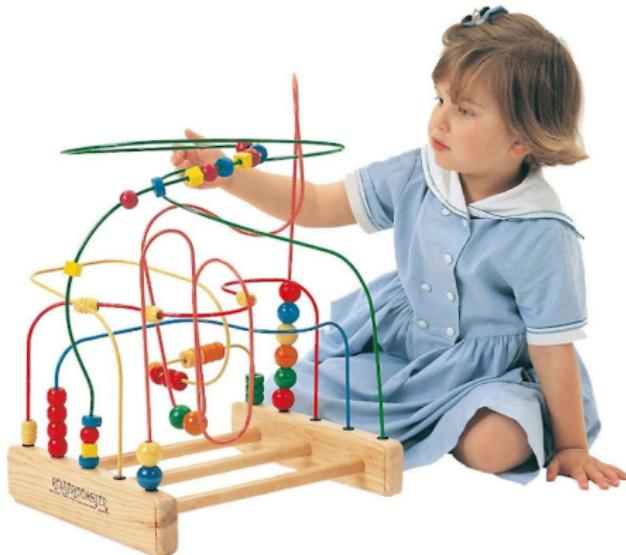


Coxeter groups and diagram algebras



Definition

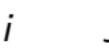
A *Coxeter system* is a group with presentation

$$\langle s_1, \dots, s_n \mid (s_i s_j)^{M_{ij}} = 1 \rangle$$

where M is a symmetric $n \times n$ matrix with $M_{ii} = 1$ and $M_{ij} \in \{2, 3, \dots\}$. n is called the *rank* of the group.

- When $i = j$, this means $s_i^2 = 1$ for all i , thus generators are idempotent and self-inverse
- When $M_{ij} = 2$, this means $(s_i s_j)^2 = 1 \Rightarrow s_i s_j = s_j s_i$, i.e. s_i and s_j commute
- When $M_{ij} \geq 3$, s_i and s_j are said to have a *braid relation*:
 - ▶ $(s_i s_j)^3 = 1$ implies $s_i s_j s_i = s_j s_i s_j$
 - ▶ $(s_i s_j)^4 = 1$ implies $s_i s_j s_i s_j = s_j s_i s_j s_i$

Instead of a matrix M we can use a edge-labeled graph on $\{1, \dots, n\}$ where

-  means s_i, s_j commute
-  means s_i, s_j have a 3-braid relation
-  means s_i, s_j have a k -braid relation

Coxeter system of type A_n :



i.e., adjacent generators have a 3-braid relation and all other pairs commute.

We represent elements of a Coxeter group by *fully reduced* words in the generators.

Theorem (Matsumoto)

For any element $w \in W$ of a Coxeter group, all reduced words for W have the same length. Furthermore, any reduced word is related to any other by a sequence of braid relations.

In type A_4 :

- $s_3 s_1 s_2 s_1 s_4 \rightarrow s_3 s_2 s_1 s_2 s_4 \rightarrow s_3 s_2 s_1 s_4 s_2 \rightarrow s_3 s_2 s_4 s_1 s_2 \rightarrow s_3 s_4 s_2 s_1 s_2 \rightarrow s_3 s_4 s_1 s_2 s_1$ are reduced words representing the same element.

The other Coxeter types we will be concerned with in this talk are closely related to A_n :

B_n :



H_n :



These look innocent, but blow up combinatorially:

- A_4 has 120 elements
- B_4 has 384 elements
- H_4 has 14,440 elements!

Given a Coxeter group $W = \langle s_1, \dots, s_n : \dots \rangle$, we can define the *Hecke algebra* of W :

Let $\mathcal{A} = \mathbb{Z}[v, v^{-1}]$ be the ring of Laurent polynomials over the integers. \mathcal{A} is a commutative ring.

Now we will define an \mathcal{A} -algebra by a certain presentation.

Definition

The *Hecke algebra* of W is the (unital, associative) \mathcal{A} -algebra with presentation

$$H(W) = \langle T_1, \dots, T_n : (\text{braid relations}), (T_i - v)(T_i + v^{-1}) = 0 \rangle$$

Note that $(T_i - v)(T_i + v^{-1}) = 0 \Leftrightarrow T_i^2 = (v - v^{-1})T_i + 1$, so we can still “get rid of any squares”.

It turns out $H(W)$ has a very simple free basis:

Theorem

$H(W)$ is free over the basis

$$\{ T_w = T_{i_1} \dots T_{i_k} : s_{i_1} \dots s_{i_k} \text{ is any reduced expression for } w \in W \}$$

Indexed by the group elements of W .

i.e., in $H(A_4)$ we have $T_{232} = T_2 T_3 T_2 = T_3 T_2 T_3 = T_{323}$ (since $s_2 s_3 s_2 = s_3 s_2 s_3$ represent the same element in W).

Typical elements look like:

- $(3v^2 - v^{-1}) T_{121} = (3v^2 - v^{-1}) T_{212}$
- $v T_1 + v^{-1} T_{24} = v T_1 + v^{-1} T_{42}$
- $T_1(v + T_{24}) = T_1(v + T_2 T_4) = v T_1 + T_1 T_2 T_4 = v T_1 + T_{124}$

In general, $H(W)$ has many bases (all necessarily indexed by the elements of W). But one is of particular interest:

Theorem (Kazdahn, Lusztig)

There exists a unique basis $\{C_w : w \in W\}$ for H satisfying some precise technical properties called the *Kazdahn-Lusztig basis*.

When a C -basis element is expanded in terms of the T -basis, the coefficients (or “structure constants”) are the *Kazdahn-Lusztig polynomials*:

$$C_w = T_w + \sum_{y < w} p_{y,w} T_y$$

In particular, $H(W)$ is generated by $C_i = T_i + v^{-1}$.

Warning

This basis is far more complicated than the T -basis! $C_{i_1 i_2 \dots i_k}$ is usually *very different* from $C_{i_1} C_{i_2} \dots C_{i_k}$! We cannot just do arithmetic/multiplication “trivially” like we could in the T -basis.

