modelled by the transition  $w_{10}$  with an arbitrary weight value  $\alpha \in \mathbb{S}$ , switching from state  $q_0$  (normal) to  $q_1$  (error). The transitions in the second column below switch back to the normal state  $q_0$ . At last, we let  $q_0$  be the only initial and final state, with  $\text{in}(q_0) = \text{out}(q_0) = 1$ , and  $\text{in}(q_1) = \text{out}(q_1) = 0$ .  $\diamond$

The *Symbolic Weighted Automata* are defined similarly as the transducers of Definition 6, by simply omitting the output symbols.

► **Definition 8.** A symbolic-weighted automaton (*swA*) over  $\Sigma$ ,  $\mathbb{S}$  and  $\bar{\Phi}$  is a tuple  $A = \langle Q, \text{in}, w_1, \text{out} \rangle$ , where  $Q$  is a finite set of states,  $\text{in} : Q \rightarrow \mathbb{S}$  (respectively  $\text{out} : Q \rightarrow \mathbb{S}$ ) are functions defining the weight for entering (respectively leaving) computation in a state, and  $w_1$  is a transition functions from  $Q \times Q$  into  $\Phi_\Sigma$ .

As above in the case of **swT**, when  $w_1(q, q') = \phi \in \Phi_\Sigma$ , we may write  $w_1(q, a, q')$  for  $\phi(a)$ . The computation of  $A$  on words  $s \in \Sigma^*$  is defined with an intermediate function  $\text{weight}_A$ , defined as follows for  $q, q' \in Q$ ,  $a \in \Sigma$ ,  $u \in \Sigma^*$ ,

$$\begin{array}{ll} \text{weight}_A(q, \varepsilon, q) & = 1 \\ \text{weight}_A(q, \varepsilon, q') & = 0 \quad \text{if } q \neq q' \\ \text{weight}_A(q, au, q') & = \bigoplus_{q'' \in Q} w_1(q, a, q'') \otimes \text{weight}_A(q'', u, q') \end{array} \tag{3}$$

and the weight value associated by  $A$  to  $s \in \Sigma^*$  is defined as follows:

$$A(s) = \bigoplus_{q, q' \in Q} \text{in}(q) \otimes \text{weight}_A(q, s, q') \otimes \text{out}(q') \tag{4}$$

The following property will be useful to the approach on symbolic weighted parsing presented in Section 5.



269 ► **Proposition 9.** *Given a swT  $T$  over  $\Sigma$ ,  $\Delta$ ,  $\mathbb{S}$  commutative, bounded and complete, and  $\bar{\Phi}$*   
 270 *effective, and a swA  $A$  over  $\Sigma$ ,  $\mathbb{S}$  and  $\bar{\Phi}$ , there exists an effectively constructible swA  $B_{T,A}$*   
 271 *over  $\Delta$ ,  $\mathbb{S}$  and  $\bar{\Phi}$ , such that for all  $t \in \Delta^*$ ,  $B_{T,A}(t) = \bigoplus_{s \in \Sigma^*} A(s) \otimes T(s, t)$ .*

272 **Proof.** Let  $T = \langle Q, \text{in}_T, \bar{w}, \text{out}_T \rangle$ , where  $\bar{w}$  contains  $w_{10}$ ,  $w_{01}$ , and  $w_{11}$ , from  $Q \times Q$  into  
 273 respectively  $\Phi_\Sigma$ ,  $\Phi_\Delta$ , and  $\Phi_{\Sigma, \Delta}$ , and let  $A = \langle P, \text{in}_A, w_1, \text{out}_A \rangle$  with  $w_1 : Q \times Q \rightarrow \Phi_\Sigma$ . The  
 274 state set of  $B_{T,A}$  will be  $Q' = P \times Q$ . The entering, leaving and transition functions of  $B_{T,A}$   
 275 will simulate synchronized computations of  $A$  and  $T$ , while reading an output word of  $\Delta^*$ .  
 276 Its state entering functions is defined for all  $p \in P$ ,  $q \in Q$  by  $\text{in}'(p, q) = \text{in}_A(p) \otimes \text{in}_T(q)$ . The  
 277 transition function  $w'_1$  will roughly perform a synchronized product of transitions defined by  
 278  $w_1$ ,  $w_{01}$  ( $T$  reading in output word and not in input word) and  $w_{11}$  ( $T$  reading in output  
 279 word and input word). Moreover,  $w'_1$  also needs to simulate transitions defined by  $w_{10}$ :  $T$   
 280 reading in input word and not in output word. Since  $B_{T,A}$  will read only in the output  
 281 word, such a transition corresponds to an  $\varepsilon$ -transition of swA, but swA have been defined  
 282 without  $\varepsilon$ -transitions. Therefore, in order to take care of this case, we perform an on-the-fly  
 283 suppression of  $\varepsilon$ -transition in the swA in construction, following the algorithm of [19].  
 284 Initially, for all  $p_1, p_2 \in P$ , and  $q_1, q_2 \in Q$ , let

$$285 \quad w'_1(\langle p_1, q_1 \rangle, \langle p_2, q_2 \rangle) = w_1(p_1, p_2) \otimes [w_{01}(q_1, q_2) \oplus \bigoplus_{\Sigma} w_{11}(q_1, q_2)].$$

