Programming Language Concepts Higher Order Functions

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Lambda Calculus

- 1930's by Alonso Church and Stephen Cole Kleene
- Mathematical foundation for computatibility and recursion
- Simplest functional paradigm language
- \(\lambda var.expr\) defines an anonymous function. Also called lambda abstraction
- expr can be any expression with other lambda abstractions and applications. Applications are one at a time.
- \bullet $(\lambda x.\lambda y.x + y)$ 3 4



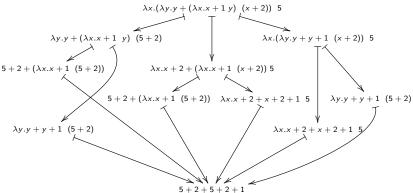
- In ' $\lambda var.expr$ ' all free occurences of var is bound by the λvar .
- Free variables of expression FV(expr)
 - $FV(name) = \{name\}$ if name is a variable
 - $FV(\lambda name.expr) = FV(expr) \{name\}$
 - $FV(M N) = FV(M) \cup FV(N)$
- lacktriangleq lpha conversion: expressions with all bound names changed to another name are equivalent:

$$\lambda f.f \ x \equiv_{\alpha} \lambda y.y \ x \equiv_{\alpha} \lambda z.z \ x$$
$$\lambda x.x + (\lambda x.x + y) \equiv_{\alpha} \lambda t.t + (\lambda x.x + y) \equiv_{\alpha} \lambda t.t + (\lambda u.u + y)$$
$$\lambda x.x + (\lambda x.x + y) \not\equiv_{\alpha} \lambda x.x + (\lambda x.x + t)$$

β Reduction

- Basic computation step, function application in λ -calculus
- lacktriangle Based on substitution. All bound occurences of λ variable parameter is substituted by the actual parameter
- $(\lambda x.M)N \mapsto_{\beta} M[x/N]$ (all x's once bound by lambda are substituted with N).
- $(\lambda x.(\lambda y.y + (\lambda x.x + 1) y)(x + 2)) 5$
- lacksquare If no further eta reduction is possible, it is called a normal form.
- There can be different reduction strategies but should reduce to same normal form. (Church Rosser property)

All possible reductions of a λ -expression. All reduce to the same normal form.



$$(f \circ g)(x) = f(g(x))$$
, $(g \circ f)(x) = g(f(x))$

- "o": Gets two unary functions and composes a new function.
- in Haskell:

```
f x = x + x
g x = x * x
compose func1 func2 x = func1 (func2 x)
t = compose f g
u = compose g f
```

$$\mathbf{u}$$
 t 3 = (3*3)+(3*3) = 18
 \mathbf{u} 3 = (3+3)*(3+3) = 36

$$\blacksquare$$
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Functions/Curry

Cartesian form vs curried form:

$$\alpha \times \beta \rightarrow \gamma \text{ vs } \alpha \rightarrow \beta \rightarrow \gamma$$

 Curry function gets a binary function in cartesian form and converts it to curried form.

```
curry f x y = f(x,y)
add (x,y) = x+y
increment = curry add 1
---
increment 5
```

- \blacksquare curry: $(\alpha \times \beta \rightarrow \gamma) \rightarrow \alpha \rightarrow \beta \rightarrow \gamma$
- Haskell library includes it as curry.



Functions/Map

```
square x = x*x
day no = case no of 1 -> "mon"; 2 -> "tue"; 3 -> "wed";
        4 -> "thu"; 5 -> "fri"; 6 -> "sat"; 7 -> "sun"
map func [] = []
map func (el:rest) = (func el):(map func rest)
----
map square [1,3,4,6]
[1,9,6,36]
map day [1,3,4,6]
["mon", "wed", "thu", "sat"]
```

- \blacksquare map: $(\alpha \to \beta) \to [\alpha] \to [\beta]$
- Gets a function and a list. Applies the function to all elements and returns a new list of results.
- Haskell library includes it as map.



Functions/Filter

```
iseven x = if \mod x = 2 == 0 then True else False
isgreater x = x > 5
filter func [] = []
filter func (el:rest) = if func el then
                                  el:(filter func rest)
                           else (filter func rest)
filter iseven [1,2,3,4,5,6,7]
[2,4,6]
filter isgreater [1,2,3,4,5,6,7]
[6,7]
```

- \blacksquare filter: $(\alpha \to Bool) \to [\alpha] \to [\alpha]$
- Gets a boolean function and a list. Returns a list with only members evaluated to True by the boolean function.
- Haskell library includes it as filter.



