Programming Languages
Variables and Storage

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Computer Engineering

Outline

Storage

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- Imperative language variables: variable has a state and value. It can be assigned to different values in same phrase.
- Two basic operations on a variable: inspect and update.

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- Cells are initially unallocated.

```
 f();
void f() {
   int x;
   ...
   x=5;
   ...
   return;
}
```

Computer memory can be considered as a collection of cells.

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- After the including block terminates, again unallocated

```
 f();
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```

Total or Selective Update

- Composite variables can be inspected and updated in total or selectively

```
struct Complex { double x,y; } a, b;
...
a=b;                    // Total update
a.x=b.y*a.x;            // Selective update
```

- ■

- Primitive variables: single cell
  Composite variables: nested cells

Array Variables

Different approaches exist in implementation of array variables:

1. Static arrays
2. Dynamic arrays
3. Flexible arrays

Static arrays

- Array size is fixed at compile time to a constant value or expression.
- C example:

```
#define MAXELS 100
int a[10];
double x[MAXELS*10][20];
```

Dynamic arrays

- Array size is defined when variable is allocated. Remains constant afterwards.
- Example: C90/GCC (not in ANSI)

```
int f(int n) {
    double a[n]; ...
}
```

- Example: C++ with templates

```
template<class T>  class Array {
      T *content;
  public:
      Array(int s) { content=new T[s]; }
      ~Array()     { delete [] content; }
};
...
Array<int>  a(10);                Array<double> b(n);
```

Flexible arrays

- Array size is completely variable. Arrays may expand or shrink at run time. Script languages like Perl, PHP, Python

- Perl example:

```perl
@a=(1,3,5);          # array size: 3
print $#a , "\n";    # output: 2 (0..2)
$a[10] = 12;         # array size 11 (intermediate elements u
$a[20] = 4;          # array size 21
print $#a , "\n";    # output: 20 (0..20)
delete $a[20];       # last  element erased,  size is 11
print $#a , "\n";    # output: 10 (0..10)
```
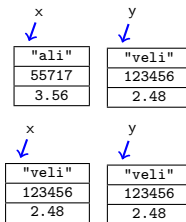
- C++ and object orient languages allow overload of [] operator to make flexible arrays possible. STL (Standard Template Library) classes in C++ like vector, map are like such flexible array implementations.

Semantic of assignment in composite variables

- Assignment by Copy vs Reference.

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- Assignment by Copy vs Reference.
- Copy: All content is copied into the other variables storage. Two copies with same values in memory.

x   y
↓   ↓

| "ali" | | "veli" |
|--------| |--------|
| 55717 | | 123456 |
| 3.56 | | 2.48 |

assignment by Copy:

x   y
↓   ↓

| "veli" | | "veli" |
|--------| |--------|
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Semantic of assignment in composite variables

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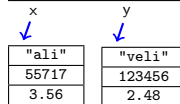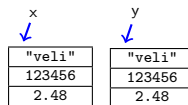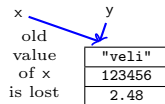- Reference: Reference of variable is copied to other variable. Two variables share the same storage and values.
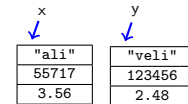
- Assignment semantics is defined by the language design
- C structures follows copy semantics. Arrays cannot be assigned. Pointers are used to implement reference semantics. C++ objects are similar.
- Java follows copy semantics for primitive types. All other types (objects) are reference semantics.
- Copy semantics is slower
- Reference semantics cause problems from storage sharing (all operations effect both variables). Deallocation of one makes the other invalid.
- Java provides copy semantic via a member function called copy(). Java garbage collector avoids invalid values (in case of deallocation)

Variable Lifetime

- Variable lifetime: The period between allocation of a variable and deallocation of a variable.
- 4 kinds of variable lifetime.
    1. Global lifetime (while program is running)
    2. Local lifetime (while declaring block is active)
    3. Heap lifetime (arbitrary)
    4. Persistent lifetime (continues after program terminates)

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lifetime of global variables

Global lifetime
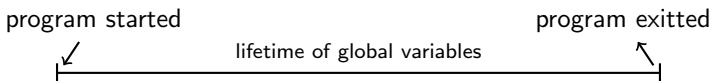
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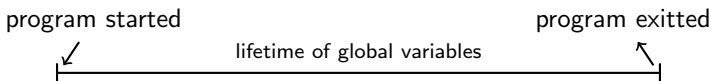
      ↙   lifetime of global variables   ↖

- What are `static` variables inside functions in C?

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- Multiple instances of same local variable may be alive at the same time in recursive functions.
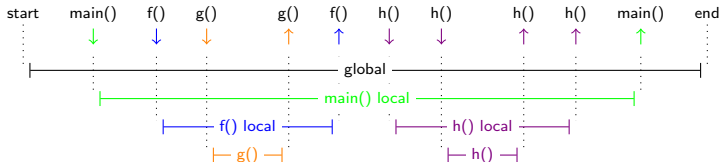
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-

