BLG 335E – Analysis of Algorithms I
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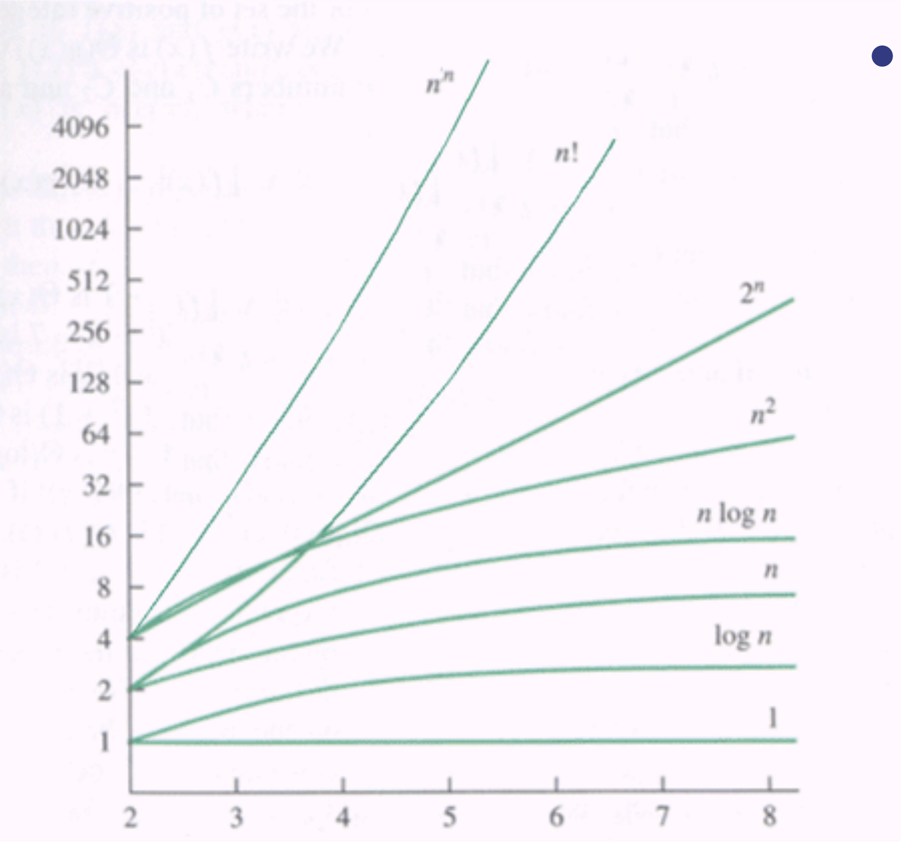
Warm-up Problem

- Order the following functions by asymptotic growth rate:
  - $n^2 + 5n + 7$
  - $\log_2 n^3$
  - 95^{17}
  - $2^{\log_2 n}$
  - n^3
  - $nlog_2n + 9n$
  - $4 \log_2 n$
  - $\log_2 n + 3n$

Warm-up Problem



- Solution:
  - 95^{17}
  - $\log_2 n^3$
  - $4 \log_2 n$
  - $2^{\log_2 n}$
  - $\log_2 n + 3n$
  - $nlog_2 n + 9n$
  - $n^2 + 5n + 7$
  - n^3

Problems

- Give tight asymptotic bounds for $T(n)$ in each of the following recurrences.

  a. $T(n) = T(n-1) + n$

  b. $T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{4}\right) + T\left(\frac{n}{8}\right) + n$

  c. $T(n) = T\left(\frac{9n}{10}\right) + n$

  d. $T(n) = 16T\left(\frac{n}{4}\right) + n^2$

  e. $T(n) = 7T\left(\frac{n}{2}\right) + n^2$

Problems

- Give tight asymptotic bounds for $T(n)$ in each of the following recurrences.