286 Iterate the following for all  $p_1 \in P$  and  $q_1, q_2 \in Q$ : for all  $p_2 \in P$  and  $q_3 \in Q$ ,

$$287 \quad w'_1(\langle p_1, q_1 \rangle, \langle p_2, q_3 \rangle) \oplus = \bigoplus_{\Sigma} w_{10}(q_1, q_2) \otimes w'_1(\langle p_1, q_2 \rangle, \langle p_2, q_3 \rangle)$$

proof correctness

$$288 \quad \text{and } \text{out}'(p_1, q_1) \oplus = \bigoplus_{\Sigma} w_{10}(q_1, q_2) \otimes \text{out}'(p_1, q_2) \quad \blacktriangleleft$$

289 The construction time and size for  $B_{T,A}$  are  $O(\|T\|^3 \cdot \|A\|^2)$ , where the sizes  $\|T\|$  and  $\|A\|$   
 290 are their number of states.

revise with nb of tr. 290  
and states

291 ► **Corollary 10.** *Given a swT  $T$  over  $\Sigma$ ,  $\Delta$ ,  $\mathbb{S}$  commutative, bounded and complete, and  $\bar{\Phi}$*   
 292 *effective, and  $s \in \Sigma^+$ , there exists an effectively constructible swA  $B_{T,s}$  over  $\Delta$ ,  $\mathbb{S}$  and  $\bar{\Phi}$ ,*  
 293 *such that for all  $t \in \Delta^*$ ,  $B_{T,s}(t) = T(s, t)$ .*

## 294 4 SW Visibly Pushdown Automata

295 The model presented in this section generalizes Symbolic VPA [6] from Boolean semirings to  
 296 arbitrary semiring weight domains. It will compute on nested words over infinite alphabets,  
 297 associating to every such word a weight value. Nested words are able to describe structures  
 298 of labeled trees, and in the context of parsing, they will be useful to represent AST.

299 Let  $\Omega$  be a countable alphabet that we assume partitioned into three subsets  $\Omega_i$ ,  $\Omega_c$ ,  $\Omega_r$ ,  
 300 whose elements are respectively called *internal*, *call* and *return* symbols. Let  $\langle \mathbb{S}, \oplus, \otimes, \mathbb{1} \rangle$   
 301 be a commutative and complete semiring and let  $\bar{\Phi} = \langle \Phi_i, \Phi_c, \Phi_r, \Phi_{ci}, \Phi_{cc}, \Phi_{cr} \rangle$  be a label  
 302 theory over  $\mathbb{S}$  where  $\Phi_i$ ,  $\Phi_c$ ,  $\Phi_r$  and  $\Phi_{cx}$  (with  $x \in \{i, c, r\}$ ) stand respectively for  $\Phi_{\Omega_i}$ ,  $\Phi_{\Omega_c}$ ,  
 303  $\Phi_{\Omega_r}$  and  $\Phi_{\Omega_c, \Omega_x}$ .

304 ► **Definition 11.** *A Symbolic Weighted Visibly Pushdown Automata (sw-VPA) over  $\Omega =$*   
 305  *$\Omega_i \uplus \Omega_c \uplus \Omega_r$ ,  $\mathbb{S}$  and  $\bar{\Phi}$  is a tuple  $A = \langle Q, P, \text{in}, \bar{w}, \text{out} \rangle$ , where  $Q$  is a finite set of states,  $P$*   
 306 *is a finite set of stack symbols,  $\text{in} : Q \rightarrow \mathbb{S}$  (respectively  $\text{out} : Q \rightarrow \mathbb{S}$ ) are functions defining*



the weight for entering (respectively leaving) a state, and  $\bar{w}$  is a sextuplet composed of the transition functions :  $w_i : Q \times P \times Q \rightarrow \Phi_{ci}$ ,  $w_i^e : Q \times Q \rightarrow \Phi_i$ ,  $w_c : Q \times P \times Q \times P \rightarrow \Phi_{cc}$ ,  $w_c^e : Q \times P \times Q \rightarrow \Phi_c$ ,  $w_r : Q \times P \times Q \rightarrow \Phi_{cr}$ ,  $w_r^e : Q \times Q \rightarrow \Phi_r$ .

Similarly as in Section 3, we extend the above transition functions as follows for all  $q, q' \in Q$ ,  $p \in P$ ,  $a \in \Omega_i$ ,  $c \in \Omega_c$ ,  $r \in \Omega_r$ , overloading their names:

$$\begin{array}{lll}
 w_i : Q \times \Omega_c \times P \times \Omega_i \times Q \rightarrow \mathbb{S} & w_i(q, c, p, a, q') = \eta_{ci}(c, a) & \text{where } \eta_{ci} = w_i(q, p, q'), \\
 w_i^e : Q \times \Omega_i \times Q \rightarrow \mathbb{S} & w_i^e(q, a, q') = \phi_i(a) & \text{where } \phi_i = w_i^e(q, q'), \\
 w_c : Q \times \Omega_c \times P \times \Omega_c \times P \times Q \rightarrow \mathbb{S} & w_c(q, c, p, c', p', q') = \eta_{cc}(c, c') & \text{where } \eta_{cc} = w_c(q, p, p', q'), \\
 w_c^e : Q \times \Omega_c \times P \times Q \rightarrow \mathbb{S} & w_c^e(q, c, p, q') = \phi_c(c) & \text{where } \phi_c = w_c^e(q, p, q'), \\
 w_r : Q \times \Omega_c \times P \times \Omega_r \times Q \rightarrow \mathbb{S} & w_r(q, c, p, r, q') = \eta_{cr}(c, r) & \text{where } \eta_{cr} = w_r(q, p, q'), \\
 w_r^e : Q \times \Omega_r \times Q \rightarrow \mathbb{S} & w_r^e(q, r, q') = \phi_r(r) & \text{where } \phi_r = w_r^e(q, q').
 \end{array}$$