```
double x;
t h(int n) {
  int a;
  if (n<1) return 1
  else return h(n-1);

d g() {
  int x;
  int b;

 f() {
  double z;
  ...
  g();
  ...

 main() {
  double k;
  f();
  ...
  h(1);
  ...;
  return 0;
```
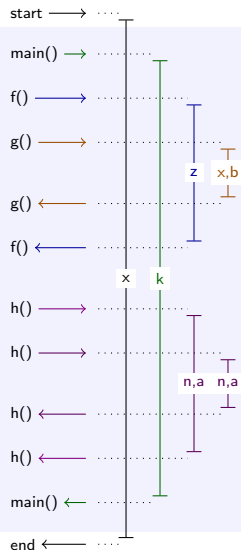
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  ```
  double *p;
  p=malloc(sizeof(double));
  *p=3.4; ...
  
  free(p);
  ```

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- p and *p are different variables p has pointer type and usually a local or global lifetime, *p is heap variable.
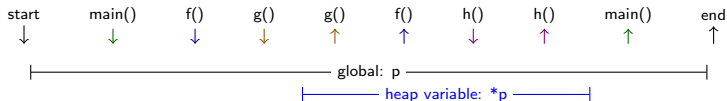- heap variable lifetime can start or end at anytime.

```
double *p;
int h() { ...
}
void g() { ...
    p=malloc(sizeof(double));
}
int f() { ...
    g(); ...
}
int main() { ...
    f();      ...
    h();      ...;
    free(p); ...
}
```

start	main()	f()	g()	g()	f()	h()	h()	main()	end
↓	↓	↓	↓	↑	↑	↓	↑	↑	↑

|————————————————————————— global: p —————————————————————————|

|————— heap variable: *p —————|

Dangling Reference

- dangling reference: trying to access a variable whose lifetime is ended and already deallocated.

```c
char *p, *q;

p=malloc(10);
q=p;
...
free(q);
printf("%s",p);
```

```c
char *f() {
    char a[]="ali";
    ....
    return a;
}
....
char *p;
p=f();
printf("%s",p);
```

- both p's are deallocated or ended lifetime variable, thus dangling reference

- sometimes operating system tolerates dangling references. Sometimes generates run-time erros like "protection fault", "segmentation fault" are generated.

Garbage variables

- garbage variables: The variables with lifetime still continue but there is no way to access.

```
                                        void f() {
                                            char *p;
char *p, *q;                                p=malloc(10); ...
...                                         return
p=malloc(10);                           }
p=q;                                    ....
...                                     f();
```

- When the pointer value is lost or lifetime of the pointer is over, heap variable is unaccessible. (*p in examples)

Garbage collection

- A solution to dangling reference and garbage problem: PL does management of heap variable deallocation automatically. This is called garbage collection. (Java, Lisp, ML, Haskell, most functional languages)
- no call like `free()` or `delete` exists.
- Language runtime needs to:
  - Keep a reference counter on each reference, initially 1.
  - Increment counter on each new assignment
  - Decrement counter at the end of the reference lifetime
  - Decrement counter at the overwritten/lost references
  - Do all these operations recursively on composite values.
  - When reference count gets 0, deallocate the heap variable

- Garbage collector deallocates heap variables having a reference count 0.
- Since it may delay execution of tasks, GC is not immediately done.
- GC usually works in a separate thread, in low priority, works when CPU is idle.
- Another but too restrictive solution to garbage: Reference cannot be assigned to a longer lifetime variable. local variable references cannot be assigned to global reference/pointer.

Persistent variable lifetime

- Variables with lifetime continues after program terminates: file, database, web service object,...

- Stored in secondary storage or external process.

- Only a few experimental language has transparent persistence. Persistence achieved via IO instructions
  C files: `fopen()`, `fseek()`, `fread()`, `fwrite()`

- In object oriented languages; serialization: Converting object into a binary image that can be written on disk or sent over the network.

- This way objects snapshot can be taken, saved, restored and object continue from where it remains.

