Database Systems Database Design

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Topics

Normalization

Introduction Normal Forms 3rd Normal Form Views

Entity/Relationship Model

Introduction E/R Diagrams

Functional Dependency

- ightharpoonup let Z be the set of all attributes of the relation R
- ▶ let $A, B \subseteq Z$
- ▶ A functionally determines $B: A \rightarrow B$ for every A value there can only be one B value
- every functional dependency is an integrity constraint

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Example Relation

R

MOVIEID	TITLE	COU	LANG	<u>ACTORID</u>	NAME	ORD
6	Usual Suspects	UK	EN	308	Gabriel Byrne	2
228	Ed Wood	US	EN	26	Johnny Depp	1
70	Being John Malkovich	US	EN	282	Cameron Diaz	2
1512	Suspiria	IT	IT	745	Udo Kier	9
70	Being John Malkovich	US	EN	503	John Malkovich	14

► assumption: the language of the movie is the language of the country where it was made

Functional Dependency Examples

- ► MOVIEID → TITLE
- ► MOVIEID → {TITLE, COUNTRY, LANGUAGE}
- ► ACTORID → NAME
- ▶ {MOVIEID, ACTORID} → ORD
- ► trivial: MOVIEID → MOVIEID
- ▶ redundant: {MOVIEID, ACTORID} → COUNTRY

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Irreducible Set

- ▶ let *S* be the set of all FDs of the relation
- ▶ let *T* ⊆ *S*
- T is an irreducible set of FDs if:
- T contains as few elements as possible, and
- ightharpoonup every FD in S can be derived from the FDs in T
- ▶ let there be only one attribute on the right hand side of FDs

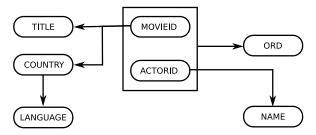
Irreducible Set Example

- ► MOVIEID → TITLE
- ► MOVIEID → COUNTRY
- ► COUNTRY → LANGUAGE
- ► ACTORID → NAME
- ► {MOVIEID, ACTORID} → ORD

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Dependency Example



Normal Forms

- ▶ normal form: set of conditions that relations must satisfy
- ▶ 1NF, 2NF, 3NF, BCNF, 4NF, 5NF
- every form adds stricter conditions to the previous form
 - every relation in 2NF is also in 1NF
 - every relation in 3NF is also in 2NF, ...
- ▶ NF of database = NF of weakest relation
- ► 1NF: attribute values have to be atomic

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Normalization

- normalization: transforming relations from one form to the next, stricter form
- ▶ transformation must be lossless

Theorem (Heath)

- ▶ let Z be the set of all attributes of the relation R
- ▶ *let* A, B, $C \subseteq Z$
- $\blacktriangleright A \rightarrow B \Rightarrow R = (A, B) \ \textit{join} \ (A, C)$

Lossless Transition Example

R1

MOVIEID	TITLE	COU	LANG
6	Usual Suspects	UK	EN
228	Ed Wood	US	EN
70	Being John Malkovich	US	EN
1512	Suspiria	IT	IT

R2

MOVIEID	ACTORID	NAME	ORD
6	308	Gabriel Byrne	2
228	26	Johnny Depp	1
70	282	Cameron Diaz	2
1512	745	Udo Kier	9
70	503	John Malkovich	14

ightharpoonup R = R1 join R2

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Lossy Transition Example

R1

MOVIEID	TITLE	COU	LANG
6	Usual Suspects	UK	EN
228	Ed Wood	US	EN
70	Being John Malkovich	US	EN
1512	Suspiria	IT	IT

R2

COU	ACTORID	NAME	ORD
UK	308	Gabriel Byrne	2
US	26	Johnny Depp	1
US	282	Cameron Diaz	2
IT	745	Udo Kier	9
US	503	John Malkovich	14

- ightharpoonup R
 eq R1 join R2
- ► {MOVIEID, ACTORID} → ORD

Anomalies

- ▶ insert: data is known but can not be inserted
- ▶ delete: deleting some data causes some other data to be lost
- ▶ update: updating data requires modifications in multiple tuples

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Anomaly Examples

• example relation

- ▶ the country of the movie "Gattaca" is known to be US, but this cannot be inserted if there is no actor in the movie
- ▶ deleting that Gabriel Byrne acts in the movie "Usual Suspects" also deletes that the movie was made in the UK
- changing the country of the movie "Being John Malkovich" requires modifications in two tuples