The intuition is the following for the above transitions.  $w_i^e$ ,  $w_c^e$ ,  $w_r^e$  describe the cases where the stack is empty.  $w_i$  and  $w_i^e$  both read an input internal symbol  $a$  and change state from  $q$  to  $q'$ , without changing the stack. Moreover,  $w_i$  reads a pair made of  $c \in \Omega_c$  and  $p \in P$  on the top of the stack ( $c$  is compared to  $a$  by the weight function  $\eta_{ci} \in \Phi_{ci}$ ).  $w_c$  and  $w_c^e$  read the input call symbol  $c'$ , push it to the stack along with  $p'$ , and change state from  $q$  to  $q'$ . Moreover,  $w_c$  reads  $c$  and  $p$  at the top of the stack ( $c$  is compared to  $c'$ ).  $w_r$  and  $w_r^e$  read the input return symbol  $r$ , and change state from  $q$  to  $q'$ . Moreover,  $w_r$  reads and pop from stack a pair made of  $c$  and  $p$ , ( $c$  is compared to  $r$ ).

Formally, the transitions of the automaton  $A$  are defined in term of an intermediate function  $\text{weight}_A$ , like in Section 3. A configuration, denoted by  $q[\gamma]$ , is here composed of a state  $q \in Q$  and a stack content  $\gamma \in \Gamma^*$ , where  $\Gamma = \Omega_c \times P$ . Hence,  $\text{weight}_A$  is a function from  $[Q \times \Gamma^*] \times \Omega^* \times [Q \times \Gamma^*]$  into  $\mathbb{S}$ . The empty stack is denoted by  $\perp$ , and the upmost symbol is the last pushed content. The following functions illustrate each of the possible cases, being : reading  $a \in \Omega_i$ , or  $c \in \Omega_c$ , or  $r \in \Omega_r$  for each possible state of the stack (empty or not), to add to  $u \in \Omega^*$ .

moved this to the beginning

intro to func

notation  $cp$  for  $\langle c, p \rangle$  ?

$$\begin{aligned}
 \text{weight}_A(q[\perp], \varepsilon, q'[\perp]) &= \mathbb{1} \text{ if } q = q' \text{ and } \mathbb{0} \text{ otherwise} & (5) \\
 \text{weight}_A\left(q \left[ \begin{array}{c} \langle c, p \rangle \\ \gamma \end{array} \right], a u, q'[\gamma']\right) &= \bigoplus_{q'' \in Q} w_i(q, c, p, a, q'') \otimes \text{weight}_A(q'' \left[ \begin{array}{c} \langle c, p \rangle \\ \gamma \end{array} \right], u, q'[\gamma']) \\
 \text{weight}_A(q[\perp], a u, q'[\gamma']) &= \bigoplus_{q'' \in Q} w_i^e(q, a, q'') \otimes \text{weight}_A(q''[\perp], u, q'[\gamma']) \\
 \text{weight}_A\left(q \left[ \begin{array}{c} \langle c, p \rangle \\ \gamma \end{array} \right], c' u, q'[\gamma']\right) &= \bigoplus_{\substack{q'' \in Q \\ p' \in P}} w_c(q, c, p, c', p', q'') \otimes \text{weight}_A(q'' \left[ \begin{array}{c} \langle c', p' \rangle \\ \langle c, p \rangle \\ \gamma \end{array} \right], u, q'[\gamma']) \\
 \text{weight}_A(q[\perp], c u, q'[\gamma']) &= \bigoplus_{\substack{q'' \in Q \\ p \in P}} w_c^e(q, c, p, q'') \otimes \text{weight}_A(q''[cp], u, q'[\gamma']) \\
 \text{weight}_A\left(q \left[ \begin{array}{c} \langle c, p \rangle \\ \gamma \end{array} \right], r u, q'[\gamma']\right) &= \bigoplus_{q'' \in Q} w_r(q, c, p, r, q'') \otimes \text{weight}_A(q''[\gamma], u, q'[\gamma']) \\
 \text{weight}_A(q[\perp], r u, q'[\gamma']) &= \bigoplus_{q'' \in Q} w_r^e(q, r, q'') \otimes \text{weight}_A(q''[\perp], u, q'[\gamma'])
 \end{aligned}$$

## XX:10 Symbolic Weighted Language Models and Parsing over Infinite Alphabets

337 The weight associated by  $A$  to  $s \in \Omega^*$  is defined according to empty stack semantics:

$$338 \quad A(s) = \bigoplus_{q, q' \in Q} \text{in}(q) \otimes \text{weight}_A(q[\perp], s, q'[\perp]) \otimes \text{out}(q'). \quad (6)$$

todo example VPA 339

340 **► Example 12.** structured words with timed symbols... intro language of music notation?  
(markup = time division, leaves = events etc)

341 Every swA  $A = \langle Q, \text{in}, w_1, \text{out} \rangle$ , over  $\Sigma, \mathbb{S}$  and  $\bar{\Phi}$  is a particular case of sw-VPA  $\langle Q, \emptyset, \text{in}, \bar{w}, \text{out} \rangle$   
342 over  $\Omega, \mathbb{S}$  and  $\bar{\Phi}$  with  $\Omega_i = \Sigma$  and  $\Omega_c = \Omega_r = \emptyset$ , and computing with an always empty stack:  
343  $w_i^e = w_1$  and all the other functions of  $\bar{w}$  are the constant  $\emptyset$ .