  a. $T(n) = T(n-1) + n$

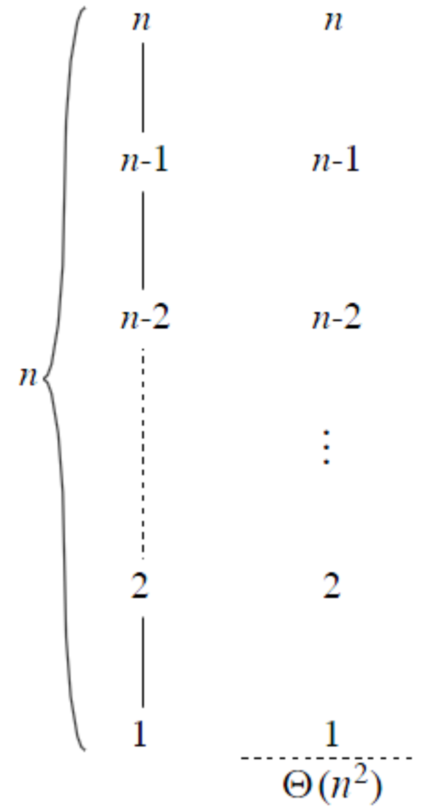
  Lower bound (Ω):

  $T(n) \geq cn^2 \; for\ some\ c > 0$

  $T(n) \geq c(n-1)^2 + n$

  $= cn^2 - 2cn + c + n \geq cn^2$

  $true\ if\ 0 < c < \dfrac{1}{2} and\ n \geq 0$



Problems

- Give tight asymptotic bounds for $T(n)$ in each of the following recurrences.

  a. $T(n) = T(n-1) + n$

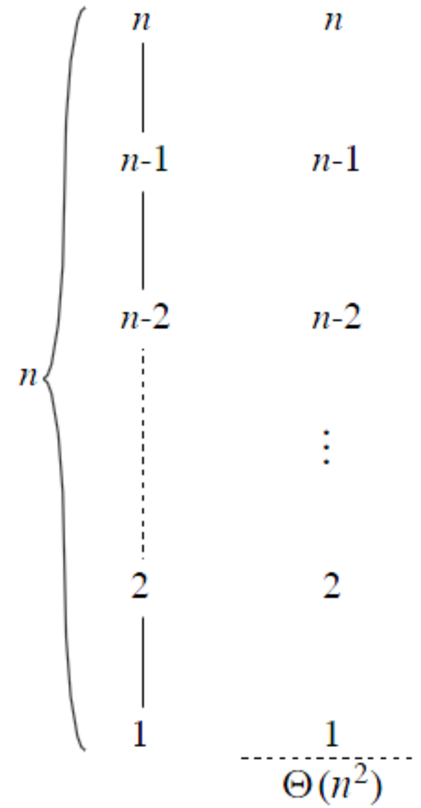
  Upper bound (*O*):

  $T(n) \leq cn^2 \; for \; some \; c > 0$

  $T(n) \leq c(n-1)^2 + n$

  $= cn^2 - 2cn + c + n \leq cn^2$

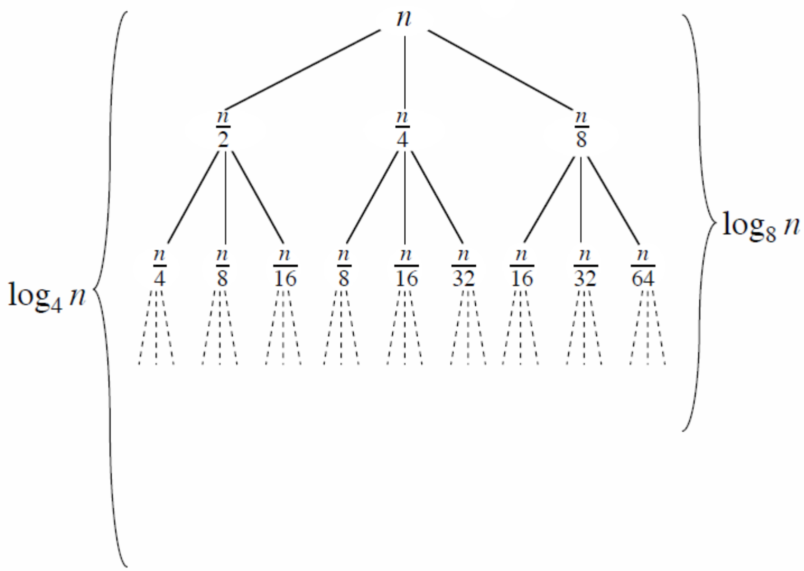
  $true \; if \; c = 1 \; and \; n \geq 1$