Memory Management

- Memory manamegement of variables involves architecture, operating system, language runtime and the compiler.
- A typical OS divides memory in sections (segments):
    - Stack section: run time stack
    - Heap section: heap variables
    - Data section: global variables
    - Code section: executable instructions, read only.
- Global variables are fixed at compile time and they are put in data section.
- Heap variables are stored in the dynamic data structures in heap section. Heap section grows and shrinks as new variables are allocated and deallocated.
- Heap section is maintained by language runtime. For C, it is `libc`.

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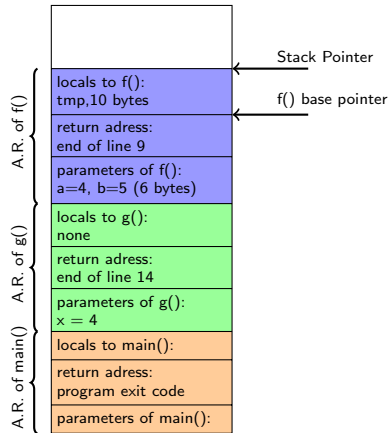   - Parameter values.

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- In a typical architecture, activation record contains:
  - Return address. Address of the next instruction after the caller.
  - Parameter values.
  - A reserved area for local variables.

```
1    int  f ( short  a ,  int  b )  {
2         char  tmp [10];
3         ...
4         return  a + b;
5    }
6    int  g ( int  x )  {
7         int  tmp,  p;
8         ...
9         tmp  =  f ( x ,  x +1);
10        ...
11        return  tmp + p;
12   }
13   int  main ()  {
14        return  g (4);
15   }
```


Stack Pointer

f() base pointer

A.R. of f()

locals to f():
tmp,10 bytes

return adress:
end of line 9

parameters of f():
a=4, b=5 (6 bytes)

A.R. of g()

locals to g():
none

return adress:
end of line 14

parameters of g():
x = 4

A.R. of main()

locals to main():

return adress:
program exit code

parameters of main():

# Function Call

- Caller side:
    1. Push parameters
    2. Push return address and jump to function code start
       (usually a single CPU instruction like `callq`)
- Function entry:
    1. Set base pointer to current stack pointer
    2. Advance stack pointer to size of local variables
- Function body can access all local variables relative to base pointer
- Function return:
    1. Set stack pointer to base pointer
    2. Pop return address and jump to return address
       (single CPU instruction like `retq`)
- Caller side after return:
    1. Recover stack pointer (remove parameters on stack)
    2. Get and use return value if exists (typically from a register)

- All locals and parameters have the same offset from base pointer
- Recursive calls execute same instructions

```
1    int h(int n) {
2        int tmp;
3        if (n <= 1) return 0;
4        else {
5            tmp = h(n-1);
6            return n+tmp;
7        }
8    }
9    int main() {
10       printf("%d\n", h(2));
11       return 0;
12   }
```
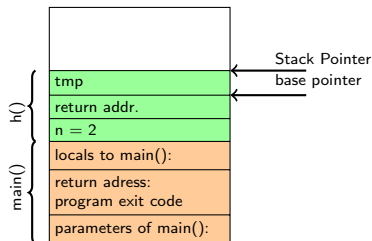
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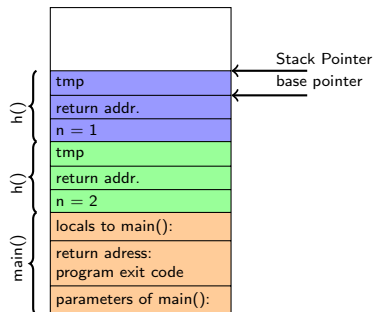
- Order of values in the activation record may differ for different languages
- Registers are used for passing primitive value parameters instead of stack
- Garbage collecting languages keep references on stack with actual variables on heap
- Languages returning nested functions as first order values require more complicated mechanisms

```python
def multiplier(a):
        def f(x):
                return a*x
        return f

twice = multiplier(2)
# multiplier/a is not alive anymore but twice=f is using it
print(twice(14))
```

Commands

Expression: program segment with a value.

Statement: program segment without a value, but alters the state.

Input, output, variable assignment, iteration...

