

BLG 335E – Analysis of Algorithms I

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Research Lab 2



Exercise 1

a) How many people must there be in a room before the probability that someone has the same birthday as you do is at least $1/2$?



Solution 1.a

Probability that someone in the room has the same birthday as me, denoted by

$P(B)$ is 1– probability that no one in the room has the same birthday as me.

$$P(B) = 1 - \left(\frac{364}{365}\right)^n$$



Solution 1.a

We wish $P(B) \geq 1/2$

$$1 - \left(\frac{364}{365}\right)^n \geq 1/2$$

$$\log(1/2) \geq \log\left(\frac{364}{365}\right)^n$$

$$-\log(2) \geq n \log\left(\frac{364}{365}\right)$$

$$\log(2) \leq n \log\left(\frac{365}{364}\right)$$

$$253 \leq n$$



Exercise 1

b) How many people must there be in a room before the probability that at least two people have a birthday on July 4 is greater than $1/2$?



Solution 1.b

Probability that at least two people have a birthday on July 4, denoted by

$P(J)$ is 1 - (probability that **exactly one person** in the room has a birthday on July 4) - (probability that **no one** in the room has a birthday on July 4).

$$P(J)$$

$$= 1 - \binom{n}{1} \left(\frac{1}{365} \right) \left(\frac{364}{365} \right)^{n-1} - \binom{n}{0} \left(\frac{1}{365} \right) \left(\frac{364}{365} \right)^n$$

$$P(J) \geq \frac{1}{2}$$

$$n \geq 613$$

Exercise 2

- Illustrate the operation of **Counting-Sort** on the array below.

$$A = [6, 0, 2, 0, 1, 3, 4, 6, 1, 3, 2]$$

1. $\text{Max}\{A[i]\}=6$

2.

0	0	0	0	0	0	0
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3.

2	2	2	2	1	0	2
---	---	---	---	---	---	---

4.

2	4	6	8	9	9	11
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Exercise 2

6	0	2	0	1	3	4	6	1	3	2
---	---	---	---	---	---	---	---	---	---	---

5.

1	2	3	4	5	6	7	8	9	10	11
					2					
					2		3			
			1		2		3			
			1		2		3			6
			1		2		3	4		6
			1		2	3	3	4		6
		1	1		2	3	3	4		6
	0	1	1		2	3	3	4		6
	0	1	1	2	2	3	3	4		6
0	0	1	1	2	2	3	3	4		6
0	0	1	1	2	2	3	3	4	6	6

0	1	2	3	4	5	6
2	4	6	8	9	9	11
2	4	5	8	9	9	11
2	4	5	7	9	9	11
2	3	5	7	9	9	11
2	3	5	7	9	9	10
2	3	5	7	8	9	10
2	3	5	6	8	9	10
2	2	5	6	8	9	10
1	2	5	6	8	9	10
1	2	4	6	8	9	10
0	2	4	6	8	9	10
0	2	4	6	8	9	9

Exercise 3

- Illustrate the operation of **Radix-Sort** on the following list of English words:

COW, DOG, SEA, RUG, ROW, MOB,
BOX, TAB, BAR, EAR, TAR, DIG, BIG,
TEA, NOW, FOX.



Exercise 3 – Solution

COW

DOG

SEA

RUG

ROW

MOB

BOX

TAB

BAR

EAR

TAR

DIG

BIG

TEA

NOW

FOX



Exercise 3 – Solution

SEA

BAR

TEA

EAR

MOB

TAR

TAB

FOX

DOG

BOX

RUG

COW

DIG

ROW

BIG

NOW



Exercise 3 – Solution

TAB

BAR

EAR

TAR

SEA

TEA

DIG

BIG

MOB

DOG

FOX

BOX

COW

ROW

NOW

RUG



Exercise 3 – Solution

BAR

BIG

BOX

COW

DIG

DOG

EAR

FOX

MOB

NOW

ROW

RUG

SEA

TAB

TAR

TEA



Exercise 4

Given input array $A[0..9]$. The array $B[0..9]$ of sorted lists or buckets.

Bucket i holds values in the interval $[i/10, (i+1)/10]$.