Now we will take a quotient of $H(W)$ by a certain ideal:

$$\mathcal{I}(W) = \langle C_{\text{long braids}} \rangle = \langle C_{ij\dots i} : s_i s_j \dots s_i \text{ is a long braid} \rangle$$

We'll call the resulting quotient $TL(W) = H(W)/\mathcal{I}(W)$ for reasons that will be explained shortly.

How can we figure out what $TL(W)$ looks like?

- Since the C_i generate $H(W)$, their equivalence classes $U_i = \pi(C_i)$ generate the quotient.
- If we can figure out what relations the U_i have we can present $TL(W)$ by generators and relations.

Let's do this "concretely" in type A_n , where adjacent numbers have a 3-braid relation, and all others commute. We have

$$\mathcal{I}(A_n) = \langle C_{i(i+1)i} : 1 \leq i < n \rangle = \langle C_{121}, C_{232}, \dots \rangle$$

Next, recalling that $C_i = T_i + v^{-1}$, you may calculate:

$$C_i^2 = (T_i + v^{-1})^2 = (v + v^{-1})T_1 + 1 + v^{-2} = (v + v^{-1})C_i$$

You can verify by another straightforward calculation that $C_i C_k = C_k C_i$ for non-adjacent, i, k .

Finally, the following calculation requires knowing what C_{121} is, but if you knew you would compute (likewise for all indices):

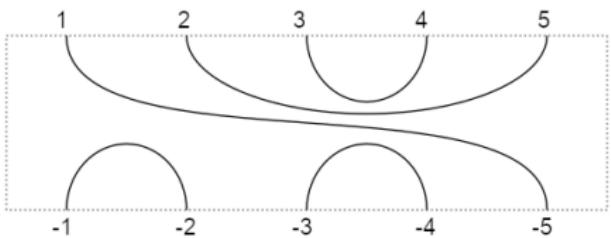
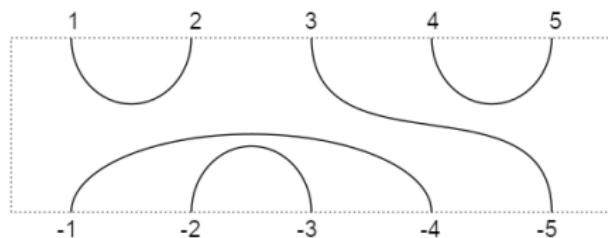
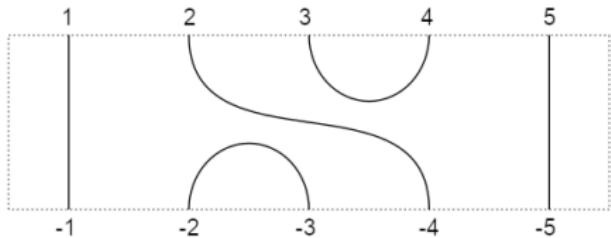
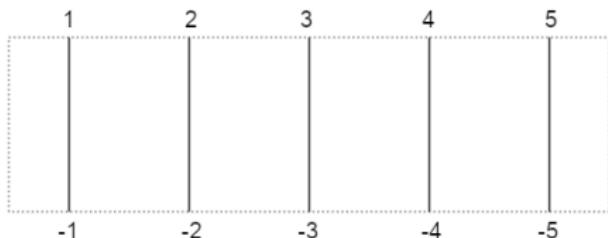
$$\begin{aligned} C_1 C_2 C_1 &= (T_1 + v^{-1})(T_2 + v^{-1})(T_1 + v^{-1}) \\ &= \color{red}{T_{121} + v^{-1}T_{21} + v^{-1}T_{12} + v^{-2}T_1 + v^{-2}T_2 + v^{-3}} + \color{blue}{T_1 + v^{-1}} \\ &= C_{121} + C_1 \\ &\equiv C_1 \quad (\text{in the quotient}) \end{aligned}$$

So to summarize, we've "showed" that $TL(A_n)$ is abstractly presented by generators and relations in the following way:

$$\begin{aligned} \langle U_1, \dots, U_n : & U_i^2 = (\nu + \nu^{-1})U_i, \\ & U_i U_k = U_k U_i \text{ for non-adjacent } i, k, \\ & U_i U_j U_i = U_i \text{ for adjacent } i, j \rangle \end{aligned}$$

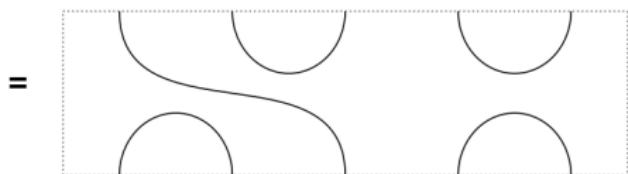
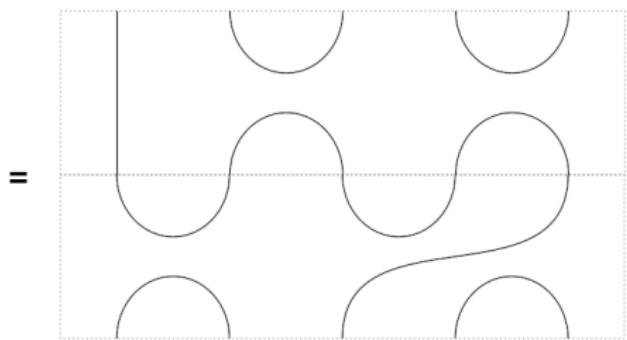
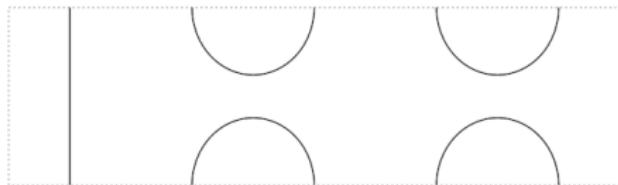
It turns out that this particular presentation reveals an isomorphism with something much more concrete...