1. Assignment
2. Procedure call
3. Block commands
4. Conditional commands
5. Iterative commands

Assignment

- C: "Var = Expr;", Pascal "Var := Expr;".
- Evaluates RHS expression and sets the value of the variable at RHS
- x = x + 1 . LHS x is a variable reference (l-value), RHS is the value
- multiple assignment: x=y=z=0;
- parallel assignment: (Perl, PHP) ($a,$b) = ($b, $a);
  ($name, $surname, $no) =
  ("Onur","Şehitoğlu",55717);
  Assignment: "reference aggregate" $\rightarrow$ "value aggregate"
- assignment with operator: x += 3; x *= 2;

Procedure call

- Procedure: user defined commands. Pascal: procedure, C: function returning void
- void *functname* (*param1* , *param2* , ..., *paramn*)
- Usage is similar to functions but call is in a statement position (on a separate line of program)

Block commands

- Composition of a block from multiple statements
- Sequential commands: `{ C₁ ; C₂; ... ; Cₙ }`
  A command is executed, after it finishes the next command is executed,…
- Commands enclosed in a block behaves like single command:"if" blocks, loop bodies,…
- Collateral commands: `{ C₁ , C₂ , ... , Cₙ }` (not C ',')!
  Commands can be executed in any order.
- The order of execution is non-deterministic. Compiler or optimizer can choose any order. If commands are independent, effectively deterministic:
  'y=3 , x=x+1 ;' vs 'x=3, x=x+1 ;'
- Can be executed in parallel.

- Concurrent commands: concurrent paradigm languages:
  $\{\ C_1\ |\ C_2\ |\ \ldots\ |\ C_n\ \}$
- All commands start concurrently in parallel. Block finishes when the last active command finishes.
- Real parallelism in multi-core/multi-processor machines.
- Transparently handled by only a few languages. Thread libraries required in languages like Java, C, C++.

```
void producer (...) {....}
void collectgarbage (...) {....}
void consumer (...) {....}
int main () {
        ...
        pthread_create (tid1 ,NULL, producer ,NULL);
        pthread_create (tid2 ,NULL, collectgarbage ,NULL);
        pthread_create (tid3 ,NULL, consumer ,NULL);
        ...
}
```

Conditional commands

- Commands to choose between alternative commands based on a condition
- in C : if (*cond*) C_1 else C_2 ;
  switch (*value*) { case L_1 : C_1 ; case L_2 : C_2 ; ...}
- if commands can be nested for multi-conditioned selection.
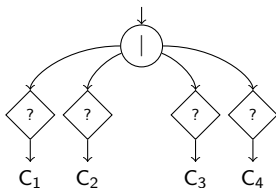- switch like commands chooses statements based on a value
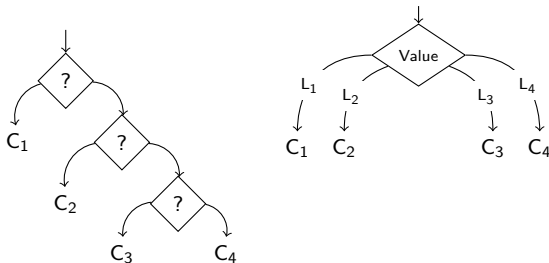
- non-deterministic conditionals: conditions are evaluated in collaterally and commands are executed if condition holds.
- hyphotetically:

  if ($cond_1$) C_1 or if ($cond_2$) C_2 or if ($cond_3$) C_3 ;

  switch (val) {
          case L_1: C_1 | case L_2: C_2 | case L_3: C_3 }

- Tests can run concurrently. First test evaluating to True wins. Others discarded.

Iterative statements

- Repeating same command or command block multiple times possibly with different data or state. Loop commands.
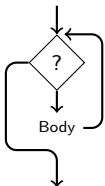
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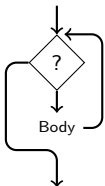
C: while (...) { ... }

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C: `while (...) { ... }`   C: `do {...} while (...);`
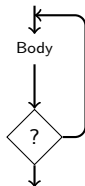
- Definite vs indefinite loops
- Indefinite iteration: Number of iterations of the loop is not known until loop finishes
- C loops are indefinite iteration loops.
- Definite iteration: Number of iterations is fixed when loop started.
- Pascal `for` loop is a definite iteration loop.
  for i:= k to m do begin .... end; has $(m - k + 1)$ iterations.
  Pascal forbids update of the loop index variable.
- List and set based iterations: PHP, Perl, Python, Shell

```
$colors=array('yellow','blue','green','red','white');
foreach ($colors as $i) {
    print $i,"␣is␣a␣color","\n";
}
```

Summary

- Variables with storage
- Variable update
- Lifetime: global, local, heap, persistent
- Commands