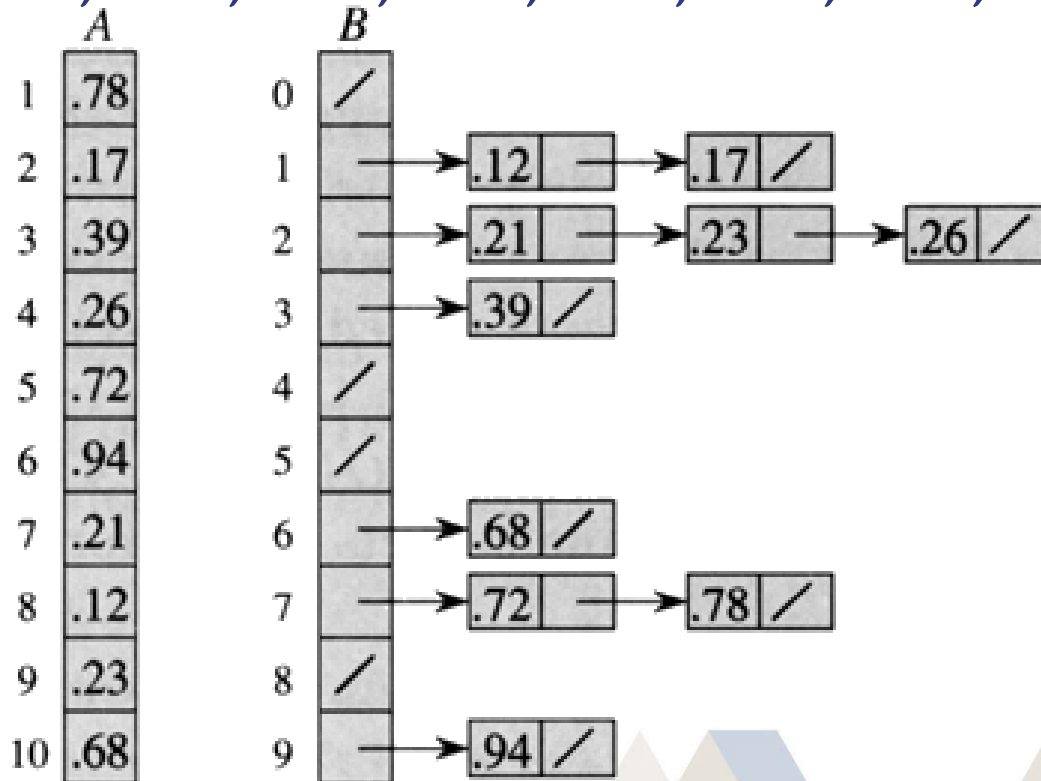
The sorted output consists of a concatenation in order of the lists first $B[0]$ then $B[1]$ then $B[2]$... and the last one is $B[9]$.



Exercise 4

- Illustrate the operation of **Bucket-Sort** on the array below.

$A = [.78, .17, .39, .26, .72, .94, .21, .12, .23, .68]$



Exercise 5

Give tight asymptotic bounds for $T(n)$ in each of the following recurrences.

$$a. \quad T(n) = 4T\left(\frac{n}{2}\right) + n^2\sqrt{n}$$
$$a = 4, \quad b = 2, \quad f(n) = n^2\sqrt{n} = n^{5/2}$$

$$n^{\log_b^a} = n^{\log_2^4} = n^2$$

$$n^{5/2} = \Omega(n^{2+1/2})$$

possibly case 3, let's check c

$$T(n) = aT(n/b) + f(n)$$

$$1 \quad f(n) = O\left(n^{\log_b a - \varepsilon}\right) \Rightarrow T(n) = \Theta\left(n^{\log_b a}\right)$$

$$2 \quad f(n) = \Theta\left(n^{\log_b a}\right) \Rightarrow T(n) = \Theta\left(n^{\log_b a} \log_2 n\right)$$

$$3 \quad f(n) = \Omega\left(n^{\log_b a + \varepsilon}\right) \text{ and } af(n/b) \leq cf(n),$$

for $\exists c \quad c < 1$ and $n > n_0$

$$\Rightarrow T(n) = \Theta(f(n))$$

Exercise 5.a

$$a \ f \left(\frac{n}{b} \right) = 4 \left(\frac{n}{2} \right)^2 \sqrt{\frac{n}{2}} = \frac{n^{5/2}}{\sqrt{2}} \leq cn^{\frac{5}{2}}$$
$$\text{for } \frac{1}{\sqrt{2}} \leq c < 1$$

Case 3 applies $T(n) = \Theta(n^2 \sqrt{n})$



Exercise 5

$$b. T(n) = 3T\left(\frac{n}{2}\right) + n \lg n$$

$$a = 2 \quad b = 3, \quad f(n) = n \lg n$$
$$n^{\log_b^a} = n^{\lg 3} \approx n^{1.585}$$

$$n \lg n = O(n^{\lg 3 - \varepsilon})$$

$$0 < \varepsilon \leq 0.58 \quad \text{case 1}$$

$$T(n) = \Theta(n^{\lg 3})$$