Consider a *non-crossing pairing* of $n + 1$ “north” nodes with $n + 1$ “south” nodes. For example, when $n = 4$:



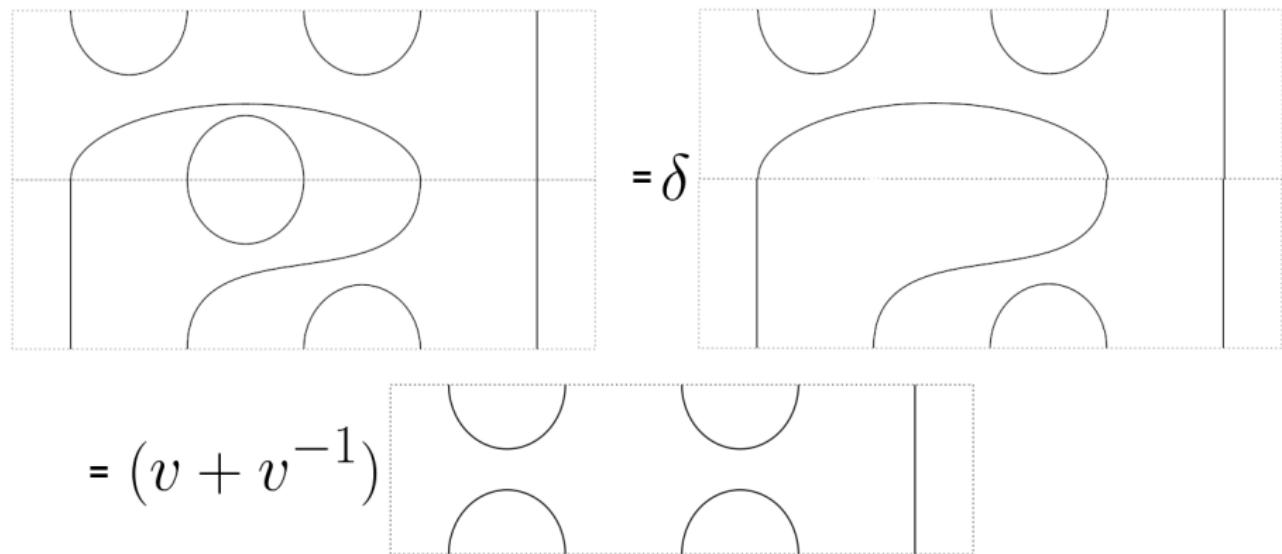
As suggested by the image, we also think of this as non-crossing 2-ary partition diagrams of the set $\{-(n+1), \dots, 1, 1, \dots, n+1\}$.

We can “multiply” these diagrams via vertical concatenation:



Form the free \mathcal{A} -algebra over these diagrams (i.e. formal \mathcal{A} -linear combinations of diagrams), and where multiplication is by vertical concatenation + “reduction rules”. Right now, the only reduction rule is

- any closed loops formed in a concatenation “come out” as a scalar multiplication by $\delta = v + v^{-1}$:



This is the *Temperley Lieb Algebra*.

Claim: The following diagrams generate all the diagrams:

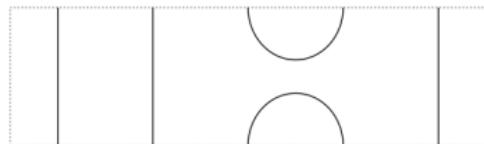
$$U_1 =$$



$$U_2 =$$



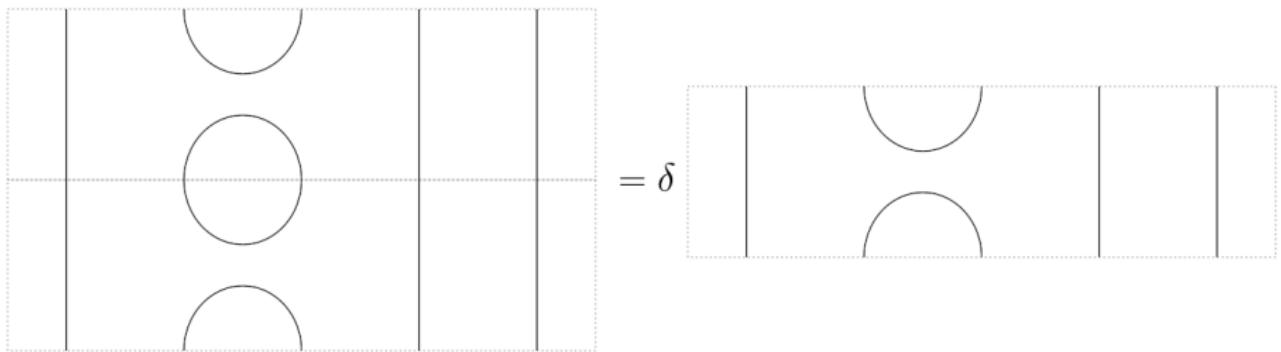
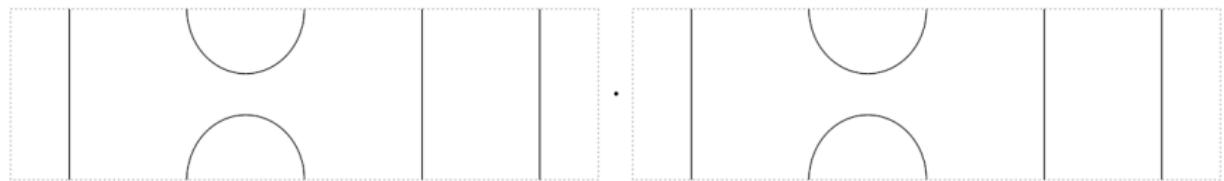
$$U_3 =$$