344 Like VPA and symbolic VPA, the class of sw-VPA is closed under the binary operators of  
345 the underlying semiring.

346 **► Proposition 13.** Let  $A_1$  and  $A_2$  be two sw-VPA over the same  $\Omega, \mathbb{S}$  and  $\bar{\Phi}$ . There  
347 exists two effectively constructible sw-VPA  $A_1 \oplus A_2$  and  $A_1 \otimes A_2$ , such that for all  $s \in \Omega^*$ ,  
348  $(A_1 \oplus A_2)(s) = A_1(s) \oplus A_2(s)$  and  $(A_1 \otimes A_2)(s) = A_1(s) \otimes A_2(s)$ .

349 **Proof.** The construction is essentially the same as in the case of the Boolean semiring [6]. ◀

total? 350

351 Let us assume that the semiring  $\mathbb{S}$  is commutative, bounded, and complete, and that  $\bar{\Phi}$  is an  
352 effective label theory. We propose a Dijkstra algorithm computing, for a sw-VPA  $A$  over  $\Omega$ ,  
353  $\mathbb{S}$  and  $\bar{\Phi}$ , the minimal weight for a word in  $\Omega^*$ . We distinguish two cases: when the stack is  
empty, and when it is not. In the case of an empty stack, let  $b_\perp : Q \times Q \rightarrow \mathbb{S}$  be such that:

$$354 \quad b_\perp(q, q') = \bigoplus_{s \in \Omega^*} \text{weight}_A(q[\perp], s, q'[\perp]). \quad (7)$$

355 Since  $\mathbb{S}$  is complete, the infinite sum in (7) is well defined, and, providing that  $\mathbb{S}$  is total,  
356 it is the minimum in  $\Omega^*$ , wrt  $\leq_\oplus$ , of the function  $s \mapsto \text{weight}_A(q[\sigma], s, q'[\sigma])$ . The term  
357  $q[\perp], s, q'[\perp]$  of this sum is the central expression in the definition (6) of  $A(s_0)$ , for the  
358 minimum  $s_0$  of the function  $\text{weight}_A$ .

$b_\top$  : mot bien par- 359  
enthésé c/r 360

360 If the stack is not empty, let  $\top$  be a fresh stack symbol which does not belong to  $\Gamma$ , and let  
 $b_\top : Q \times P \times Q \rightarrow \Phi_c$  be such that, for every two states  $q, q' \in Q$  and stack symbol  $p \in P$ :

$$361 \quad b_\top(q, p, q') : c \mapsto \bigoplus_{s \in \Omega^*} \text{weight}_A\left(q \left[ \begin{array}{c} \langle c, p \rangle \\ \top \end{array} \right], s, q' \left[ \begin{array}{c} \langle c, p \rangle \\ \top \end{array} \right] \right) \quad (8)$$

362 Intuitively, the function defined in (8) associates to  $c \in \Omega_c$  the minimum weight of a  
363 computation of  $A$  starting in state  $q$  with a stack  $\langle c, p \rangle \cdot \gamma \in \Gamma^+$  and ending in state  $q'$  with  
364 the same stack, such that the computation can not pop the pair made of  $c$  and  $p$  at the top  
365 of this stack, but may only read these symbols. Moreover,  $A$  may push another pair  $\langle c', p' \rangle$   
366 on the top of  $\langle c, p \rangle \cdot \gamma$ , following the third case of in the definition (5) of  $\text{weight}_A$ , and may  
367 pop  $\langle c', p' \rangle$  later, following the fifth case of (5) (return symbol).

368 Algorithm 1 constructs iteratively markings  $d_\perp : Q \times Q \rightarrow \mathbb{S}$  and  $d_\top : Q \times P \times Q \rightarrow \Phi_c$   
369 that converges eventually to  $b_\top$  and  $b_\perp$ .

explication Fig. 3 370  
suivant cas de (5) 371

complete \*\* 372

detail with nb tr. 373  
and states 374

375 The infinite sums in the updates of  $d$  in Algorithm 1, Figure 3 are well defined since  $\mathbb{S}$   
is complete. \*\* effectively computable by hypothesis that the label theory is effective\*\*  
The algorithm performs  $2 \cdot |Q|^2$  iterations until  $P$  is empty, and each iteration has a time  
complexity  $O(|Q|^2 \cdot |P|)$ . That gives a time complexity  $O(|Q|^4 \cdot |P|)$ . It can be reduced by  
implementing  $P$  as a priority queue, prioritized by the value returned by  $d$ .

375 The correctness of Algorithm 1 is ensured by the invariant expressed in the following  
376 lemma.

■ **Algorithm 1** Best search for sw-VPA

---

**initially** let  $\mathcal{Q} = (Q \times Q) \cup (Q \times P \times Q)$ , and let  $d_{\perp}(q_1, q_2) = d_{\top}(q_1, p, q_2) = \mathbb{1}$  if  $q_1 = q_2$  and  $d_{\perp}(q_1, q_2) = d_{\top}(q_1, p, q_2) = 0$  otherwise

**while**  $\mathcal{Q} \neq \emptyset$  **do**

**extract**  $\langle q_1, q_2 \rangle$  or  $\langle q_1, p, q_2 \rangle$  from  $\mathcal{Q}$  such that  $d_{\perp}(q_1, q_2)$ , resp.  $\bigoplus_{c \in \Omega_c} d_{\top}(q_1, p, q_2)(c)$ , is minimal in  $\mathbb{S}$  wrt  $\leq_{\oplus}$

**update**  $d_{\perp}$  with  $\langle q_1, q_2 \rangle$  or  $d_{\top}$  with  $\langle q_1, p, q_2 \rangle$  (Figure 3).