Problems

- Give tight asymptotic bounds for $T(n)$ in each of the following recurrences.

  b. $T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{4}\right) + T\left(\frac{n}{8}\right) + n$



$$n$$

$$n\left(\frac{4+2+1}{8}\right) = \frac{7}{8}n$$

$$n\left(\frac{1}{4} + \frac{2}{8} + \frac{3}{16} + \frac{2}{32} + \frac{1}{64}\right) = n\frac{16+16+12+4+1}{64} = n\frac{49}{64} = \frac{7}{8}^2 n$$

$$\vdots$$

$$\sum_{i=1}^{\log n} \left(\frac{7}{8}\right)^i n = \Theta(n)$$

Problems

- Give tight asymptotic bounds for $T(n)$ in each of the following recurrences.

$b.\ \ T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{4}\right) + T\left(\frac{n}{8}\right) + n$

Upper bound (*O*):

$$T(n) \leq \frac{cn}{2} + \frac{cn}{4} + \frac{cn}{8} + n = \frac{7cn}{8} + n \leq cn$$

$$true\ if\ c \geq 8$$

Problems

- Give tight asymptotic bounds for $T(n)$ in each of the following recurrences.

  b. $T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{4}\right) + T\left(\frac{n}{8}\right) + n$

  Lower bound (Ω):

  $$T(n) \geq \frac{cn}{2} + \frac{cn}{4} + \frac{cn}{8} + n = \frac{7cn}{8} + n \geq cn$$

  $$true\ if\ 0 < c \leq 8$$

Master Method

$$T(n) = aT(n/b) + f(n)$$

1 $f(n) = O\left(n^{\log_b a-\varepsilon}\right) \Rightarrow T(n) = \Theta\left(n^{\log_b a}\right)$

2 $f(n) = \Theta\left(n^{\log_b a}\right) \Rightarrow T(n) = \Theta\left(n^{\log_b a} \log_2 n\right)$

3 $f(n) = \Omega\left(n^{\log_b a+\varepsilon}\right)$ and $af(n/b) \le cf(n),$

$for \;\; \exists c \;\; c<1 \;\; and \;\; n>n_0$

$\Rightarrow T(n) = \Theta\left(f(n)\right)$

Problems

- Give tight asymptotic bounds for $T(n)$ in each of the following recurrences.

c. $T(n) = T\left(\frac{9n}{10}\right) + n$

$a = 1, b = \frac{10}{9}, f(n) = n = \Omega\left(\mathrm{n}^{\log_{\frac{10}{9}} 1 + 1}\right)$

$possibly\ case\ 3, let's\ check\ c$

$1\frac{9n}{10} \leq cn\ holds\ for\ c = \frac{9}{10} \leq 1$

$certainly\ case\ 3:$

$T(n) = \Theta(\mathrm{n})$

Problems

- Give tight asymptotic bounds for $T(n)$ in each of the following recurrences.

d. $T(n) = 16T\left(\frac{n}{4}\right) + n^2 \quad a = 16, b = 4, f(n) = n^2$

$$n^2 = \Theta\left(n^{\log_4 16}\right), case\ 2:$$

$$T(n) = \Theta(n^2 \log_2 n)$$

e. $T(n) = 7T\left(\frac{n}{2}\right) + n^2 \qquad a = 7, b = 2, f(n) = n^2$

$$n^2 = O\left(n^{\log_2 7-\varepsilon}\right), case\ 1:$$

$$T(n) = \Theta(n^{\log_2 7})$$