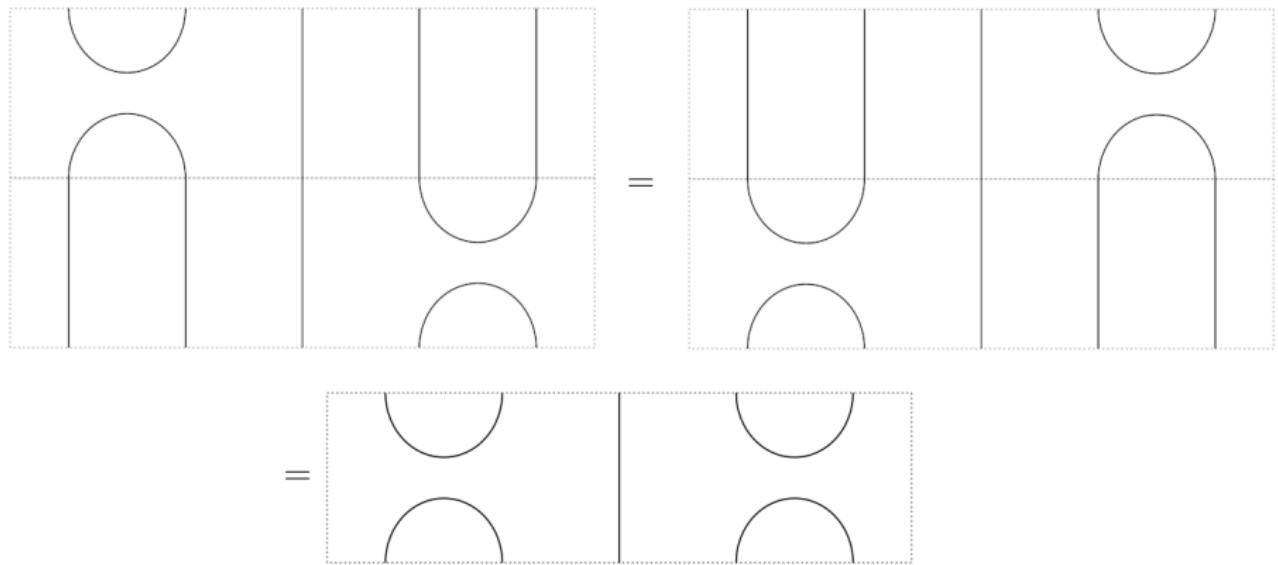
$$U_4 =$$



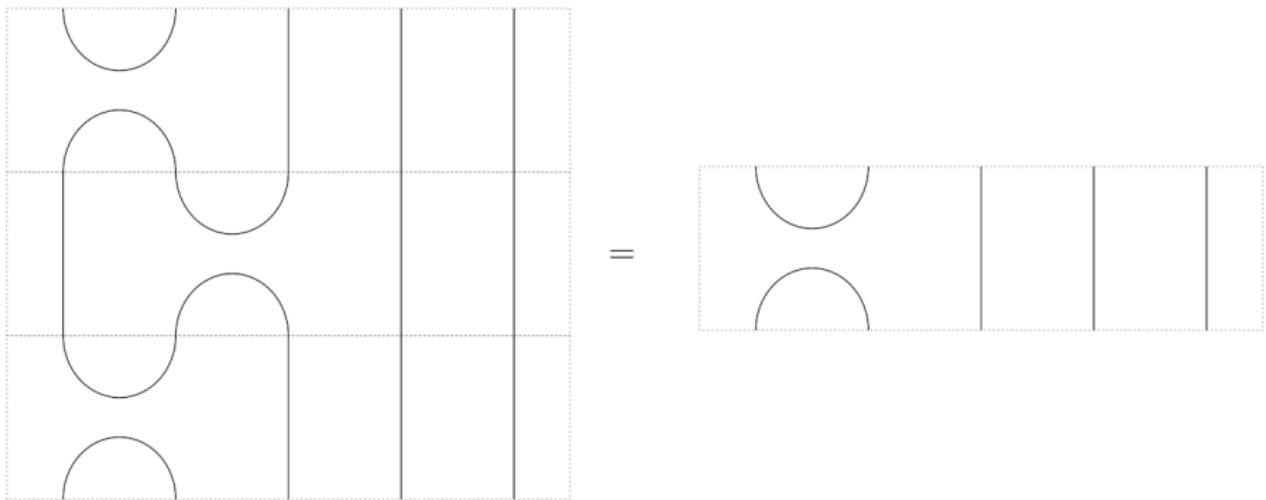
Furthermore, notice $U_i^2 = (v + v^{-1})U_i$



And $U_i U_k = U_k U_i$ for non-adjacent, U_i, U_k :



Finally, $U_i U_{(i+1)} U_i = U_i$ for adjacent i, j :



In fact, $TL(A_n)$ is isomorphic to the Temperley-Lieb algebra of diagrams with $n + 1$ nodes, which has been well known for a while.

One can form the *Generalized Temperley-Lieb algebras* $TL(W)$ by the same construction as a quotient of the Hecke algebra.

Question: Is there a diagram algebra realization for the other Generalized TL algebras?

One is known for types A, B, H, D, E and \widetilde{C}_n .

If you repeat the whole process for type B_n , you'll get the following presentation of $TL(B_n)$.

$$\begin{aligned} \langle U_1, \dots, U_n : & U_i^2 = (v + v^{-1})U_i, \\ & U_i U_k = U_k U_i \text{ for non-adjacent } i, k, \\ & U_i U_j U_i = U_i \text{ for adjacent } i, j \text{ and } \{i, j\} \neq \{1, 2\}, \\ & U_1 U_2 U_1 U_2 = 2U_1 U_2 \rangle \end{aligned}$$

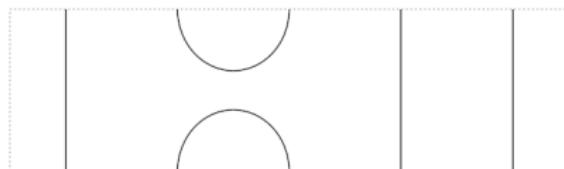
This can be realized by “decorated” Temperley-Lieb diagrams, where arcs are allowed to carry a “decoration” (according to certain rules).

