1.

Robot Mouse Races:

Performance measure: destination, speed

Environment: maze (road, barrier)
Actuators: movement, trace path
Sensors: displacement, distance

Robothespian:

Performance measure: level of human-like behavior

Environment: any object it can contact Actuators: response human-like behavior

Sensors: optical recognize, GPS, sound recognize and so on. Generally, needing all

sensors to recognize external environment change.

2(a).

Solution by HW1 q2(a).cpp (at the bottom of document)

start at: AT

expand: AM AN AS AX IT

choose: AM

expand:

choose: AN

expand: IN ON

choose: AS

expand: IS US

choose: AX

expand: OX

choose: IT

expand: IF

goal at: IN

 $path = AT \Rightarrow AN \Rightarrow IN$

2(b).

If we consider this problem as a graph, we can transit from s1 which is "ab" to s2 which is "12" by at least one step if 'a' equal to '1' or 'b' equal to '2'. Hence, if we want to transit s1 to s2 with 'a' not equal to '1' and 'b' not equal to '2', it needs at least two steps to make it.

By above conclusion, hamming distance can be considered as a heuristic for this problem because number of unmatchable characters between strings means the minimum steps to transit between 2 strings.

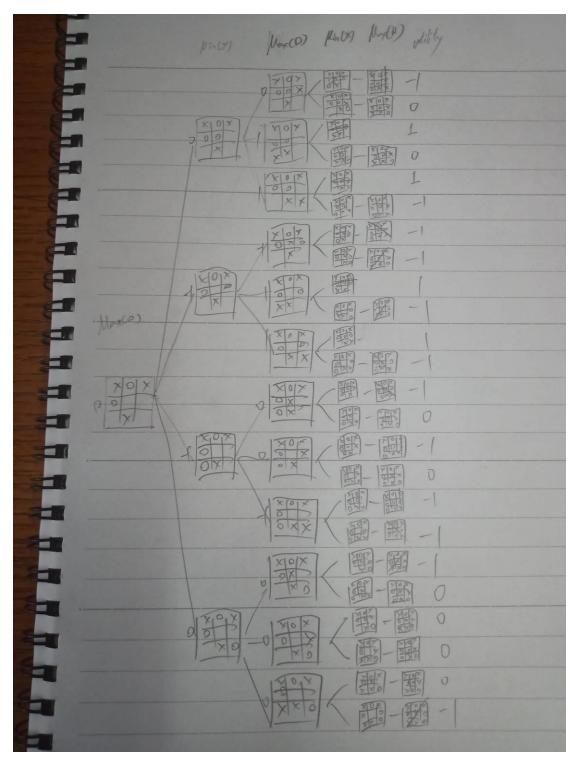
2(c).