2nd Normal Form

- ▶ 2NF: every non-key attribute depends on the primary key
- ▶ in a relation *R* that conforms to 1NF, if:
 - ightharpoonup R(A, B, C, D), primary key: (A, B)
 - $A \rightarrow D$
- ▶ to transform to 2NF:
 - R1(A, D), primary key: A
 - ► R2(A, B, C), primary key: (A, B) A is a foreign key referencing R1

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1NF-2NF Transition Example

- ► among the non-key attributes, only ORD depends on the primary key
 - ► A: {MOVIEID}
 - ► B: {ACTORID}
 - ► C: {NAME, ORD}
 - ▶ D: {TITLE, COUNTRY, LANGUAGE}
- ► R1(MOVIEID, TITLE, COUNTRY, LANGUAGE) primary key: MOVIEID
- ► R2(MOVIEID, ACTORID, NAME, ORD) primary key: {MOVIEID, ACTORID} MOVIEID is a foreign key referencing R1

1NF-2NF Transition Example

- ▶ R2 still not in 2NF: ACTORID → NAME
 - ► A: {ACTORID}
 - ► B: {MOVIEID}
 - ► *C*: {ORD}
 - ► *D*: {NAME}
- ► R3(ACTORID, NAME) primary key: ACTORID
- ► R4(MOVIEID, ACTORID, ORD)
 primary key: {MOVIEID, ACTORID}
 ACTORID is a foreign key referencing R3

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2NF Relation Examples

R1

MOVIEID	TITLE	COU	LANG
6	Usual Suspects	UK	EN
228	Ed Wood	US	EN
70	Being John Malkovich	US	EN
1512	Suspiria	ΙΤ	IT

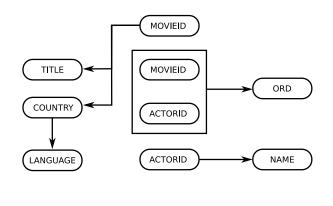
R3

ACTORID	NAME
308	Gabriel Byrne
26	Johnny Depp
282	Cameron Diaz
745	Udo Kier
503	John Malkovich

R4

MOVIEID	<u>ACTORID</u>	ORD
6	308	2
228	26	1
70	282	2
1512	745	9
70	503	14

Dependency Example



2NF Corrected Anomalies

- ▶ if the country of the movie "Gattaca" is known to be US, this can be inserted to R1
- ▶ if Gabriel Byrne is deleted from the movie "Usual Suspects", the country of the movie is still kept in R1
- ► changing the country of the movie "Being John Malkovich" requires updating only one tuple in R1

2NF Remaining Anomalies

- ▶ it is known that movies made in Brazil are in Portuguese but this can not be inserted if there is no movie made in Brazil
- ▶ deleting the movie "Suspiria" also deletes that the language of movies made in Italy is Italian
- ► changing the language of the movies made in the US requires two tuples to be updated

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3rd Normal Form

- ▶ 3NF: non-key attributes do not depend on any attributes other than the primary key
- ightharpoonup in a relation R that conforms to 2NF, if:
 - ightharpoonup R(A, B, C, D), primary key: A
 - ightharpoonup C
 ightharpoonup D
- ▶ to transform to 3NF:
 - ightharpoonup R1(C,D), primary key: C
 - ► R2(A, B, C), primary key: A C is a foreign key referencing R1

2NF-3NF Transition Example

- ▶ R1: COUNTRY → LANGUAGE
 - ► A: {MOVIEID}
 - ► *B*: {TITLE}
 - ► C: {COUNTRY}
 - ► D: {LANGUAGE}
- ► R5(COUNTRY, LANGUAGE) primary key: COUNTRY
- ► R6(MOVIEID, TITLE, COUNTRY)

 primary key: MOVIEID

 COUNTRY is a foreign key referencing R5

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3NF Relation Examples

R6

MOVIEID	TITLE	COU
6	Usual Suspects	UK
228	Ed Wood	US
70	Being John Malkovich	US
1512	Suspiria	ΙΤ

R5

11.5			
<u>COU</u>	LANG		
UK	EN		
US	EN		
IT	ΙΤ		

R3

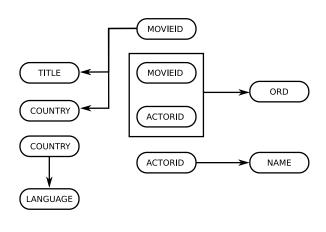
ACTORID	NAME
308	Gabriel Byrne
26	Johnny Depp
282	Cameron Diaz
745	Udo Kier
503	John Malkovich