$TL(B_4)$ is generated by

$$U_1 = 2$$



$$U_2 =$$



$$U_3 =$$



$$U_4 =$$



Multiplication is still concatenation + “reduction rules”, but now there are more reduction rules:

- Replace any instance of 2 decorations with 1 decoration.
- An undecorated loop comes out as δ .
- A decorated loop comes out as $\delta/2$.

$$\text{---} = \delta$$

$$\text{---} = \delta/2$$

$$\begin{array}{c} | \\ \bullet \\ | \end{array} = \begin{array}{c} | \\ \bullet \end{array}$$

This time we'll only check the new relation $U_1 U_2 U_1 U_2 = 2U_1 U_2$.

$$\begin{array}{c} \text{Diagram 1: } \\ \text{Diagram 2: } \\ \text{Diagram 3: } \\ \text{Diagram 4: } \end{array} = 4 \quad \begin{array}{c} \text{Diagram 5: } \\ \text{Diagram 6: } \end{array} = 4$$

The diagram illustrates the relation $U_1 U_2 U_1 U_2 = 2U_1 U_2$ using string diagrams. The left side shows four diagrams labeled 1 through 4. Diagram 1 is a single vertical line with a dot at height 2. Diagram 2 is a vertical line with a loop above it, and a dot at height 1. Diagram 3 is a vertical line with a loop below it, and a dot at height 1. Diagram 4 is a vertical line with a loop below it, and a dot at height 0. The right side shows two pairs of diagrams, each pair summing to 4. The first pair (Diagrams 5 and 6) consists of a vertical line with a loop above it and a dot at height 2, plus a vertical line with a loop below it and a dot at height 2. The second pair (Diagrams 5 and 6) consists of a vertical line with a loop above it and a dot at height 1, plus a vertical line with a loop below it and a dot at height 1.

If you do this for type H_n , you'll get the following presentation for $TL(H_n)$:

$$\langle U_1, \dots, U_n : U_i^2 = (v + v^{-1})U_i,$$

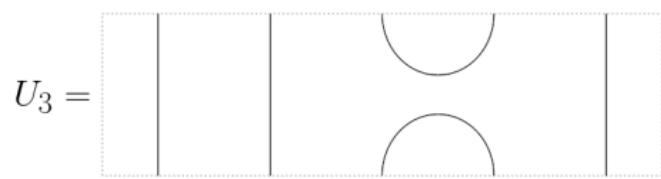
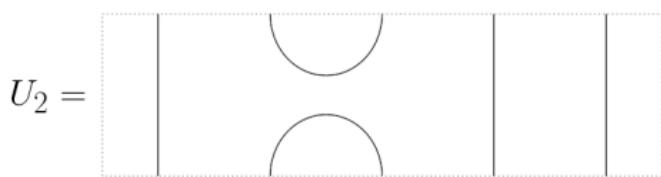
$$U_i U_k = U_k U_i \text{ for non-adjacent } i, k,$$

$$U_i U_j U_i = U_i \text{ for adjacent } i, j \text{ and } \{i, j\} \neq \{1, 2\},$$

$$U_1 U_2 U_1 U_2 U_1 = 3U_1 U_2 U_1 - U_1 \rangle$$

The diagram realization is the same as the previous one with different reduction rules.

$TL(H_4)$ is generated by



Reduction rules for H_n are

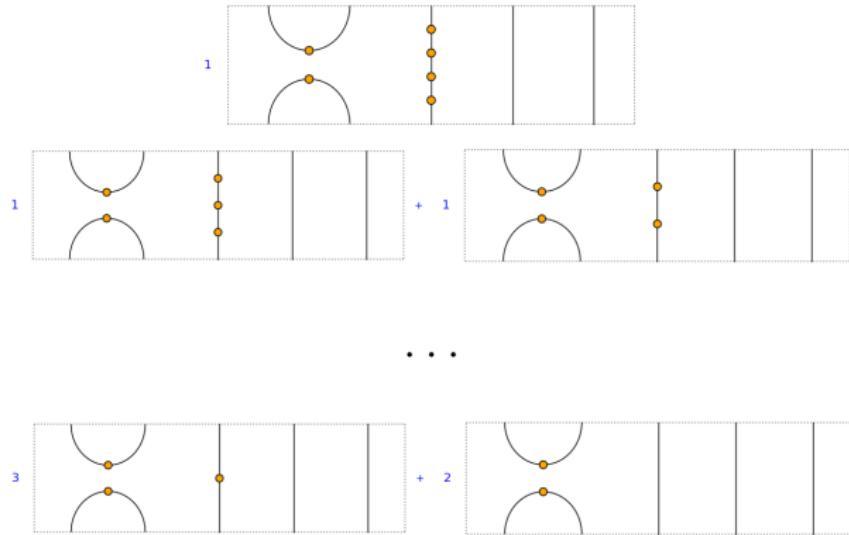
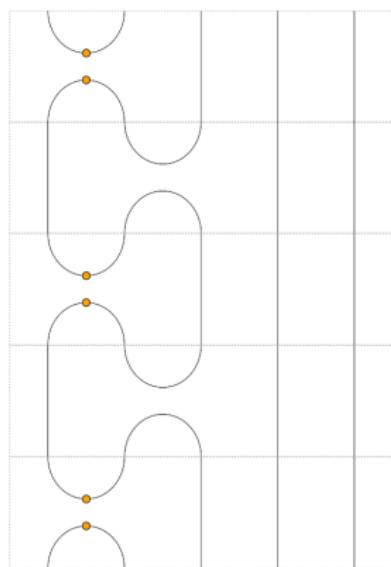
- If a diagram has an edge with 2 decorations, split it into two copies that have 1 and 0 decorations on that edge, respectively.
- An undecorated loop comes out as δ
- A decorated loop comes out as 0 (i.e., the whole diagram is gone)

$$\text{---} = \delta$$

$$\text{---} = 0$$

$$\begin{array}{c} | \\ \bullet \\ | \end{array} = \begin{array}{c} | \\ \bullet \\ | \end{array} + \begin{array}{c} | \end{array}$$

This time, I'll just show a computation of $U_1 U_2 U_1 U_2 U_1$, and you can check™ that it's the same as what you get from $3U_1 U_2 U_1 - U_1$.



Ultimately, all such diagram realizations are of the same “scheme”, where the basis elements are decorated Temperley-Lieb diagrams and multiplication is by concatenation + “decoration reduction rules”.

Recall that $H(W)$ has bases indexed by the elements of W . It turns out:

Theorem

Bases of $TL(W)$ are indexed by the *fully-commutative* elements of W ; $w \in W$ is *fully commutative* if any reduced word for w can be transformed into any other reduced word by using only *commutation* relations (no longer braid relations); equivalently, no reduced word for w contains a long braid.