Solution by HW1_q2(c).cpp (at the bottom of document)

```
start at: AT
expand: AM AN AS AX IT
choose: AN
expand: IN ON
goal at: IN
AT \Rightarrow AN \Rightarrow IN
3.
Init(): D_X = \{0, 1, ..., 9\}, D_Y = \{0, 1, ..., 9\}, D_Z = \{0, 1, ..., 9\}
Arc checking Loop(time = 1; ; time++):
     time = 1:
     Arc(Y~X): D_X = \{0, 1, ..., 9\}, D_Y = \{0, 1, 2, 3\}, D_Z = \{0, 1, ..., 9\}
     Arc(X \sim Y): D_X = \{0, 1, 4, 9\}, D_Y = \{0, 1, 2, 3\}, D_Z = \{0, 1, ..., 9\}
     Arc(Z\sim X): D_X = \{0, 1, 4, 9\}, D_Y = \{0, 1, 2, 3\}, D_Z = \{0, 1\}
     Arc(X~Z): D_X = \{0, 1\}, D_Y = \{0, 1, 2, 3\}, D_Z = \{0, 1\}
     time = 2:
     Arc(Y~X): D_X = \{0, 1\}, D_Y = \{0, 1\}, D_Z = \{0, 1\}
     Arc(X \sim Y): D_X = \{0, 1\}, D_Y = \{0, 1\}, D_Z = \{0, 1\}
     Arc(Z \sim X): D_X = \{0, 1\}, D_Y = \{0, 1\}, D_Z = \{0, 1\}
     Arc(X \sim Z): D_X = \{0, 1\}, D_Y = \{0, 1\}, D_Z = \{0, 1\}
     time = 3:
     Arc(Y~X): D_X = \{0, 1\}, D_Y = \{0, 1\}, D_Z = \{0, 1\}
     Exit because 3 domains are halting
D_X = \{0, 1\}, D_Y = \{0, 1\}, D_Z = \{0, 1\} \Rightarrow (X, Y, Z) = (0, 0, 0), (1, 1, 1)
4(a).
append 3 more variable b1, b2, b3 where denotes second digit borrows b1 from first
digit, third digit borrows b2 from second digit, last digit borrows b3 from third digit.
0.100I + 10V - 200O = 10U + 10N + R
     simplify from
     1000F + 100I + 10V + E - 1000F - 100O - 10U - R = 100O + 10N + E
1. F, I, V, E, O, U, R, N \in {0, 1, ..., 9}
2. F≠0 (first digit can't be zero)
3. F \neq I \neq V \neq E \neq O \neq U \neq R \neq N
4. b1, b2, b3 \in \{0, 1\}
5. Digit constrain
     F - b1 - F = 0 \Rightarrow b1 = 0
     I - O + 10b1 - b2 = O \Rightarrow I - O - b2 = O
     V - U + 10b2 - b3 = N
     E + 10b3 - R = E
```

```
4(b).
arrange constrain:
     R = 0, b3 = 0
           consider last digit, it may be
           E - R = E => R = 0
           or 10 + E - R = E \Rightarrow R = 10 \Rightarrow conflict to R's domain
     100I + 10V - 200O = 10U + 10N (constrain1)
     F, I, V, E, O, U, N \in \{1, 2, ..., 9\} (domain)
     F \neq I \neq V \neq E \neq O \neq U \neq N (constrain2)
     b2 \in \{0, 1\} \text{ (domain)}
     b1 = 0
     I = 2O + b2 (constrain3)
     V + 10b2 = N + U (constrain4)
explicit var: R = b1 = b3 = 0
implicit var: F, I, V, E, O, U, N, b2
Var(MRV, degree): b2(2, 2), I(9,3), V(9,3), O(9,3), U(9,3), N(9,3), F(9, 1), E(9,1)
Iteration1:
     choose b2 = 0;
