

DISCRETE MATHEMATICS

2006 – 2007 Spring Term

RECITATION 1

16.02.2007

Text book – Exercise 2.2 – Questions: 4,14

4. Simplify following statement

$$[[[(p \wedge q) \wedge r] \vee [(p \wedge r) \wedge \neg r] \vee \neg q] \rightarrow s]$$

$$p \rightarrow q = \neg p \vee q$$

$$\begin{aligned} &= [[[(p \wedge q) \wedge r] \vee [p \wedge (r \wedge \neg r)] \vee \neg q] \rightarrow s] \\ &= [[[(p \wedge q) \wedge r] \vee [p \wedge \mathbf{F}] \vee \neg q] \rightarrow s] \\ &= [[[(p \wedge q) \wedge r] \vee \mathbf{F}] \vee \neg q] \rightarrow s \\ &= [(p \wedge q) \wedge r] \vee \neg q \rightarrow s \\ &= [(p \wedge q \wedge r) \vee \neg q] \rightarrow s \\ &= [(p \vee \neg q) \wedge (q \vee \neg q) \wedge (r \vee \neg q)] \rightarrow s \\ &= [(p \vee \neg q) \wedge \mathbf{T} \wedge (r \vee \neg q)] \rightarrow s \\ &= [(p \vee \neg q) \wedge (r \vee \neg q)] \rightarrow s \\ &= \neg [(p \vee \neg q) \wedge (r \vee \neg q)] \vee s \\ &= \neg [(p \wedge r) \vee \neg q] \vee s \\ &= [\neg(p \wedge r) \wedge q] \vee s \end{aligned}$$

14. a) Show that $p \rightarrow [q \rightarrow (p \wedge q)]$ is a tautology

$$\begin{aligned} &= \neg p \vee [\neg q \vee (p \wedge q)] \\ &= \neg p \vee [(\neg q \vee p) \wedge (\neg q \vee q)] \\ &= \neg p \vee [(\neg q \vee p) \wedge \mathbf{T}] \\ &= \neg p \vee (\neg q \vee p) \\ &= (\neg p \vee p) \vee \neg q \\ &= \mathbf{T} \vee \neg q \\ &= \mathbf{T} \end{aligned}$$

b) Show that $(p \vee q) \rightarrow [q \rightarrow q]$ is a tautology using the result in the previous question.

Rewrite the last result we obtained like $p' \rightarrow [q' \rightarrow (p' \wedge q')]$

In our current statement, we can apply substitution like

$$p' = (p \vee q)$$

$$q' = q$$

If statement $(p' \wedge q')$ equals q , then we can deduce that our initial statement is a tautology

$$(p' \wedge q') = (p \vee q) \wedge q = (p \wedge q) \vee (q \wedge q) = (p \wedge q) \vee q$$

Using law of absorption

$$(p \wedge q) \vee q = q$$

$p \vee (p \wedge q) \leftrightarrow p$	Law of
$p \wedge (p \vee q) \leftrightarrow p$	Absorption

c) Test if $(p \vee q) \rightarrow [q \rightarrow (p \wedge q)]$ is tautology.

$$\begin{aligned} &= \neg(p \vee q) \vee [\neg q \vee (p \wedge q)] \\ &= (\neg p \wedge \neg q) \vee [(\neg q \vee p) \wedge (\neg q \vee q)] \\ &= (\neg p \wedge \neg q) \vee [(\neg q \vee p) \wedge \mathbf{T}] \\ &= (\neg p \wedge \neg q) \vee (\neg q \vee p) \\ &= (\neg p \wedge \neg q) \vee p \vee \neg q \\ &= [(\neg p \vee p) \wedge (\neg q \vee p)] \vee \neg q \\ &= [\mathbf{T} \wedge (\neg q \vee p)] \vee \neg q \\ &= (\neg q \vee p) \vee \neg q \\ &= \neg q \vee \neg q \vee p \\ &= \neg q \vee p \quad \text{NOT TAUTOLOGY!} \end{aligned}$$

Text book – Exercise 2.3 – Questions: 4,8

4. Deduce results from the following verbal statements using Modus Ponens or Modus Tollens

- a)** \rightarrow If Tolga can't fix the computer Tahir will arrive and check the computer
 \rightarrow Tolga wasn't able to fix the computer.

Modus Ponens	
$\frac{p \rightarrow q \quad p}{\therefore q}$	$(p \rightarrow q) \wedge p \Rightarrow q$

∴ Tahir arrives and checks the computer. **Modus Ponens**

- b) → If Ayça has solved the problem she should've get the answer as 137.
→ Ayça hasn't reached 137.

