

Question1)

$G = (V, E, s, t, c)$

Source = s

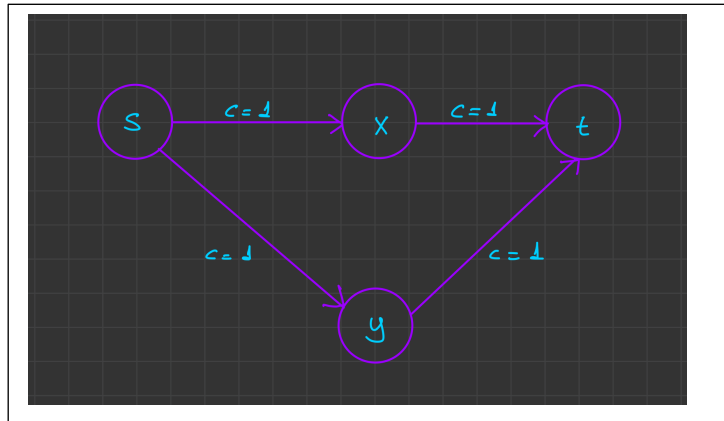
Sink = t

Capacity function = c

Vertices $V = \{s, t, x, y\}$

Edges E and capacities c :

- $(s, x), c=1$
- $(x, t), c=1$
- $(s, y), c=1$
- $(y, t), c=1$



Consider the flow path that gives max flow 1: $s \rightarrow x \rightarrow t$

Flow function f :

- $f(s, x)=1$
- $f(x, t)=1$
- $f(s, y)=0$
- $f(y, t)=0$

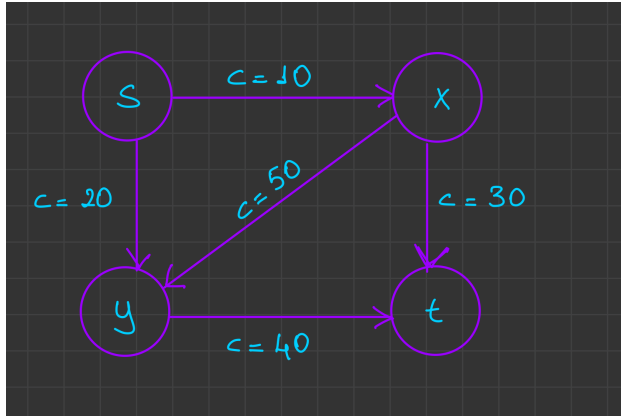
Consider the flow path that gives max flow 1: $s \rightarrow y \rightarrow t$

Flow function f :

- $f(s, x)=0$
- $f(x, t)=0$
- $f(s, y)=1$
- $f(y, t)=1$

Question2)

Theorem A is **incorrect**, below is a **counter-example** in which edges have unique capacity :



Flow function f1:

- $f1(s,x) = 10$
- **$f1(x,t) = 0$**
- $f1(s,y) = 20$
- **$f1(y,t) = 30$**
- $f1(x,y) = 10$

→ Maximum flow 30

Flow function f2:

- $f2(s,x) = 10$
- **$f2(x,t) = 10$**
- $f2(s,y) = 20$
- **$f2(y,t) = 20$**
- $f2(x,y) = 0$

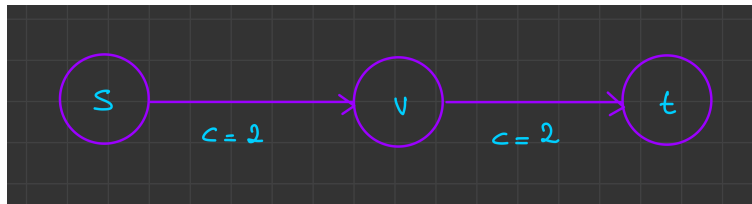
→ Maximum flow 30

Question3)

F is **not guaranteed** to be a flow on G.

The capacity constraint, which is a must for a flow in a network, may not hold for F.
Meaning that the combined flow F may exceed the capacity of an edge in the network.

Counter example:



1st node s → source
2nd node v
3rd node t → sink
capacity c

Flow function 1:

- $f_1(s,v) = 2$
- $f_1(v,t) = 2$

Flow function 2:

- $f_2(s,v) = 1$
- $f_2(v,t) = 1$

Combination of the above flow functions f1 and f2:

- $F(s,v) = f_1(s,v) + f_2(s,v) = 3$
- $F(v,t) = f_1(v,t) + f_2(v,t) = 3$

F exceeds the capacity of edges as it has flow of 3 on both edges.

The capacity constrained is violated by F, this viaolation shows that F is not a valid flow on G.