     explicit var: R = 0, b1 = b2 = b3 = 0
     stack = \{\$,b2=1\} // \$ is bottom of stack
     constrain1: 100I + 10V - 200O = 10U + 10N
     constrain2: F≠I≠V≠E≠O≠U≠N
     constrain3: I = 2O \Rightarrow I \in \{2, 4, 6, 8\}
     constrain4: V = N + U
     Var(MRV, degree): I(4,3), V(9,3), O(9,3), U(9,3), N(9,3), F(9, 1), E(9,1)
Iteration2:
     choose I = 8;
     explicit var: I = 8, R = 0, b1 = b2 = b3 = 0
     stack = \{\$, b2=1, I=2, I=4, I=6\} // \$ is bottom of stack
     constrain3: 8 = 2O \Rightarrow O \in \{4\} \Rightarrow O = 4 becomes explicit (SOLVE)
     constrain1: 800 + 10V - 800 = 10U + 10N \Rightarrow 10V = 10U + 10N
           \Rightarrow V = N + U \Rightarrow constrain4 (SOLVE)
     constrain2: F \neq V \neq E \neq U \neq N \neq 4, 8 \Rightarrow F, V, E, U, N \in \{1, 2, ..., 9\} - \{4, 8\}
     constrain4: V = N + U
     Var(MRV, degree): V(7,2), U(7,2), N(7,2), F(7, 1), E(7,1)
Iteration3:
     choose V = 9;
     explicit var: I = 8, V = 9, O = 4, R = 0, b1 = b2 = b3 = 0
```

```
stack = \{\$,b2=1,I=2,I=4,I=6,V=1,V=2,V=3,V=5,V=6,V=7\}
     constrain4: 9 = N + U \Rightarrow N, U \in \{2, 3, 6, 7\}
     constrain2: F \neq E \neq U \neq N \neq 4, 8, 9 \Rightarrow F, E, U, N \in \{1, 2, 3, 5, 6, 7\}
     Var(MRV, degree): U(4,2), N(4,2), F(6, 1), E(6,1)
Iteration4:
     choose U = 7:
     explicit var: I = 8, V = 9, O = 4, U = 7, R = 0, b1 = b2 = b3 = 0
     stack = \{\$, b2=1, I=2, I=4, I=6, V=1, V=2, V=3, V=5, V=6, V=7, U=2, U=3, U=6\}
     constrain4: 2 = N \Rightarrow N \in \{2\} \Rightarrow N = 2 becomes explicit (SOLVE)
     constrain2: F \neq E \neq 2, 4, 7, 8, 9 \Rightarrow F, E \in \{1, 3, 5, 6\}
     Var(MRV, degree): F(4, 1), E(4,1)
Iteration4:
     choose F = 6;
     explicit var: F = 6, I = 8, V = 9, O = 4, U = 7, R = 0, N = 2, b1 = b2 = b3 = 0
     stack =
{$,b2=1,I=2,I=4,I=6,V=1,V=2,V=3,V=5,V=6,V=7,U=2,U=3,U=6,F=1,F=3,F=5}
     constrain2: E \neq 2, 3, 4, 7, 8, 9 \Rightarrow E \in \{1, 3, 5\}
     Var(MRV, degree): E(3,1)
Iteration5:
     choose E = 5:
     explicit var: F = 6, I = 8, V = 9, E = 5, O = 4, U = 7, R = 0, N = 2, b1 = b2 = b3 = 0
0
     stack =
{$,b2=1,I=2,I=4,I=6,V=1,V=2,V=3,V=5,V=6,V=7,U=2,U=3,U=6,F=1,F=3,F=5,E=1,
E=5
     constrain2: (EMPTY) => (SOLVE)
     (No unsolved constrain)
Solution: F = 6, I = 8, V = 9, E = 5, O = 4, U = 7, R = 0, N = 2, b1 = b2 = b3 = 0
This solution denotes to 6895 - 6470 = 425
5(a).
```