In particular, any basis of $TL(W)$ is in bijection with $W_c =$ the fully commutative elements of W .

As far as bases of $TL(W)$ go, there are particular ones of interest. Certain algebras over $\mathcal{A} = \mathbb{Z}[v, v^{-1}]$ may have a *canonical basis* or *IC-basis*, a term coming from representation theory (but the existence or uniqueness is not guaranteed)

An *example* of such a canonical basis is the Kazhdan-Lusztig basis of the $H(W)$, but it is not the case that a canonical basis for $TL(W)$ can simply be obtained from the projection of the Kazhdan-Lusztig basis of $H(W)$.

Question: Can the canonical basis be understood in terms of diagrams? If so, can we describe the bijection between fully commutative elements and their “canonical diagrams”?

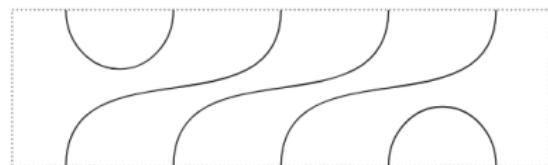
In cases A , B , and H , the bijection is non-trivial but can be described combinatorially. In cases A and H , the canonical diagram basis can be described combinatorially as “admissible diagrams” satisfying certain restrictions. The same is true in case B , but it requires some additional creativity.

Some canonical diagrams in $TL(A_4)$ with their corresponding FC word:
<https://math.colorado.edu/~chme3268/diagrams>)

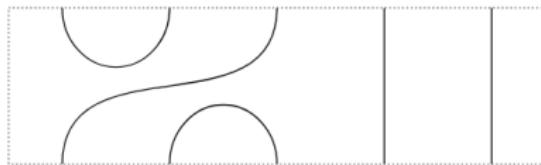
(4)



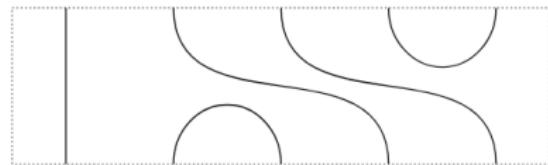
(1, 2, 3, 4)



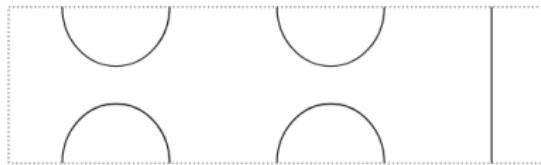
(1, 2)



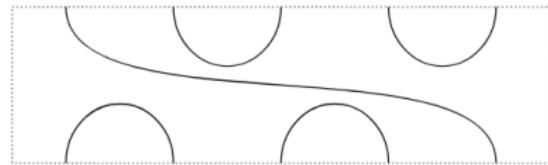
(4, 3, 2)



(1, 3)

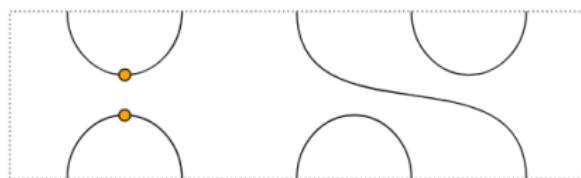


(2, 4, 1, 3)

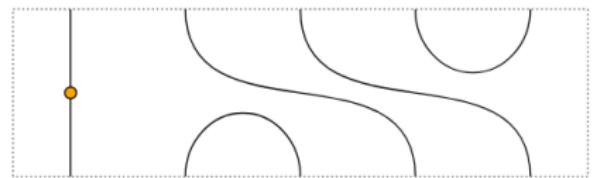


Some canonical diagrams in $TL(H_4)$:

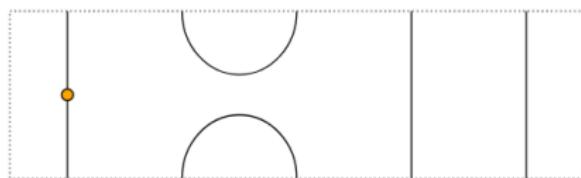
(1, 4, 3)



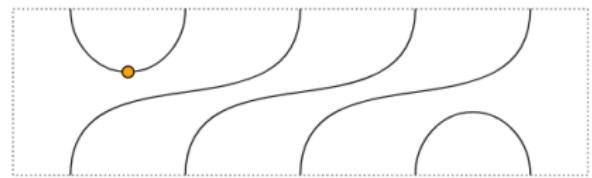
(4, 3, 2, 1, 2)



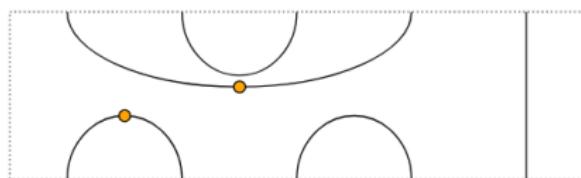
(2, 1, 2)



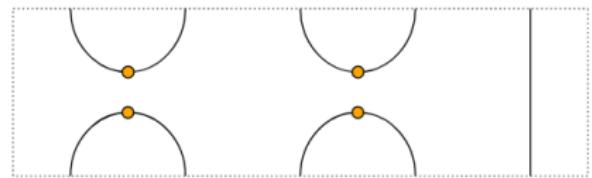
(1, 2, 1, 2, 3, 4)



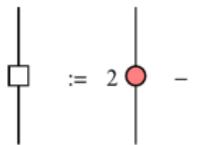
(2, 1, 3)