---

For all  $q_0, q_3 \in Q$ ,

$$\begin{aligned}
 d_{\top}(q_1, p, q_3) &\oplus= d_{\top}(q_1, p, q_2) \otimes \bigoplus_{\Omega_i} w_i(q_2, p, q_3) \\
 d_{\perp}(q_1, p, q_3) &\oplus= d_{\perp}(q_1, q_2) \otimes \bigoplus_{\Omega_i} w_i^e(q_2, q_3) \\
 d_{\top}(q_0, p, q_3) &\oplus= \bigoplus_{\Omega_c}^2 [(w_c(q_0, p, p', q_1) \otimes_2 d_{\top}(q_1, p', q_2)) \otimes_2 \bigoplus_{\Omega_r} w_r(q_2, p', q_3)] \\
 d_{\perp}(q_0, q_3) &\oplus= \bigoplus_{\Omega_c} (w_c^e(q_0, p, q_1) \otimes d_{\top}(q_1, p, q_2) \otimes \bigoplus_{\Omega_r} w_r(q_2, p, q_3)) \\
 d_{\perp}(q_1, q_3) &\oplus= d_{\perp}(q_1, q_2) \otimes \bigoplus_{\Omega_r} w_r^e(q_2, q_3) \\
 d_{\top}(q_1, p, q_3) &\oplus= d_{\top}(q_1, p, q_2) \otimes d_{\top}(q_2, p, q_3), \text{ if } \langle q_2, \top, q_3 \rangle \notin P \\
 d_{\perp}(q_1, q_3) &\oplus= d_{\perp}(q_1, q_2) \otimes d_{\perp}(q_2, q_3), \text{ if } \langle q_2, \perp, q_3 \rangle \notin P
 \end{aligned}$$

■ **Figure 3** Update  $d_{\perp}$  with  $\langle q_1, q_2 \rangle$  or  $d_{\top}$  with  $\langle q_1, p, q_2 \rangle$ .

377 ► **Lemma 14.** For all  $\langle q_1, q_2 \rangle \notin \mathcal{Q}$ ,  $d_{\perp}(q_1, q_2) = b_{\perp}(q_1, q_2)/$

378 The proof is by contradiction, assuming a counter-example minimal in the length of the  
379 witness word.

380 ► **Lemma 15.** For all  $\langle q_1, p, q_2 \rangle \notin \mathcal{Q}$ ,  $d_{\top}(q_1, p, q_2) = b_{\top}(q_1, p, q_2)$ ,

381 For computing the minimal weight of a computation of  $A$ , we use the fact that, at the  
382 termination of Algorithm 1,  $\bigoplus_{s \in \Omega^*} A(s) = \bigoplus_{q, q' \in Q} \text{in}(q) \otimes d_{\perp}(q, q') \otimes \text{out}(q')$ .

383 In order to obtain effectively a witness (word of  $\Omega^*$  with a computation of  $A$  of minimal  
384 weight), we require the additional property of convexity of weight functions.

385 ► **Proposition 16.** For a sw-VPA  $A$  over  $\Omega$ ,  $\mathbb{S}$  commutative, bounded, total and complete,  
386 and  $\bar{\Phi}$  effective, one can construct in PTIME a word  $t \in \Omega^*$  such that  $A(t)$  is minimal wrt  
387 the natural ordering for  $\mathbb{S}$ .

## 388 5 Symbolic Weighted Parsing

389 Let us now apply the models and results of the previous sections to the problem of parsing  
390 over infinite alphabet. Let  $\Sigma$  be a countable input alphabet, and  $\Omega = \Omega_i \uplus \Omega_c \uplus \Omega_r$  be a  
391 countable output alphabet. Let  $\langle \mathbb{S}, \oplus, \mathbb{0}, \otimes, \mathbb{1} \rangle$  be a commutative, bounded, and complete  
392 semiring and let  $\bar{\Phi}$  be an effective label theory over  $\mathbb{S}$ , containing  $\Phi_{\Sigma}$ ,  $\Phi_{\Sigma, \Omega_i}$ , as well as  $\Phi_i$ ,  
393  $\Phi_c$ ,  $\Phi_r$ ,  $\Phi_{cr}$  (following the notations of Section 4). We assume given the following input:

total?