5(b).

9 nodes, first expansion will return 0, every value smaller or equal to 0 will drop all expansion.

6(a).

A

В

 $P \Rightarrow Q: \neg P \lor Q$

```
L \wedge M \Rightarrow P: \neg L \vee \neg M \vee P
L \wedge B \Rightarrow M: \neg L \vee \neg B \vee M
A \wedge B \Rightarrow L: \neg A \vee \neg B \vee L
A \wedge P \Longrightarrow L: \neg A \vee \neg P \vee L
6(b).
(1)A
(2) B
(3) \neg P \lor Q
(4) \neg L \lor \neg M \lor P
(5) \neg L \lor \neg B \lor M
(6) \neg A \lor \neg B \lor L
(7) \neg A \lor \neg P \lor L
(8) \neg Q // complement of proof target
(9) L by (1)(2)(6)
(10) M by (2)(5)(9)
(11) P by (4)(9)(10)
(12) Q by (3)(11)
(13) False by (8)(12)
Get contradiction, which means Q can't be False.
We can say Q is True by this proof.
Appendix:
HW1_q2(a).cpp:
#include <iostream>
#include <vector>
#include <algorithm>
#include <queue>
using namespace std;
class Frontier{
public:
      int pos,pos_p;
      Frontier(int p,int p_p){
            pos = p;
            pos_p = p_p;
      }
};
int main(int argc,char const *argv[]){
```

```
string sa[] =
{"AN","AM","AS","AT","AX","BE","BY","GO","HE","HI","IT","IS","IN","IF","ME
","MY","NO","OF","OH","OK","ON","OR","OX","SO","TO","UP","US","WE"};
     vector<string> v(sa,sa + sizeof(sa) / sizeof(string));
     sort(v.begin(), v.end());
     vector<int> parent(v.size(),-1);
     vector<bool> inq(v.size(),false);
     int start, goal;
     for(int i = 0; i < v.size(); i++){
          if(v[i] == "AT")
               start = i;
          else if(v[i] == "IN")
               goal = i;
     }
     queue<Frontier> q;
     q.push(Frontier(start,9999)); // no parent
     inq[start] = true;
     while(!q.empty()){
          Frontier node = q.front();
          q.pop();
          parent[node.pos] = node.pos_p;
          string s = v[node.pos];
          cout << "choose: " << s << '\n';
          if(s == v[goal])
               break;
          cout << "expand: ";</pre>
          for(int i = 0;i < v.size();i++)
               if(!inq[i] \&\& (s[0] == v[i][0] || s[1] == v[i][1])){
                    inq[i] = true;
                    cout << v[i] << ' ';
                    q.push(Frontier(i,node.pos));
          cout << '\n';
     }
     vector<int> path;
     int pos = goal;
     while(true){
          if(pos == 9999)
```

```
break;
          path.push_back(pos);
          pos = parent[pos];
     }
     cout << v[path[path.size() - 1]];</pre>
     for(int i = path.size() - 2; i >= 0; i--)
          cout << " => " << v[path[i]];
     cout << '\n';
     return 0;
}
HW1_q2(c).cpp:
#include <iostream>
#include <vector>
#include <algorithm>
#include <queue>
using namespace std;
int h(string s1,string s2){ // heuristic function
     int n = s1.size();
     for(int i = 0;i < s1.length();i++)
          if(s1[i] == s2[i])
               n--;
     return n;
}
class Frontier{
public:
     int pos,pos_p,fval,gval;
     Frontier(int p,int p p,int f){
          pos = p;
          pos_p = p_p;
          fval = f;
     }
     friend bool operator>(const Frontier & obj1,const Frontier & obj2){
          return obj1.fval > obj2.fval;
     }
};
int main(int argc,char const *argv[]){
```

```
string sa[] =
{"AN","AM","AS","AT","AX","BE","BY","GO","HE","HI","IT","IS","IN","IF","ME
","MY","NO","OF","OH","OK","ON","OR","OX","SO","TO","UP","US","WE"};
     vector<string> v(sa,sa + sizeof(sa) / sizeof(string));
     sort(v.begin(), v.end());
     vector<int> parent(v.size(),-1);
     vector<bool> inq(v.size(),false);
     int start, goal;
     for(int i = 0; i < v.size(); i++){
          if(v[i] == "AT")
               start = i;
          else if(v[i] == "IN")
               goal = i;
     }
     priority queue<Frontier, vector<Frontier>, greater<Frontier>> q;
     q.push(Frontier(start,9999,h(v[start],v[goal]))); // no parent
     inq[start] = true;
     while(!q.empty()){
          Frontier node = q.top();
          q.pop();
          parent[node.pos] = node.pos_p;
          string s = v[node.pos];
          cout << "choose: " << s << '\n';
          if(s == v[goal])
               break;
          cout << "expand: ";</pre>
          for(int i = 0;i < v.size();i++)
               if(!inq[i] \ \&\& \ (s[0] == v[i][0] \ \| \ s[1] == v[i][1])) \{
                    inq[i] = true;
                    cout << v[i] << ' ';
                    q.push(Frontier(i,node.pos,h(v[i],v[goal])));
          cout << '\n';
     }
     vector<int> path;
     int pos = goal;
     while(true){
          if(pos == 9999)
```

```
break;
    path.push_back(pos);
    pos = parent[pos];
}
cout << v[path[path.size() - 1]];
for(int i = path.size() - 2;i >= 0;i--)
    cout << " => " << v[path[i]];
cout << '\n';
return 0;
}</pre>
```