∴ Ayça wasn't able to solve the problem. **Modus Tollens**

Modus Tollens	
$p \rightarrow q$ $\neg q$	
<hr/>	$(p \rightarrow q) \wedge \neg q \Rightarrow \neg p$
$\therefore \neg p$	

8. Explain the following inferences.

$(\neg p \vee q) \rightarrow r$
 $r \rightarrow (s \vee t)$
 $\neg s \wedge \neg u$
 $\neg u \rightarrow \neg t$

∴ p

Steps	Rules
1) $\neg s \wedge \neg u$	Given
2) $\neg u$	(1) and simplification
3) $\neg u \rightarrow \neg t$	Given
4) $\neg t$	(2), (3) and Modus Ponens
5) $\neg s$	(1) and simplification
6) $\neg s \wedge \neg t$	(4), (5) (Rule of Conjunction)
7) $r \rightarrow (s \vee t)$	Given
8) $\neg(s \vee t) \rightarrow \neg r$	(1) and contrapositiveness
9) $(\neg s \wedge \neg t) \rightarrow \neg r$	De Morgan
10) $\neg r$	(6),(9)and Modus Ponens
11) $(\neg p \vee q) \rightarrow r$	Given
12) $\neg r \rightarrow \neg(\neg p \vee q)$	(11) and contrapositiveness
13) $\neg r \rightarrow (p \wedge \neg q)$	De Morgan
14) $p \wedge \neg q$	(10),(13) and Modus Ponens
15) ∴ p	(14) and simplification

Basitleştirme	
$p \wedge q$	
<hr/>	$p \wedge q \Rightarrow p$
$\therefore p$	
Rule of Conjunctive Simplification	

p
q
<hr/>
∴ p ∧ q
Rule of Conjunction

$p \rightarrow r \leftrightarrow \neg r \rightarrow \neg p \leftrightarrow \neg p \vee r$
Contrapositive

10. Show the correctness of the following statements

b) $[p \wedge (p \rightarrow q) \wedge (\neg q \vee r)] \rightarrow r$

$$\begin{array}{l} p \\ (p \rightarrow q) \\ (\neg q \vee r) \\ \hline \therefore r \end{array}$$

	Steps	Rules
1)	$p \rightarrow q$	Given
2)	q	Modus Ponens
3)	$\neg q \vee r$	Given
4)	r	Disjunctive Syllogism

Ayrııcı Kıyas	
$p \vee q$ $\neg p$	$(p \vee q) \wedge \neg p \Rightarrow q$
$\therefore q$	
Rule of Disjunctive Syllogism	

Varsayımlı Kıyas	
$p \rightarrow q$ $q \rightarrow r$	$(p \rightarrow q) \wedge (q \rightarrow r) \Rightarrow (p \rightarrow r)$
$\therefore p \rightarrow r$	
Law of the Syllogism	

c)

	Steps	Rules
$p \rightarrow q$	1) $p \rightarrow q$	Given
$\neg q$	2) $\neg q$	Given
$\neg r$	3) $\neg p$	Modus Tollens
$\therefore \neg(p \vee r)$	4) $\neg r$	Given
	5) $\neg p \wedge \neg r$	Rule of Conjunction
	6) $\neg(p \vee r)$	DeMorgan

d)

	Steps	Rules
$p \rightarrow q$	1) $r \rightarrow \neg q$	Given
$r \rightarrow \neg q$	2) $q \rightarrow \neg r$	Contrapositive (1)
r	3) $p \rightarrow q$	Given
$\therefore \neg p$	4) $p \rightarrow \neg r$	Law of the Syllogism
	5) r	Given
	6) $\neg p$	Modus Tollens

f)

$$\begin{array}{l}
 p \wedge q \\
 p \rightarrow (r \wedge q) \\
 r \rightarrow (s \vee t) \\
 \hline
 \neg s \\
 \hline
 \therefore t
 \end{array}$$

	Steps	Rules
1)	$p \wedge q$	Given
2)	p	Simplification
3)	$r \wedge q$	Modus Ponens
4)	r	Simplification
5)	$s \vee t$	Modus Ponens
6)	$\neg s$	Given
7)	t	Rule of Disjunctive Syllogism

g)

$$\begin{array}{l}
 p \rightarrow (q \rightarrow r) \\
 p \vee s \\
 t \rightarrow q \\
 \hline
 \neg s \\
 \hline
 \therefore \neg r \rightarrow \neg t
 \end{array}$$

	Steps	Rules
1)	$p \vee s$	Given
2)	$\neg s$	Given
3)	p	Rule of Disjunctive Syllogism
4)	$q \rightarrow r$	Modus Ponens
5)	$t \rightarrow q$	Given
6)	$t \rightarrow r$	Law of the Syllogism
7)	$\neg r \rightarrow \neg t$	Contrapositive (6)