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- 394 – a swT  $T$  over  $\Sigma$ ,  $\Omega_i$ ,  $\mathbb{S}$ , and  $\bar{\Phi}$ , defining a measure  $T : \Sigma^* \times \Omega_i^* \rightarrow \mathbb{S}$ ,
  - 395 – a sw-VPA  $A$  over  $\Omega$ ,  $\mathbb{S}$ , and  $\bar{\Phi}$ , defining a measure  $A : \Omega^* \rightarrow \mathbb{S}$ ,
  - 396 – an input word  $s \in \Sigma^*$ .
- 397 For all  $s \in \Sigma^*$  and  $t \in \Omega^*$ , let  $d(s, t) = T(s, t|_{\Omega_i})$ , where  $t|_{\Omega_i} \in \Omega_i^*$  is the projection of  $t$
- 398 onto  $\Omega_i$ , obtained from  $t$  by removing all symbols in  $\Omega \setminus \Omega_i$ . *Symbolic weighted parsing* is the
- 399 problem, given the above input, to find  $t \in \Omega^*$  minimizing  $d(s, t) \otimes A(t)$  wrt  $\leq_\oplus$ , i.e. s.t.

$$400 \quad d(s, t) \otimes A(t) = \bigoplus_{t' \in \mathcal{T}(\Omega)} d(s, t') \otimes A(t') \quad (9)$$

401 Following the terminology of [21], **sw**-parsing is the problem of computing the distance (9)

402 between the input  $s$  and the output weighted language of  $A$ , and returning a witness  $t$ . Every

403 labeled tree  $t$  can be linearized into a nested word  $\text{lin}(t) \in \Omega^*$ , assuming e.g. that  $\Omega_i$  contain

404 the symbols labelling the leaves (symbols of rank 0) and  $\Omega_c$  and  $\Omega_r$  contain respectively

405 one left and right parenthesis  $\langle_b$  and  $\rangle_b$  for each symbol  $b$  labelling inner nodes (symbols of

406 rank  $> 0$ ). With this representation, the projection  $\text{lin}(t)|_{\Omega_i}$  is then the sequence of leaves

407 of  $t$ , enumerated in a *dfs*-traversal. We show in Appendix A how to convert a (sw) tree

408 automaton  $A$  into a sw-VPA computing  $A(\text{lin}(t))$  for every tree  $t$ . That also holds, for the

409 set of ASTs of a weighted CF-grammar. Therefore, **sw**-parsing generalizes the problem of

410 searching, the best derivation of a weighted CF-grammar that yields a given input, sometimes

411 referred as *weighted parsing*, see e.g. [13] and [23] for a more general weighted parsing

412 framework. The latter indeed corresponds to the particular case where the alphabet is finite,

413  $T$  computes identity and  $A$  is obtained from the weighted CF grammar.

414 ► **Proposition 17.** *The problem of Symbolic Weighted parsing can be solved in PTIME in*

415 *the size of the input swT  $T$ , sw-VPA  $A$  and input word  $s$ , and the computation time of the*

416 *functions of the label theory.*

417 **Proof.** (sketch) We follow a *Bar-Hillel* construction, also called parsing by intersection.

418 We first extend the swT  $T$  over  $\Sigma$ ,  $\Omega_i$ ,  $\mathbb{S}$ , and  $\bar{\Phi}$ , into a swT  $T'$  over  $\Sigma$  and  $\Omega$  (and the same

419 semiring and label theory), such that for all  $s \in \Sigma^*$ , and  $u \in \Omega^*$ ,  $T'(s, u) = T(s, u|_{\Omega_i})$ . The

420 transducer  $T'$  simply skips every symbol  $b \in \Omega \setminus \Omega_i$ , by the addition to the transition of

421  $T$ , of new transitions of the form  $w_{01}(q, \varepsilon, b, q')$ . Then, given an input word  $s \in \Sigma^*$ , using

422 Corolary 10, we compute the swA  $B_{T',s}$ , such that for all  $t \in \Omega^*$ ,  $B_{T',s}(t) = d(s, t)$ .

423 Next, we compute the sw-VPA  $B_{T',s} \otimes A$ , using Proposition 13. It remains to compute a

424 best nested-word  $w \in \Omega^*$  using the best-search procedure of Proposition 16. ◀

2 lines Application  
to Automated Mu-  
sic Transcription:  
implementation  $\neq$   
but same principle,  
on-the-fly automata  
construction during  
best search, for effi-  
ciency.

## Conclusion

427 We have introduced weighted language models (SW transducers and visibly pushdown

428 automata) computing over infinite alphabets, and applied them to the problem of parsing

429 with infinitely many possible input symbols (typically timed events). This approach extends

430 conventional parsing and weighted parsing by computing a derivation tree modulo a generic

431 distance between words, defined by a SW transducer given in input. This enables to consider

432 finer word relationships than strict equality, opening possibilities of quantitative analysis via

433 this method.

TODO future work

434 Ongoing and future work include

435 – The study of other theoretical properties of SW models, such as the extension of the best

436 search algorithm from 1-best to  $n$ -best [17], and to  $k$ -closed semirings [20] (instead of *bounded*,

437 which corresponds to 0-closed).

- 438 – ...there is room to improve the complexity bounds for the algorithms ... modular approach
- 439 with oracles ...
- 440 – present here an offline algorithm for best search, semi-online implementation for AMT
- 441 (bar-by-bar approach) with an on-the-fly automata construction.

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## A

 Nested-Words and Parse-Trees

The hierarchical structure of nested-words, defined with the *call* and *return* markup symbols suggest a correspondence with trees. The lifting of this correspondence to languages, of tree automata and VPA, has been discussed in [1], and [4] for the weighted case. In this section, we describe a correspondence between the symbolic-weighted extensions of tree automata and VPA.