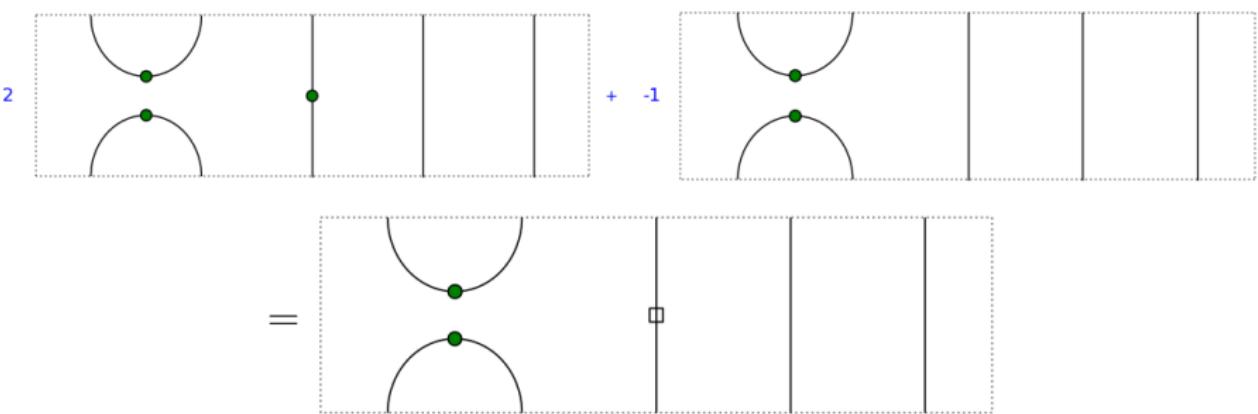
(1, 2, 1, 3, 2, 1)



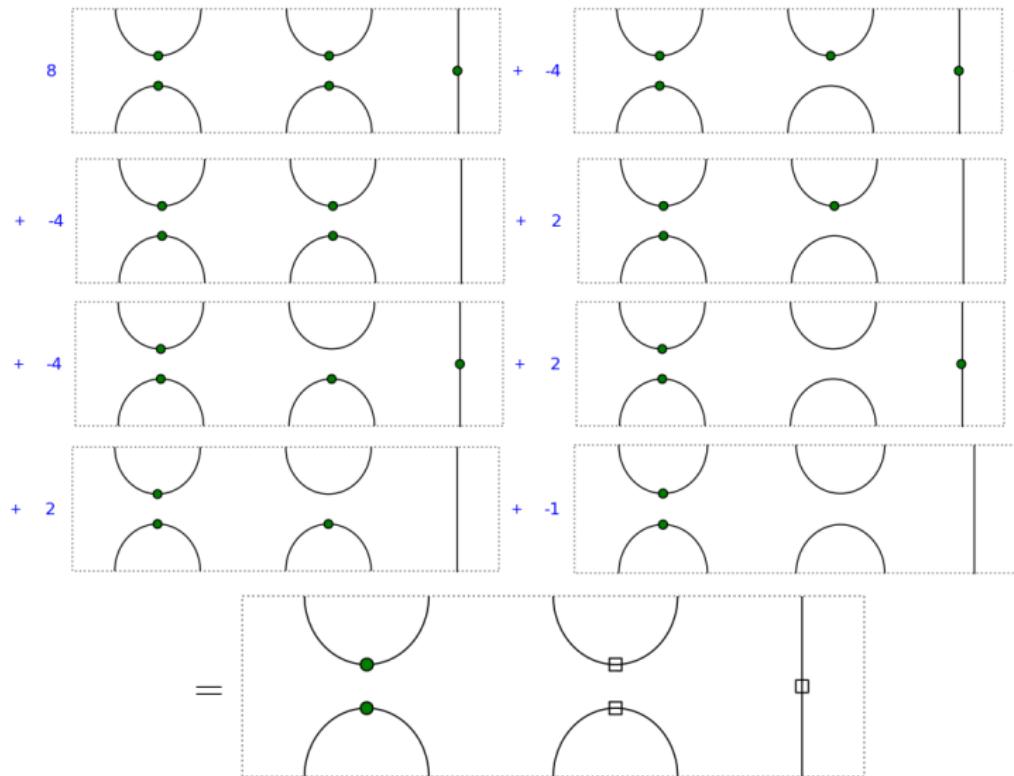
In type B , in order to describe the canonical basis, we need to invent a notational shorthand: the “square” decoration:



Like the reduction rule in type H , this takes place on “whole diagrams”; for example:

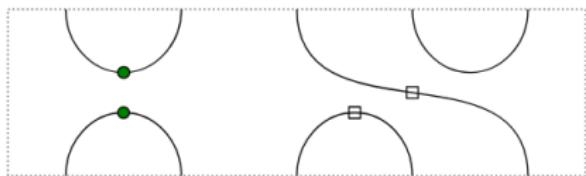


This notational device can compound to express unwieldy linear combinations as single diagrams:

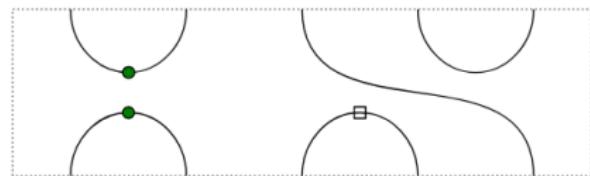


Using the square we can much more easily describe the canonical diagrams in type B :

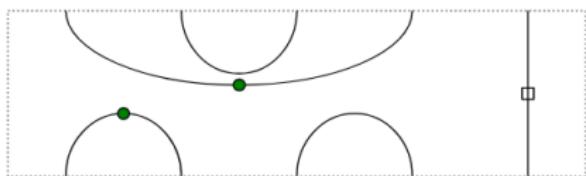
$(1, 4, 2, 1, 3, 2, 1)$



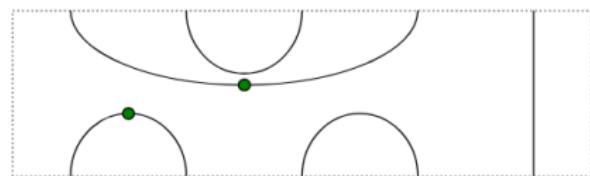
$(1, 4, 3, 2, 1)$



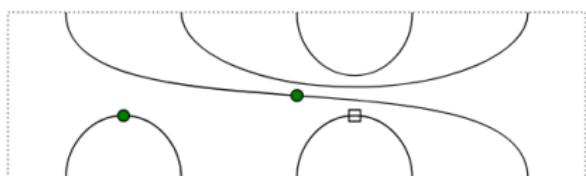
$(2, 1, 3, 2, 4, 1, 3)$



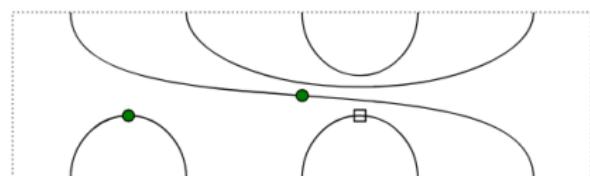
$(2, 1, 3)$



$(3, 2, 4, 1, 3, 2, 1)$



$(3, 2, 4, 1, 3, 2, 1)$



There is a certain involution on $W_c =$ the fully commutative elements of W that has a very complicated definition.