Functions/Reduce (Fold Right)

```
sum \times y = x+y
product \times y = x*y
reduce func s [] = s
reduce func s (el:rest) = func el (reduce func s rest)
reduce sum 0 [1,2,3,4]
10
                                  // 1+2+3+4+0
reduce product 1 [1,2,3,4]
                                  // 1*2*3*4*1
24
```

- \blacksquare reduce: $(\alpha \to \beta \to \beta) \to \beta \to [\alpha] \to \beta$
- Gets a binary function, a list and a seed element. Applies function to all elements right to left with a single value. reduce $f s [a_1, a_2, ..., a_n] = f a_1 (f a_2 (.... (f a_n s)))$
- Haskell library includes it as foldr.



- Sum of a numbers in a list: listsum = reduce sum 0
- Product of a numbers in a list: listproduct = reduce product 1
- Sum of squares of a list: squaresum x = reduce sum 0 (map square x)

Functions/Fold Left

- foldl: $(\alpha \to \beta \to \alpha) \to \alpha \to [\beta] \to \alpha$
- Reduce operation, left associative.: $reduce\ f\ s\ [a_1,a_2,...,a_n] = f\ (f\ (f\ ...(f\ s\ a_1)\ a_2\ ...))\ a_n$
- Haskell library includes it as fold1.

Functions/Iterate

- \blacksquare iterate: $(\alpha \to \alpha) \to \alpha \to int \to \alpha$
- Applies same function for given number of times, starting with the initial seed value. iterate f s $n = f^n$ s $= \underbrace{f(f(f...(f s)))}$

Functions/Value Iteration (for)

```
for func s m n =
    if m>n then s
    else for func (func s m) (m+1) n

for sum 0 1 4
10     // sum (sum (sum 0 1) 2) 3) 4
for product 1 1 4
24     // product (product (product 1 1) 2) 3) 4
```

- for: $(\alpha \to int \to \alpha) \to \alpha \to int \to int \to \alpha$
- Applies a binary integer function to a range of integers in order.

for
$$f \, s \, m \, n = f(f \, (f \, (f \, (f \, s \, m) \, (m+1)) \, (m+2)) \, ...) \, n$$

- multiply (with summation): multiply x = iterate (sum x) x
- integer power operation (Haskell '^'):
 power x = iterate (product x) x
- sum of values in range 1 to n: seriessum = for sum 0 1
- Factorial operation: factorial = for product 1 1

Higher Order Functions in C

C allows similar definitions based on function pointers. Example: bsearch() and qsort() funtions in C library.

```
typedef struct Person { char name[30]; int no;} person;
int cmpnmbs(void *a, void *b) {
    person *ka=(person *)a; person *kb=(person *)b;
   return ka->no - kb->no;
int cmpnames(void *a, void *b) {
    person *ka=(person *)a; person *kb=(person *)b;
    return strncmp(ka->name,kb->name,30);
int main() {     int i:
    person list[]={{"veli",4},{"ali",12},{"ayse",8},
                  {"osman",6},{"fatma",1},{"mehmet",3}};
    qsort(list ,6,sizeof(person),cmpnmbs);
    qsort(list ,6,sizeof(person),cmpnames);
    . . .
```

Fibonacci

Fibonacci series: 1 1 2 3 5 8 13 21 .. fib(0) = 1; fib(1) = 1; fib(n) = fib(n-1) + fib(n-2)

```
fib n = let f (x,y) = (y,x+y)

(a,b) = iterate f (0,1) n

in b

fib 5  // f(f(f(f(0,1))))

8  //(0,1)->(1,1)->(1,2)->(2,3)->(3,5)->(5,8)
```

Sorting

Quicksort:

- 1 First element of the list is x and rest is xs
- 2 select smaller elements of xs from x, sort them and put before x.
- 3 select greater elements of xs from x, sort them and put after x.

Taking the reverse

- First element is x rest is xs
- Reverse the xs, append x at the end

Loose time for appending x at the end at each step (N times append of size N).

- Fast version, extra parameter (initially empty list) added:

```
reverse1 [] = []
reverse1 (x:xs) = (reverse1 xs) ++ [x]
reverse2 \times = reverse2, \times [] where
         reverse2 , [] \times = \times
         reverse2 ' (x:xs) y = reverse2 ' xs (x:y)
reverse1 [1..10000]
                               // slow
                               // fast
reverse2 [1..10000]
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 - When recursion at the deepest, return the extra parameter.

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reverse1 [] = []
reverse1 (x:xs) = (reverse1 xs) ++ [x]
reverse2 \times = reverse2 \times 1
         reverse2 , [] \times = \times
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