Let  $\Omega$  be a countable ranked alphabet, such that every symbol  $a \in \Omega$  has a rank  $\text{rk}(a) \in [0..M]$  where  $M$  is a fixed natural number. We denote by  $\Omega_k$  the subset of all symbols  $a$  of  $\Omega$  with  $\text{rk}(a) = k$ , where  $0 \leq k \leq M$ , and  $\Omega_{>0} = \Omega \setminus \Omega_0$ . The free  $\Omega$ -algebra of finite, ordered,  $\Omega$ -labeled trees is denoted by  $\mathcal{T}(\Omega)$ . It is the smallest set such that  $\Omega_0 \subset \mathcal{T}(\Omega)$  and for all  $1 \leq k \leq M$ , all  $a \in \Omega_k$ , and all  $t_1, \dots, t_k \in \mathcal{T}(\Omega)$ ,  $a(t_1, \dots, t_k) \in \mathcal{T}(\Omega)$ . Let us assume a commutative semiring  $\mathbb{S}$  and a label theory  $\bar{\Phi}$  over  $\mathbb{S}$  containing one set  $\Phi_{\Omega_k}$  for each  $k \in [0..M]$ .

**Definition 18.** A symbolic-weighted tree automaton (*swTA*) over  $\Omega$ ,  $\mathbb{S}$ , and  $\bar{\Phi}$  is a triplet  $A = \langle Q, \text{in}, \bar{w} \rangle$  where  $Q$  is a finite set of states,  $\text{in} : Q \rightarrow \Phi_{\Omega}$  is the starting weight function, and  $\bar{w}$  is a tuple of transition functions containing, for each  $k \in [0..M]$ , the functions  $w_k : Q \times Q^k \rightarrow \Phi_{\Omega_{>0}, \Omega_k}$  and  $w_k^e : Q \times Q^k \rightarrow \Phi_{\Omega_k}$ .

We define a transition function  $w : Q \times (\Omega_{>0} \cup \{\varepsilon\}) \times \Omega \times \bigcup_{k=0}^M Q^k \rightarrow \mathbb{S}$  by:

$$\begin{aligned} w(q_0, a, b, q_1 \dots q_k) &= \eta(a, b) & \text{where } \eta &= w_k(q_0, q_1 \dots q_k) \\ w(q_0, \varepsilon, b, q_1 \dots q_k) &= \phi(b) & \text{where } \phi &= w_k^e(q_0, q_1 \dots q_k). \end{aligned}$$

where  $q_1 \dots q_k$  is  $\varepsilon$  if  $k = 0$ . The first case deals with a strict subtree, with a parent node labeled by  $a$ , and the second case is for a root tree.

Every *swTA* defines a mapping from trees of  $\mathcal{T}(\Omega)$  into  $\mathbb{S}$ , based on the following intermediate function  $\text{weight}_A : Q \times (\Omega \cup \{\varepsilon\}) \times \mathcal{T}(\Omega) \rightarrow \mathbb{S}$

$$\text{weight}_A(q_0, a, t) = \bigoplus_{q_1 \dots q_k \in Q^k} w(q_0, a, b, q_1 \dots q_k) \otimes \bigotimes_{i=1}^k \text{weight}_A(q_i, b, t_i) \quad (10)$$

where  $q_0 \in Q$ ,  $a \in \Omega_{>0} \cup \{\varepsilon\}$  and  $t = b(t_1, \dots, t_k) \in \mathcal{T}(\Omega)$ ,  $0 \leq k \leq M$ .

Finally, the weight associated by  $A$  to  $t \in \mathcal{T}(\Omega)$  is

$$A(t) = \bigoplus_{q \in Q} \text{in}(q) \otimes \text{weight}_A(q, \varepsilon, t) \quad (11)$$

Intuitively,  $w(q_0, a, b, q_1 \dots q_k)$  can be seen as the weight of a production rule  $q_0 \rightarrow b(q_1, \dots, q_k)$  of a regular tree grammar [5], that replaces the non-terminal symbol  $q_0$  by  $b(q_1, \dots, q_k)$ , provided that the parent of  $q_0$  is labeled by  $a$  (or  $q_0$  is the root node if  $a = \varepsilon$ ). The above production rule can also be seen as a rule of a weighted CF grammar, of the form  $[a, b] q_0 := q_1 \dots q_k$  if  $k > 0$ , and  $[a] q_0 := b$  if  $k = 0$ . In the first case,  $b$  is a label of the rule, and in the second case, it is a terminal symbol. And in both cases,  $a$  is a constraint on the label of rule applied on the parent node in the derivation tree. This features of observing the parent's label are useful in the case of infinite alphabet, where it is not possible to memorize a label with the states. The weight of a labeled derivation tree  $t$  of the weighted CF grammar associated to  $A$  as above, is  $\text{weight}_A(q, t)$ , when  $q$  is the start non-terminal. We shall now establish a correspondence between such derivation tree  $t$  and some word describing a linearization of  $t$ , in a way that  $\text{weight}_A(q, t)$  can be computed by a *sw-VPA*.



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539 Let  $\hat{\Omega}$  be the countable (unranked) alphabet obtained from  $\Omega$  by:  $\hat{\Omega} = \Omega_i \uplus \Omega_c \uplus \Omega_r$ , with  
 540  $\Omega_i = \Omega_0$ ,  $\Omega_c = \{ \langle a \mid a \in \Omega_{>0} \rangle \}$ ,  $\Omega_r = \{ a \mid a \in \Omega_{>0} \}$ .