Definition

The *Mathas-Lusztig involution* $\lambda : W \rightarrow W$ can be defined on $w \in W$ as follows: Let T_{w_0} be the T -basis element in $H(W)$ corresponding to the longest element $w_0 \in W$. Compute $T_{w_0} C_w$ and expand the result in terms of the C -basis:

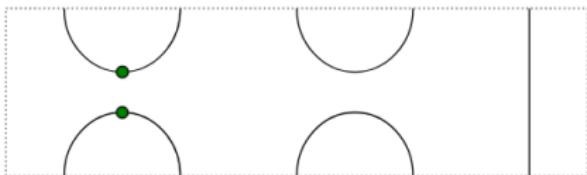
$$T_{w_0} C_w = \sum_{y \in W} \alpha_{y,w} C_y$$

There exists a unique element y in the same *left Kazdahn-Lusztig cell* as w such that $\alpha_{y,w} \neq 0$, and $\lambda(w) = y$.

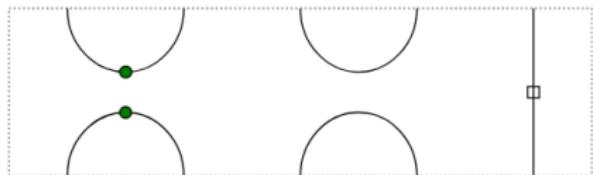
Can this involution be better understood under the bijection with diagrams?

One conjecture supported with evidence is that in type B , when the ML involution is *not* the identity (it may often be), it appears to have an interpretation as “toggling a square” on the diagram:

(1, 3)



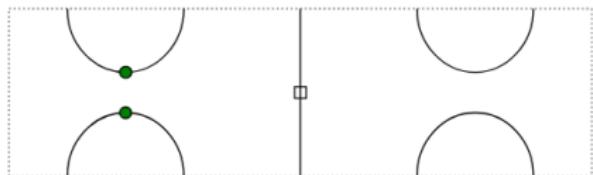
(1, 3, 2, 4, 1, 3)



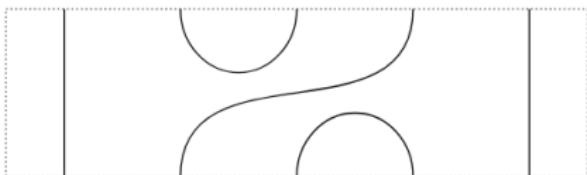
(1, 4)



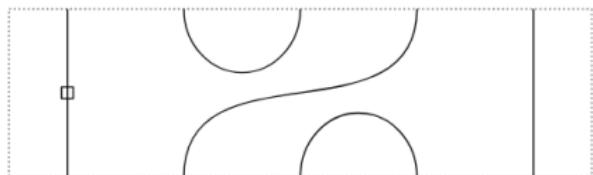
(1, 4, 2, 1)



(2, 3)



(2, 1, 2, 3)



Just for fun, I've included the decorations and reduction rules for type \tilde{C} , but we won't talk about it.

$$(1) \quad \begin{array}{c} \bullet \\ | \\ \bullet \end{array} = \begin{array}{c} \blacktriangle \\ | \end{array};$$

$$(2) \quad \begin{array}{c} \circ \\ | \\ \circ \end{array} = \begin{array}{c} \blacktriangle \\ | \end{array};$$

$$(3) \quad \begin{array}{c} \blacktriangle \\ | \\ \bullet \end{array} = \begin{array}{c} \bullet \\ | \\ \blacktriangle \end{array} = 2 \begin{array}{c} \bullet \\ | \end{array};$$

$$(4) \quad \begin{array}{c} \blacktriangle \\ | \\ \circ \end{array} = \begin{array}{c} \circ \\ | \\ \blacktriangle \end{array} = 2 \begin{array}{c} \circ \\ | \end{array};$$

$$(5) \quad \text{Diagram} = \text{Diagram} \blacktriangle = \text{Diagram} \blacktriangle = \delta;$$



Dana C. Ernst. “A diagrammatic representation of an affine $\mathbb{C}\$$ Temperley–Lieb algebra”. In: *arXiv:0905.4457 [math]* (May 2009). arXiv: 0905.4457 version: 1. URL: <http://arxiv.org/abs/0905.4457> (visited on 09/24/2020).



R. M. Green. “Cellular algebras arising from Hecke algebras of type H_n ”. In: *arXiv:q-alg/9712019* (Dec. 1997). arXiv: q-alg/9712019 version: 1. URL: <http://arxiv.org/abs/q-alg/9712019> (visited on 09/24/2020).



R. M. Green. “Decorated tangles and canonical bases”. In: *arXiv:math/0108076* (Aug. 2001). arXiv: math/0108076 version: 1. URL: <http://arxiv.org/abs/math/0108076> (visited on 09/24/2020).



R. M. Green. “Generalized Temperley-Lieb algebras and decorated tangles”. In: *arXiv:q-alg/9712018* (Dec. 1997). arXiv: q-alg/9712018 version: 1. URL: <http://arxiv.org/abs/q-alg/9712018> (visited on 09/24/2020).



R. M. Green and J. Losonczy. “Canonical Bases for Hecke Algebra Quotients”. In: *arXiv:math/9904045* (Apr. 1999). arXiv: math/9904045 version: 1. URL: <http://arxiv.org/abs/math/9904045> (visited on 09/24/2020).