R4

MOVIEID	ACTORID	ORD
6	308	2
228	26	1
70	282	2
1512	745	9
70	503	14

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Dependency Example



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3NF Corrected Anomalies

- ▶ if movies made in Brazil are in known to be in Portuguese, this can be inserted into R5
- ▶ if the movie "Suspiria" is deleted, R5 still keeps that movies made in Italy are in Italian
- ► changing the language of the movies made in the US requires modifying only one tuple in R5

Boyce-Codd Normal Form

- ▶ BCNF: all functional dependencies must be on candidate keys
- ► consider dependencies between key attributes

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BCNF Example

- ▶ let movie title be unique
- candidate keys:
 - ► {MOVIEID, ACTORID}
 - ► {TITLE, ACTORID}
- ▶ non-conforming functional dependencies:
 - ► MOVIEID → TITLE
 - ightharpoonup TITLE ightharpoonup MOVIEID

Views

- presenting a derived table like a base table: view
- ► isolating users and application programs from changes in database structure
- creating a view:

```
CREATE VIEW view_name AS
   SELECT ...
```

▶ SELECT will be executed every time the view is used

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View Example

▶ identifiers, titles and years of new movies

```
CREATE VIEW NEW_MOVIE AS
SELECT ID, TITLE, YR FROM MOVIE
WHERE (YR > 1995)
```

```
SELECT * FROM NEW_MOVIE
```

Updating Views

- ▶ any change will have to performed on the base tables
- rules need to be defined
- creating a rule:

```
CREATE RULE rule_name AS
  ON event TO view_name
  [ WHERE condition ]
  DO [ INSTEAD ] sql_statement
```

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View Rule Example

▶ modify the title of a new movie

```
UPDATE NEW_MOVIE SET TITLE = ...
WHERE (ID = ...)
```

▶ rule for updating the base table

```
CREATE RULE UPDATE_TITLE AS

ON UPDATE TO NEW_MOVIE

DO INSTEAD

UPDATE MOVIE SET TITLE = new.TITLE

WHERE (ID = old.ID)
```

References

Required Reading: Date

- ► Chapter 11: Functional Dependencies
- ► Chapter 12: Further Normalization I: 1NF, 2NF, 3NF, BCNF
- ► Chapter 10: Views

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Entity/Relationship Model

- ▶ modeling approach (Chen 1976)
- components:
 - entities
 - properties
 - relationships

Entities

- entity: set of "things" with the same attributes
- elements of the set are *instances* of the entity
- ► strong: can exist by itself
- ▶ weak: existence depends on another entity

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Entity Examples

▶ entity: movie, person

▶ person instance: Johnny Depp

▶ strong entity: person

▶ weak entity: movie

Properties

- property: data describing entities
- ► simple / composite
- key
- ► single / multiple valued
- empty
- ▶ base / derived

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Property Examples

▶ property: title, country, language

▶ simple: first name, last name

► composite: full name

▶ base: date of birth

▶ derived: age

Relationships

- ► relationship: connection between entities
- participant: entities in the relationship
- ► *degree*: number of participants
- ► total: all instances of the entity participate in the relationship (otherwise partial)

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Relationship Types

- one-to-one
 - e.g. capital city relationship between countries and cities
- one-to-many
 - e.g. management relationship between employees and projects
- many-to-many
 - e.g. enrollment relationship between students and courses

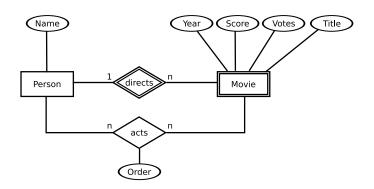
Entity/Relationship Diagrams

- entity: box
 - weak: double lines
- property: circle
 - derived: dashed lines
 - multi-valued: double lines
 - composite: sub-circles
- relationship: diamond
 - between weak and strong: double lines
 - total: connection double lines
 - ▶ 1 or n depending on the type of the relationship

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Entity/Relationship Diagram Example

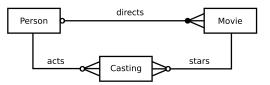


Entity/Relationship Diagrams

- properties excluded if diagram too large
- ▶ alternative representation for relationship: line
- ▶ name of relationship above the line
- ▶ fork at the end of line if "many"
- ▶ circle at the end of line: empty if partial, filled if total
- ▶ no many-to-many relationships
- ▶ add an entity and use two one-to-many relationships

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Entity/Relationship Diagram Example



Applying to Design

- every entity a relation
- ▶ every property an attribute
- ▶ every many-to-many relationship a relation
 - ► foreign keys to participating entities
- ► for every one-to-many relationship a foreign key from the "many" side to the "one" side

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References

Required Reading: Date

► Chapter 14: Semantic Modeling

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