541 We associate to  $\hat{\Omega}$  a label theory  $\hat{\Phi}$  like in Section 4, and we define a linearization of trees of  
 542  $\mathcal{T}(\Omega)$  into words of  $\hat{\Omega}^*$  as follows:

543  $\text{lin}(a) = a$  for all  $a \in \Omega_0$ ,

544  $\text{lin}(b(t_1, \dots, t_k)) = \langle b \text{ lin}(t_1) \dots \text{lin}(t_k) b \rangle$  when  $b \in \Omega_k$  for  $1 \leq k \leq M$ .

545 ► **Proposition 19.** *For all swTA  $A$  over  $\Omega$ ,  $\mathbb{S}$  commutative, and  $\bar{\Phi}$ , there exists an effectively*  
 546 *constructible sw-VPA  $A'$  over  $\hat{\Omega}$ ,  $\mathbb{S}$  and  $\hat{\Phi}$  such that for all  $t \in \mathcal{T}(\Omega)$ ,  $A'(\text{lin}(t)) = A(t)$ .*

547 **Proof.** Let  $A = \langle Q, \text{in}, \bar{w} \rangle$  where  $\bar{w}$  is presented as above by a function We build  $A' =$   
 548  $\langle Q', P', \text{in}', \bar{w}', \text{out}' \rangle$ , where  $Q' = \bigcup_{k=0}^M Q^k$  is the set of sequences of state symbols of  $A$ , of  
 549 length at most  $M$ , including the empty sequence denoted by  $\varepsilon$ , and where  $P' = Q'$  and  $\bar{w}$  is  
 550 defined by:

$$\begin{array}{lll}
 \mathbf{w}_i(q_0 \bar{u}, \langle c, \bar{p}, a, \bar{u} \rangle) & = & \mathbf{w}(q_0, c, a, \varepsilon) \quad \text{for all } c \in \Omega_{>0}, a \in \Omega_0 \\
 \mathbf{w}_i^e(q_0 \bar{u}, a, \bar{u}) & = & \mathbf{w}(q_0, \varepsilon, a, \varepsilon) \quad \text{for all } a \in \Omega_0 \\
 \mathbf{w}_c(q_0 \bar{u}, \langle c, \bar{p}, \langle d, \bar{u}, \bar{q} \rangle \rangle) & = & \mathbf{w}(q_0, c, d, \bar{q}) \quad \text{for all } c, d \in \Omega_{>0} \\
 551 \quad \mathbf{w}_c^e(q_0 \bar{u}, \langle c, \bar{u}, \bar{q} \rangle) & = & \mathbf{w}(q_0, \varepsilon, c, \bar{q}) \quad \text{for all } c \in \Omega_{>0} \\
 \mathbf{w}_r(\varepsilon, \langle c, \bar{p}, c \rangle, \bar{p}) & = & \mathbb{1} \quad \text{for all } c \in \Omega_{>0} \\
 \mathbf{w}_r^e(\bar{u}, c, \bar{q}) & = & \mathbb{0} \quad \text{for all } c \in \Omega_{>0}
 \end{array}$$

552 All cases not matched by one of the above equations have a weight  $\mathbb{0}$ , for instance  $\mathbf{w}_r(\bar{u}, \langle c, \bar{p}, d \rangle, \bar{q}) =$   
 553  $\mathbb{0}$  if  $c \neq d$  or  $\bar{u} \neq \varepsilon$  or  $\bar{q} \neq \bar{p}$ . ◀

554 **Todo list**

555	<input type="checkbox"/> register: skip refs and details, add Mikolaj recent . . . . .	1
556	<input type="checkbox"/> chap. intersection in [15] . . . . .	3
557	<input type="checkbox"/> expressiveness: VPA have restricted equality test. comparable to pebble automata?	
558	→ conclusion . . . . .	3
559	<input type="checkbox"/> The results are established for a general class of semirings. They can be instantiated	
560	for concrete cases . . . . .	3
561	<input type="checkbox"/> partial application is needed? . . . . .	5
562	<input type="checkbox"/> notion of diagram of functions akin BDD for transitions in practice . . . . .	5
563	<input type="checkbox"/> mv appendix? . . . . .	5
564	<input type="checkbox"/> precise/restrict complexity . . . . .	5
565	<input type="checkbox"/> added $u$ and $v$ . . . . .	6
566	<input type="checkbox"/> unique → similar . . . . .	6
567	<input type="checkbox"/> similar → single . . . . .	6
568	<input type="checkbox"/> modif. . . . .	7
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570	<input type="checkbox"/> proof correctness . . . . .	8
571	<input type="checkbox"/> revise with nb of tr. and states . . . . .	8
572	<input type="checkbox"/> moved this to the beginning . . . . .	9
573	<input type="checkbox"/> intro to func . . . . .	9
574	<input type="checkbox"/> notation $cp$ for $\langle c, p \rangle$ ? . . . . .	9
575	<input type="checkbox"/> todo example VPA . . . . .	10
576	<input type="checkbox"/> total? . . . . .	10
577	<input type="checkbox"/> 2 cases of $b$ . . . . .	10
578	<input type="checkbox"/> $b_{\top}$ : mot bien parenthésé $c/r$ . . . . .	10
579	<input type="checkbox"/> explication Fig. 3 suivant cas de (5) . . . . .	10
580	<input type="checkbox"/> complete ** . . . . .	10
581	<input type="checkbox"/> detail with nb tr. and states . . . . .	10
582	<input type="checkbox"/> total? . . . . .	11
583	<input type="checkbox"/> 2 lines Application to Automated Music Transcription: implementation $\neq$ but same	
584	principle, on-the-fly automata construction during best search, for efficiency. . . .	12